

Machine Words in Isabelle/HOL

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Abstract

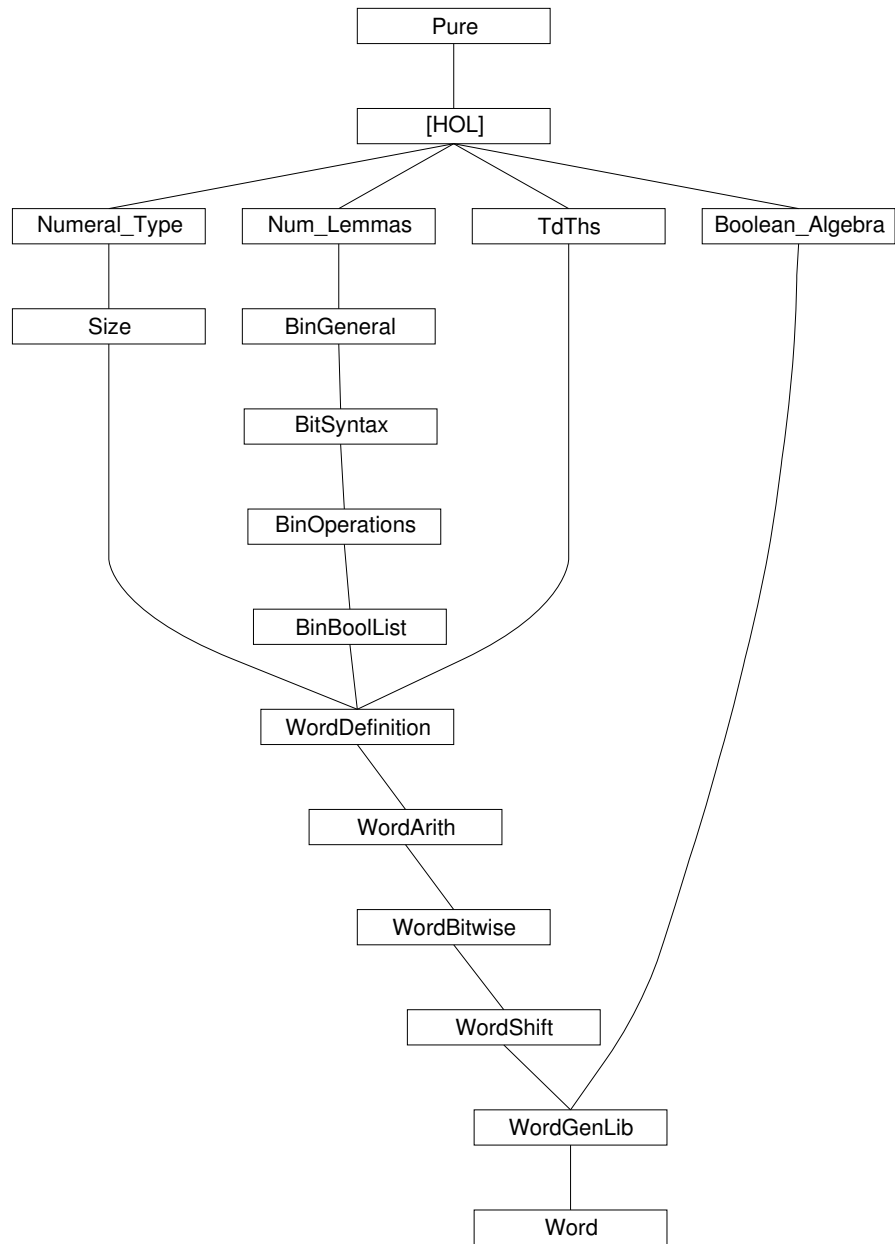
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
imports Main
begin
```

1.1 Preliminary lemmas

```
lemma (in type-definition) univ:
  UNIV = Abs ‘ A
proof
  show Abs ‘ A  $\subseteq$  UNIV by (rule subset-UNIV)
  show UNIV  $\subseteq$  Abs ‘ A
  proof
    fix x :: 'b
    have x = Abs (Rep x) by (rule Rep-inverse [symmetric])
    moreover have Rep x  $\in$  A by (rule Rep)
    ultimately show x  $\in$  Abs ‘ A by (rule image-eqI)
  qed
qed
```

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  by (simp add: univ card-image inj-on-def Abs-inject)
```

1.2 Cardinalities of types

```
syntax -type-card :: type  $\Rightarrow$  nat ((1CARD/(1'(-))))
```

```
translations CARD('t)  $\Rightarrow$  CONST card (CONST UNIV :: 't set)
```

```
typed-print-translation <<
let
  fun card-univ-tr' show-sorts - [Const (@{const-syntax UNIV}, Type(-, [T, -]))]
=
  Syntax.const @{syntax-const -type-card} $ Syntax.term-of-typ show-sorts T;
in [(@{const-syntax card}, card-univ-tr')]
end
>>
```

```
lemma card-unit [simp]: CARD(unit) = 1
  unfolding UNIV-unit by simp
```

```
lemma card-bool [simp]: CARD(bool) = 2
  unfolding UNIV-bool by simp
```

```
lemma card-prod [simp]: CARD('a  $\times$  'b) = CARD('a::finite) * CARD('b::finite)
  unfolding UNIV-Times-UNIV [symmetric] by (simp only: card-cartesian-product)
```

```
lemma card-sum [simp]: CARD('a + 'b) = CARD('a::finite) + CARD('b::finite)
  unfolding UNIV-Plus-UNIV [symmetric] by (simp only: finite card-Plus)
```

```

lemma card-option [simp]:  $CARD('a \text{ option}) = Suc \ CARD('a::finite)$ 
  unfolding UNIV-option-conv
  apply (subgoal-tac (None::'a option)  $\notin$  range Some)
  apply (simp add: card-image)
  apply fast
  done

```

```

lemma card-set [simp]:  $CARD('a \text{ set}) = 2 ^ CARD('a::finite)$ 
  unfolding Pow-UNIV [symmetric]
  by (simp only: card-Pow finite numeral-2-eq-2)

```

```

lemma card-nat [simp]:  $CARD(nat) = 0$ 
  by (simp add: infinite-UNIV-nat card-eq-0-iff)

```

1.3 Classes with at least 1 and 2

Class *finite* already captures “at least 1”

```

lemma zero-less-card-finite [simp]:  $0 < CARD('a::finite)$ 
  unfolding neq0-conv [symmetric] by simp

```

```

lemma one-le-card-finite [simp]:  $Suc \ 0 \leq CARD('a::finite)$ 
  by (simp add: less-Suc-eq-le [symmetric])

```

Class for cardinality “at least 2”

```

class card2 = finite +
  assumes two-le-card:  $2 \leq CARD('a)$ 

```

```

lemma one-less-card:  $Suc \ 0 < CARD('a::card2)$ 
  using two-le-card [where  $'a = 'a$ ] by simp

```

```

lemma one-less-int-card:  $1 < int \ CARD('a::card2)$ 
  using one-less-card [where  $'a = 'a$ ] by simp

```

1.4 Numeral Types

```

typedef (open) num0 = UNIV :: nat set ..
typedef (open) num1 = UNIV :: unit set ..

```

```

typedef (open)  $'a \text{ bit0} = \{0 \ ..< 2 * int \ CARD('a::finite)\}$ 
proof
  show  $0 \in \{0 \ ..< 2 * int \ CARD('a)\}$ 
    by simp
qed

```

```

typedef (open)  $'a \text{ bit1} = \{0 \ ..< 1 + 2 * int \ CARD('a::finite)\}$ 
proof
  show  $0 \in \{0 \ ..< 1 + 2 * int \ CARD('a)\}$ 
    by simp

