

Matrix

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June 21, 2010

1 Various algebraic structures combined with a lattice

```
theory Lattice-Algebras  
imports Complex-Main  
begin
```

```
class semilattice-inf-ab-group-add = ordered-ab-group-add + semilattice-inf  
begin
```

```
lemma add-inf-distrib-left:  
   $a + \inf b\ c = \inf (a + b)\ (a + c)$   
<proof>
```

```
lemma add-inf-distrib-right:  
   $\inf a\ b + c = \inf (a + c)\ (b + c)$   
<proof>
```

```
end
```

```
class semilattice-sup-ab-group-add = ordered-ab-group-add + semilattice-sup  
begin
```

```
lemma add-sup-distrib-left:  
   $a + \sup b\ c = \sup (a + b)\ (a + c)$   
<proof>
```

```
lemma add-sup-distrib-right:  
   $\sup a\ b + c = \sup (a+c)\ (b+c)$   
<proof>
```

```
end
```

```
class lattice-ab-group-add = ordered-ab-group-add + lattice  
begin
```

```
subclass semilattice-inf-ab-group-add <proof>
```

subclass *semilattice-sup-ab-group-add* $\langle \text{proof} \rangle$

lemmas *add-sup-inf-distrib* = *add-inf-distrib-right add-inf-distrib-left add-sup-distrib-right add-sup-distrib-left*

lemma *inf-eq-neg-sup*: $\text{inf } a \ b = - \ \text{sup } (-a) \ (-b)$
 $\langle \text{proof} \rangle$

lemma *sup-eq-neg-inf*: $\text{sup } a \ b = - \ \text{inf } (-a) \ (-b)$
 $\langle \text{proof} \rangle$

lemma *neg-inf-eq-sup*: $- \ \text{inf } a \ b = \text{sup } (-a) \ (-b)$
 $\langle \text{proof} \rangle$

lemma *neg-sup-eq-inf*: $- \ \text{sup } a \ b = \text{inf } (-a) \ (-b)$
 $\langle \text{proof} \rangle$

lemma *add-eq-inf-sup*: $a + b = \text{sup } a \ b + \text{inf } a \ b$
 $\langle \text{proof} \rangle$

1.1 Positive Part, Negative Part, Absolute Value

definition

nprt :: 'a \Rightarrow 'a **where**
nprt *x* = $\text{inf } x \ 0$

definition

pprt :: 'a \Rightarrow 'a **where**
pprt *x* = $\text{sup } x \ 0$

lemma *pprt-neg*: $\text{pprt } (-x) = - \ \text{nprt } x$
 $\langle \text{proof} \rangle$

lemma *nprt-neg*: $\text{nprt } (-x) = - \ \text{pprt } x$
 $\langle \text{proof} \rangle$

lemma *prts*: $a = \text{pprt } a + \text{nprt } a$
 $\langle \text{proof} \rangle$

lemma *zero-le-pprt[simp]*: $0 \leq \text{pprt } a$
 $\langle \text{proof} \rangle$

lemma *nprt-le-zero[simp]*: $\text{nprt } a \leq 0$
 $\langle \text{proof} \rangle$

lemma *le-eq-neg*: $a \leq -b \longleftrightarrow a + b \leq 0$ (**is** ?l = ?r)
 $\langle \text{proof} \rangle$

lemma *pprt-0[simp]*: $\text{pprt } 0 = 0$ $\langle \text{proof} \rangle$

lemma *nprrt-0*[*simp*]: $\text{nprrt } 0 = 0 \langle \text{proof} \rangle$

lemma *pprrt-eq-id* [*simp*, *no-atp*]: $0 \leq x \implies \text{pprrt } x = x \langle \text{proof} \rangle$

lemma *nprrt-eq-id* [*simp*, *no-atp*]: $x \leq 0 \implies \text{nprrt } x = x \langle \text{proof} \rangle$

lemma *pprrt-eq-0* [*simp*, *no-atp*]: $x \leq 0 \implies \text{pprrt } x = 0 \langle \text{proof} \rangle$

lemma *nprrt-eq-0* [*simp*, *no-atp*]: $0 \leq x \implies \text{nprrt } x = 0 \langle \text{proof} \rangle$

lemma *sup-0-imp-0*: $\text{sup } a \text{ } (- a) = 0 \implies a = 0 \langle \text{proof} \rangle$

lemma *inf-0-imp-0*: $\text{inf } a \text{ } (-a) = 0 \implies a = 0 \langle \text{proof} \rangle$

lemma *inf-0-eq-0* [*simp*, *no-atp*]: $\text{inf } a \text{ } (- a) = 0 \longleftrightarrow a = 0 \langle \text{proof} \rangle$

lemma *sup-0-eq-0* [*simp*, *no-atp*]: $\text{sup } a \text{ } (- a) = 0 \longleftrightarrow a = 0 \langle \text{proof} \rangle$

lemma *zero-le-double-add-iff-zero-le-single-add* [*simp*]:
 $0 \leq a + a \longleftrightarrow 0 \leq a \langle \text{proof} \rangle$

lemma *double-zero* [*simp*]:
 $a + a = 0 \longleftrightarrow a = 0 \langle \text{proof} \rangle$

lemma *zero-less-double-add-iff-zero-less-single-add* [*simp*]:
 $0 < a + a \longleftrightarrow 0 < a \langle \text{proof} \rangle$

lemma *double-add-le-zero-iff-single-add-le-zero* [*simp*]:
 $a + a \leq 0 \longleftrightarrow a \leq 0 \langle \text{proof} \rangle$

lemma *double-add-less-zero-iff-single-less-zero* [*simp*]:
 $a + a < 0 \longleftrightarrow a < 0 \langle \text{proof} \rangle$

declare *neg-inf-eq-sup* [*simp*] *neg-sup-eq-inf* [*simp*]

lemma *le-minus-self-iff*: $a \leq - a \longleftrightarrow a \leq 0$

<proof>

lemma *minus-le-self-iff*: $- a \leq a \longleftrightarrow 0 \leq a$
<proof>

lemma *zero-le-iff-zero-nprt*: $0 \leq a \longleftrightarrow \text{nprt } a = 0$
<proof>

lemma *le-zero-iff-zero-pprt*: $a \leq 0 \longleftrightarrow \text{pprt } a = 0$
<proof>

lemma *le-zero-iff-pprt-id*: $0 \leq a \longleftrightarrow \text{pprt } a = a$
<proof>

lemma *zero-le-iff-nprt-id*: $a \leq 0 \longleftrightarrow \text{nprt } a = a$
<proof>

lemma *pprt-mono* [*simp*, *no-atp*]: $a \leq b \implies \text{pprt } a \leq \text{pprt } b$
<proof>

lemma *nprt-mono* [*simp*, *no-atp*]: $a \leq b \implies \text{nprt } a \leq \text{nprt } b$
<proof>

end

lemmas *add-sup-inf-distribs* = *add-inf-distrib-right add-inf-distrib-left add-sup-distrib-right add-sup-distrib-left*

class *lattice-ab-group-add-abs* = *lattice-ab-group-add + abs +*
 assumes *abs-lattice*: $|a| = \text{sup } a \ (- a)$
begin

lemma *abs-prts*: $|a| = \text{pprt } a - \text{nprt } a$
<proof>

subclass *ordered-ab-group-add-abs*
<proof>

end

lemma *sup-eq-if*:
 fixes $a :: 'a :: \{\text{lattice-ab-group-add}, \text{linorder}\}$
 shows $\text{sup } a \ (- a) = (\text{if } a < 0 \text{ then } - a \text{ else } a)$
<proof>

lemma *abs-if-lattice*:
 fixes $a :: 'a :: \{\text{lattice-ab-group-add-abs}, \text{linorder}\}$
 shows $|a| = (\text{if } a < 0 \text{ then } - a \text{ else } a)$

$\langle proof \rangle$

lemma *estimate-by-abs*:

$a + b \leq (c :: 'a :: lattice-ab-group-add-abs) \implies a \leq c + abs\ b$
 $\langle proof \rangle$

class *lattice-ring* = *ordered-ring* + *lattice-ab-group-add-abs*
begin

subclass *semilattice-inf-ab-group-add* $\langle proof \rangle$

subclass *semilattice-sup-ab-group-add* $\langle proof \rangle$

end

lemma *abs-le-mult*: $abs\ (a * b) \leq (abs\ a) * (abs\ (b :: 'a :: lattice-ring))$
 $\langle proof \rangle$

instance *lattice-ring* \subseteq *ordered-ring-abs*
 $\langle proof \rangle$

lemma *mult-le-prts*:

assumes

$a1 \leq (a :: 'a :: lattice-ring)$

$a \leq a2$

$b1 \leq b$

$b \leq b2$

shows

$a * b \leq ppert\ a2 * ppert\ b2 + ppert\ a1 * npert\ b2 + npert\ a2 * ppert\ b1 + npert\ a1$
 $* npert\ b1$
 $\langle proof \rangle$

lemma *mult-ge-prts*:

assumes

$a1 \leq (a :: 'a :: lattice-ring)$

$a \leq a2$

$b1 \leq b$

$b \leq b2$

shows

$a * b \geq npert\ a1 * ppert\ b2 + npert\ a2 * npert\ b2 + ppert\ a1 * ppert\ b1 + ppert\ a2$
 $* npert\ b1$
 $\langle proof \rangle$

instance *int* :: *lattice-ring*
 $\langle proof \rangle$

instance *real* :: *lattice-ring*
 $\langle proof \rangle$

end

```

theory Matrix
imports Main Lattice-Algebras
begin

types 'a infmatrix = nat  $\Rightarrow$  nat  $\Rightarrow$  'a

definition nonzero-positions :: (nat  $\Rightarrow$  nat  $\Rightarrow$  'a::zero)  $\Rightarrow$  nat  $\times$  nat  $\Rightarrow$  bool
where
  nonzero-positions A = {pos. A (fst pos) (snd pos)  $\sim$  0}

typedef 'a matrix = {(f::(nat  $\Rightarrow$  nat  $\Rightarrow$  'a::zero)). finite (nonzero-positions f)}
  <proof>

declare Rep-matrix-inverse[simp]

lemma finite-nonzero-positions : finite (nonzero-positions (Rep-matrix A))
  <proof>

definition nrows :: ('a::zero) matrix  $\Rightarrow$  nat where
  nrows A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max
    ((image fst) (nonzero-positions (Rep-matrix A))))

definition ncols :: ('a::zero) matrix  $\Rightarrow$  nat where
  ncols A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max ((image
    snd) (nonzero-positions (Rep-matrix A))))

lemma nrows:
  assumes hyp: nrows A  $\leq$  m
  shows (Rep-matrix A m n) = 0
  <proof>

definition transpose-infmatrix :: 'a infmatrix  $\Rightarrow$  'a infmatrix where
  transpose-infmatrix A j i == A i j

definition transpose-matrix :: ('a::zero) matrix  $\Rightarrow$  'a matrix where
  transpose-matrix == Abs-matrix o transpose-infmatrix o Rep-matrix

declare transpose-infmatrix-def[simp]

lemma transpose-infmatrix-twice[simp]: transpose-infmatrix (transpose-infmatrix
  A) = A
  <proof>

lemma transpose-infmatrix: transpose-infmatrix (% j i. P j i) = (% j i. P i j)
  <proof>

```

lemma *transpose-infmatrix-closed[simp]*: $\text{Rep-matrix } (\text{Abs-matrix } (\text{transpose-infmatrix } (\text{Rep-matrix } x))) = \text{transpose-infmatrix } (\text{Rep-matrix } x)$
 $\langle \text{proof} \rangle$

lemma *infmatrixforward*: $(x::'a \text{ infmatrix}) = y \implies \forall a b. x a b = y a b$ $\langle \text{proof} \rangle$

lemma *transpose-infmatrix-inject*: $(\text{transpose-infmatrix } A = \text{transpose-infmatrix } B) = (A = B)$
 $\langle \text{proof} \rangle$

lemma *transpose-matrix-inject*: $(\text{transpose-matrix } A = \text{transpose-matrix } B) = (A = B)$
 $\langle \text{proof} \rangle$

lemma *transpose-matrix[simp]*: $\text{Rep-matrix}(\text{transpose-matrix } A) j i = \text{Rep-matrix } A i j$
 $\langle \text{proof} \rangle$

lemma *transpose-transpose-id[simp]*: $\text{transpose-matrix } (\text{transpose-matrix } A) = A$
 $\langle \text{proof} \rangle$

lemma *nrows-transpose[simp]*: $\text{nrows } (\text{transpose-matrix } A) = \text{ncols } A$
 $\langle \text{proof} \rangle$

lemma *ncols-transpose[simp]*: $\text{ncols } (\text{transpose-matrix } A) = \text{nrows } A$
 $\langle \text{proof} \rangle$

lemma *ncols*: $\text{ncols } A \leq n \implies \text{Rep-matrix } A m n = 0$
 $\langle \text{proof} \rangle$

lemma *ncols-le*: $(\text{ncols } A \leq n) = (! j i. n \leq i \longrightarrow (\text{Rep-matrix } A j i) = 0)$ (**is** $- = ?st$)
 $\langle \text{proof} \rangle$

lemma *less-ncols*: $(n < \text{ncols } A) = (? j i. n \leq i \ \& \ (\text{Rep-matrix } A j i) \neq 0)$
 $\langle \text{proof} \rangle$

lemma *le-ncols*: $(n \leq \text{ncols } A) = (\forall m. (\forall j i. m \leq i \longrightarrow (\text{Rep-matrix } A j i) = 0) \longrightarrow n \leq m)$
 $\langle \text{proof} \rangle$

lemma *nrows-le*: $(\text{nrows } A \leq n) = (! j i. n \leq j \longrightarrow (\text{Rep-matrix } A j i) = 0)$ (**is** $?s$)
 $\langle \text{proof} \rangle$

lemma *less-nrows*: $(m < \text{nrows } A) = (? j i. m \leq j \ \& \ (\text{Rep-matrix } A j i) \neq 0)$
 $\langle \text{proof} \rangle$

lemma *le-nrows*: $(n \leq \text{nrows } A) = (\forall m. (\forall j i. m \leq j \longrightarrow (\text{Rep-matrix } A j i) = 0) \longrightarrow n \leq m)$

$i) = 0) \longrightarrow n \leq m)$
 $\langle \text{proof} \rangle$

lemma *nrows-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies m < \text{nrows } A$
 $\langle \text{proof} \rangle$

lemma *ncols-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies n < \text{ncols } A$
 $\langle \text{proof} \rangle$

lemma *finite-natarray1*: $\text{finite } \{x. x < (n::\text{nat})\}$
 $\langle \text{proof} \rangle$

lemma *finite-natarray2*: $\text{finite } \{\text{pos}. (\text{fst pos}) < (m::\text{nat}) \ \& \ (\text{snd pos}) < (n::\text{nat})\}$
 $\langle \text{proof} \rangle$

lemma *RepAbs-matrix*:
assumes $\text{aem}: ?m. !j \ i. m \leq j \longrightarrow x \ j \ i = 0$ **(is ?em)** **and** $\text{aen}: ?n. !j \ i. (n \leq i \longrightarrow x \ j \ i = 0)$ **(is ?en)**
shows $(\text{Rep-matrix } (\text{Abs-matrix } x)) = x$
 $\langle \text{proof} \rangle$

definition *apply-infmatrix* :: $('a \Rightarrow 'b) \Rightarrow 'a \ \text{infmatrix} \Rightarrow 'b \ \text{infmatrix}$ **where**
 $\text{apply-infmatrix } f == \% A. (\% j \ i. f \ (A \ j \ i))$

definition *apply-matrix* :: $('a \Rightarrow 'b) \Rightarrow ('a::\text{zero}) \ \text{matrix} \Rightarrow ('b::\text{zero}) \ \text{matrix}$ **where**
 $\text{apply-matrix } f == \% A. \text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))$

definition *combine-infmatrix* :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \ \text{infmatrix} \Rightarrow 'b \ \text{infmatrix} \Rightarrow 'c \ \text{infmatrix}$ **where**
 $\text{combine-infmatrix } f == \% A \ B. (\% j \ i. f \ (A \ j \ i) \ (B \ j \ i))$

definition *combine-matrix* :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a::\text{zero}) \ \text{matrix} \Rightarrow ('b::\text{zero}) \ \text{matrix} \Rightarrow ('c::\text{zero}) \ \text{matrix}$ **where**
 $\text{combine-matrix } f == \% A \ B. \text{Abs-matrix } (\text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B))$

lemma *expand-apply-infmatrix[simp]*: $\text{apply-infmatrix } f \ A \ j \ i = f \ (A \ j \ i)$
 $\langle \text{proof} \rangle$

lemma *expand-combine-infmatrix[simp]*: $\text{combine-infmatrix } f \ A \ B \ j \ i = f \ (A \ j \ i) \ (B \ j \ i)$
 $\langle \text{proof} \rangle$

definition *commutative* :: $('a \Rightarrow 'a \Rightarrow 'b) \Rightarrow \text{bool}$ **where**
 $\text{commutative } f == !x \ y. f \ x \ y = f \ y \ x$

definition *associative* :: $('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{bool}$ **where**
 $\text{associative } f == !x \ y \ z. f \ (f \ x \ y) \ z = f \ x \ (f \ y \ z)$