```

qed

lemma *card-num0* [*simp*]: $CARD\ (num0) = 0$
unfolding *type-definition.card* [*OF type-definition-num0*]
by *simp*

lemma *card-num1* [*simp*]: $CARD(num1) = 1$
unfolding *type-definition.card* [*OF type-definition-num1*]
by (*simp only: card-unit*)

lemma *card-bit0* [*simp*]: $CARD('a\ bit0) = 2 * CARD('a::finite)$
unfolding *type-definition.card* [*OF type-definition-bit0*]
by *simp*

lemma *card-bit1* [*simp*]: $CARD('a\ bit1) = Suc\ (2 * CARD('a::finite))$
unfolding *type-definition.card* [*OF type-definition-bit1*]
by *simp*

instance *num1* :: *finite*

proof

show *finite* (*UNIV::num1 set*)
unfolding *type-definition.univ* [*OF type-definition-num1*]
using *finite* **by** (*rule finite-imageI*)

qed

instance *bit0* :: (*finite*) *card2*

proof

show *finite* (*UNIV::'a bit0 set*)
unfolding *type-definition.univ* [*OF type-definition-bit0*]
by *simp*
show $2 \leq CARD('a\ bit0)$
by *simp*

qed

instance *bit1* :: (*finite*) *card2*

proof

show *finite* (*UNIV::'a bit1 set*)
unfolding *type-definition.univ* [*OF type-definition-bit1*]
by *simp*
show $2 \leq CARD('a\ bit1)$
by *simp*

qed

1.5 Locale for modular arithmetic subtypes

locale *mod-type* =

fixes *n* :: *int*

and *Rep* :: '*a*::{*zero,one,plus,times,uminus,minus*} \Rightarrow *int*

and *Abs* :: *int* \Rightarrow '*a*::{*zero,one,plus,times,uminus,minus*}

```

assumes type: type-definition Rep Abs {0.. $n$ }
and size1:  $1 < n$ 
and zero-def:  $0 = \text{Abs } 0$ 
and one-def:  $1 = \text{Abs } 1$ 
and add-def:  $x + y = \text{Abs } ((\text{Rep } x + \text{Rep } y) \bmod n)$ 
and mult-def:  $x * y = \text{Abs } ((\text{Rep } x * \text{Rep } y) \bmod n)$ 
and diff-def:  $x - y = \text{Abs } ((\text{Rep } x - \text{Rep } y) \bmod n)$ 
and minus-def:  $-x = \text{Abs } ((-\text{Rep } x) \bmod n)$ 
begin

lemma size0:  $0 < n$ 
using size1 by simp

lemmas definitions =
  zero-def one-def add-def mult-def minus-def diff-def

lemma Rep-less-n:  $\text{Rep } x < n$ 
by (rule type-definition.Rep [OF type, simplified, THEN conjunct2])

lemma Rep-le-n:  $\text{Rep } x \leq n$ 
by (rule Rep-less-n [THEN order-less-imp-le])

lemma Rep-inject-sym:  $x = y \longleftrightarrow \text{Rep } x = \text{Rep } y$ 
by (rule type-definition.Rep-inject [OF type, symmetric])

lemma Rep-inverse:  $\text{Abs } (\text{Rep } x) = x$ 
by (rule type-definition.Rep-inverse [OF type])

lemma Abs-inverse:  $m \in \{0.. $n$ \} \implies \text{Rep } (\text{Abs } m) = m$ 
by (rule type-definition.Abs-inverse [OF type])

lemma Rep-Abs-mod:  $\text{Rep } (\text{Abs } (m \bmod n)) = m \bmod n$ 
by (simp add: Abs-inverse pos-mod-conj [OF size0])

lemma Rep-Abs-0:  $\text{Rep } (\text{Abs } 0) = 0$ 
by (simp add: Abs-inverse size0)

lemma Rep-0:  $\text{Rep } 0 = 0$ 
by (simp add: zero-def Rep-Abs-0)

lemma Rep-Abs-1:  $\text{Rep } (\text{Abs } 1) = 1$ 
by (simp add: Abs-inverse size1)

lemma Rep-1:  $\text{Rep } 1 = 1$ 
by (simp add: one-def Rep-Abs-1)

lemma Rep-mod:  $\text{Rep } x \bmod n = \text{Rep } x$ 
apply (rule-tac  $x=x$  in type-definition.Abs-cases [OF type])
apply (simp add: type-definition.Abs-inverse [OF type])

```


apply (*simp add: mod-pos-pos-trivial*)
done

lemmas *Rep-simps* =
Rep-inject-sym Rep-inverse Rep-Abs-mod Rep-mod Rep-Abs-0 Rep-Abs-1

lemma *comm-ring-1*: *OFCLASS('a, comm-ring-1-class)*
apply (*intro-classes, unfold definitions*)
apply (*simp-all add: Rep-simps zmod-simps field-simps*)
done

end

locale *mod-ring* = *mod-type* +
constrains *n* :: *int*
and *Rep* :: '*a*::{*number-ring*} \Rightarrow *int*
and *Abs* :: *int* \Rightarrow '*a*::{*number-ring*}
begin

lemma *of-nat-eq*: *of-nat k = Abs (int k mod n)*
apply (*induct k*)
apply (*simp add: zero-def*)
apply (*simp add: Rep-simps add-def one-def zmod-simps add-ac*)
done

lemma *of-int-eq*: *of-int z = Abs (z mod n)*
apply (*cases z rule: int-diff-cases*)
apply (*simp add: Rep-simps of-nat-eq diff-def zmod-simps*)
done

lemma *Rep-number-of*:
Rep (number-of w) = number-of w mod n
by (*simp add: number-of-eq of-int-eq Rep-Abs-mod*)

lemma *iszero-number-of*:
iszero (number-of w::'a) \longleftrightarrow number-of w mod n = 0
by (*simp add: Rep-simps number-of-eq of-int-eq iszero-def zero-def*)

lemma *cases*:
assumes *1*: $\bigwedge z. \llbracket (x::'a) = \text{of-int } z; 0 \leq z; z < n \rrbracket \Longrightarrow P$
shows *P*
apply (*cases x rule: type-definition.Abs-cases [OF type]*)
apply (*rule-tac z=y in 1*)
apply (*simp-all add: of-int-eq mod-pos-pos-trivial*)
done

lemma *induct*:
 $(\bigwedge z. \llbracket 0 \leq z; z < n \rrbracket \Longrightarrow P (\text{of-int } z)) \Longrightarrow P (x::'a)$
by (*cases x rule: cases simp*)

end

1.6 Number ring instances

Unfortunately a number ring instance is not possible for *num1*, since 0 and 1 are not distinct.

instantiation *num1* :: {*comm-ring*,*comm-monoid-mult*,*number*}
begin

lemma *num1-eq-iff*: (*x*::*num1*) = (*y*::*num1*) \longleftrightarrow *True*
 by (*induct x*, *induct y*) *simp*

instance proof
qed (*simp-all add: num1-eq-iff*)

end

instantiation
bit0 and *bit1* :: (*finite*) {*zero*,*one*,*plus*,*times*,*uminus*,*minus*}
begin

definition *Abs-bit0'* :: *int* \Rightarrow '*a bit0* **where**
Abs-bit0' x = *Abs-bit0* (*x mod int CARD('a bit0)*)

definition *Abs-bit1'* :: *int* \Rightarrow '*a bit1* **where**
Abs-bit1' x = *Abs-bit1* (*x mod int CARD('a bit1)*)

definition 0 = *Abs-bit0 0*

definition 1 = *Abs-bit0 1*

definition *x + y* = *Abs-bit0' (Rep-bit0 x + Rep-bit0 y)*

definition *x * y* = *Abs-bit0' (Rep-bit0 x * Rep-bit0 y)*

definition *x - y* = *Abs-bit0' (Rep-bit0 x - Rep-bit0 y)*

definition $-x$ = *Abs-bit0' (- Rep-bit0 x)*

definition 0 = *Abs-bit1 0*

definition 1 = *Abs-bit1 1*

definition *x + y* = *Abs-bit1' (Rep-bit1 x + Rep-bit1 y)*

definition *x * y* = *Abs-bit1' (Rep-bit1 x * Rep-bit1 y)*

definition *x - y* = *Abs-bit1' (Rep-bit1 x - Rep-bit1 y)*

definition $-x$ = *Abs-bit1' (- Rep-bit1 x)*

instance ..

end

interpretation *bit0*:
mod-type int CARD('a::finite bit0)
Rep-bit0 :: '*a::finite bit0* \Rightarrow *int*

```

      Abs-bit0 :: int ⇒ 'a::finite bit0
apply (rule mod-type.intro)
apply (simp add: int-mult type-definition-bit0)
apply (rule one-less-int-card)
apply (rule zero-bit0-def)
apply (rule one-bit0-def)
apply (rule plus-bit0-def [unfolded Abs-bit0'-def])
apply (rule times-bit0-def [unfolded Abs-bit0'-def])
apply (rule minus-bit0-def [unfolded Abs-bit0'-def])
apply (rule uminus-bit0-def [unfolded Abs-bit0'-def])
done

```

interpretation bit1:

```

      mod-type int CARD('a::finite bit1)
      Rep-bit1 :: 'a::finite bit1 ⇒ int
      Abs-bit1 :: int ⇒ 'a::finite bit1
apply (rule mod-type.intro)
apply (simp add: int-mult type-definition-bit1)
apply (rule one-less-int-card)
apply (rule zero-bit1-def)
apply (rule one-bit1-def)
apply (rule plus-bit1-def [unfolded Abs-bit1'-def])
apply (rule times-bit1-def [unfolded Abs-bit1'-def])
apply (rule minus-bit1-def [unfolded Abs-bit1'-def])
apply (rule uminus-bit1-def [unfolded Abs-bit1'-def])
done

```

```

instance bit0 :: (finite) comm-ring-1
  by (rule bit0.comm-ring-1)+

```

```

instance bit1 :: (finite) comm-ring-1
  by (rule bit1.comm-ring-1)+

```

instantiation bit0 and bit1 :: (finite) number-ring
begin

```

definition (number-of w :: - bit0) = of-int w

```

```

definition (number-of w :: - bit1) = of-int w

```

instance proof

```

qed (rule number-of-bit0-def number-of-bit1-def)+

```

end

interpretation bit0:

```

      mod-ring int CARD('a::finite bit0)
      Rep-bit0 :: 'a::finite bit0 ⇒ int
      Abs-bit0 :: int ⇒ 'a::finite bit0

```

..

interpretation *bit1*:

mod-ring int CARD('a::finite bit1)
Rep-bit1 :: 'a::finite bit1 \Rightarrow int
Abs-bit1 :: int \Rightarrow 'a::finite bit1

..

Set up cases, induction, and arithmetic

lemmas *bit0-cases* [*case-names of-int*, *cases type: bit0*] = *bit0.cases***lemmas** *bit1-cases* [*case-names of-int*, *cases type: bit1*] = *bit1.cases***lemmas** *bit0-induct* [*case-names of-int*, *induct type: bit0*] = *bit0.induct***lemmas** *bit1-induct* [*case-names of-int*, *induct type: bit1*] = *bit1.induct***lemmas** *bit0-iszero-number-of* [*simp*] = *bit0.iszero-number-of***lemmas** *bit1-iszero-number-of* [*simp*] = *bit1.iszero-number-of*

1.7 Syntax

syntax*-NumeralType* :: *num-const* \Rightarrow *type* (-)*-NumeralType0* :: *type* (0)*-NumeralType1* :: *type* (1)**translations***(type)* 1 == *(type)* *num1**(type)* 0 == *(type)* *num0***parse-translation** $\langle\langle$ *let**fun mk-bintype n =**let*

fun mk-bit 0 = Syntax.const @{type-syntax bit0}
| mk-bit 1 = Syntax.const @{type-syntax bit1};

fun bin-of n =

if n = 1 then Syntax.const @{type-syntax num1}
else if n = 0 then Syntax.const @{type-syntax num0}
else if n = ~1 then raise TERM (negative type numeral, [])
else

*let val (q, r) = Integer.div-mod n 2;**in mk-bit r \$ bin-of q end;**in bin-of n end;**fun numeral-tr (*-NumeralType*) [Const (str, -)] =**mk-bintype (the (Int.fromString str))**| numeral-tr (*-NumeralType*) ts = raise TERM (numeral-tr, ts);**in [(@{syntax-const -NumeralType}, numeral-tr)] end;*

»

print-translation «

let

fun int-of [] = 0

| *int-of* (b :: bs) = b + 2 * *int-of* bs;

fun bin-of (Const (@{type-syntax num0}, -)) = []

| *bin-of* (Const (@{type-syntax num1}, -)) = [1]

| *bin-of* (Const (@{type-syntax bit0}, -) \$ bs) = 0 :: *bin-of* bs

| *bin-of* (Const (@{type-syntax bit1}, -) \$ bs) = 1 :: *bin-of* bs

| *bin-of* t = raise TERM (*bin-of*, [t]);

fun bit-tr' b [t] =

let

val rev-digs = b :: *bin-of* t handle TERM - => raise Match

val i = *int-of* rev-digs;

val num = *string-of-int* (abs i);

in

Syntax.const @{syntax-const -NumeralType} \$ *Syntax.free* num

end

| *bit-tr' b* - = raise Match;

in [(@{type-syntax bit0}, *bit-tr' 0*), (@{type-syntax bit1}, *bit-tr' 1*)] *end*;

»

1.8 Examples

lemma CARD(0) = 0 **by** simp

lemma CARD(17) = 17 **by** simp

lemma 8 * 11 ^ 3 - 6 = (2::5) **by** simp

end

2 Size: The len classes

theory Size

imports Numeral-Type

begin

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in *Numeral-Type*.

class len0 =

fixes len-of :: 'a itself \Rightarrow nat

Some theorems are only true on words with length greater 0.

class len = len0 +

```

assumes len-gt-0 [iff]:  $0 < \text{len-of TYPE } ('a)$ 

instantiation num0 and num1 :: len0
begin

definition
  len-num0:  $\text{len-of } (x::\text{num0 } \text{itself}) = 0$ 

definition
  len-num1:  $\text{len-of } (x::\text{num1 } \text{itself}) = 1$ 

instance ..

end

instantiation bit0 and bit1 :: (len0) len0
begin

definition
  len-bit0:  $\text{len-of } (x::'a::\text{len0 } \text{bit0 } \text{itself}) = 2 * \text{len-of TYPE } ('a)$ 

definition
  len-bit1:  $\text{len-of } (x::'a::\text{len0 } \text{bit1 } \text{itself}) = 2 * \text{len-of TYPE } ('a) + 1$ 

instance ..

end

lemmas len-of-numeral-defs [simp] = len-num0 len-num1 len-bit0 len-bit1

instance num1 :: len by (intro-classes) simp
instance bit0 :: (len) len by (intro-classes) simp
instance bit1 :: (len0) len by (intro-classes) simp

— Examples:
lemma len-of TYPE(17) = 17 by simp
lemma len-of TYPE(0) = 0 by simp

— not simplified:
lemma len-of TYPE('a::len0) = x
  oops

end

```

3 Num-Lemmas: Useful Numerical Lemmas

theory *Num-Lemmas*

```

imports Main Parity
begin

lemma contentsI:  $y = \{x\} \implies \text{contents } y = x$ 
  unfolding contents-def by auto — FIXME move

lemmas split-split = prod.split [unfolded prod-case-split]
lemmas split-split-asm = prod.split-asm [unfolded prod-case-split]
lemmas split-splits = split-split split-split-asm

lemmas funpow-0 = funpow.simps(1)
lemmas funpow-Suc = funpow.simps(2)

lemma nonemptyE:  $S \sim \{\} \implies (!x. x : S \implies R) \implies R$  by auto

lemma gt-or-eq-0:  $0 < y \vee 0 = (y::\text{nat})$  by arith

declare iszero-0 [iff]

lemmas xtr1 = xtrans(1)
lemmas xtr2 = xtrans(2)
lemmas xtr3 = xtrans(3)
lemmas xtr4 = xtrans(4)
lemmas xtr5 = xtrans(5)
lemmas xtr6 = xtrans(6)
lemmas xtr7 = xtrans(7)
lemmas xtr8 = xtrans(8)

lemmas nat-simps = diff-add-inverse2 diff-add-inverse
lemmas nat-iffs = le-add1 le-add2

lemma sum-imp-diff:  $j = k + i \implies j - i = (k :: \text{nat})$  by arith

lemma nobm1:
   $0 < (\text{number-of } w :: \text{nat}) \implies$ 
     $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Int.pred } w)$ 
  apply (unfold nat-number-of-def One-nat-def nat-1 [symmetric] pred-def)
  apply (simp add: number-of-eq nat-diff-distrib [symmetric])
  done

lemma zless2:  $0 < (2 :: \text{int})$  by arith

lemmas zless2p [simp] = zless2 [THEN zero-less-power]
lemmas zle2p [simp] = zless2p [THEN order-less-imp-le]

lemmas pos-mod-sign2 = zless2 [THEN pos-mod-sign [where b = 2::int]]
lemmas pos-mod-bound2 = zless2 [THEN pos-mod-bound [where b = 2::int]]

— the inverse(s) of number-of