To reason about associativity and commutativity of operations on matrices, let's take a step back and look at the general situation: Assume that we have sets A and B with $B \subset A$ and an abstraction $u : A \rightarrow B$. This abstraction has to fulfill $u(b) = b$ for all $b \in B$, but is arbitrary otherwise. Each function $f : A \times A \rightarrow A$ now induces a function $f' : B \times B \rightarrow B$ by $f' = u \circ f$. It is obvious that commutativity of f implies commutativity of f' : $f'xy = u(fxy) = u(fyx) = f'yx$.

lemma *combine-infmatrix-commute*:

commutative f \implies commutative (combine-infmatrix f)
<proof>

lemma *combine-matrix-commute*:

commutative f \implies commutative (combine-matrix f)
<proof>

On the contrary, given an associative function f we cannot expect f' to be associative. A counterexample is given by $A = \mathbb{Z}$, $B = \{-1, 0, 1\}$, as f we take addition on \mathbb{Z} , which is clearly associative. The abstraction is given by $u(a) = 0$ for $a \notin B$. Then we have

$$f'(f'11) - 1 = u(f(u(f11)) - 1) = u(f(u2) - 1) = u(f0 - 1) = -1,$$

but on the other hand we have

$$f'1(f'1 - 1) = u(f1(u(f1 - 1))) = u(f10) = 1.$$

A way out of this problem is to assume that $f(A \times A) \subset A$ holds, and this is what we are going to do:

lemma *nonzero-positions-combine-infmatrix[simp]*: *$f \ 0 \ 0 = 0 \implies \text{nonzero-positions (combine-infmatrix f A B)} \subseteq (\text{nonzero-positions A}) \cup (\text{nonzero-positions B})$*
<proof>

lemma *finite-nonzero-positions-Rep[simp]*: *finite (nonzero-positions (Rep-matrix A))*
<proof>

lemma *combine-infmatrix-closed [simp]*:

$f \ 0 \ 0 = 0 \implies \text{Rep-matrix (Abs-matrix (combine-infmatrix f (Rep-matrix A) (Rep-matrix B)))} = \text{combine-infmatrix f (Rep-matrix A) (Rep-matrix B)}$
<proof>

We need the next two lemmas only later, but it is analog to the above one, so we prove them now:

lemma *nonzero-positions-apply-infmatrix[simp]*: *$f \ 0 = 0 \implies \text{nonzero-positions (apply-infmatrix f A)} \subseteq \text{nonzero-positions A}$*
<proof>

lemma *apply-infmatrix-closed* [simp]:

$f\ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{apply-infmatrix } f\ (\text{Rep-matrix } A))) =$
 $\text{apply-infmatrix } f\ (\text{Rep-matrix } A)$
 ⟨proof⟩

lemma *combine-infmatrix-assoc*[simp]: $f\ 0\ 0 = 0 \implies \text{associative } f \implies \text{associative}$
 $(\text{combine-infmatrix } f)$
 ⟨proof⟩

lemma *comb*: $f = g \implies x = y \implies f\ x = g\ y$
 ⟨proof⟩

lemma *combine-matrix-assoc*: $f\ 0\ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-matrix}$
 $f)$
 ⟨proof⟩

lemma *Rep-apply-matrix*[simp]: $f\ 0 = 0 \implies \text{Rep-matrix } (\text{apply-matrix } f\ A)\ j\ i =$
 $f\ (\text{Rep-matrix } A\ j\ i)$
 ⟨proof⟩

lemma *Rep-combine-matrix*[simp]: $f\ 0\ 0 = 0 \implies \text{Rep-matrix } (\text{combine-matrix } f$
 $A\ B)\ j\ i = f\ (\text{Rep-matrix } A\ j\ i)\ (\text{Rep-matrix } B\ j\ i)$
 ⟨proof⟩

lemma *combine-nrows-max*: $f\ 0\ 0 = 0 \implies \text{nrows } (\text{combine-matrix } f\ A\ B) \leq$
 $\text{max } (\text{nrows } A)\ (\text{nrows } B)$
 ⟨proof⟩

lemma *combine-ncols-max*: $f\ 0\ 0 = 0 \implies \text{ncols } (\text{combine-matrix } f\ A\ B) \leq \text{max}$
 $(\text{ncols } A)\ (\text{ncols } B)$
 ⟨proof⟩

lemma *combine-nrows*: $f\ 0\ 0 = 0 \implies \text{nrows } A \leq q \implies \text{nrows } B \leq q \implies$
 $\text{nrows}(\text{combine-matrix } f\ A\ B) \leq q$
 ⟨proof⟩

lemma *combine-ncols*: $f\ 0\ 0 = 0 \implies \text{ncols } A \leq q \implies \text{ncols } B \leq q \implies$
 $\text{ncols}(\text{combine-matrix } f\ A\ B) \leq q$
 ⟨proof⟩

definition *zero-r-neutral* :: $('a \Rightarrow 'b::\text{zero} \Rightarrow 'a) \Rightarrow \text{bool}$ **where**
 $\text{zero-r-neutral } f == ! a. f\ a\ 0 = a$

definition *zero-l-neutral* :: $('a::\text{zero} \Rightarrow 'b \Rightarrow 'b) \Rightarrow \text{bool}$ **where**
 $\text{zero-l-neutral } f == ! a. f\ 0\ a = a$

definition *zero-closed* :: $(('a::\text{zero} \Rightarrow ('b::\text{zero} \Rightarrow ('c::\text{zero})) \Rightarrow \text{bool}) \Rightarrow \text{bool}$ **where**
 $\text{zero-closed } f == (!x. f\ x\ 0 = 0) \ \&\ (!y. f\ 0\ y = 0)$

consts *foldseq* :: ('a ⇒ 'a ⇒ 'a) ⇒ (nat ⇒ 'a) ⇒ nat ⇒ 'a

primrec

foldseq *f* *s* 0 = *s* 0

foldseq *f* *s* (Suc *n*) = *f* (*s* 0) (*foldseq* *f* (% *k*. *s*(Suc *k*)) *n*)

consts *foldseq-transposed* :: ('a ⇒ 'a ⇒ 'a) ⇒ (nat ⇒ 'a) ⇒ nat ⇒ 'a

primrec

foldseq-transposed *f* *s* 0 = *s* 0

foldseq-transposed *f* *s* (Suc *n*) = *f* (*foldseq-transposed* *f* *s* *n*) (*s* (Suc *n*))

lemma *foldseq-assoc* : *associative f* ⇒ *foldseq f* = *foldseq-transposed f*

⟨*proof*⟩

lemma *foldseq-distr*: [[*associative f*; *commutative f*]] ⇒ *foldseq f* (% *k*. *f* (*u* *k*) (*v* *k*)) *n* = *f* (*foldseq f* *u* *n*) (*foldseq f* *v* *n*)

⟨*proof*⟩

theorem [[*associative f*; *associative g*; ∀ *a b c d*. *g* (*f* *a* *b*) (*f* *c* *d*) = *f* (*g* *a* *c*) (*g* *b* *d*); ? *x y*. (*f* *x*) ≠ (*f* *y*); ? *x y*. (*g* *x*) ≠ (*g* *y*); *f* *x* *x* = *x*; *g* *x* *x* = *x*]] ⇒ *f*=*g* | (! *y*. *f* *y* *x* = *y*) | (! *y*. *g* *y* *x* = *y*)

⟨*proof*⟩

lemma *foldseq-zero*:

assumes *fz*: *f* 0 0 = 0 **and** *sz*: ! *i*. *i* ≤ *n* ⇒ *s* *i* = 0

shows *foldseq f* *s* *n* = 0

⟨*proof*⟩

lemma *foldseq-significant-positions*:

assumes *p*: ! *i*. *i* ≤ *N* ⇒ *S* *i* = *T* *i*

shows *foldseq f* *S* *N* = *foldseq f* *T* *N*

⟨*proof*⟩

lemma *foldseq-tail*:

assumes *M* ≤ *N*

shows *foldseq f* *S* *N* = *foldseq f* (% *k*. (if *k* < *M* then (*S* *k*) else (*foldseq f* (% *k*. *S*(*k*+*M*)) (*N*-*M*)))) *M*

⟨*proof*⟩

lemma *foldseq-zerotail*:

assumes

fz: *f* 0 0 = 0

and *sz*: ! *i*. *n* ≤ *i* ⇒ *s* *i* = 0

and *nm*: *n* ≤ *m*

shows

foldseq f *s* *n* = *foldseq f* *s* *m*

⟨*proof*⟩

lemma *foldseq-zerotail2*:

assumes ! $x. f\ x\ 0 = x$
and ! $i. n < i \longrightarrow s\ i = 0$
and $nm: n \leq m$
shows $foldseq\ f\ s\ n = foldseq\ f\ s\ m$
 $\langle proof \rangle$

lemma *foldseq-zerostart*:
 ! $x. f\ 0\ (f\ 0\ x) = f\ 0\ x \implies ! i. i \leq n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$
 $\langle proof \rangle$

lemma *foldseq-zerostart2*:
 ! $x. f\ 0\ x = x \implies ! i. i < n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ n = s\ n$
 $\langle proof \rangle$

lemma *foldseq-almostzero*:
assumes $f0x: ! x. f\ 0\ x = x$ **and** $fx0: ! x. f\ x\ 0 = x$ **and** $s0: ! i. i \neq j \longrightarrow s\ i = 0$
shows $foldseq\ f\ s\ n = (if\ (j \leq n)\ then\ (s\ j)\ else\ 0)$
 $\langle proof \rangle$

lemma *foldseq-distr-unary*:
assumes !! $a\ b. g\ (f\ a\ b) = f\ (g\ a)\ (g\ b)$
shows $g(foldseq\ f\ s\ n) = foldseq\ f\ (\% x. g(s\ x))\ n$
 $\langle proof \rangle$

definition *mult-matrix-n* :: $nat \Rightarrow (('a::zero) \Rightarrow ('b::zero) \Rightarrow ('c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$ **where**
 $mult-matrix-n\ n\ fmul\ fadd\ A\ B == Abs-matrix(\% j\ i. foldseq\ fadd\ (\% k. fmul\ (Rep-matrix\ A\ j\ k)\ (Rep-matrix\ B\ k\ i))\ n)$

definition *mult-matrix* :: $((('a::zero) \Rightarrow ('b::zero) \Rightarrow ('c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$ **where**
 $mult-matrix\ fmul\ fadd\ A\ B == mult-matrix-n\ (\max\ (ncols\ A)\ (nrows\ B))\ fmul\ fadd\ A\ B$

lemma *mult-matrix-n*:
assumes $ncols\ A \leq n$ (**is** ?An) $nrows\ B \leq n$ (**is** ?Bn) $fadd\ 0\ 0 = 0$ $fmul\ 0\ 0 = 0$
shows $c:mult-matrix\ fmul\ fadd\ A\ B = mult-matrix-n\ n\ fmul\ fadd\ A\ B$
 $\langle proof \rangle$

lemma *mult-matrix-nm*:
assumes $ncols\ A \leq n$ $nrows\ B \leq n$ $ncols\ A \leq m$ $nrows\ B \leq m$ $fadd\ 0\ 0 = 0$ $fmul\ 0\ 0 = 0$
shows $mult-matrix-n\ n\ fmul\ fadd\ A\ B = mult-matrix-n\ m\ fmul\ fadd\ A\ B$
 $\langle proof \rangle$

definition *r-distributive* :: $('a \Rightarrow 'b \Rightarrow 'b) \Rightarrow ('b \Rightarrow 'b \Rightarrow 'b) \Rightarrow bool$ **where**
 $r-distributive\ fmul\ fadd == ! a\ u\ v. fmul\ a\ (fadd\ u\ v) = fadd\ (fmul\ a\ u)\ (fmul\ a\ v)$

$v)$

definition *l-distributive* :: ($'a \Rightarrow 'b \Rightarrow 'a$) \Rightarrow ($'a \Rightarrow 'a \Rightarrow 'a$) \Rightarrow *bool* **where**
 l-distributive *fmul fadd* == ! $a\ u\ v.$ *fmul* (*fadd* $u\ v$) a = *fadd* (*fmul* $u\ a$) (*fmul* $v\ a$)

definition *distributive* :: ($'a \Rightarrow 'a \Rightarrow 'a$) \Rightarrow ($'a \Rightarrow 'a \Rightarrow 'a$) \Rightarrow *bool* **where**
 distributive *fmul fadd* == *l-distributive* *fmul fadd* & *r-distributive* *fmul fadd*

lemma *max1*: !! $a\ x\ y.$ ($a::nat$) $\leq x \implies a \leq \max\ x\ y$ \langle *proof* \rangle

lemma *max2*: !! $b\ x\ y.$ ($b::nat$) $\leq y \implies b \leq \max\ x\ y$ \langle *proof* \rangle

lemma *r-distributive-matrix*:

assumes

r-distributive *fmul fadd*

associative *fadd*

commutative *fadd*

fadd $0\ 0 = 0$

 ! $a.$ *fmul* $a\ 0 = 0$

 ! $a.$ *fmul* $0\ a = 0$

shows *r-distributive* (*mult-matrix* *fmul fadd*) (*combine-matrix* *fadd*)

\langle *proof* \rangle

lemma *l-distributive-matrix*:

assumes

l-distributive *fmul fadd*

associative *fadd*

commutative *fadd*

fadd $0\ 0 = 0$

 ! $a.$ *fmul* $a\ 0 = 0$

 ! $a.$ *fmul* $0\ a = 0$

shows *l-distributive* (*mult-matrix* *fmul fadd*) (*combine-matrix* *fadd*)

\langle *proof* \rangle

instantiation *matrix* :: (*zero*) *zero*

begin

definition *zero-matrix-def* [*code del*]: $0 = \text{Abs-matrix } (\lambda j\ i.\ 0)$

instance \langle *proof* \rangle

end

lemma *Rep-zero-matrix-def*[*simp*]: *Rep-matrix* $0\ j\ i = 0$
 \langle *proof* \rangle

lemma *zero-matrix-def-nrows*[*simp*]: *nrows* $0 = 0$
 \langle *proof* \rangle

lemma *zero-matrix-def-ncols*[simp]: $\text{ncols } 0 = 0$
 ⟨proof⟩

lemma *combine-matrix-zero-l-neutral*: $\text{zero-l-neutral } f \implies \text{zero-l-neutral } (\text{combine-matrix } f)$
 ⟨proof⟩

lemma *combine-matrix-zero-r-neutral*: $\text{zero-r-neutral } f \implies \text{zero-r-neutral } (\text{combine-matrix } f)$
 ⟨proof⟩

lemma *mult-matrix-zero-closed*: $\llbracket \text{fadd } 0 \ 0 = 0; \text{zero-closed } \text{fmul} \rrbracket \implies \text{zero-closed}$
 (*mult-matrix fmul fadd*)
 ⟨proof⟩

lemma *mult-matrix-n-zero-right*[simp]: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies$
mult-matrix-n n fmul fadd A 0 = 0
 ⟨proof⟩

lemma *mult-matrix-n-zero-left*[simp]: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies$
mult-matrix-n n fmul fadd 0 A = 0
 ⟨proof⟩

lemma *mult-matrix-zero-left*[simp]: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies \text{mult-matrix}$
fmul fadd 0 A = 0
 ⟨proof⟩

lemma *mult-matrix-zero-right*[simp]: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies \text{mult-matrix}$
fmul fadd A 0 = 0
 ⟨proof⟩

lemma *apply-matrix-zero*[simp]: $f \ 0 = 0 \implies \text{apply-matrix } f \ 0 = 0$
 ⟨proof⟩

lemma *combine-matrix-zero*: $f \ 0 \ 0 = 0 \implies \text{combine-matrix } f \ 0 \ 0 = 0$
 ⟨proof⟩

lemma *transpose-matrix-zero*[simp]: $\text{transpose-matrix } 0 = 0$
 ⟨proof⟩

lemma *apply-zero-matrix-def*[simp]: $\text{apply-matrix } (\% \ x. \ 0) \ A = 0$
 ⟨proof⟩

definition *singleton-matrix* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow ('a::\text{zero}) \Rightarrow 'a \text{ matrix}$ **where**
singleton-matrix j i a == Abs-matrix(% m n. if j = m & i = n then a else 0)

definition *move-matrix* :: $('a::\text{zero}) \text{ matrix} \Rightarrow \text{int} \Rightarrow \text{int} \Rightarrow 'a \text{ matrix}$ **where**
move-matrix A y x == Abs-matrix(% j i. if (neg ((int j) - y)) | (neg ((int i) - x))
 then 0 else Rep-matrix A (nat ((int j) - y)) (nat ((int i) - x)))

definition *take-rows* :: ('a::zero) matrix \Rightarrow nat \Rightarrow 'a matrix **where**
take-rows A r == Abs-matrix(% j i. if (j < r) then (Rep-matrix A j i) else 0)

definition *take-columns* :: ('a::zero) matrix \Rightarrow nat \Rightarrow 'a matrix **where**
take-columns A c == Abs-matrix(% j i. if (i < c) then (Rep-matrix A j i) else 0)

definition *column-of-matrix* :: ('a::zero) matrix \Rightarrow nat \Rightarrow 'a matrix **where**
column-of-matrix A n == take-columns (move-matrix A 0 (- int n)) 1

definition *row-of-matrix* :: ('a::zero) matrix \Rightarrow nat \Rightarrow 'a matrix **where**
row-of-matrix A m == take-rows (move-matrix A (- int m) 0) 1

lemma *Rep-singleton-matrix[simp]*: Rep-matrix (singleton-matrix j i e) m n = (if j = m & i = n then e else 0)
 <proof>

lemma *apply-singleton-matrix[simp]*: f 0 = 0 \implies apply-matrix f (singleton-matrix j i x) = (singleton-matrix j i (f x))
 <proof>

lemma *singleton-matrix-zero[simp]*: singleton-matrix j i 0 = 0
 <proof>

lemma *nrows-singleton[simp]*: nrows(singleton-matrix j i e) = (if e = 0 then 0 else Suc j)
 <proof>

lemma *ncols-singleton[simp]*: ncols(singleton-matrix j i e) = (if e = 0 then 0 else Suc i)
 <proof>

lemma *combine-singleton*: f 0 0 = 0 \implies combine-matrix f (singleton-matrix j i a) (singleton-matrix j i b) = singleton-matrix j i (f a b)
 <proof>

lemma *transpose-singleton[simp]*: transpose-matrix (singleton-matrix j i a) = singleton-matrix i j a
 <proof>

lemma *Rep-move-matrix[simp]*:
 Rep-matrix (move-matrix A y x) j i =
 (if (neg ((int j) - y)) | (neg ((int i) - x)) then 0 else Rep-matrix A (nat((int j) - y))
 (nat((int i) - x)))
 <proof>

lemma *move-matrix-0-0[simp]*: move-matrix A 0 0 = A
 <proof>

lemma *move-matrix-ortho*: $\text{move-matrix } A \ j \ i = \text{move-matrix } (\text{move-matrix } A \ j \ 0) \ 0 \ i$
 <proof>

lemma *transpose-move-matrix[simp]*:
 $\text{transpose-matrix } (\text{move-matrix } A \ x \ y) = \text{move-matrix } (\text{transpose-matrix } A) \ y \ x$
 <proof>

lemma *move-matrix-singleton[simp]*: $\text{move-matrix } (\text{singleton-matrix } u \ v \ x) \ j \ i =$
 $(\text{if } (j + \text{int } u < 0) \mid (i + \text{int } v < 0) \text{ then } 0 \text{ else } (\text{singleton-matrix } (\text{nat } (j + \text{int } u)) \ (\text{nat } (i + \text{int } v)) \ x))$
 <proof>

lemma *Rep-take-columns[simp]*:
 $\text{Rep-matrix } (\text{take-columns } A \ c) \ j \ i =$
 $(\text{if } i < c \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$
 <proof>

lemma *Rep-take-rows[simp]*:
 $\text{Rep-matrix } (\text{take-rows } A \ r) \ j \ i =$
 $(\text{if } j < r \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$
 <proof>

lemma *Rep-column-of-matrix[simp]*:
 $\text{Rep-matrix } (\text{column-of-matrix } A \ c) \ j \ i = (\text{if } i = 0 \text{ then } (\text{Rep-matrix } A \ j \ c) \text{ else } 0)$
 <proof>

lemma *Rep-row-of-matrix[simp]*:
 $\text{Rep-matrix } (\text{row-of-matrix } A \ r) \ j \ i = (\text{if } j = 0 \text{ then } (\text{Rep-matrix } A \ r \ i) \text{ else } 0)$
 <proof>

lemma *column-of-matrix*: $\text{ncols } A \leq n \implies \text{column-of-matrix } A \ n = 0$
 <proof>

lemma *row-of-matrix*: $\text{nrows } A \leq n \implies \text{row-of-matrix } A \ n = 0$
 <proof>

lemma *mult-matrix-singleton-right[simp]*:
assumes
 $! x. \text{fmul } x \ 0 = 0$
 $! x. \text{fmul } 0 \ x = 0$
 $! x. \text{fadd } 0 \ x = x$
 $! x. \text{fadd } x \ 0 = x$
shows $(\text{mult-matrix } \text{fmul } \text{fadd } A \ (\text{singleton-matrix } j \ i \ e)) = \text{apply-matrix } (\% \ x. \text{fmul } x \ e) \ (\text{move-matrix } (\text{column-of-matrix } A \ j) \ 0 \ (\text{int } i))$
 <proof>

lemma *mult-matrix-ext*:

assumes

eprem:

? *e*. (! *a b*. *a* ≠ *b* ⟶ *fmul a e* ≠ *fmul b e*)

and *fprems*:

! *a*. *fmul 0 a* = 0

! *a*. *fmul a 0* = 0

! *a*. *fadd a 0* = *a*

! *a*. *fadd 0 a* = *a*

and *contraprems*:

mult-matrix fmul fadd A = *mult-matrix fmul fadd B*

shows

A = *B*

⟨*proof*⟩

definition *foldmatrix* :: (*'a* ⇒ *'a* ⇒ *'a*) ⇒ (*'a* ⇒ *'a* ⇒ *'a*) ⇒ (*'a infmatrix*) ⇒ *nat* ⇒ *nat* ⇒ *'a* **where**

foldmatrix f g A m n == *foldseq-transposed g (% j. foldseq f (A j) n) m*

definition *foldmatrix-transposed* :: (*'a* ⇒ *'a* ⇒ *'a*) ⇒ (*'a* ⇒ *'a* ⇒ *'a*) ⇒ (*'a infmatrix*) ⇒ *nat* ⇒ *nat* ⇒ *'a* **where**

foldmatrix-transposed f g A m n == *foldseq g (% j. foldseq-transposed f (A j) n) m*

lemma *foldmatrix-transpose*:

assumes

! *a b c d*. *g(f a b) (f c d)* = *f (g a c) (g b d)*

shows

foldmatrix f g A m n = *foldmatrix-transposed g f (transpose-infmatrix A) n m*

⟨*proof*⟩

lemma *foldseq-foldseq*:

assumes

associative f

associative g

! *a b c d*. *g(f a b) (f c d)* = *f (g a c) (g b d)*

shows

foldseq g (% j. foldseq f (A j) n) m = *foldseq f (% j. foldseq g ((transpose-infmatrix A) j) m) n*

⟨*proof*⟩

lemma *mult-n-nrows*:

assumes

! *a*. *fmul 0 a* = 0

! *a*. *fmul a 0* = 0

fadd 0 0 = 0

shows *nrows (mult-matrix-n n fmul fadd A B)* ≤ *nrows A*

⟨*proof*⟩

lemma *mult-n-ncols:*

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows $\text{ncols } (\text{mult-matrix-n } n \ \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$

$\langle \text{proof} \rangle$

lemma *mult-nrows:*

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows $\text{nrows } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{nrows } A$

$\langle \text{proof} \rangle$

lemma *mult-ncols:*

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows $\text{ncols } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$

$\langle \text{proof} \rangle$

lemma *nrows-move-matrix-le:* $\text{nrows } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{nrows } A)) + j)$

$\langle \text{proof} \rangle$

lemma *ncols-move-matrix-le:* $\text{ncols } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{ncols } A)) + i)$

$\langle \text{proof} \rangle$

lemma *mult-matrix-assoc:*

assumes

$! a. \text{fmul1 } 0 \ a = 0$

$! a. \text{fmul1 } a \ 0 = 0$

$! a. \text{fmul2 } 0 \ a = 0$

$! a. \text{fmul2 } a \ 0 = 0$

$\text{fadd1 } 0 \ 0 = 0$

$\text{fadd2 } 0 \ 0 = 0$

$! a \ b \ c \ d. \text{fadd2 } (\text{fadd1 } a \ b) (\text{fadd1 } c \ d) = \text{fadd1 } (\text{fadd2 } a \ c) (\text{fadd2 } b \ d)$

associative fadd1

associative fadd2

$! a \ b \ c. \text{fmul2 } (\text{fmul1 } a \ b) \ c = \text{fmul1 } a \ (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul2 } (\text{fadd1 } a \ b) \ c = \text{fadd1 } (\text{fmul2 } a \ c) (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul1 } c \ (\text{fadd2 } a \ b) = \text{fadd2 } (\text{fmul1 } c \ a) (\text{fmul1 } c \ b)$

shows $\text{mult-matrix } \text{fmul2 } \text{fadd2 } (\text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ B) \ C = \text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ (\text{mult-matrix } \text{fmul2 } \text{fadd2 } B \ C)$

$\langle \text{proof} \rangle$

lemma

assumes

$! a. \text{fmul1 } 0 \ a = 0$

$! a. \text{fmul1 } a \ 0 = 0$

$! a. \text{fmul2 } 0 \ a = 0$

$! a. \text{fmul2 } a \ 0 = 0$

$\text{fadd1 } 0 \ 0 = 0$

$\text{fadd2 } 0 \ 0 = 0$

$! a \ b \ c \ d. \text{fadd2 } (\text{fadd1 } a \ b) (\text{fadd1 } c \ d) = \text{fadd1 } (\text{fadd2 } a \ c) (\text{fadd2 } b \ d)$

associative fadd1

associative fadd2

$! a \ b \ c. \text{fmul2 } (\text{fmul1 } a \ b) \ c = \text{fmul1 } a \ (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul2 } (\text{fadd1 } a \ b) \ c = \text{fadd1 } (\text{fmul2 } a \ c) (\text{fmul2 } b \ c)$

$! a \ b \ c. \text{fmul1 } c \ (\text{fadd2 } a \ b) = \text{fadd2 } (\text{fmul1 } c \ a) (\text{fmul1 } c \ b)$

shows

$(\text{mult-matrix } \text{fmul1 } \text{fadd1 } A) \circ (\text{mult-matrix } \text{fmul2 } \text{fadd2 } B) = \text{mult-matrix } \text{fmul2}$
 $\text{fadd2 } (\text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ B)$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-assoc-simple:*

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

associative fadd

commutative fadd

associative fmul

distributive fmul fadd

shows $\text{mult-matrix } \text{fmul } \text{fadd } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \ C = \text{mult-matrix } \text{fmul}$
 $\text{fadd } A \ (\text{mult-matrix } \text{fmul } \text{fadd } B \ C)$
 $\langle \text{proof} \rangle$

lemma *transpose-apply-matrix:* $f \ 0 = 0 \implies \text{transpose-matrix } (\text{apply-matrix } f \ A)$
 $= \text{apply-matrix } f \ (\text{transpose-matrix } A)$
 $\langle \text{proof} \rangle$

lemma *transpose-combine-matrix:* $f \ 0 \ 0 = 0 \implies \text{transpose-matrix } (\text{combine-matrix}$
 $f \ A \ B) = \text{combine-matrix } f \ (\text{transpose-matrix } A) \ (\text{transpose-matrix } B)$
 $\langle \text{proof} \rangle$

lemma *Rep-mult-matrix:*

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows

$\text{Rep-matrix}(\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \ j \ i =$
 $\text{foldseq } \text{fadd } (\% \ k. \text{fmul } (\text{Rep-matrix } A \ j \ k) (\text{Rep-matrix } B \ k \ i)) \ (\text{max } (\text{ncols } A))$

(*nrows B*))
 <proof>

lemma *transpose-mult-matrix*:

assumes

! *a*. *fmul 0 a = 0*

! *a*. *fmul a 0 = 0*

fadd 0 0 = 0

! *x y*. *fmul y x = fmul x y*

shows

transpose-matrix (mult-matrix fmul fadd A B) = mult-matrix fmul fadd (transpose-matrix B) (transpose-matrix A)

<proof>

lemma *column-transpose-matrix*: *column-of-matrix (transpose-matrix A) n = transpose-matrix (row-of-matrix A n)*

<proof>

lemma *take-columns-transpose-matrix*: *take-columns (transpose-matrix A) n = transpose-matrix (take-rows A n)*

<proof>

instantiation *matrix* :: (*{zero, ord}*) *ord*
begin

definition

le-matrix-def: $A \leq B \longleftrightarrow (\forall j\ i. \text{Rep-matrix } A\ j\ i \leq \text{Rep-matrix } B\ j\ i)$

definition

less-def: $A < (B :: 'a\ matrix) \longleftrightarrow A \leq B \wedge \neg B \leq A$

instance <proof>

end

instance *matrix* :: (*{zero, order}*) *order*
 <proof>

lemma *le-apply-matrix*:

assumes

f 0 = 0

! *x y*. $x \leq y \longrightarrow f\ x \leq f\ y$

(*a*::('a::{ord, zero}) *matrix*) $\leq b$

shows

apply-matrix f a \leq *apply-matrix f b*

<proof>

lemma *le-combine-matrix*:

assumes

$f\ 0\ 0 = 0$
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$
 $A \leq B$
 $C \leq D$
shows
 $combine_matrix\ f\ A\ C \leq combine_matrix\ f\ B\ D$
 $\langle proof \rangle$

lemma *le-left-combine-matrix*:
assumes
 $f\ 0\ 0 = 0$
 $! a\ b\ c. a \leq b \longrightarrow f\ c\ a \leq f\ c\ b$
 $A \leq B$
shows
 $combine_matrix\ f\ C\ A \leq combine_matrix\ f\ C\ B$
 $\langle proof \rangle$

lemma *le-right-combine-matrix*:
assumes
 $f\ 0\ 0 = 0$
 $! a\ b\ c. a \leq b \longrightarrow f\ a\ c \leq f\ b\ c$
 $A \leq B$
shows
 $combine_matrix\ f\ A\ C \leq combine_matrix\ f\ B\ C$
 $\langle proof \rangle$

lemma *le-transpose-matrix*: $(A \leq B) = (transpose_matrix\ A \leq transpose_matrix\ B)$
 $\langle proof \rangle$

lemma *le-foldseq*:
assumes
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$
 $! i. i \leq n \longrightarrow s\ i \leq t\ i$
shows
 $foldseq\ f\ s\ n \leq foldseq\ f\ t\ n$
 $\langle proof \rangle$

lemma *le-left-mult*:
assumes
 $! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$
 $! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ c\ a \leq fmul\ c\ b$
 $! a. fmul\ 0\ a = 0$
 $! a. fmul\ a\ 0 = 0$
 $fadd\ 0\ 0 = 0$
 $0 \leq C$
 $A \leq B$
shows
 $mult_matrix\ fmul\ fadd\ C\ A \leq mult_matrix\ fmul\ fadd\ C\ B$

$\langle \text{proof} \rangle$

lemma *le-right-mult*:

assumes

$! a \ b \ c \ d. a \leq b \ \& \ c \leq d \longrightarrow \text{fadd } a \ c \leq \text{fadd } b \ d$

$! c \ a \ b. 0 \leq c \ \& \ a \leq b \longrightarrow \text{fmul } a \ c \leq \text{fmul } b \ c$

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

$0 \leq C$

$A \leq B$

shows

$\text{mult-matrix fmul fadd } A \ C \leq \text{mult-matrix fmul fadd } B \ C$

$\langle \text{proof} \rangle$

lemma *spec2*: $! j \ i. P \ j \ i \Longrightarrow P \ j \ i \ \langle \text{proof} \rangle$

lemma *neg-imp*: $(\neg Q \longrightarrow \neg P) \Longrightarrow P \longrightarrow Q \ \langle \text{proof} \rangle$

lemma *singleton-matrix-le[simp]*: $(\text{singleton-matrix } j \ i \ a \leq \text{singleton-matrix } j \ i \ b) = (a \leq (b::\text{'a}::\text{order}))$

$\langle \text{proof} \rangle$

lemma *singleton-le-zero[simp]*: $(\text{singleton-matrix } j \ i \ x \leq 0) = (x \leq (0::\text{'a}::\{\text{order}, \text{zero}\}))$