```

lemma *nmod2*: $n \bmod (2::int) = 0 \mid n \bmod 2 = 1$ **by** *arith*

lemma *emep1*:

even n ==> even d ==> 0 <= d ==> (n + 1) mod (d :: int) = (n mod d) + 1

apply (*simp add: add-commute*)

apply (*safe dest!: even-equiv-def [THEN iffD1]*)

apply (*subst pos-zmod-mult-2*)

apply *arith*

apply (*simp add: mod-mult-mult1*)

done

lemmas *eme1p* = *emep1* [*simplified add-commute*]

lemma *le-diff-eq'*: $(a \leq c - b) = (b + a \leq (c::int))$ **by** *arith*

lemma *less-diff-eq'*: $(a < c - b) = (b + a < (c::int))$ **by** *arith*

lemma *diff-le-eq'*: $(a - b \leq c) = (a \leq b + (c::int))$ **by** *arith*

lemma *diff-less-eq'*: $(a - b < c) = (a < b + (c::int))$ **by** *arith*

lemmas *m1mod2k* = *zless2p* [*THEN zmod-minus1*]

lemmas *m1mod22k* = *mult-pos-pos* [*OF zless2 zless2p, THEN zmod-minus1*]

lemmas *p1mod22k'* = *zless2p* [*THEN order-less-imp-le, THEN pos-zmod-mult-2*]

lemmas *z1pmod2'* = *zero-le-one* [*THEN pos-zmod-mult-2, simplified*]

lemmas *z1pdiv2'* = *zero-le-one* [*THEN pos-zdiv-mult-2, simplified*]

lemma *p1mod22k*:

$(2 * b + 1) \bmod (2 * 2 ^ n) = 2 * (b \bmod 2 ^ n) + (1::int)$

by (*simp add: p1mod22k' add-commute*)

lemma *z1pmod2*:

$(2 * b + 1) \bmod 2 = (1::int)$ **by** *arith*

lemma *z1pdiv2*:

$(2 * b + 1) \div 2 = (b::int)$ **by** *arith*

lemmas *zdiv-le-dividend* = *xtr3* [*OF div-by-1 [symmetric] zdiv-mono2, simplified int-one-le-iff-zero-less, simplified, standard*]

lemma *axbbyy*:

$a + m + m = b + n + n ==> (a = 0 \mid a = 1) ==> (b = 0 \mid b = 1) ==>$

$a = b \ \& \ m = (n :: int)$ **by** *arith*

lemma *axrmod2*:

$(1 + x + x) \bmod 2 = (1 :: int) \ \& \ (0 + x + x) \bmod 2 = (0 :: int)$ **by** *arith*

lemma *axrdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$ **by** *arith*

lemmas *iszero-minus* = *trans* [*THEN trans*,
OF iszero-def neg-equal-0-iff-equal iszero-def [symmetric], standard]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel [symmetric] add-commute*,
standard]

lemmas *add-diff-cancel2* = *add-commute* [*THEN diff-eq-eq [THEN iffD2], standard*]

lemma *zmod-uminus*: $-(a :: \text{int}) \text{ mod } b \text{ mod } b = -a \text{ mod } b$
by (*simp add : zmod-zminus1-eq-if*)

lemma *zmod-zsub-distrib*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b \text{ mod } c) \text{ mod } c$
apply (*unfold diff-int-def*)
apply (*rule trans [OF - mod-add-eq [symmetric]]*)
apply (*simp add: zmod-uminus mod-add-eq [symmetric]*)
done

lemma *zmod-zsub-right-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a - b \text{ mod } c) \text{ mod } c$
apply (*unfold diff-int-def*)
apply (*rule trans [OF - mod-add-right-eq [symmetric]]*)
apply (*simp add : zmod-uminus mod-add-right-eq [symmetric]*)
done

lemma *zmod-zsub-left-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b) \text{ mod } c$
by (*rule mod-add-left-eq [where b = - b, simplified diff-int-def [symmetric]]*)

lemma *zmod-zsub-self* [*simp*]:
 $((b :: \text{int}) - a) \text{ mod } a = b \text{ mod } a$
by (*simp add: zmod-zsub-right-eq*)

lemma *zmod-zmult1-eq-rev*:
 $b * a \text{ mod } c = b \text{ mod } c * a \text{ mod } (c :: \text{int})$
apply (*simp add: mult-commute*)
apply (*subst zmod-zmult1-eq*)
apply *simp*
done

lemmas *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]
zmod-zsub-left-eq zmod-zsub-right-eq mod-add-left-eq
mod-add-right-eq zmod-zmult1-eq zmod-zmult1-eq-rev

lemma *mod-plus-right*:
 $((a + x) \text{ mod } m = (b + x) \text{ mod } m) = (a \text{ mod } m = b \text{ mod } (m :: \text{nat}))$
apply (*induct x*)
apply (*simp-all add: mod-Suc*)
apply *arith*

done

lemma *nat-minus-mod*: $(n - n \bmod m) \bmod m = (0 :: nat)$
by (*induct* *n*) (*simp-all* *add* : *mod-Suc*)

lemmas *nat-minus-mod-plus-right* = *trans* [*OF* *nat-minus-mod mod-0* [*symmetric*],
THEN *mod-plus-right* [*THEN* *iffD2*], *standard*, *simplified*]

lemmas *push-mods'* = *mod-add-eq* [*standard*]
mod-mult-eq [*standard*] *zmod-zsub-distrib* [*standard*]
zmod-uminus [*symmetric*, *standard*]

lemmas *push-mods* = *push-mods'* [*THEN* *eq-reflection*, *standard*]

lemmas *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN* *eq-reflection*, *standard*]

lemmas *mod-simps* =
mod-mult-self2-is-0 [*THEN* *eq-reflection*]
mod-mult-self1-is-0 [*THEN* *eq-reflection*]
mod-mod-trivial [*THEN* *eq-reflection*]

lemma *nat-mod-eq*:
 $!!b. b < n ==> a \bmod n = b \bmod n ==> a \bmod n = (b :: nat)$
by (*induct* *a*) *auto*

lemmas *nat-mod-eq'* = *refl* [*THEN* [*2*] *nat-mod-eq*]

lemma *nat-mod-lem*:
 $(0 :: nat) < n ==> b < n = (b \bmod n = b)$
apply *safe*
apply (*erule* *nat-mod-eq'*)
apply (*erule* *subst*)
apply (*erule* *mod-less-divisor*)
done

lemma *mod-nat-add*:
 $(x :: nat) < z ==> y < z ==>$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
apply (*rule* *nat-mod-eq*)
apply *auto*
apply (*rule* *trans*)
apply (*rule* *le-mod-geq*)
apply *simp*
apply (*rule* *nat-mod-eq'*)
apply *arith*
done

lemma *mod-nat-sub*:
 $(x :: nat) < z ==> (x - y) \bmod z = x - y$
by (*rule* *nat-mod-eq'*) *arith*

lemma *int-mod-lem*:

$(0 :: \text{int}) < n \implies (0 \leq b \ \& \ b < n) = (b \bmod n = b)$
apply *safe*
apply (*erule* (1) *mod-pos-pos-trivial*)
apply (*erule-tac* [!] *subst*)
apply *auto*
done

lemma *int-mod-eq*:

$(0 :: \text{int}) \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$
by *clarsimp* (*rule mod-pos-pos-trivial*)

lemmas *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

lemma *int-mod-le*: $0 \leq a \implies 0 < (n :: \text{int}) \implies a \bmod n \leq a$

apply (*cases* $a < n$)
apply (*auto dest: mod-pos-pos-trivial pos-mod-bound* [**where** $a=a$])
done

lemma *int-mod-le'*: $0 \leq b - n \implies 0 < (n :: \text{int}) \implies b \bmod n \leq b - n$

by (*rule int-mod-le* [**where** $a = b - n$ **and** $n = n$, *simplified*])

lemma *int-mod-ge*: $a < n \implies 0 < (n :: \text{int}) \implies a \leq a \bmod n$

apply (*cases* $0 \leq a$)
apply (*drule* (1) *mod-pos-pos-trivial*)
apply *simp*
apply (*rule order-trans* [*OF* - *pos-mod-sign*])
apply *simp*
apply *assumption*
done

lemma *int-mod-ge'*: $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$

by (*rule int-mod-ge* [**where** $a = b + n$ **and** $n = n$, *simplified*])

lemma *mod-add-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
by (*auto intro: int-mod-eq*)

lemma *mod-sub-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
by (*auto intro: int-mod-eq*)

lemmas *zmde* = *zmod-zdiv-equality* [*THEN* *diff-eq-eq* [*THEN* *iffD2*], *symmetric*]

lemmas *mcl* = *mult-cancel-left* [*THEN* *iffD1*, *THEN* *make-pos-rule*]

```

lemma zdiv-mult-self:  $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$ 
  apply (rule mcl)
  prefer 2
  apply (erule asm-rl)
  apply (simp add: zmde ring-distrib)
done