$\langle \text{proof} \rangle$

lemma *singleton-ge-zero[simp]*: $(0 \leq \text{singleton-matrix } j \ i \ x) = ((0::\text{'a}::\{\text{order}, \text{zero}\}) \leq x)$

$\langle \text{proof} \rangle$

lemma *move-matrix-le-zero[simp]*: $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (\text{move-matrix } A \ j \ i \leq 0) = (A \leq (0::\text{'a}::\{\text{order}, \text{zero}\}) \text{ matrix})$

$\langle \text{proof} \rangle$

lemma *move-matrix-zero-le[simp]*: $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (0 \leq \text{move-matrix } A \ j \ i) = ((0::\text{'a}::\{\text{order}, \text{zero}\}) \text{ matrix} \leq A)$

$\langle \text{proof} \rangle$

lemma *move-matrix-le-move-matrix-iff[simp]*: $0 \leq j \Longrightarrow 0 \leq i \Longrightarrow (\text{move-matrix } A \ j \ i \leq \text{move-matrix } B \ j \ i) = (A \leq (B::\text{'a}::\{\text{order}, \text{zero}\}) \text{ matrix})$

$\langle \text{proof} \rangle$

instantiation *matrix* :: $(\{\text{lattice}, \text{zero}\}) \text{ lattice}$

begin

definition [*code del*]: $\text{inf} = \text{combine-matrix inf}$

definition [*code del*]: $\text{sup} = \text{combine-matrix sup}$

instance

```

    <proof>

end

instantiation matrix :: ({plus, zero}) plus
begin

definition
  plus-matrix-def [code del]:  $A + B = \text{combine-matrix } (op +) A B$ 

instance <proof>

end

instantiation matrix :: ({uminus, zero}) uminus
begin

definition
  minus-matrix-def [code del]:  $- A = \text{apply-matrix } uminus A$ 

instance <proof>

end

instantiation matrix :: ({minus, zero}) minus
begin

definition
  diff-matrix-def [code del]:  $A - B = \text{combine-matrix } (op -) A B$ 

instance <proof>

end

instantiation matrix :: ({plus, times, zero}) times
begin

definition
  times-matrix-def [code del]:  $A * B = \text{mult-matrix } (op *) (op +) A B$ 

instance <proof>

end

instantiation matrix :: ({lattice, uminus, zero}) abs
begin

definition
  abs-matrix-def [code del]:  $abs (A :: 'a \text{ matrix}) = sup A (- A)$ 

```

```

instance ⟨proof⟩

end

instance matrix :: (monoid-add) monoid-add
⟨proof⟩

instance matrix :: (comm-monoid-add) comm-monoid-add
⟨proof⟩

instance matrix :: (group-add) group-add
⟨proof⟩

instance matrix :: (ab-group-add) ab-group-add
⟨proof⟩

instance matrix :: (ordered-ab-group-add) ordered-ab-group-add
⟨proof⟩

instance matrix :: (lattice-ab-group-add) semilattice-inf-ab-group-add ⟨proof⟩
instance matrix :: (lattice-ab-group-add) semilattice-sup-ab-group-add ⟨proof⟩

instance matrix :: (semiring-0) semiring-0
⟨proof⟩

instance matrix :: (ring) ring ⟨proof⟩

instance matrix :: (ordered-ring) ordered-ring
⟨proof⟩

instance matrix :: (lattice-ring) lattice-ring
⟨proof⟩

lemma Rep-matrix-add[simp]:
  Rep-matrix ((a::('a::monoid-add)matrix)+b) j i = (Rep-matrix a j i) + (Rep-matrix
  b j i)
  ⟨proof⟩

lemma Rep-matrix-mult: Rep-matrix ((a::('a::semiring-0) matrix) * b) j i =
  foldseq (op +) (% k. (Rep-matrix a j k) * (Rep-matrix b k i)) (max (ncols a)
  (nrows b))
  ⟨proof⟩

lemma apply-matrix-add: ! x y. f (x+y) = (f x) + (f y) ⟹ f 0 = (0::'a)
  ⟹ apply-matrix f ((a::('a::monoid-add) matrix) + b) = (apply-matrix f a) +
  (apply-matrix f b)
  ⟨proof⟩

```


lemma *singleton-matrix-add*: $\text{singleton-matrix } j \ i \ ((a::\text{monoid-add})+b) = (\text{singleton-matrix } j \ i \ a) + (\text{singleton-matrix } j \ i \ b)$
 <proof>

lemma *nrows-mult*: $\text{nrows } ((A::('a::\text{semiring-0}) \text{ matrix}) * B) \leq \text{nrows } A$
 <proof>

lemma *ncols-mult*: $\text{ncols } ((A::('a::\text{semiring-0}) \text{ matrix}) * B) \leq \text{ncols } B$
 <proof>

definition

$\text{one-matrix} :: \text{nat} \Rightarrow ('a::\{\text{zero,one}\}) \text{ matrix}$ **where**
 $\text{one-matrix } n = \text{Abs-matrix } (\% j \ i. \text{ if } j = i \ \& \ j < n \text{ then } 1 \text{ else } 0)$

lemma *Rep-one-matrix[simp]*: $\text{Rep-matrix } (\text{one-matrix } n) \ j \ i = (\text{if } (j = i \ \& \ j < n) \text{ then } 1 \text{ else } 0)$
 <proof>

lemma *nrows-one-matrix[simp]*: $\text{nrows } ((\text{one-matrix } n) :: ('a::\text{zero-neq-one}) \text{ matrix}) = n$ (**is** ?r = -)
 <proof>

lemma *ncols-one-matrix[simp]*: $\text{ncols } ((\text{one-matrix } n) :: ('a::\text{zero-neq-one}) \text{ matrix}) = n$ (**is** ?r = -)
 <proof>

lemma *one-matrix-mult-right[simp]*: $\text{ncols } A \leq n \implies (A::('a::\{\text{semiring-1}\}) \text{ matrix}) * (\text{one-matrix } n) = A$
 <proof>

lemma *one-matrix-mult-left[simp]*: $\text{nrows } A \leq n \implies (\text{one-matrix } n) * A = (A::('a::\text{semiring-1}) \text{ matrix})$
 <proof>

lemma *transpose-matrix-mult*: $\text{transpose-matrix } ((A::('a::\text{comm-ring}) \text{ matrix}) * B) = (\text{transpose-matrix } B) * (\text{transpose-matrix } A)$
 <proof>

lemma *transpose-matrix-add*: $\text{transpose-matrix } ((A::('a::\text{monoid-add}) \text{ matrix}) + B) = \text{transpose-matrix } A + \text{transpose-matrix } B$
 <proof>

lemma *transpose-matrix-diff*: $\text{transpose-matrix } ((A::('a::\text{group-add}) \text{ matrix}) - B) = \text{transpose-matrix } A - \text{transpose-matrix } B$
 <proof>

lemma *transpose-matrix-minus*: $\text{transpose-matrix } (-(A::('a::\text{group-add}) \text{ matrix})) = - \text{transpose-matrix } (A::'a \text{ matrix})$
 <proof>

definition *right-inverse-matrix* :: ('a::{ring-1}) matrix \Rightarrow 'a matrix \Rightarrow bool **where**
right-inverse-matrix A X == (A * X = one-matrix (max (nrows A) (ncols X)))
 \wedge nrows X \leq ncols A

definition *left-inverse-matrix* :: ('a::{ring-1}) matrix \Rightarrow 'a matrix \Rightarrow bool **where**
left-inverse-matrix A X == (X * A = one-matrix (max(nrows X) (ncols A))) \wedge
ncols X \leq nrows A

definition *inverse-matrix* :: ('a::{ring-1}) matrix \Rightarrow 'a matrix \Rightarrow bool **where**
inverse-matrix A X == (right-inverse-matrix A X) \wedge (left-inverse-matrix A X)

lemma *right-inverse-matrix-dim*: right-inverse-matrix A X \implies nrows A = ncols X
 \langle proof \rangle

lemma *left-inverse-matrix-dim*: left-inverse-matrix A Y \implies ncols A = nrows Y
 \langle proof \rangle

lemma *left-right-inverse-matrix-unique*:
assumes left-inverse-matrix A Y right-inverse-matrix A X
shows X = Y
 \langle proof \rangle

lemma *inverse-matrix-inject*: \llbracket inverse-matrix A X; inverse-matrix A Y $\rrbracket \implies$ X = Y
 \langle proof \rangle

lemma *one-matrix-inverse*: inverse-matrix (one-matrix n) (one-matrix n)
 \langle proof \rangle

lemma *zero-imp-mult-zero*: (a::'a::semiring-0) = 0 \mid b = 0 \implies a * b = 0
 \langle proof \rangle

lemma *Rep-matrix-zero-imp-mult-zero*:
! j i k. (Rep-matrix A j k = 0) \mid (Rep-matrix B k i) = 0 \implies A * B =
(0::('a::lattice-ring) matrix)
 \langle proof \rangle

lemma *add-nrows*: nrows (A::('a::monoid-add) matrix) \leq u \implies nrows B \leq u
 \implies nrows (A + B) \leq u
 \langle proof \rangle

lemma *move-matrix-row-mult*: move-matrix ((A::('a::semiring-0) matrix) * B) j
0 = (move-matrix A j 0) * B
 \langle proof \rangle

lemma *move-matrix-col-mult*: move-matrix ((A::('a::semiring-0) matrix) * B) 0
i = A * (move-matrix B 0 i)

⟨proof⟩

lemma *move-matrix-add*: $((\text{move-matrix } (A + B) \ j \ i) :: ('a :: \text{monoid-add}) \ \text{matrix}))$
= $(\text{move-matrix } A \ j \ i) + (\text{move-matrix } B \ j \ i)$
⟨proof⟩

lemma *move-matrix-mult*: $\text{move-matrix } ((A :: ('a :: \text{semiring-0}) \ \text{matrix}) * B) \ j \ i =$
 $(\text{move-matrix } A \ j \ 0) * (\text{move-matrix } B \ 0 \ i)$
⟨proof⟩

definition *scalar-mult* :: $('a :: \text{ring}) \Rightarrow 'a \ \text{matrix} \Rightarrow 'a \ \text{matrix}$ **where**
scalar-mult $a \ m == \text{apply-matrix } (op * a) \ m$

lemma *scalar-mult-zero[simp]*: $\text{scalar-mult } y \ 0 = 0$
⟨proof⟩

lemma *scalar-mult-add*: $\text{scalar-mult } y \ (a + b) = (\text{scalar-mult } y \ a) + (\text{scalar-mult } y \ b)$
⟨proof⟩

lemma *Rep-scalar-mult[simp]*: $\text{Rep-matrix } (\text{scalar-mult } y \ a) \ j \ i = y * (\text{Rep-matrix } a \ j \ i)$
⟨proof⟩

lemma *scalar-mult-singleton[simp]*: $\text{scalar-mult } y \ (\text{singleton-matrix } j \ i \ x) = \text{singleton-matrix } j \ i \ (y * x)$
⟨proof⟩

lemma *Rep-minus[simp]*: $\text{Rep-matrix } (-(A :: (\text{group-add})) \ x \ y) = - (\text{Rep-matrix } A \ x \ y)$
⟨proof⟩

lemma *Rep-abs[simp]*: $\text{Rep-matrix } (\text{abs } (A :: (\text{lattice-ab-group-add})) \ x \ y) = \text{abs } (\text{Rep-matrix } A \ x \ y)$
⟨proof⟩

end

theory *SparseMatrix*
imports *Matrix*
begin

types
 $'a \ \text{svec} = (\text{nat} * 'a) \ \text{list}$
 $'a \ \text{spmat} = ('a \ \text{svec}) \ \text{svec}$

definition *sparse-row-vector* :: $('a :: \text{ab-group-add}) \ \text{svec} \Rightarrow 'a \ \text{matrix}$ **where**

sparse-row-vector-def: $\text{sparse-row-vector } arr = \text{foldl } (\% \ m \ x. \ m + (\text{singleton-matrix } 0 \ (\text{fst } x) \ (\text{snd } x))) \ 0 \ arr$

definition *sparse-row-matrix* :: ('a::ab-group-add) *spmat* \Rightarrow 'a *matrix* **where**
sparse-row-matrix-def: $\text{sparse-row-matrix } arr = \text{foldl } (\% \ m \ r. \ m + (\text{move-matrix } (\text{sparse-row-vector } (\text{snd } r)) \ (\text{int } (\text{fst } r)) \ 0)) \ 0 \ arr$

code-datatype *sparse-row-vector* *sparse-row-matrix*

lemma *sparse-row-vector-empty* [simp]: $\text{sparse-row-vector } [] = 0$
 <proof>

lemma *sparse-row-matrix-empty* [simp]: $\text{sparse-row-matrix } [] = 0$
 <proof>

lemmas [code] = *sparse-row-vector-empty* [symmetric]

lemma *foldl-distrstart*: $! \ a \ x \ y. \ (f \ (g \ x \ y) \ a = g \ x \ (f \ y \ a)) \Longrightarrow \ (\text{foldl } f \ (g \ x \ y) \ l = g \ x \ (\text{foldl } f \ y \ l))$
 <proof>

lemma *sparse-row-vector-cons*[simp]:
 $\text{sparse-row-vector } (a \# \ arr) = (\text{singleton-matrix } 0 \ (\text{fst } a) \ (\text{snd } a)) + (\text{sparse-row-vector } arr)$
 <proof>

lemma *sparse-row-vector-append*[simp]:
 $\text{sparse-row-vector } (a \ @ \ b) = (\text{sparse-row-vector } a) + (\text{sparse-row-vector } b)$
 <proof>

lemma *nrows-spvec*[simp]: $\text{nrows } (\text{sparse-row-vector } x) \leq (\text{Suc } 0)$
 <proof>

lemma *sparse-row-matrix-cons*: $\text{sparse-row-matrix } (a \# \ arr) = ((\text{move-matrix } (\text{sparse-row-vector } (\text{snd } a)) \ (\text{int } (\text{fst } a)) \ 0)) + \text{sparse-row-matrix } arr$
 <proof>

lemma *sparse-row-matrix-append*: $\text{sparse-row-matrix } (arr \ @ \ brr) = (\text{sparse-row-matrix } arr) + (\text{sparse-row-matrix } brr)$
 <proof>

primrec *sorted-spvec* :: 'a *spvec* \Rightarrow *bool* **where**
 $\text{sorted-spvec } [] = \text{True}$
 $| \text{sorted-spvec-step: sorted-spvec } (a \# \ as) = (\text{case as of } [] \Rightarrow \text{True} \mid b \# \ bs \Rightarrow ((\text{fst } a < \text{fst } b) \ \& \ (\text{sorted-spvec } as)))$

primrec *sorted-spmat* :: 'a *spmat* \Rightarrow *bool* **where**
 $\text{sorted-spmat } [] = \text{True}$
 $| \text{sorted-spmat } (a \# \ as) = ((\text{sorted-spvec } (\text{snd } a)) \ \& \ (\text{sorted-spmat } as))$

declare *sorted-spvec.simps* [*simp del*]

lemma *sorted-spvec-empty*[*simp*]: *sorted-spvec* [] = *True*
 ⟨*proof*⟩

lemma *sorted-spvec-cons1*: *sorted-spvec* (a#as) \implies *sorted-spvec* as
 ⟨*proof*⟩

lemma *sorted-spvec-cons2*: *sorted-spvec* (a#b#t) \implies *sorted-spvec* (a#t)
 ⟨*proof*⟩

lemma *sorted-spvec-cons3*: *sorted-spvec*(a#b#t) \implies *fst* a < *fst* b
 ⟨*proof*⟩

lemma *sorted-sparse-row-vector-zero*[*rule-format*]: $m \leq n \implies$ *sorted-spvec* ((n,a)#arr)
 \longrightarrow *Rep-matrix* (*sparse-row-vector* arr) *j* *m* = 0
 ⟨*proof*⟩

lemma *sorted-sparse-row-matrix-zero*[*rule-format*]: $m \leq n \implies$ *sorted-spvec* ((n,a)#arr)
 \longrightarrow *Rep-matrix* (*sparse-row-matrix* arr) *m* *j* = 0
 ⟨*proof*⟩

primrec *minus-spvec* :: ('a::ab-group-add) *spvec* \Rightarrow 'a *spvec* **where**
minus-spvec [] = []
 | *minus-spvec* (a#as) = (*fst* a, -(*snd* a))#(*minus-spvec* as)

primrec *abs-spvec* :: ('a::lattice-ab-group-add-abs) *spvec* \Rightarrow 'a *spvec* **where**
abs-spvec [] = []
 | *abs-spvec* (a#as) = (*fst* a, *abs* (*snd* a))#(*abs-spvec* as)

lemma *sparse-row-vector-minus*:
sparse-row-vector (*minus-spvec* v) = - (*sparse-row-vector* v)
 ⟨*proof*⟩

instance *matrix* :: (*lattice-ab-group-add-abs*) *lattice-ab-group-add-abs*
 ⟨*proof*⟩

lemma *sparse-row-vector-abs*:
sorted-spvec (v :: 'a::lattice-ring *spvec*) \implies *sparse-row-vector* (*abs-spvec* v) = *abs*
 (*sparse-row-vector* v)
 ⟨*proof*⟩

lemma *sorted-spvec-minus-spvec*:
sorted-spvec v \implies *sorted-spvec* (*minus-spvec* v)
 ⟨*proof*⟩

lemma *sorted-spvec-abs-spvec*:
sorted-spvec v \implies *sorted-spvec* (*abs-spvec* v)

$\langle \text{proof} \rangle$

definition

$\text{smult-spvec } y = \text{map } (\% a. (\text{fst } a, y * \text{snd } a))$

lemma $\text{smult-spvec-empty}[simp]: \text{smult-spvec } y \ [] = []$

$\langle \text{proof} \rangle$

lemma $\text{smult-spvec-cons}: \text{smult-spvec } y (a \# \text{arr}) = (\text{fst } a, y * (\text{snd } a)) \# (\text{smult-spvec } y \ \text{arr})$

$\langle \text{proof} \rangle$

fun $\text{addmult-spvec} :: ('a::\text{ring}) \Rightarrow 'a \ \text{spvec} \Rightarrow 'a \ \text{spvec} \Rightarrow 'a \ \text{spvec}$ **where**

$\text{addmult-spvec } y \ \text{arr} \ [] = \text{arr} \mid$

$\text{addmult-spvec } y \ [] \ \text{brr} = \text{smult-spvec } y \ \text{brr} \mid$

$\text{addmult-spvec } y \ ((i,a) \# \text{arr}) \ ((j,b) \# \text{brr}) = ($

$\text{if } i < j \text{ then } ((i,a) \# (\text{addmult-spvec } y \ \text{arr} \ ((j,b) \# \text{brr})))$

$\text{else (if } (j < i) \text{ then } ((j, y * b) \# (\text{addmult-spvec } y \ ((i,a) \# \text{arr}) \ \text{brr}))$

$\text{else } ((i, a + y*b) \# (\text{addmult-spvec } y \ \text{arr} \ \text{brr})))$

lemma $\text{addmult-spvec-empty1}[simp]: \text{addmult-spvec } y \ [] \ a = \text{smult-spvec } y \ a$

$\langle \text{proof} \rangle$

lemma $\text{addmult-spvec-empty2}[simp]: \text{addmult-spvec } y \ a \ [] = a$

$\langle \text{proof} \rangle$

lemma $\text{sparse-row-vector-map}: (! x y. f (x+y) = (f x) + (f y)) \Longrightarrow (f :: 'a \Rightarrow ('a::\text{lattice-ring}))$

$0 = 0 \Longrightarrow$

$\text{sparse-row-vector } (\text{map } (\% x. (\text{fst } x, f (\text{snd } x))) \ a) = \text{apply-matrix } f \ (\text{sparse-row-vector } a)$

$\langle \text{proof} \rangle$

lemma $\text{sparse-row-vector-smult}: \text{sparse-row-vector } (\text{smult-spvec } y \ a) = \text{scalar-mult}$

$y \ (\text{sparse-row-vector } a)$

$\langle \text{proof} \rangle$

lemma $\text{sparse-row-vector-addmult-spvec}: \text{sparse-row-vector } (\text{addmult-spvec } (y :: 'a::\text{lattice-ring})$

$a \ b) =$

$(\text{sparse-row-vector } a) + (\text{scalar-mult } y \ (\text{sparse-row-vector } b))$

$\langle \text{proof} \rangle$

lemma $\text{sorted-smult-spvec}: \text{sorted-spvec } a \Longrightarrow \text{sorted-spvec } (\text{smult-spvec } y \ a)$

$\langle \text{proof} \rangle$

lemma $\text{sorted-spvec-addmult-spvec-helper}: \llbracket \text{sorted-spvec } (\text{addmult-spvec } y \ ((a, b) \# \text{arr}) \ \text{brr}); aa < a; \text{sorted-spvec } ((a, b) \# \text{arr});$

$\text{sorted-spvec } ((aa, ba) \# \text{brr}) \rrbracket \Longrightarrow \text{sorted-spvec } ((aa, y * ba) \# \text{addmult-spvec } y \ ((a, b) \# \text{arr}) \ \text{brr})$

$\langle \text{proof} \rangle$

lemma *sorted-spvec-addmult-spvec-helper2*:

$\llbracket \text{sorted-spvec } (\text{addmult-spvec } y \text{ arr } ((aa, ba) \# brr)); a < aa; \text{sorted-spvec } ((a, b) \# \text{arr}); \text{sorted-spvec } ((aa, ba) \# brr) \rrbracket$
 $\implies \text{sorted-spvec } ((a, b) \# \text{addmult-spvec } y \text{ arr } ((aa, ba) \# brr))$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-addmult-spvec-helper3[rule-format]*:

$\text{sorted-spvec } (\text{addmult-spvec } y \text{ arr } brr) \longrightarrow \text{sorted-spvec } ((aa, b) \# \text{arr}) \longrightarrow$
 $\text{sorted-spvec } ((aa, ba) \# brr)$
 $\longrightarrow \text{sorted-spvec } ((aa, b + y * ba) \# (\text{addmult-spvec } y \text{ arr } brr))$
 $\langle \text{proof} \rangle$

lemma *sorted-addmult-spvec*: $\text{sorted-spvec } a \implies \text{sorted-spvec } b \implies \text{sorted-spvec } (\text{addmult-spvec } y \text{ a } b)$

$\langle \text{proof} \rangle$

fun *mult-spvec-spmat* :: $('a::\text{lattice-ring}) \text{ spvec} \Rightarrow 'a \text{ spvec} \Rightarrow 'a \text{ smat} \Rightarrow 'a \text{ spvec}$
where

$\text{mult-spvec-spmat } c \ [] \ brr = c \mid$
 $\text{mult-spvec-spmat } c \text{ arr } [] = c \mid$
 $\text{mult-spvec-spmat } c \ ((i,a)\#\text{arr}) \ ((j,b)\#brr) = ($
 $\text{if } (i < j) \text{ then } \text{mult-spvec-spmat } c \text{ arr } ((j,b)\#brr)$
 $\text{else if } (j < i) \text{ then } \text{mult-spvec-spmat } c \ ((i,a)\#\text{arr}) \ brr$
 $\text{else } \text{mult-spvec-spmat } (\text{addmult-spvec } a \ c \ b) \text{ arr } brr)$

lemma *sparse-row-mult-spvec-spmat[rule-format]*: $\text{sorted-spvec } (a::('a::\text{lattice-ring}) \text{ spvec}) \longrightarrow \text{sorted-spvec } B \longrightarrow$

$\text{sparse-row-vector } (\text{mult-spvec-spmat } c \ a \ B) = (\text{sparse-row-vector } c) + (\text{sparse-row-vector } a) * (\text{sparse-row-matrix } B)$
 $\langle \text{proof} \rangle$

lemma *sorted-mult-spvec-spmat[rule-format]*:

$\text{sorted-spvec } (c::('a::\text{lattice-ring}) \text{ spvec}) \longrightarrow \text{sorted-spmat } B \longrightarrow \text{sorted-spvec } (\text{mult-spvec-spmat } c \ a \ B)$
 $\langle \text{proof} \rangle$

consts

$\text{mult-spmat} :: ('a::\text{lattice-ring}) \text{ smat} \Rightarrow 'a \text{ smat} \Rightarrow 'a \text{ smat}$

primrec

$\text{mult-spmat } [] \ A = []$
 $\text{mult-spmat } (a\#as) \ A = (\text{fst } a, \text{mult-spvec-spmat } [] \ (\text{snd } a) \ A)\#(\text{mult-spmat } as \ A)$

lemma *sparse-row-mult-spmat*:

$\text{sorted-spmat } A \implies \text{sorted-spvec } B \implies$

sparse-row-matrix (mult-spmat A B) = (*sparse-row-matrix* A) * (*sparse-row-matrix* B)
 ⟨proof⟩

lemma *sorted-spvec-mult-spmat*[rule-format]:
sorted-spvec (A::('a::lattice-ring) spmat) \longrightarrow *sorted-spvec* (mult-spmat A B)
 ⟨proof⟩

lemma *sorted-spmat-mult-spmat*:
sorted-spmat (B::('a::lattice-ring) spmat) \implies *sorted-spmat* (mult-spmat A B)
 ⟨proof⟩

fun *add-spvec* :: ('a::lattice-ab-group-add) spvec \Rightarrow 'a spvec \Rightarrow 'a spvec **where**

add-spvec arr [] = arr |
add-spvec [] brr = brr |
add-spvec ((i,a)#arr) ((j,b)#brr) = (
 if i < j then (i,a)#(*add-spvec* arr ((j,b)#brr))
 else if (j < i) then (j,b) # *add-spvec* ((i,a)#arr) brr
 else (i, a+b) # *add-spvec* arr brr)

lemma *add-spvec-empty1*[simp]: *add-spvec* [] a = a
 ⟨proof⟩

lemma *sparse-row-vector-add*: *sparse-row-vector* (*add-spvec* a b) = (*sparse-row-vector* a) + (*sparse-row-vector* b)
 ⟨proof⟩

fun *add-spmat* :: ('a::lattice-ab-group-add) spmat \Rightarrow 'a spmat \Rightarrow 'a spmat **where**

add-spmat [] bs = bs |
add-spmat as [] = as |
add-spmat ((i,a)#as) ((j,b)#bs) = (
 if i < j then
 (i,a) # *add-spmat* as ((j,b)#bs)
 else if j < i then
 (j,b) # *add-spmat* ((i,a)#as) bs
 else
 (i, *add-spvec* a b) # *add-spmat* as bs)

lemma *add-spmat-Nil2*[simp]: *add-spmat* as [] = as
 ⟨proof⟩

lemma *sparse-row-add-spmat*: *sparse-row-matrix* (*add-spmat* A B) = (*sparse-row-matrix* A) + (*sparse-row-matrix* B)
 ⟨proof⟩

lemmas [code] = *sparse-row-add-spmat* [symmetric]

lemmas [code] = sparse-row-vector-add [symmetric]

lemma sorted-add-spvec-helper1[rule-format]: add-spvec ((a,b)#arr) brr = (ab, bb) # list \longrightarrow (ab = a | (brr \neq [] & ab = fst (hd brr)))
 <proof>

lemma sorted-add-spmat-helper1[rule-format]: add-spmat ((a,b)#arr) brr = (ab, bb) # list \longrightarrow (ab = a | (brr \neq [] & ab = fst (hd brr)))
 <proof>

lemma sorted-add-spvec-helper: add-spvec arr brr = (ab, bb) # list \implies ((arr \neq [] & ab = fst (hd arr)) | (brr \neq [] & ab = fst (hd brr)))
 <proof>

lemma sorted-add-spmat-helper: add-spmat arr brr = (ab, bb) # list \implies ((arr \neq [] & ab = fst (hd arr)) | (brr \neq [] & ab = fst (hd brr)))
 <proof>

lemma add-spvec-commute: add-spvec a b = add-spvec b a
 <proof>

lemma add-spmat-commute: add-spmat a b = add-spmat b a
 <proof>

lemma sorted-add-spvec-helper2: add-spvec ((a,b)#arr) brr = (ab, bb) # list \implies aa < a \implies sorted-spvec ((aa, ba) # brr) \implies aa < ab
 <proof>

lemma sorted-add-spmat-helper2: add-spmat ((a,b)#arr) brr = (ab, bb) # list \implies aa < a \implies sorted-spmat ((aa, ba) # brr) \implies aa < ab
 <proof>

lemma sorted-spvec-add-spvec[rule-format]: sorted-spvec a \longrightarrow sorted-spvec b \longrightarrow sorted-spvec (add-spvec a b)
 <proof>

lemma sorted-spvec-add-spmat[rule-format]: sorted-spvec A \longrightarrow sorted-spvec B \longrightarrow sorted-spmat (add-spmat A B)
 <proof>

lemma sorted-spmat-add-spmat[rule-format]: sorted-spmat A \implies sorted-spmat B \implies sorted-spmat (add-spmat A B)
 <proof>

fun le-spvec :: ('a::lattice-ab-group-add) spvec \Rightarrow 'a spvec \Rightarrow bool **where**

le-spvec [] [] = True |
 le-spvec ((-,a)#as) [] = (a <= 0 & le-spvec as []) |
 le-spvec [] ((-,b)#bs) = (0 <= b & le-spvec [] bs) |

```

le-spvec ((i,a)#as) ((j,b)#bs) = (
  if (i < j) then a <= 0 & le-spvec as ((j,b)#bs)
  else if (j < i) then 0 <= b & le-spvec ((i,a)#as) bs
  else a <= b & le-spvec as bs)

```

fun *le-spmat* :: ('a::lattice-ab-group-add) *spmat* \Rightarrow 'a *spmat* \Rightarrow bool **where**

```

le-spmat [] [] = True |
le-spmat ((i,a)#as) [] = (le-spvec a [] & le-spmat as []) |
le-spmat [] ((j,b)#bs) = (le-spvec [] b & le-spmat [] bs) |
le-spmat ((i,a)#as) ((j,b)#bs) = (
  if i < j then (le-spvec a [] & le-spmat as ((j,b)#bs))
  else if j < i then (le-spvec [] b & le-spmat ((i,a)#as) bs)
  else (le-spvec a b & le-spmat as bs))

```

definition *disj-matrices* :: ('a::zero) *matrix* \Rightarrow 'a *matrix* \Rightarrow bool **where**

```

disj-matrices A B == (! j i. (Rep-matrix A j i  $\neq$  0)  $\longrightarrow$  (Rep-matrix B j i =
0)) & (! j i. (Rep-matrix B j i  $\neq$  0)  $\longrightarrow$  (Rep-matrix A j i = 0))

```

declare [[*simp-depth-limit* = 6]]

lemma *disj-matrices-contr1*: *disj-matrices* A B \Longrightarrow Rep-matrix A j i \neq 0 \Longrightarrow
Rep-matrix B j i = 0
<proof>

lemma *disj-matrices-contr2*: *disj-matrices* A B \Longrightarrow Rep-matrix B j i \neq 0 \Longrightarrow
Rep-matrix A j i = 0
<proof>

lemma *disj-matrices-add*: *disj-matrices* A B \Longrightarrow *disj-matrices* C D \Longrightarrow *disj-matrices*
A D \Longrightarrow *disj-matrices* B C \Longrightarrow
(A + B <= C + D) = (A <= C & B <= (D::('a::lattice-ab-group-add) matrix))
<proof>

lemma *disj-matrices-zero1*[*simp*]: *disj-matrices* 0 B
<proof>

lemma *disj-matrices-zero2*[*simp*]: *disj-matrices* A 0
<proof>

lemma *disj-matrices-commute*: *disj-matrices* A B = *disj-matrices* B A
<proof>

lemma *disj-matrices-add-le-zero*: *disj-matrices* A B \Longrightarrow
(A + B <= 0) = (A <= 0 & (B::('a::lattice-ab-group-add) matrix) <= 0)
<proof>

lemma *disj-matrices-add-zero-le*: *disj-matrices* A B \Longrightarrow

$(0 \leq A + B) = (0 \leq A \ \& \ 0 \leq (B::('a::lattice-ab-group-add) \text{ matrix}))$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-x-le*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$
 $(A \leq B + C) = (A \leq C \ \& \ 0 \leq (B::('a::lattice-ab-group-add) \text{ matrix}))$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-le-x*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$
 $(B + A \leq C) = (A \leq C \ \& \ (B::('a::lattice-ab-group-add) \text{ matrix}) \leq 0)$
 $\langle \text{proof} \rangle$

lemma *disj-sparse-row-singleton*: $i \leq j \implies \text{sorted-spvec}((j,y)\#v) \implies \text{disj-matrices}$
 $(\text{sparse-row-vector } v) (\text{singleton-matrix } 0 \ i \ x)$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-x-add*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$
 $(B+C) (A::('a::lattice-ab-group-add) \text{ matrix})$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-x*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$
 $(B+C) (A::('a::lattice-ab-group-add) \text{ matrix})$
 $\langle \text{proof} \rangle$

lemma *disj-singleton-matrices[simp]*: $\text{disj-matrices } (\text{singleton-matrix } j \ i \ x) (\text{singleton-matrix}$
 $u \ v \ y) = (j \neq u \mid i \neq v \mid x = 0 \mid y = 0)$
 $\langle \text{proof} \rangle$