```

```

lemma eqne:  $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$ 
  unfolding equiv-def refl-on-def quotient-def Image-def by auto

```

```

lemmas Rep-Integ-ne = Integ.Rep-Integ
  [THEN equiv-intrel [THEN eqne, simplified Integ-def [symmetric]], standard]

```

```

lemmas riq = Integ.Rep-Integ [simplified Integ-def]
lemmas intrel-refl = refl [THEN equiv-intrel-iff [THEN iffD1], standard]
lemmas Rep-Integ-equiv = quotient-eq-iff
  [OF equiv-intrel riq riq, simplified Integ.Rep-Integ-inject, standard]
lemmas Rep-Integ-same =
  Rep-Integ-equiv [THEN intrel-refl [THEN rev-iffD2], standard]

```

```

lemma RI-int:  $(a, 0) : \text{Rep-Integ } (\text{int } a)$ 
  unfolding int-def by auto

```

```

lemmas RI-intrel [simp] = UNIV-I [THEN quotientI,
  THEN Integ.Abs-Integ-inverse [simplified Integ-def], standard]

```

```

lemma RI-minus:  $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (-x)$ 
  apply (rule-tac z=x in eq-Abs-Integ)
  apply (clarsimp simp: minus)
done

```

```

lemma RI-add:
   $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$ 
   $(a + c, b + d) : \text{Rep-Integ } (x + y)$ 
  apply (rule-tac z=x in eq-Abs-Integ)
  apply (rule-tac z=y in eq-Abs-Integ)
  apply (clarsimp simp: add)
done

```

```

lemma mem-same:  $a : S \implies a = b \implies b : S$ 
  by fast

```

```

lemma RI-eq-diff':  $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$ 
  apply (unfold diff-def)
  apply (rule mem-same)
  apply (rule RI-minus RI-add RI-int) +
  apply simp

```

done

lemma *RI-eq-diff*: $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$
apply *safe*
apply (*rule Rep-Integ-same*)
prefer 2
apply (*erule asm-rl*)
apply (*rule RI-eq-diff'*)
done

lemma *mod-power-lem*:
 $a > 1 \implies a^n \bmod a^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int})^n)$
apply *clarsimp*
apply *safe*
apply (*simp add: dvd-eq-mod-eq-0 [symmetric]*)
apply (*erule le-iff-add [THEN iffD1]*)
apply (*force simp: zpower-zadd-distrib*)
apply (*rule mod-pos-pos-trivial*)
apply (*simp*)
apply (*rule power-strict-increasing*)
apply *auto*
done

lemma *min-pm* [*simp*]: $\min a b + (a - b) = (a :: \text{nat})$ **by** *arith*

lemmas *min-pm1* [*simp*] = *trans [OF add-commute min-pm]*

lemma *rev-min-pm* [*simp*]: $\min b a + (a - b) = (a :: \text{nat})$ **by** *arith*

lemmas *rev-min-pm1* [*simp*] = *trans [OF add-commute rev-min-pm]*

lemma *pl-pl-rels*:
 $a + b = c + d \implies$
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$ **by** *arith*

lemmas *pl-pl-rels'* = *add-commute [THEN [2] trans, THEN pl-pl-rels]*

lemma *minus-eq*: $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$ **by** *arith*

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$ **by** *arith*

lemmas *pl-pl-mm'* = *add-commute [THEN [2] trans, THEN pl-pl-mm]*

lemma *min-minus* [*simp*]: $\min m (m - k) = (m - k :: \text{nat})$ **by** *arith*

lemmas *min-minus'* [*simp*] = *trans [OF min-max.inf-commute min-minus]*

lemma *nat-no-eq-iff*:
 $(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$

```

    (number-of b = (number-of c :: nat)) = (b = c)
  apply (unfold nat-number-of-def)
  apply safe
  apply (drule (2) eq-nat-nat-iff [THEN iffD1])
  apply (simp add: number-of-eq)
  done

lemmas dne = box-equals [OF div-mod-equality add-0-right add-0-right]
lemmas dtle = xtr3 [OF dne [symmetric] le-add1]
lemmas th2 = order-trans [OF order-refl [THEN [2] mult-le-mono] dtle]

lemma td-gal:
  0 < c ==> (a >= b * c) = (a div c >= (b :: nat))
  apply safe
  apply (erule (1) xtr4 [OF div-le-mono div-mult-self-is-m])
  apply (erule th2)
  done

lemmas td-gal-lt = td-gal [simplified not-less [symmetric], simplified]

lemma div-mult-le: (a :: nat) div b * b <= a
  apply (cases b)
  prefer 2
  apply (rule order-refl [THEN th2])
  apply auto
  done

lemmas sdl = split-div-lemma [THEN iffD1, symmetric]

lemma given-quot: f > (0 :: nat) ==> (f * l + (f - 1)) div f = l
  by (rule sdl, assumption) (simp (no-asm))

lemma given-quot-alt: f > (0 :: nat) ==> (l * f + f - Suc 0) div f = l
  apply (frule given-quot)
  apply (rule trans)
  prefer 2
  apply (erule asm-rl)
  apply (rule-tac f=%n. n div f in arg-cong)
  apply (simp add : mult-ac)
  done

lemma diff-mod-le: (a::nat) < d ==> b dvd d ==> a - a mod b <= d - b
  apply (unfold dvd-def)
  apply clarify
  apply (case-tac k)
  apply clarsimp
  apply clarify
  apply (cases b > 0)
  apply (drule mult-commute [THEN xtr1])

```

```

apply (frule (1) td-gal-lt [THEN iffD1])
apply (clarsimp simp: le-simps)
apply (rule mult-div-cancel [THEN [2] xtr4])
apply (rule mult-mono)
apply auto
done

lemma less-le-mult':
   $w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: int)$ 
apply (rule mult-right-mono)
apply (rule zless-imp-add1-zle)
apply (erule (1) mult-right-less-imp-less)
apply assumption
done

lemmas less-le-mult = less-le-mult' [simplified left-distrib, simplified]

lemmas less-le-mult-minus = iffD2 [OF le-diff-eq less-le-mult,
  simplified left-diff-distrib, standard]

lemma lrlem':
  assumes d:  $(i :: nat) \leq j \vee m < j'$ 
  assumes R1:  $i * k \leq j * k \implies R$ 
  assumes R2:  $Suc\ m * k' \leq j' * k' \implies R$ 
  shows R using d
  apply safe
  apply (rule R1, erule mult-le-mono1)
  apply (rule R2, erule Suc-le-eq [THEN iffD2 [THEN mult-le-mono1]])
  done

lemma lrlem:  $(0 :: nat) < sc \implies$ 
   $(sc - n + (n + lb * n) <= m * n) = (sc + lb * n <= m * n)$ 
  apply safe
  apply arith
  apply (case-tac sc >= n)
  apply arith
  apply (insert linorder-le-less-linear [of m lb])
  apply (erule-tac k=n and k'=n in lrlem')
  apply arith
  apply simp
  done

lemma gen-minus:  $0 < n \implies f\ n = f\ (Suc\ (n - 1))$ 
  by auto

lemma mpl-lem:  $j <= (i :: nat) \implies k < j \implies i - j + k < i$  by arith

lemma nonneg-mod-div:
   $0 <= a \implies 0 <= b \implies 0 <= (a \bmod b :: int) \ \& \ 0 <= a \div b$ 

```

```

apply (cases b = 0, clarsimp)
apply (auto intro: pos-imp-zdiv-nonneg-iff [THEN iffD2])
done

end

```

4 BinGeneral: Basic Definitions for Binary Integers

```

theory BinGeneral
imports Num-Lemmas
begin

```

4.1 Further properties of numerals

```

datatype bit = B0 | B1

```

definition

```

  Bit :: int  $\Rightarrow$  bit  $\Rightarrow$  int (infixl BIT 90) where
    k BIT b = (case b of B0  $\Rightarrow$  0 | B1  $\Rightarrow$  1) + k + k

```

```

lemma BIT-B0-eq-Bit0 [simp]: w BIT B0 = Int.Bit0 w
unfolding Bit-def Bit0-def by simp

```

```

lemma BIT-B1-eq-Bit1 [simp]: w BIT B1 = Int.Bit1 w
unfolding Bit-def Bit1-def by simp