lemma *disj-move-sparse-vec-mat[simplified disj-matrices-commute]*:
 $j \leq a \implies \text{sorted-spvec}((a,c)\#as) \implies \text{disj-matrices } (\text{move-matrix } (\text{sparse-row-vector}$
 $b) (\text{int } j) \ i) (\text{sparse-row-matrix } as)$
 $\langle \text{proof} \rangle$

lemma *disj-move-sparse-row-vector-twice*:
 $j \neq u \implies \text{disj-matrices } (\text{move-matrix } (\text{sparse-row-vector } a) \ j \ i) (\text{move-matrix}$
 $(\text{sparse-row-vector } b) \ u \ v)$
 $\langle \text{proof} \rangle$

lemma *le-spvec-iff-sparse-row-le[rule-format]*: $(\text{sorted-spvec } a) \longrightarrow (\text{sorted-spvec}$
 $b) \longrightarrow (\text{le-spvec } a \ b) = (\text{sparse-row-vector } a \leq \text{sparse-row-vector } b)$
 $\langle \text{proof} \rangle$

lemma *le-spvec-empty2-sparse-row[rule-format]*: $\text{sorted-spvec } b \longrightarrow \text{le-spvec } b \ [] =$
 $(\text{sparse-row-vector } b \leq 0)$
 $\langle \text{proof} \rangle$

lemma *le-spvec-empty1-sparse-row[rule-format]*: $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } []$
 $b = (0 \leq \text{sparse-row-vector } b))$
 $\langle \text{proof} \rangle$

lemma *le-spmat-iff-sparse-row-le*[*rule-format*]: (*sorted-spvec A*) \longrightarrow (*sorted-spmat A*) \longrightarrow (*sorted-spvec B*) \longrightarrow (*sorted-spmat B*) \longrightarrow
le-spmat A B = (*sparse-row-matrix A* <= *sparse-row-matrix B*)
 <proof>

declare [[*simp-depth-limit* = 999]]

primrec *abs-spmat* :: ('a::lattice-ring) *spmat* \Rightarrow 'a *spmat* **where**
abs-spmat [] = [] |
abs-spmat (a#as) = (fst a, *abs-spvec* (snd a))#(*abs-spmat* as)

primrec *minus-spmat* :: ('a::lattice-ring) *spmat* \Rightarrow 'a *spmat* **where**
minus-spmat [] = [] |
minus-spmat (a#as) = (fst a, *minus-spvec* (snd a))#(*minus-spmat* as)

lemma *sparse-row-matrix-minus*:
sparse-row-matrix (*minus-spmat A*) = - (*sparse-row-matrix A*)
 <proof>

lemma *Rep-sparse-row-vector-zero*: $x \neq 0 \implies \text{Rep-matrix } (\text{sparse-row-vector } v)$
 $x \cdot y = 0$
 <proof>

lemma *sparse-row-matrix-abs*:
sorted-spvec A \implies *sorted-spmat A* \implies *sparse-row-matrix* (*abs-spmat A*) = *abs*
 (*sparse-row-matrix A*)
 <proof>

lemma *sorted-spvec-minus-spmat*: *sorted-spvec A* \implies *sorted-spvec* (*minus-spmat A*)
 <proof>

lemma *sorted-spvec-abs-spmat*: *sorted-spvec A* \implies *sorted-spvec* (*abs-spmat A*)
 <proof>

lemma *sorted-spmat-minus-spmat*: *sorted-spmat A* \implies *sorted-spmat* (*minus-spmat A*)
 <proof>

lemma *sorted-spmat-abs-spmat*: *sorted-spmat A* \implies *sorted-spmat* (*abs-spmat A*)
 <proof>

definition *diff-spmat* :: ('a::lattice-ring) *spmat* \Rightarrow 'a *spmat* \Rightarrow 'a *spmat* **where**
diff-spmat A B == *add-spmat A* (*minus-spmat B*)

lemma *sorted-spmat-diff-spmat*: *sorted-spmat A* \implies *sorted-spmat B* \implies *sorted-spmat*
 (*diff-spmat A B*)
 <proof>

lemma *sorted-spvec-diff-spmat*: *sorted-spvec A* \implies *sorted-spvec B* \implies *sorted-spvec*
(diff-spmat A B)
 ⟨proof⟩

lemma *sparse-row-diff-spmat*: *sparse-row-matrix (diff-spmat A B)* = (*sparse-row-matrix*
A) - (*sparse-row-matrix B*)
 ⟨proof⟩

definition *sorted-sparse-matrix* :: 'a *spmat* \Rightarrow *bool* **where**
sorted-sparse-matrix A == (*sorted-spvec A*) & (*sorted-spmat A*)

lemma *sorted-sparse-matrix-imp-spvec*: *sorted-sparse-matrix A* \implies *sorted-spvec A*
 ⟨proof⟩

lemma *sorted-sparse-matrix-imp-spmat*: *sorted-sparse-matrix A* \implies *sorted-spmat*
A
 ⟨proof⟩

lemmas *sorted-sp-simps* =
sorted-spvec.simps
sorted-spmat.simps
sorted-sparse-matrix-def

lemma *bool1*: (\neg *True*) = *False* ⟨proof⟩
lemma *bool2*: (\neg *False*) = *True* ⟨proof⟩
lemma *bool3*: ((*P::bool*) \wedge *True*) = *P* ⟨proof⟩
lemma *bool4*: (*True* \wedge (*P::bool*)) = *P* ⟨proof⟩
lemma *bool5*: ((*P::bool*) \wedge *False*) = *False* ⟨proof⟩
lemma *bool6*: (*False* \wedge (*P::bool*)) = *False* ⟨proof⟩
lemma *bool7*: ((*P::bool*) \vee *True*) = *True* ⟨proof⟩
lemma *bool8*: (*True* \vee (*P::bool*)) = *True* ⟨proof⟩
lemma *bool9*: ((*P::bool*) \vee *False*) = *P* ⟨proof⟩
lemma *bool10*: (*False* \vee (*P::bool*)) = *P* ⟨proof⟩
lemmas *boolarith* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10*

lemma *if-case-eq*: (*if b then x else y*) = (*case b of True => x | False => y*)
 ⟨proof⟩

consts
pprt-spvec :: ('a::{\i{lattice-ab-group-add}}) *spvec* \Rightarrow 'a *spvec*
nprrt-spvec :: ('a::{\i{lattice-ab-group-add}}) *spvec* \Rightarrow 'a *spvec*
pprt-spmat :: ('a::{\i{lattice-ab-group-add}}) *spmat* \Rightarrow 'a *spmat*
nprrt-spmat :: ('a::{\i{lattice-ab-group-add}}) *spmat* \Rightarrow 'a *spmat*

primrec
pprt-spvec [] = []
pprt-spvec (a#as) = (fst a, pprt (snd a)) # (pprt-spvec as)

primrec

$\text{npert-spvec } [] = []$
 $\text{npert-spvec } (a \# as) = (\text{fst } a, \text{npert } (\text{snd } a)) \# (\text{npert-spvec } as)$

primrec

$\text{pprt-spmat } [] = []$
 $\text{pprt-spmat } (a \# as) = (\text{fst } a, \text{pprt-spmat } (\text{snd } a)) \# (\text{pprt-spmat } as)$

primrec

$\text{npert-spmat } [] = []$
 $\text{npert-spmat } (a \# as) = (\text{fst } a, \text{npert-spmat } (\text{snd } a)) \# (\text{npert-spmat } as)$

lemma *pprt-add: disj-matrices* $A \ (B::('a::\text{lattice-ring}) \text{ matrix}) \implies \text{pprt } (A+B) =$
 $\text{pprt } A + \text{pprt } B$
 $\langle \text{proof} \rangle$

lemma *npert-add: disj-matrices* $A \ (B::('a::\text{lattice-ring}) \text{ matrix}) \implies \text{npert } (A+B) =$
 $\text{npert } A + \text{npert } B$
 $\langle \text{proof} \rangle$

lemma *pprt-singleton[simp]:* $\text{pprt } (\text{singleton-matrix } j \ i \ (x::('a::\text{lattice-ring}))) = \text{singleton-matrix}$
 $j \ i \ (\text{pprt } x)$
 $\langle \text{proof} \rangle$

lemma *npert-singleton[simp]:* $\text{npert } (\text{singleton-matrix } j \ i \ (x::('a::\text{lattice-ring}))) = \text{singleton-matrix}$
 $j \ i \ (\text{npert } x)$
 $\langle \text{proof} \rangle$

lemma *less-imp-le:* $a < b \implies a \leq (b::('a::\text{order})) \langle \text{proof} \rangle$

lemma *sparse-row-vector-pprt:* $\text{sorted-spmat } (v::('a::\text{lattice-ring}) \text{ spvec}) \implies \text{sparse-row-vector}$
 $(\text{pprt-spmat } v) = \text{pprt } (\text{sparse-row-vector } v)$
 $\langle \text{proof} \rangle$

lemma *sparse-row-vector-npert:* $\text{sorted-spmat } (v::('a::\text{lattice-ring}) \text{ spvec}) \implies \text{sparse-row-vector}$
 $(\text{npert-spmat } v) = \text{npert } (\text{sparse-row-vector } v)$
 $\langle \text{proof} \rangle$

lemma *pprt-move-matrix:* $\text{pprt } (\text{move-matrix } (A::('a::\text{lattice-ring}) \text{ matrix}) \ j \ i) =$
 $\text{move-matrix } (\text{pprt } A) \ j \ i$
 $\langle \text{proof} \rangle$

lemma *npert-move-matrix:* $\text{npert } (\text{move-matrix } (A::('a::\text{lattice-ring}) \text{ matrix}) \ j \ i) =$
 $\text{move-matrix } (\text{npert } A) \ j \ i$
 $\langle \text{proof} \rangle$

lemma *sparse-row-matrix-pprt: sorted-spvec* ($m :: 'a::lattice-ring\ spmat$) \implies *sorted-spmat*
 $m \implies$ *sparse-row-matrix* (*pprt-spmat* m) = *pprt* (*sparse-row-matrix* m)
 ⟨*proof*⟩

lemma *sparse-row-matrix-nprt: sorted-spvec* ($m :: 'a::lattice-ring\ spmat$) \implies *sorted-spmat*
 $m \implies$ *sparse-row-matrix* (*nprrt-spmat* m) = *nprrt* (*sparse-row-matrix* m)
 ⟨*proof*⟩

lemma *sorted-pprt-spvec: sorted-spvec* $v \implies$ *sorted-spvec* (*pprt-spvec* v)
 ⟨*proof*⟩

lemma *sorted-nprt-spvec: sorted-spvec* $v \implies$ *sorted-spvec* (*nprrt-spvec* v)
 ⟨*proof*⟩

lemma *sorted-spvec-pprt-spmat: sorted-spvec* $m \implies$ *sorted-spvec* (*pprt-spmat* m)
 ⟨*proof*⟩

lemma *sorted-spvec-nprt-spmat: sorted-spvec* $m \implies$ *sorted-spvec* (*nprrt-spmat* m)
 ⟨*proof*⟩

lemma *sorted-spmat-pprt-spmat: sorted-spmat* $m \implies$ *sorted-spmat* (*pprt-spmat* m)
 ⟨*proof*⟩

lemma *sorted-spmat-nprt-spmat: sorted-spmat* $m \implies$ *sorted-spmat* (*nprrt-spmat* m)
 ⟨*proof*⟩

definition *mult-est-spmat* :: ($'a::lattice-ring$) $spmat \Rightarrow 'a\ spmat \Rightarrow 'a\ spmat \Rightarrow$
 $'a\ spmat \Rightarrow 'a\ spmat$ **where**
mult-est-spmat $r1\ r2\ s1\ s2 ==$
add-spmat (*mult-spmat* (*pprt-spmat* $s2$) (*pprt-spmat* $r2$)) (*add-spmat* (*mult-spmat*
 (*pprt-spmat* $s1$) (*nprrt-spmat* $r2$))
 (*add-spmat* (*mult-spmat* (*nprrt-spmat* $s2$) (*pprt-spmat* $r1$)) (*mult-spmat* (*nprrt-spmat*
 $s1$) (*nprrt-spmat* $r1$))))

lemmas *sparse-row-matrix-op-simps* =
sorted-sparse-matrix-imp-spmat sorted-sparse-matrix-imp-spvec
sparse-row-add-spmat sorted-spvec-add-spmat sorted-spmat-add-spmat
sparse-row-diff-spmat sorted-spvec-diff-spmat sorted-spmat-diff-spmat
sparse-row-matrix-minus sorted-spvec-minus-spmat sorted-spmat-minus-spmat
sparse-row-mult-spmat sorted-spvec-mult-spmat sorted-spmat-mult-spmat
sparse-row-matrix-abs sorted-spvec-abs-spmat sorted-spmat-abs-spmat
le-spmat-iff-sparse-row-le
sparse-row-matrix-pprt sorted-spvec-pprt-spmat sorted-spmat-pprt-spmat
sparse-row-matrix-nprt sorted-spvec-nprt-spmat sorted-spmat-nprt-spmat

lemma *zero-eq-Numeral0*: ($0::number-ring$) = *N numeral0* ⟨*proof*⟩

```

lemmas sparse-row-matrix-arith-simps [simplified zero-eq-Numeral0] =
  mult-spmat.simps mult-spvec-spmat.simps
  addmult-spvec.simps
  smult-spvec-empty smult-spvec-cons
  add-spmat.simps add-spvec.simps
  minus-spmat.simps minus-spvec.simps
  abs-spmat.simps abs-spvec.simps
  diff-spmat-def
  le-spmat.simps le-spvec.simps
  pprt-spmat.simps pprt-spvec.simps
  nprr-spmat.simps nprr-spvec.simps
  mult-est-spmat-def

```

end

```

theory LP
imports Main Lattice-Algebras
begin

```

```

lemma linprog-dual-estimate:
  assumes
     $A * x \leq (b :: 'a :: lattice-ring)$ 
     $0 \leq y$ 
     $abs (A - A') \leq \delta A$ 
     $b \leq b'$ 
     $abs (c - c') \leq \delta c$ 
     $abs x \leq r$ 
  shows
     $c * x \leq y * b' + (y * \delta A + abs (y * A' - c') + \delta c) * r$ 
  <proof>

```

```

lemma le-ge-imp-abs-diff-1:
  assumes
     $A1 \leq (A :: 'a :: lattice-ring)$ 
     $A \leq A2$ 
  shows  $abs (A - A1) \leq A2 - A1$ 
  <proof>

```

```

lemma mult-le-prts:
  assumes
     $a1 \leq (a :: 'a :: lattice-ring)$ 
     $a \leq a2$ 
     $b1 \leq b$ 
     $b \leq b2$ 

```



```

shows
  a * b <= pprr a2 * pprr b2 + pprr a1 * nprr b2 + nprr a2 * pprr b1 + nprr a1
  * nprr b1
<proof>

lemma mult-le-dual-prts:
assumes
  A * x ≤ (b::'a::lattice-ring)
  0 ≤ y
  A1 ≤ A
  A ≤ A2
  c1 ≤ c
  c ≤ c2
  r1 ≤ x
  x ≤ r2
shows
  c * x ≤ y * b + (let s1 = c1 - y * A2; s2 = c2 - y * A1 in pprr s2 * pprr r2
  + pprr s1 * nprr r2 + nprr s2 * pprr r1 + nprr s1 * nprr r1)
  (is - <= - + ?C)
<proof>

end

```

2 Floating Point Representation of the Reals

```

theory ComputeFloat
imports Complex-Main Lattice-Algebras
uses ~~/src/Tools/float.ML (~~/src/HOL/Tools/float-arith.ML)
begin

```

```

definition
  pow2 :: int ⇒ real where
  pow2 a = (if (0 ≤ a) then (2^(nat a)) else (inverse (2^(nat (-a)))))

```

```

definition
  float :: int * int ⇒ real where
  float x = real (fst x) * pow2 (snd x)

```

```

lemma pow2-0[simp]: pow2 0 = 1
<proof>

```

```

lemma pow2-1[simp]: pow2 1 = 2
<proof>

```

```

lemma pow2-neg: pow2 x = inverse (pow2 (-x))
<proof>

```

```

lemma pow2-add1: pow2 (1 + a) = 2 * (pow2 a)
<proof>

```

lemma *pow2-add*: $\text{pow2 } (a+b) = (\text{pow2 } a) * (\text{pow2 } b)$
 $\langle \text{proof} \rangle$

lemma *float* $(a, e) + \text{float } (b, e) = \text{float } (a + b, e)$
 $\langle \text{proof} \rangle$

definition

int-of-real :: $\text{real} \Rightarrow \text{int}$ **where**
int-of-real $x = (\text{SOME } y. \text{real } y = x)$

definition

real-is-int :: $\text{real} \Rightarrow \text{bool}$ **where**
real-is-int $x = (\text{EX } (u::\text{int}). x = \text{real } u)$

lemma *real-is-int-def2*: $\text{real-is-int } x = (x = \text{real } (\text{int-of-real } x))$
 $\langle \text{proof} \rangle$

lemma *float-transfer*: $\text{real-is-int } ((\text{real } a) * (\text{pow2 } c)) \Longrightarrow \text{float } (a, b) = \text{float } (\text{int-of-real } ((\text{real } a) * (\text{pow2 } c)), b - c)$
 $\langle \text{proof} \rangle$

lemma *pow2-int*: $\text{pow2 } (\text{int } c) = 2^c$
 $\langle \text{proof} \rangle$

lemma *float-transfer-nat*: $\text{float } (a, b) = \text{float } (a * 2^c, b - \text{int } c)$
 $\langle \text{proof} \rangle$

lemma *real-is-int-real[simp]*: $\text{real-is-int } (\text{real } (x::\text{int}))$
 $\langle \text{proof} \rangle$

lemma *int-of-real-real[simp]*: $\text{int-of-real } (\text{real } x) = x$
 $\langle \text{proof} \rangle$

lemma *real-int-of-real[simp]*: $\text{real-is-int } x \Longrightarrow \text{real } (\text{int-of-real } x) = x$
 $\langle \text{proof} \rangle$

lemma *real-is-int-add-int-of-real*: $\text{real-is-int } a \Longrightarrow \text{real-is-int } b \Longrightarrow (\text{int-of-real } (a+b)) = (\text{int-of-real } a) + (\text{int-of-real } b)$
 $\langle \text{proof} \rangle$

lemma *real-is-int-add[simp]*: $\text{real-is-int } a \Longrightarrow \text{real-is-int } b \Longrightarrow \text{real-is-int } (a+b)$
 $\langle \text{proof} \rangle$

lemma *int-of-real-sub*: $\text{real-is-int } a \Longrightarrow \text{real-is-int } b \Longrightarrow (\text{int-of-real } (a-b)) = (\text{int-of-real } a) - (\text{int-of-real } b)$
 $\langle \text{proof} \rangle$

lemma *real-is-int-sub[simp]*: $\text{real-is-int } a \Longrightarrow \text{real-is-int } b \Longrightarrow \text{real-is-int } (a-b)$

$\langle \text{proof} \rangle$

lemma *real-is-int-rep*: *real-is-int* $x \implies ?! (a::\text{int}). \text{real } a = x$
 $\langle \text{proof} \rangle$

lemma *int-of-real-mult*:
 assumes *real-is-int* a *real-is-int* b
 shows (*int-of-real* ($a*b$)) = (*int-of-real* a) * (*int-of-real* b)
 $\langle \text{proof} \rangle$

lemma *real-is-int-mult[simp]*: *real-is-int* $a \implies \text{real-is-int } b \implies \text{real-is-int } (a*b)$
 $\langle \text{proof} \rangle$

lemma *real-is-int-0[simp]*: *real-is-int* ($0::\text{real}$)
 $\langle \text{proof} \rangle$

lemma *real-is-int-1[simp]*: *real-is-int* ($1::\text{real}$)
 $\langle \text{proof} \rangle$

lemma *real-is-int-n1*: *real-is-int* ($-1::\text{real}$)
 $\langle \text{proof} \rangle$

lemma *real-is-int-number-of[simp]*: *real-is-int* ((*number-of* :: $\text{int} \Rightarrow \text{real}$) x)
 $\langle \text{proof} \rangle$

lemma *int-of-real-0[simp]*: *int-of-real* ($0::\text{real}$) = ($0::\text{int}$)
 $\langle \text{proof} \rangle$

lemma *int-of-real-1[simp]*: *int-of-real* ($1::\text{real}$) = ($1::\text{int}$)
 $\langle \text{proof} \rangle$

lemma *int-of-real-number-of[simp]*: *int-of-real* (*number-of* b) = *number-of* b
 $\langle \text{proof} \rangle$

lemma *float-transfer-even*: *even* $a \implies \text{float } (a, b) = \text{float } (a \text{ div } 2, b+1)$
 $\langle \text{proof} \rangle$

lemma *int-div-zdiv*: *int* ($a \text{ div } b$) = (*int* a) *div* (*int* b)
 $\langle \text{proof} \rangle$

lemma *int-mod-zmod*: *int* ($a \text{ mod } b$) = (*int* a) *mod* (*int* b)
 $\langle \text{proof} \rangle$

lemma *abs-div-2-less*: $a \neq 0 \implies a \neq -1 \implies \text{abs}((a::\text{int}) \text{ div } 2) < \text{abs } a$
 $\langle \text{proof} \rangle$

function *norm-float* :: $\text{int} \Rightarrow \text{int} \Rightarrow \text{int} \times \text{int}$ **where**
 norm-float a b = (if $a \neq 0 \wedge \text{even } a$ then *norm-float* ($a \text{ div } 2$) ($b + 1$)
 else if $a = 0$ then ($0, 0$) else (a, b))

$\langle proof \rangle$

termination $\langle proof \rangle$

lemma *norm-float*: $float\ x = float\ (split\ norm\text{-}float\ x)$
 $\langle proof \rangle$

lemma *float-add-l0*: $float\ (0, e) + x = x$
 $\langle proof \rangle$

lemma *float-add-r0*: $x + float\ (0, e) = x$
 $\langle proof \rangle$

lemma *float-add*:
 $float\ (a1, e1) + float\ (a2, e2) =$
 $(if\ e1 \leq e2\ then\ float\ (a1 + a2 * 2^{nat(e2-e1)}, e1)$
 $else\ float\ (a1 * 2^{nat(e1-e2)} + a2, e2))$
 $\langle proof \rangle$

lemma *float-add-assoc1*:
 $(x + float\ (y1, e1)) + float\ (y2, e2) = (float\ (y1, e1) + float\ (y2, e2)) + x$
 $\langle proof \rangle$

lemma *float-add-assoc2*:
 $(float\ (y1, e1) + x) + float\ (y2, e2) = (float\ (y1, e1) + float\ (y2, e2)) + x$
 $\langle proof \rangle$

lemma *float-add-assoc3*:
 $float\ (y1, e1) + (x + float\ (y2, e2)) = (float\ (y1, e1) + float\ (y2, e2)) + x$
 $\langle proof \rangle$

lemma *float-add-assoc4*:
 $float\ (y1, e1) + (float\ (y2, e2) + x) = (float\ (y1, e1) + float\ (y2, e2)) + x$
 $\langle proof \rangle$

lemma *float-mult-l0*: $float\ (0, e) * x = float\ (0, 0)$
 $\langle proof \rangle$

lemma *float-mult-r0*: $x * float\ (0, e) = float\ (0, 0)$
 $\langle proof \rangle$

definition
 $lbound :: real \Rightarrow real$
where
 $lbound\ x = min\ 0\ x$

definition
 $ubound :: real \Rightarrow real$
where

$ubound\ x = max\ 0\ x$

lemma *lbound*: $lbound\ x \leq x$
<proof>

lemma *ubound*: $x \leq ubound\ x$
<proof>

lemma *float-mult*:
 $float\ (a1,\ e1) * float\ (a2,\ e2) =$
 $(float\ (a1 * a2,\ e1 + e2))$
<proof>

lemma *float-minus*:
 $-(float\ (a,b)) = float\ (-a,\ b)$
<proof>

lemma *zero-less-pow2*:
 $0 < pow2\ x$
<proof>

lemma *zero-le-float*:
 $(0 \leq float\ (a,b)) = (0 \leq a)$
<proof>

lemma *float-le-zero*:
 $(float\ (a,b) \leq 0) = (a \leq 0)$
<proof>

lemma *float-abs*:
 $abs\ (float\ (a,b)) = (if\ 0 \leq a\ then\ (float\ (a,b))\ else\ (float\ (-a,b)))$
<proof>

lemma *float-zero*:
 $float\ (0,\ b) = 0$
<proof>

lemma *float-pprt*:
 $pprt\ (float\ (a,\ b)) = (if\ 0 \leq a\ then\ (float\ (a,b))\ else\ (float\ (0,\ b)))$
<proof>

lemma *pprt-lbound*: $pprt\ (lbound\ x) = float\ (0,\ 0)$
<proof>

lemma *nprrt-ubound*: $nprrt\ (ubound\ x) = float\ (0,\ 0)$
<proof>

lemma *float-nprtt*:
 $nprrt\ (float\ (a,\ b)) = (if\ 0 \leq a\ then\ (float\ (0,b))\ else\ (float\ (a,\ b)))$

$\langle proof \rangle$

lemma *norm-0-1*: $(0::\text{number-ring}) = \text{Numeral0} \ \& \ (1::\text{number-ring}) = \text{Numeral1}$
 $\langle proof \rangle$

lemma *add-left-zero*: $0 + a = (a::'a::\text{comm-monoid-add})$
 $\langle proof \rangle$

lemma *add-right-zero*: $a + 0 = (a::'a::\text{comm-monoid-add})$
 $\langle proof \rangle$

lemma *mult-left-one*: $1 * a = (a::'a::\text{semiring-1})$
 $\langle proof \rangle$

lemma *mult-right-one*: $a * 1 = (a::'a::\text{semiring-1})$
 $\langle proof \rangle$

lemma *int-pow-0*: $(a::\text{int})^{\text{Numeral0}} = 1$
 $\langle proof \rangle$

lemma *int-pow-1*: $(a::\text{int})^{\text{Numeral1}} = a$
 $\langle proof \rangle$

lemma *zero-eq-Numeral0-nring*: $(0::'a::\text{number-ring}) = \text{Numeral0}$
 $\langle proof \rangle$

lemma *one-eq-Numeral1-nring*: $(1::'a::\text{number-ring}) = \text{Numeral1}$
 $\langle proof \rangle$

lemma *zero-eq-Numeral0-nat*: $(0::\text{nat}) = \text{Numeral0}$
 $\langle proof \rangle$

lemma *one-eq-Numeral1-nat*: $(1::\text{nat}) = \text{Numeral1}$
 $\langle proof \rangle$

lemma *zpower-Pls*: $(z::\text{int})^{\text{Numeral0}} = \text{Numeral1}$
 $\langle proof \rangle$

lemma *zpower-Min*: $(z::\text{int})^{((-1)::\text{nat})} = \text{Numeral1}$
 $\langle proof \rangle$

lemma *fst-cong*: $a=a' \implies \text{fst } (a,b) = \text{fst } (a',b)$
 $\langle proof \rangle$

lemma *snd-cong*: $b=b' \implies \text{snd } (a,b) = \text{snd } (a,b')$
 $\langle proof \rangle$

lemma *lift-bool*: $x \implies x = \text{True}$
 $\langle proof \rangle$

lemma *nlift-bool*: $\sim x \implies x = \text{False}$
 $\langle \text{proof} \rangle$

lemma *not-false-eq-true*: $(\sim \text{False}) = \text{True}$ $\langle \text{proof} \rangle$

lemma *not-true-eq-false*: $(\sim \text{True}) = \text{False}$ $\langle \text{proof} \rangle$

lemmas *binarith* =
normalize-bin-simps
pred-bin-simps succ-bin-simps
add-bin-simps minus-bin-simps mult-bin-simps

lemma *int-eq-number-of-eq*:
 $((\text{number-of } v)::\text{int}) = (\text{number-of } w) \iff \text{iszero } ((\text{number-of } (v + \text{uminus } w))::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-iszero-number-of-Pls*: $\text{iszero } (\text{Numeral0}::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-nonzero-number-of-Min*: $\sim(\text{iszero } ((-1)::\text{int}))$
 $\langle \text{proof} \rangle$

lemma *int-iszero-number-of-Bit0*: $\text{iszero } ((\text{number-of } (\text{Int.Bit0 } w))::\text{int}) \iff \text{iszero } ((\text{number-of } w)::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-iszero-number-of-Bit1*: $\neg \text{iszero } ((\text{number-of } (\text{Int.Bit1 } w))::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-less-number-of-eq-neg*: $((\text{number-of } x)::\text{int}) < \text{number-of } y \iff \text{neg } ((\text{number-of } (x + (\text{uminus } y)))::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-not-neg-number-of-Pls*: $\neg (\text{neg } (\text{Numeral0}::\text{int}))$
 $\langle \text{proof} \rangle$

lemma *int-neg-number-of-Min*: $\text{neg } (-1::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-neg-number-of-Bit0*: $\text{neg } ((\text{number-of } (\text{Int.Bit0 } w))::\text{int}) = \text{neg } ((\text{number-of } w)::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-neg-number-of-Bit1*: $\text{neg } ((\text{number-of } (\text{Int.Bit1 } w))::\text{int}) = \text{neg } ((\text{number-of } w)::\text{int})$
 $\langle \text{proof} \rangle$

lemma *int-le-number-of-eq*: $((\text{number-of } x)::\text{int}) \leq \text{number-of } y \iff (\neg \text{neg } ((\text{number-of } (x - \text{number-of } y))::\text{int}))$

$(y + (\text{uminus } x))::\text{int})$
 $\langle \text{proof} \rangle$

lemmas *intarithrel* =
int-eq-number-of-eq
lift-bool[OF int-iszero-number-of-Pls] nlift-bool[OF int-nonzero-number-of-Min]
int-iszero-number-of-Bit0
lift-bool[OF int-iszero-number-of-Bit1] int-less-number-of-eq-neg nlift-bool[OF int-not-neg-number-of-Pls]
lift-bool[OF int-neg-number-of-Min]
int-neg-number-of-Bit0 int-neg-number-of-Bit1 int-le-number-of-eq

lemma *int-number-of-add-sym*: $((\text{number-of } v)::\text{int}) + \text{number-of } w = \text{number-of}$
 $(v + w)$
 $\langle \text{proof} \rangle$

lemma *int-number-of-diff-sym*: $((\text{number-of } v)::\text{int}) - \text{number-of } w = \text{number-of}$
 $(v + (\text{uminus } w))$
 $\langle \text{proof} \rangle$

lemma *int-number-of-mult-sym*: $((\text{number-of } v)::\text{int}) * \text{number-of } w = \text{number-of}$
 $(v * w)$
 $\langle \text{proof} \rangle$

lemma *int-number-of-minus-sym*: $-(\text{number-of } v)::\text{int} = \text{number-of } (\text{uminus } v)$
 $\langle \text{proof} \rangle$

lemmas *intarith* = *int-number-of-add-sym int-number-of-minus-sym int-number-of-diff-sym*
int-number-of-mult-sym

lemmas *natarith* = *add-nat-number-of diff-nat-number-of mult-nat-number-of eq-nat-number-of*
less-nat-number-of

lemmas *powerarith* = *nat-number-of zpower-number-of-even*
zpower-number-of-odd[simplified zero-eq-Numeral0-nring one-eq-Numeral1-nring]
zpower-Pls zpower-Min

lemmas *floatarith[simplified norm-0-1]* = *float-add float-add-l0 float-add-r0 float-mult*
float-mult-l0 float-mult-r0
float-minus float-abs zero-le-float float-pprt float-nprt pprrt-lbound nprrt-ubound

lemmas *arith* = *binarith intarith intarithrel natarith powerarith floatarith not-false-eq-true*
not-true-eq-false