```

```

lemmas BIT-simps = BIT-B0-eq-Bit0 BIT-B1-eq-Bit1

```

```

hide-const (open) B0 B1

```

```

lemma Min-ne-Pls [iff]:
  Int.Min  $\sim$  Int.Pl
unfolding Min-def Pls-def by auto

```

```

lemmas Pls-ne-Min [iff] = Min-ne-Pls [symmetric]

```

```

lemmas PlsMin-defs [intro!] =
  Pls-def Min-def Pls-def [symmetric] Min-def [symmetric]

```

```

lemmas PlsMin-simps [simp] = PlsMin-defs [THEN Eq-TrueI]

```

```

lemma number-of-False-cong:
  False  $\implies$  number-of x = number-of y
by (rule FalseE)

```


lemma *BIT-eq*: $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$
apply (*unfold Bit-def*)
apply (*simp* (*no-asm-use*) *split*: *bit.split-asm*)
apply *simp-all*
apply (*drule-tac f=even in arg-cong, clarsimp*) +
done

lemmas *BIT-eqE* [*elim!*] = *BIT-eq* [*THEN conjE, standard*]

lemma *BIT-eq-iff* [*simp*]:
 $(u \text{ BIT } b = v \text{ BIT } c) = (u = v \ \wedge \ b = c)$
by (*rule iffI*) *auto*

lemmas *BIT-eqI* [*intro!*] = *conjI* [*THEN BIT-eq-iff* [*THEN iffD2*]]

lemma *less-Bits*:
 $(v \text{ BIT } b < w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$
unfolding *Bit-def* **by** (*auto split*: *bit.split*)

lemma *le-Bits*:
 $(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim = \text{bit.B1} \mid c \sim = \text{bit.B0}))$
unfolding *Bit-def* **by** (*auto split*: *bit.split*)

lemma *no-no* [*simp*]: *number-of* (*number-of* *i*) = *i*
unfolding *number-of-eq* **by** *simp*

lemma *Bit-B0*:
 $k \text{ BIT } \text{bit.B0} = k + k$
by (*unfold Bit-def*) *simp*

lemma *Bit-B1*:
 $k \text{ BIT } \text{bit.B1} = k + k + 1$
by (*unfold Bit-def*) *simp*

lemma *Bit-B0-2t*: $k \text{ BIT } \text{bit.B0} = 2 * k$
by (*rule trans, rule Bit-B0*) *simp*

lemma *Bit-B1-2t*: $k \text{ BIT } \text{bit.B1} = 2 * k + 1$
by (*rule trans, rule Bit-B1*) *simp*

lemma *B-mod-2'*:
 $X = 2 \implies (w \text{ BIT } \text{bit.B1}) \bmod X = 1 \ \& \ (w \text{ BIT } \text{bit.B0}) \bmod X = 0$
apply (*simp* (*no-asm*) *only*: *Bit-B0 Bit-B1*)
apply (*simp add*: *z1pmod2*)
done

lemma *B1-mod-2* [*simp*]: $(\text{Int.Bit1 } w) \bmod 2 = 1$
unfolding *numeral-simps number-of-is-id* **by** (*simp add*: *z1pmod2*)

lemma *B0-mod-2* [simp]: $(\text{Int.Bit0 } w) \bmod 2 = 0$
unfolding *numeral-simps number-of-is-id* **by** *simp*

lemma *neB1E* [elim!]:
assumes *ne*: $y \neq \text{bit.B1}$
assumes *y*: $y = \text{bit.B0} \implies P$
shows *P*
apply (*rule y*)
apply (*cases y rule: bit.exhaust, simp*)
apply (*simp add: ne*)
done

lemma *bin-ex-rl*: $EX\ w\ b.\ w\ \text{BIT}\ b = \text{bin}$
apply (*unfold Bit-def*)
apply (*cases even bin*)
apply (*clarsimp simp: even-equiv-def*)
apply (*auto simp: odd-equiv-def split: bit.split*)
done

lemma *bin-exhaust*:
assumes *Q*: $\bigwedge x\ b.\ \text{bin} = x\ \text{BIT}\ b \implies Q$
shows *Q*
apply (*insert bin-ex-rl [of bin]*)
apply (*erule exE*)
apply (*rule Q*)
apply *force*
done

4.2 Destructors for binary integers

definition *bin-rl* :: $\text{int} \Rightarrow \text{int} \times \text{bit}$ **where**
[*code del*]: *bin-rl* *w* = (*THE* (*r*, *l*). $w = r\ \text{BIT}\ l$)

lemma *bin-rl-char*: $(\text{bin-rl } w = (r, l)) = (r\ \text{BIT}\ l = w)$
apply (*unfold bin-rl-def*)
apply *safe*
apply (*cases w rule: bin-exhaust*)
apply *auto*
done

definition
bin-rest-def [*code del*]: *bin-rest* *w* = *fst* (*bin-rl* *w*)

definition
bin-last-def [*code del*]: *bin-last* *w* = *snd* (*bin-rl* *w*)

primrec *bin-nth* **where**
 Z : *bin-nth* *w* 0 = (*bin-last* *w* = *bit.B1*)

| *Suc*: $\text{bin-nth } w \ (\text{Suc } n) = \text{bin-nth } (\text{bin-rest } w) \ n$

lemma *bin-rl*: $\text{bin-rl } w = (\text{bin-rest } w, \text{bin-last } w)$
unfolding *bin-rest-def bin-last-def* **by** *auto*

lemma *bin-rl-simps* [*simp*]:
 $\text{bin-rl } \text{Int.Pl}s = (\text{Int.Pl}s, \text{bit.B0})$
 $\text{bin-rl } \text{Int.Min} = (\text{Int.Min}, \text{bit.B1})$
 $\text{bin-rl } (\text{Int.Bit0 } r) = (r, \text{bit.B0})$
 $\text{bin-rl } (\text{Int.Bit1 } r) = (r, \text{bit.B1})$
 $\text{bin-rl } (r \text{ BIT } b) = (r, b)$
unfolding *bin-rl-char* **by** *simp-all*

declare *bin-rl-simps*(1–4) [*code*]

lemmas *bin-rl-simp* [*simp*] = *iffD1* [*OF bin-rl-char bin-rl*]

lemma *bin-abs-lem*:
 $\text{bin} = (w \text{ BIT } b) \implies \sim \text{bin} = \text{Int.Min} \dashrightarrow \sim \text{bin} = \text{Int.Pl}s \dashrightarrow$
 $\text{nat } (\text{abs } w) < \text{nat } (\text{abs } \text{bin})$
apply (*clarsimp simp add: bin-rl-char*)
apply (*unfold Pls-def Min-def Bit-def*)
apply (*cases b*)
apply (*clarsimp, arith*)
apply (*clarsimp, arith*)
done

lemma *bin-induct*:
assumes *PPls*: $P \ \text{Int.Pl}s$
and *PMin*: $P \ \text{Int.Min}$
and *PBit*: $\forall \text{bin bit}. P \ \text{bin} \implies P \ (\text{bin BIT bit})$
shows $P \ \text{bin}$
apply (*rule-tac P=P and a=bin and f1=nat o abs*
in *wf-measure [THEN wf-induct]*)
apply (*simp add: measure-def inv-image-def*)
apply (*case-tac x rule: bin-exhaust*)
apply (*frule bin-abs-lem*)
apply (*auto simp add : PPls PMin PBit*)
done

lemma *numeral-induct*:
assumes *Pls*: $P \ \text{Int.Pl}s$
assumes *Min*: $P \ \text{Int.Min}$
assumes *Bit0*: $\bigwedge w. \llbracket P \ w; w \neq \text{Int.Pl}s \rrbracket \implies P \ (\text{Int.Bit0 } w)$
assumes *Bit1*: $\bigwedge w. \llbracket P \ w; w \neq \text{Int.Min} \rrbracket \implies P \ (\text{Int.Bit1 } w)$
shows $P \ x$
apply (*induct x rule: bin-induct*)
apply (*rule Pls*)
apply (*rule Min*)

```

apply (case-tac bit)
apply (case-tac bin = Int.Pls)
  apply simp
apply (simp add: Bit0)
apply (case-tac bin = Int.Min)
  apply simp
apply (simp add: Bit1)
done

```

```

lemma bin-rest-simps [simp]:
  bin-rest Int.Pls = Int.Pls
  bin-rest Int.Min = Int.Min
  bin-rest (Int.Bit0 w) = w
  bin-rest (Int.Bit1 w) = w
  bin-rest (w BIT b) = w
unfolding bin-rest-def by auto