$\langle ML \rangle$

end


```

theory Compute-Oracle imports Pure
uses am.ML am-compiler.ML am-interpreter.ML am-ghc.ML am-sml.ML report.ML
compute.ML linker.ML
begin

end
theory ComputeHOL
imports Complex-Main ~~/src/Tools/Compute-Oracle/Compute-Oracle
begin

lemma Trueprop-eq-eq: Trueprop  $X == (X == \text{True})$   $\langle \text{proof} \rangle$ 
lemma meta-eq-trivial:  $x == y \implies x == y$   $\langle \text{proof} \rangle$ 
lemma meta-eq-imp-eq:  $x == y \implies x = y$   $\langle \text{proof} \rangle$ 
lemma eq-trivial:  $x = y \implies x = y$   $\langle \text{proof} \rangle$ 
lemma bool-to-true:  $x :: \text{bool} \implies x == \text{True}$   $\langle \text{proof} \rangle$ 
lemma transmeta-1:  $x = y \implies y == z \implies x = z$   $\langle \text{proof} \rangle$ 
lemma transmeta-2:  $x == y \implies y = z \implies x = z$   $\langle \text{proof} \rangle$ 
lemma transmeta-3:  $x == y \implies y == z \implies x = z$   $\langle \text{proof} \rangle$ 


lemma If-True:  $\text{If } \text{True} = (\lambda x y. x)$   $\langle \text{proof} \rangle$ 
lemma If-False:  $\text{If } \text{False} = (\lambda x y. y)$   $\langle \text{proof} \rangle$ 

lemmas compute-if = If-True If-False


lemma bool1:  $(\neg \text{True}) = \text{False}$   $\langle \text{proof} \rangle$ 
lemma bool2:  $(\neg \text{False}) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool3:  $(P \wedge \text{True}) = P$   $\langle \text{proof} \rangle$ 
lemma bool4:  $(\text{True} \wedge P) = P$   $\langle \text{proof} \rangle$ 
lemma bool5:  $(P \wedge \text{False}) = \text{False}$   $\langle \text{proof} \rangle$ 
lemma bool6:  $(\text{False} \wedge P) = \text{False}$   $\langle \text{proof} \rangle$ 
lemma bool7:  $(P \vee \text{True}) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool8:  $(\text{True} \vee P) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool9:  $(P \vee \text{False}) = P$   $\langle \text{proof} \rangle$ 
lemma bool10:  $(\text{False} \vee P) = P$   $\langle \text{proof} \rangle$ 
lemma bool11:  $(\text{True} \longrightarrow P) = P$   $\langle \text{proof} \rangle$ 
lemma bool12:  $(P \longrightarrow \text{True}) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool13:  $(\text{True} \longrightarrow P) = P$   $\langle \text{proof} \rangle$ 
lemma bool14:  $(P \longrightarrow \text{False}) = (\neg P)$   $\langle \text{proof} \rangle$ 
lemma bool15:  $(\text{False} \longrightarrow P) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool16:  $(\text{False} = \text{False}) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool17:  $(\text{True} = \text{True}) = \text{True}$   $\langle \text{proof} \rangle$ 
lemma bool18:  $(\text{False} = \text{True}) = \text{False}$   $\langle \text{proof} \rangle$ 
lemma bool19:  $(\text{True} = \text{False}) = \text{False}$   $\langle \text{proof} \rangle$ 

```

lemmas *compute-bool* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10 bool11 bool12 bool13 bool14 bool15 bool16 bool17 bool18 bool19*

lemma *compute-fst*: $\text{fst } (x,y) = x$ *<proof>*
lemma *compute-snd*: $\text{snd } (x,y) = y$ *<proof>*
lemma *compute-pair-eq*: $((a, b) = (c, d)) = (a = c \wedge b = d)$ *<proof>*

lemma *prod-case-simp*: $\text{prod-case } f \ (x,y) = f \ x \ y$ *<proof>*

lemmas *compute-pair* = *compute-fst compute-snd compute-pair-eq prod-case-simp*

lemma *compute-the*: $\text{the } (\text{Some } x) = x$ *<proof>*
lemma *compute-None-Some-eq*: $(\text{None} = \text{Some } x) = \text{False}$ *<proof>*
lemma *compute-Some-None-eq*: $(\text{Some } x = \text{None}) = \text{False}$ *<proof>*
lemma *compute-None-None-eq*: $(\text{None} = \text{None}) = \text{True}$ *<proof>*
lemma *compute-Some-Some-eq*: $(\text{Some } x = \text{Some } y) = (x = y)$ *<proof>*

definition

option-case-compute :: $'b \ \text{option} \Rightarrow 'a \Rightarrow ('b \Rightarrow 'a) \Rightarrow 'a$

where

option-case-compute *opt a f* = *option-case a f opt*

lemma *option-case-compute*: $\text{option-case} = (\lambda \ a \ f \ \text{opt}. \ \text{option-case-compute } \text{opt } a \ f)$
<proof>

lemma *option-case-compute-None*: $\text{option-case-compute } \text{None} = (\lambda \ a \ f. \ a)$
<proof>

lemma *option-case-compute-Some*: $\text{option-case-compute } (\text{Some } x) = (\lambda \ a \ f. \ f \ x)$
<proof>

lemmas *compute-option-case* = *option-case-compute option-case-compute-None option-case-compute-Some*

lemmas *compute-option* = *compute-the compute-None-Some-eq compute-Some-None-eq compute-None-None-eq compute-Some-Some-eq compute-option-case*

lemma *length-cons*: $\text{length } (x \# xs) = 1 + (\text{length } xs)$
<proof>

lemma *length-nil*: $\text{length } [] = 0$

$\langle proof \rangle$

lemmas *compute-list-length* = *length-nil length-cons*

definition

list-case-compute :: 'b list \Rightarrow 'a \Rightarrow ('b \Rightarrow 'b list \Rightarrow 'a) \Rightarrow 'a

where

list-case-compute l a f = *list-case* a f l

lemma *list-case-compute*: *list-case* = (λ (a::'a) f (l::'b list). *list-case-compute* l a f)

$\langle proof \rangle$

lemma *list-case-compute-empty*: *list-case-compute* ([]::'b list) = (λ (a::'a) f. a)

$\langle proof \rangle$

lemma *list-case-compute-cons*: *list-case-compute* (u#v) = (λ (a::'a) f. (f (u::'b) v))

$\langle proof \rangle$

lemmas *compute-list-case* = *list-case-compute list-case-compute-empty list-case-compute-cons*

lemma *compute-list-nth*: ((x#xs) ! n) = (if n = 0 then x else (xs ! (n - 1)))

$\langle proof \rangle$

lemmas *compute-list* = *compute-list-case compute-list-length compute-list-nth*

lemmas *compute-let* = *Let-def*

lemmas *compute-hol* = *compute-if compute-bool compute-pair compute-option compute-list compute-let*

$\langle ML \rangle$

end

```

theory ComputeNumeral
imports ComputeHOL ComputeFloat
begin

```

```

lemmas bitnorm = normalize-bin-simps

```

```

lemma neg1: neg Int.Pls = False <proof>
lemma neg2: neg Int.Min = True <proof>
lemma neg3: neg (Int.Bit0 x) = neg x <proof>
lemma neg4: neg (Int.Bit1 x) = neg x <proof>
lemmas bitneg = neg1 neg2 neg3 neg4

```

```

lemma iszero1: iszero Int.Pls = True <proof>
lemma iszero2: iszero Int.Min = False <proof>
lemma iszero3: iszero (Int.Bit0 x) = iszero x <proof>
lemma iszero4: iszero (Int.Bit1 x) = False <proof>
lemmas bitiszero = iszero1 iszero2 iszero3 iszero4

```

definition

```

  lezero x == (x ≤ 0)
lemma lezero1: lezero Int.Pls = True <proof>
lemma lezero2: lezero Int.Min = True <proof>
lemma lezero3: lezero (Int.Bit0 x) = lezero x <proof>
lemma lezero4: lezero (Int.Bit1 x) = neg x <proof>
lemmas bitlezero = lezero1 lezero2 lezero3 lezero4

```

```

lemmas biteq = eq-bin-simps

```

```

lemmas bitless = less-bin-simps

```

```

lemmas bitle = le-bin-simps

```

```

lemmas bitsucc = succ-bin-simps

```

```

lemmas bitpred = pred-bin-simps

```

```

lemmas bituminus = minus-bin-simps

```

lemmas *bitadd = add-bin-simps*

lemma *mult-Pls-right*: $x * \text{Int.Pl} = \text{Int.Pl}$ *<proof>*

lemma *mult-Min-right*: $x * \text{Int.Min} = - x$ *<proof>*

lemma *multb0x*: $(\text{Int.Bit0 } x) * y = \text{Int.Bit0 } (x * y)$ *<proof>*

lemma *multxb0*: $x * (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x * y)$ *<proof>*

lemma *multb1*: $(\text{Int.Bit1 } x) * (\text{Int.Bit1 } y) = \text{Int.Bit1 } (\text{Int.Bit0 } (x * y) + x + y)$ *<proof>*

lemmas *bitmul = mult-Pls mult-Min mult-Pls-right mult-Min-right multb0x multxb0 multb1*

lemmas *bitarith = bitnorm bitiszero bitneg bitlezero biteq bitless bitle bitsucc bitpred bituminus bitadd bitmul*

definition

nat-norm-number-of $(x::\text{nat}) == x$

lemma *nat-norm-number-of*: *nat-norm-number-of* $(\text{number-of } w) = (\text{if } \text{lezero } w \text{ then } 0 \text{ else } \text{number-of } w)$ *<proof>*

lemma *natnorm0*: $(0::\text{nat}) = \text{number-of } (\text{Int.Pl})$ *<proof>*

lemma *natnorm1*: $(1::\text{nat}) = \text{number-of } (\text{Int.Bit1 } \text{Int.Pl})$ *<proof>*

lemmas *natnorm = natnorm0 natnorm1 nat-norm-number-of*

lemma *natsuc*: *Suc* $(\text{number-of } x) = (\text{if } \text{neg } x \text{ then } 1 \text{ else } \text{number-of } (\text{Int.succ } x))$ *<proof>*

lemma *natadd*: $\text{number-of } x + ((\text{number-of } y)::\text{nat}) = (\text{if } \text{neg } x \text{ then } (\text{number-of } y) \text{ else } (\text{if } \text{neg } y \text{ then } \text{number-of } x \text{ else } (\text{number-of } (x + y))))$ *<proof>*

lemma *natsub*: $(\text{number-of } x) - ((\text{number-of } y)::\text{nat}) = (\text{if } \text{neg } x \text{ then } 0 \text{ else } (\text{if } \text{neg } y \text{ then } \text{number-of } x \text{ else } (\text{nat-norm-number-of } (\text{number-of } (x + (- y))))))$ *<proof>*

lemma *natmul*: $(\text{number-of } x) * ((\text{number-of } y)::\text{nat}) = (\text{if } \text{neg } x \text{ then } 0 \text{ else } (\text{if } \text{neg } y \text{ then } 0 \text{ else } \text{number-of } (x * y)))$ *<proof>*

lemma *nateq*: $((\text{number-of } x)::\text{nat}) = (\text{number-of } y) = ((\text{lezero } x \wedge \text{lezero } y) \vee (x = y))$

<proof>

lemma *natless*: (((*number-of* *x*)::nat) < (*number-of* *y*)) = ((*x* < *y*) ∧ (¬ (*lezero* *y*)))
<proof>

lemma *natle*: (((*number-of* *x*)::nat) ≤ (*number-of* *y*)) = (*y* < *x* → *lezero* *x*)
<proof>

fun *natfac* :: nat ⇒ nat

where

natfac *n* = (if *n* = 0 then 1 else *n* * (*natfac* (*n* - 1)))

lemmas *compute-natarith* = *bitarith natnorm natsuc natadd natsub natmul nateq*
natless natle natfac.simps

lemma *number-eq*: (((*number-of* *x*)::'a::{*number-ring*, *linordered-idom*}) = (*number-of* *y*)) = (*x* = *y*)
<proof>

lemma *number-le*: (((*number-of* *x*)::'a::{*number-ring*, *linordered-idom*}) ≤ (*number-of* *y*)) = (*x* ≤ *y*)
<proof>

lemma *number-less*: (((*number-of* *x*)::'a::{*number-ring*, *linordered-idom*}) < (*number-of* *y*)) = (*x* < *y*)
<proof>

lemma *number-diff*: ((*number-of* *x*)::'a::{*number-ring*, *linordered-idom*}) - *number-of* *y* = *number-of* (*x* + (- *y*))
<proof>

lemmas *number-norm* = *number-of-Pls[symmetric] numeral-1-eq-1[symmetric]*

lemmas *compute-numberarith* = *number-of-minus[symmetric] number-of-add[symmetric]*
number-diff number-of-mult[symmetric] number-norm number-eq number-le number-less

lemma *compute-real-of-nat-number-of*: real ((*number-of* *v*)::nat) = (if *neg* *v* then 0 else *number-of* *v*)
<proof>

lemma *compute-nat-of-int-number-of*: nat ((*number-of* *v*)::int) = (*number-of* *v*)
<proof>

lemmas *compute-num-conversions* = *compute-real-of-nat-number-of compute-nat-of-int-number-of*
real-number-of

lemmas *zpowerarith* = *zpower-number-of-even*
zpower-number-of-odd[simplified zero-eq-Numeral0-nring one-eq-Numeral1-nring]

zpower-Pls zpower-Min

lemma *adjust*: $\text{adjust } b \ (q, r) = (\text{if } 0 \leq r - b \text{ then } (2 * q + 1, r - b) \text{ else } (2 * q, r))$
 ⟨proof⟩

lemma *negateSnd*: $\text{negateSnd } (q, r) = (q, -r)$
 ⟨proof⟩

lemma *divmod*: $\text{divmod-int } a \ b = (\text{if } 0 \leq a \text{ then}$
 $\text{if } 0 \leq b \text{ then } \text{posDivAlg } a \ b$
 $\text{else if } a = 0 \text{ then } (0, 0)$
 $\text{else } \text{negateSnd } (\text{negDivAlg } (-a) (-b))$
 else
 $\text{if } 0 < b \text{ then } \text{negDivAlg } a \ b$
 $\text{else } \text{negateSnd } (\text{posDivAlg } (-a) (-b)))$
 ⟨proof⟩

lemmas *compute-div-mod* = *div-int-def mod-int-def divmod adjust negateSnd pos-DivAlg.simps negDivAlg.simps*

lemma *even-Pls*: $\text{even } (\text{Int.Pl}) = \text{True}$
 ⟨proof⟩

lemma *even-Min*: $\text{even } (\text{Int.Min}) = \text{False}$
 ⟨proof⟩

lemma *even-B0*: $\text{even } (\text{Int.Bit0 } x) = \text{True}$
 ⟨proof⟩

lemma *even-B1*: $\text{even } (\text{Int.Bit1 } x) = \text{False}$
 ⟨proof⟩

lemma *even-number-of*: $\text{even } ((\text{number-of } w)::\text{int}) = \text{even } w$
 ⟨proof⟩

lemmas *compute-even* = *even-Pls even-Min even-B0 even-B1 even-number-of*

lemmas *compute-numeral* = *compute-if compute-let compute-pair compute-bool*
 compute-natarith compute-numberarith max-def min-def
compute-num-conversions zpowerarith compute-div-mod compute-even

end

```

theory Cplex
imports SparseMatrix LP ComputeFloat ComputeNumeral
uses Cplex-tools.ML CplexMatrixConverter.ML FloatSparseMatrixBuilder.ML
      fspmlp.ML (matrixlp.ML)
begin

lemma spm-mult-le-dual-prts:
  assumes
    sorted-sparse-matrix A1
    sorted-sparse-matrix A2
    sorted-sparse-matrix c1
    sorted-sparse-matrix c2
    sorted-sparse-matrix y
    sorted-sparse-matrix r1
    sorted-sparse-matrix r2
    sorted-spvec b
    le-spmat [] y
    sparse-row-matrix A1  $\leq$  A
    A  $\leq$  sparse-row-matrix A2
    sparse-row-matrix c1  $\leq$  c
    c  $\leq$  sparse-row-matrix c2
    sparse-row-matrix r1  $\leq$  x
    x  $\leq$  sparse-row-matrix r2
    A * x  $\leq$  sparse-row-matrix (b::('a::lattice-ring) spmat)
  shows
    c * x  $\leq$  sparse-row-matrix (add-spmat (mult-spmat y b)
      (let s1 = diff-spmat c1 (mult-spmat y A2); s2 = diff-spmat c2 (mult-spmat y
A1) in
        add-spmat (mult-spmat (pprt-spmat s2) (pprt-spmat r2)) (add-spmat (mult-spmat
(pprt-spmat s1) (nprrt-spmat r2))
          (add-spmat (mult-spmat (nprrt-spmat s2) (pprt-spmat r1)) (mult-spmat (nprrt-spmat
s1) (nprrt-spmat r1))))))
    <proof>

lemma spm-mult-le-dual-prts-no-let:
  assumes
    sorted-sparse-matrix A1
    sorted-sparse-matrix A2
    sorted-sparse-matrix c1
    sorted-sparse-matrix c2
    sorted-sparse-matrix y
    sorted-sparse-matrix r1
    sorted-sparse-matrix r2
    sorted-spvec b
    le-spmat [] y
    sparse-row-matrix A1  $\leq$  A

```



```

 $A \leq \text{sparse-row-matrix } A2$ 
 $\text{sparse-row-matrix } c1 \leq c$ 
 $c \leq \text{sparse-row-matrix } c2$ 
 $\text{sparse-row-matrix } r1 \leq x$ 
 $x \leq \text{sparse-row-matrix } r2$ 
 $A * x \leq \text{sparse-row-matrix } (b::('a::\text{lattice-ring}) \text{ spmat})$ 
shows
 $c * x \leq \text{sparse-row-matrix } (\text{add-spmat } (\text{mult-spmat } y \ b)$ 
 $(\text{mult-est-spmat } r1 \ r2 \ (\text{diff-spmat } c1 \ (\text{mult-spmat } y \ A2)) \ (\text{diff-spmat } c2 \ (\text{mult-spmat}$ 
 $y \ A1))))$ 
 $\langle \text{proof} \rangle$ 

 $\langle ML \rangle$ 

end

```