```

```

declare bin-rest-simps(1-4) [code]

```

```

lemma bin-last-simps [simp]:
  bin-last Int.Pls = bit.B0
  bin-last Int.Min = bit.B1
  bin-last (Int.Bit0 w) = bit.B0
  bin-last (Int.Bit1 w) = bit.B1
  bin-last (w BIT b) = b
unfolding bin-last-def by auto

```

```

declare bin-last-simps(1-4) [code]

```

```

lemma bin-r-l-extras [simp]:
  bin-last 0 = bit.B0
  bin-last (- 1) = bit.B1
  bin-last -1 = bit.B1
  bin-last 1 = bit.B1
  bin-rest 1 = 0
  bin-rest 0 = 0
  bin-rest (- 1) = - 1
  bin-rest -1 = -1
apply (unfold number-of-Min)
apply (unfold Pls-def [symmetric] Min-def [symmetric])
apply (unfold numeral-1-eq-1 [symmetric])
apply (auto simp: number-of-eq)
done

```

```

lemma bin-last-mod:
  bin-last w = (if w mod 2 = 0 then bit.B0 else bit.B1)
apply (case-tac w rule: bin-exhaust)
apply (case-tac b)
apply auto

```

done

```
lemma bin-rest-div:
  bin-rest w = w div 2
  apply (case-tac w rule: bin-exhaust)
  apply (rule trans)
  apply clarsimp
  apply (rule refl)
  apply (drule trans)
  apply (rule Bit-def)
  apply (simp add: z1pdiv2 split: bit.split)
done
```

```
lemma Bit-div2 [simp]: (w BIT b) div 2 = w
  unfolding bin-rest-div [symmetric] by auto
```

```
lemma Bit0-div2 [simp]: (Int.Bit0 w) div 2 = w
  using Bit-div2 [where b=bit.B0] by simp
```

```
lemma Bit1-div2 [simp]: (Int.Bit1 w) div 2 = w
  using Bit-div2 [where b=bit.B1] by simp
```

```
lemma bin-nth-lem [rule-format]:
  ALL y. bin-nth x = bin-nth y --> x = y
  apply (induct x rule: bin-induct)
  apply safe
  apply (erule rev-mp)
  apply (induct-tac y rule: bin-induct)
  apply (safe del: subset-antisym)
  apply (drule-tac x=0 in fun-cong, force)
  apply (erule notE, rule ext,
    drule-tac x=Suc x in fun-cong, force)
  apply (drule-tac x=0 in fun-cong, force)
  apply (erule rev-mp)
  apply (induct-tac y rule: bin-induct)
  apply (safe del: subset-antisym)
  apply (drule-tac x=0 in fun-cong, force)
  apply (erule notE, rule ext,
    drule-tac x=Suc x in fun-cong, force)
  apply (drule-tac x=0 in fun-cong, force)
  apply (case-tac y rule: bin-exhaust)
  apply clarify
  apply (erule allE)
  apply (erule impE)
  prefer 2
  apply (erule BIT-eqI)
  apply (drule-tac x=0 in fun-cong, force)
  apply (rule ext)
  apply (drule-tac x=Suc ?x in fun-cong, force)
```

done

lemma *bin-nth-eq-iff*: $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$
by (*auto elim: bin-nth-lem*)

lemmas *bin-eqI* = *ext* [*THEN bin-nth-eq-iff* [*THEN iffD1*], *standard*]

lemma *bin-nth-Pls* [*simp*]: $\sim \text{bin-nth Int.Pls } n$
by (*induct n*) *auto*

lemma *bin-nth-Min* [*simp*]: $\text{bin-nth Int.Min } n$
by (*induct n*) *auto*

lemma *bin-nth-0-BIT*: $\text{bin-nth } (w \text{ BIT } b) \ 0 = (b = \text{bit.B1})$
by *auto*

lemma *bin-nth-Suc-BIT*: $\text{bin-nth } (w \text{ BIT } b) (\text{Suc } n) = \text{bin-nth } w \ n$
by *auto*

lemma *bin-nth-minus* [*simp*]: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) \ n = \text{bin-nth } w \ (n - 1)$
by (*cases n*) *auto*

lemma *bin-nth-minus-Bit0* [*simp*]:
 $0 < n \implies \text{bin-nth } (\text{Int.Bit0 } w) \ n = \text{bin-nth } w \ (n - 1)$
using *bin-nth-minus* [**where** *b=bit.B0*] **by** *simp*

lemma *bin-nth-minus-Bit1* [*simp*]:
 $0 < n \implies \text{bin-nth } (\text{Int.Bit1 } w) \ n = \text{bin-nth } w \ (n - 1)$
using *bin-nth-minus* [**where** *b=bit.B1*] **by** *simp*

lemmas *bin-nth-0* = *bin-nth.simps(1)*

lemmas *bin-nth-Suc* = *bin-nth.simps(2)*

lemmas *bin-nth-simps* =
bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus
bin-nth-minus-Bit0 bin-nth-minus-Bit1

4.3 Recursion combinator for binary integers

lemma *brlem*: $(\text{bin} = \text{Int.Min}) = (- \text{bin} + \text{Int.pred } 0 = 0)$
unfolding *Min-def pred-def* **by** *arith*

function

bin-rec :: $'a \Rightarrow 'a \Rightarrow (\text{int} \Rightarrow \text{bit} \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{int} \Rightarrow 'a$

where

bin-rec *f1 f2 f3 bin* = (*if bin = Int.Pls then f1*
else if bin = Int.Min then f2
else case bin-rl bin of (w, b) => f3 w b (bin-rec f1 f2 f3 w))

by *pat-completeness auto*

termination

apply (*relation measure* (*nat o abs o snd o snd o snd*))
apply *simp*
apply (*simp add: Pls-def brlem*)
apply (*clarsimp simp: bin-rl-char pred-def*)
apply (*frule thin-rl [THEN refl [THEN bin-abs-lem [rule-format]]]*)
apply (*unfold Pls-def Min-def number-of-eq*)
prefer 2
apply (*erule asm-rl*)
apply *auto*
done

declare *bin-rec.simps* [*simp del*]

lemma *bin-rec-PM*:

f = bin-rec f1 f2 f3 ==> f Int.Pls = f1 & f Int.Min = f2
by (*auto simp add: bin-rec.simps*)

lemma *bin-rec-Pls*: *bin-rec f1 f2 f3 Int.Pls = f1*

by (*simp add: bin-rec.simps*)

lemma *bin-rec-Min*: *bin-rec f1 f2 f3 Int.Min = f2*

by (*simp add: bin-rec.simps*)

lemma *bin-rec-Bit0*:

f3 Int.Pls bit.B0 f1 = f1 ==>
bin-rec f1 f2 f3 (Int.Bit0 w) = f3 w bit.B0 (bin-rec f1 f2 f3 w)
by (*simp add: bin-rec-Pls bin-rec.simps [of - - Int.Bit0 w]*)

lemma *bin-rec-Bit1*:

f3 Int.Min bit.B1 f2 = f2 ==>
bin-rec f1 f2 f3 (Int.Bit1 w) = f3 w bit.B1 (bin-rec f1 f2 f3 w)
by (*simp add: bin-rec-Min bin-rec.simps [of - - Int.Bit1 w]*)

lemma *bin-rec-Bit*:

f = bin-rec f1 f2 f3 ==> f3 Int.Pls bit.B0 f1 = f1 ==>
f3 Int.Min bit.B1 f2 = f2 ==> f (w BIT b) = f3 w b (f w)
by (*cases b, simp add: bin-rec-Bit0, simp add: bin-rec-Bit1*)

lemmas *bin-rec-simps = refl [THEN bin-rec-Bit] bin-rec-Pls bin-rec-Min*
bin-rec-Bit0 bin-rec-Bit1

4.4 Truncating binary integers

definition

bin-sign-def [*code del*] : *bin-sign = bin-rec Int.Pls Int.Min (%w b s. s)*

```

lemma bin-sign-simps [simp]:
  bin-sign Int.Pls = Int.Pls
  bin-sign Int.Min = Int.Min
  bin-sign (Int.Bit0 w) = bin-sign w
  bin-sign (Int.Bit1 w) = bin-sign w
  bin-sign (w BIT b) = bin-sign w
unfolding bin-sign-def by (auto simp: bin-rec-simps)

declare bin-sign-simps(1-4) [code]

lemma bin-sign-rest [simp]:
  bin-sign (bin-rest w) = (bin-sign w)
by (cases w rule: bin-exhaust) auto

consts
  bintrunc :: nat => int => int
primrec
  Z : bintrunc 0 bin = Int.Pls
  Suc : bintrunc (Suc n) bin = bintrunc n (bin-rest bin) BIT (bin-last bin)

consts
  sbintrunc :: nat => int => int
primrec
  Z : sbintrunc 0 bin =
    (case bin-last bin of bit.B1 => Int.Min | bit.B0 => Int.Pls)
  Suc : sbintrunc (Suc n) bin = sbintrunc n (bin-rest bin) BIT (bin-last bin)

lemma sign-bintr:
  !!w. bin-sign (bintrunc n w) = Int.Pls
by (induct n) auto

lemma bintrunc-mod2p:
  !!w. bintrunc n w = (w mod 2 ^ n :: int)
apply (induct n, clarsimp)
apply (simp add: bin-last-mod bin-rest-div Bit-def zmod-zmult2-eq
  cong: number-of-False-cong)
done

lemma sbintrunc-mod2p:
  !!w. sbintrunc n w = ((w + 2 ^ n) mod 2 ^ (Suc n) - 2 ^ n :: int)
apply (induct n)
apply clarsimp
apply (subst mod-add-left-eq)
apply (simp add: bin-last-mod)
apply (simp add: number-of-eq)
apply clarsimp
apply (simp add: bin-last-mod bin-rest-div Bit-def
  cong: number-of-False-cong)
apply (clarsimp simp: mod-mult-mult1 [symmetric])

```



```

      zmod-zdiv-equality [THEN diff-eq-eq [THEN iffD2 [THEN sym]]])
  apply (rule trans [symmetric, OF - emep1])
  apply auto
  apply (auto simp: even-def)
done

```

4.5 Simplifications for (s)bintrunc

```

lemma bit-bool:
  (b = (b' = bit.B1)) = (b' = (if b then bit.B1 else bit.B0))
  by (cases b') auto

```

```

lemmas bit-bool1 [simp] = refl [THEN bit-bool [THEN iffD1], symmetric]

```

```

lemma bin-sign-lem:
  !!bin. (bin-sign (sbintrunc n bin) = Int.Min) = bin-nth bin n
  apply (induct n)
  apply (case-tac bin rule: bin-exhaust, case-tac b, auto)+
done

```

```

lemma nth-bintr:
  !!w m. bin-nth (bintrunc m w) n = (n < m & bin-nth w n)
  apply (induct n)
  apply (case-tac m, auto)[1]
  apply (case-tac m, auto)[1]
done

```

```

lemma nth-sbintr:
  !!w m. bin-nth (sbintrunc m w) n =
    (if n < m then bin-nth w n else bin-nth w m)
  apply (induct n)
  apply (case-tac m, simp-all split: bit.splits)[1]
  apply (case-tac m, simp-all split: bit.splits)[1]
done

```

```

lemma bin-nth-Bit:
  bin-nth (w BIT b) n = (n = 0 & b = bit.B1 | (EX m. n = Suc m & bin-nth w m))
  by (cases n) auto

```

```

lemma bin-nth-Bit0:
  bin-nth (Int.Bit0 w) n = (EX m. n = Suc m & bin-nth w m)
  using bin-nth-Bit [where b=bit.B0] by simp

```

```

lemma bin-nth-Bit1:
  bin-nth (Int.Bit1 w) n = (n = 0 | (EX m. n = Suc m & bin-nth w m))
  using bin-nth-Bit [where b=bit.B1] by simp

```

```

lemma bintrunc-bintrunc-l:

```

$n \leq m \implies (\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } n w)$
by (rule bin-eqI) (auto simp add : nth-bintr)

lemma *sbintrunc-sbintrunc-l*:
 $n \leq m \implies (\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } n w)$
by (rule bin-eqI) (auto simp: nth-sbintr)

lemma *bintrunc-bintrunc-ge*:
 $n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w)$
by (rule bin-eqI) (auto simp: nth-bintr)

lemma *bintrunc-bintrunc-min* [simp]:
 $\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$
apply (rule bin-eqI)
apply (auto simp: nth-bintr)
done

lemma *sbintrunc-sbintrunc-min* [simp]:
 $\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$
apply (rule bin-eqI)
apply (auto simp: nth-sbintr min-max.inf-absorb1 min-max.inf-absorb2)
done

lemmas *bintrunc-Pls* =
 bintrunc.Suc [where bin=Int.Pls, simplified bin-last-simps bin-rest-simps, standard]

lemmas *bintrunc-Min* [simp] =
 bintrunc.Suc [where bin=Int.Min, simplified bin-last-simps bin-rest-simps, standard]

lemmas *bintrunc-BIT* [simp] =
 bintrunc.Suc [where bin= $w \text{ BIT } b$, simplified bin-last-simps bin-rest-simps, standard]

lemma *bintrunc-Bit0* [simp]:
 $\text{bintrunc } (\text{Suc } n) (\text{Int.Bit0 } w) = \text{Int.Bit0 } (\text{bintrunc } n w)$
using *bintrunc-BIT* [where $b=\text{bit.B0}$] **by** simp

lemma *bintrunc-Bit1* [simp]:
 $\text{bintrunc } (\text{Suc } n) (\text{Int.Bit1 } w) = \text{Int.Bit1 } (\text{bintrunc } n w)$
using *bintrunc-BIT* [where $b=\text{bit.B1}$] **by** simp

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*
bintrunc-Bit0 bintrunc-Bit1

lemmas *sbintrunc-Suc-Pls* =
 sbintrunc.Suc [where bin=Int.Pls, simplified bin-last-simps bin-rest-simps, standard]

lemmas *sbintrunc-Suc-Min* =
sbintrunc.Suc [where *bin*=*Int.Min*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Suc-BIT* [simp] =
sbintrunc.Suc [where *bin*=*w BIT b*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemma *sbintrunc-Suc-Bit0* [simp]:
sbintrunc (Suc n) (Int.Bit0 w) = Int.Bit0 (sbintrunc n w)
using *sbintrunc-Suc-BIT* [where *b*=*bit.B0*] **by** *simp*

lemma *sbintrunc-Suc-Bit1* [simp]:
sbintrunc (Suc n) (Int.Bit1 w) = Int.Bit1 (sbintrunc n w)
using *sbintrunc-Suc-BIT* [where *b*=*bit.B1*] **by** *simp*

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls sbintrunc-Suc-Min sbintrunc-Suc-BIT*
sbintrunc-Suc-Bit0 sbintrunc-Suc-Bit1

lemmas *sbintrunc-Pls* =
sbintrunc.Z [where *bin*=*Int.Pl*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-Min* =
sbintrunc.Z [where *bin*=*Int.Min*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-0-BIT-B0* [simp] =
sbintrunc.Z [where *bin*=*w BIT bit.B0*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-0-BIT-B1* [simp] =
sbintrunc.Z [where *bin*=*w BIT bit.B1*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemma *sbintrunc-0-Bit0* [simp]: *sbintrunc 0 (Int.Bit0 w) = Int.Pl*
using *sbintrunc-0-BIT-B0* **by** *simp*

lemma *sbintrunc-0-Bit1* [simp]: *sbintrunc 0 (Int.Bit1 w) = Int.Min*
using *sbintrunc-0-BIT-B1* **by** *simp*

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls sbintrunc-Min sbintrunc-0-BIT-B0 sbintrunc-0-BIT-B1
sbintrunc-0-Bit0 sbintrunc-0-Bit1

lemmas *bintrunc-simps* = *bintrunc.Z bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps sbintrunc-Sucs*

lemma *bintrunc-minus*:

$0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$

by *auto*

lemma *sbintrunc-minus*:

$0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$

by *auto*

lemmas *bintrunc-minus-simps* =

bintrunc-Sucs [THEN [2] *bintrunc-minus* [symmetric, THEN trans], standard]

lemmas *sbintrunc-minus-simps* =

sbintrunc-Sucs [THEN [2] *sbintrunc-minus* [symmetric, THEN trans], standard]

lemma *bintrunc-n-Pls* [simp]:

$\text{bintrunc } n \ \text{Int.Pl} = \text{Int.Pl}$

by (induct *n*) *auto*

lemma *sbintrunc-n-PM* [simp]:

$\text{sbintrunc } n \ \text{Int.Pl} = \text{Int.Pl}$

$\text{sbintrunc } n \ \text{Int.Min} = \text{Int.Min}$

by (induct *n*) *auto*

lemmas *thobini1* = *arg-cong* [where $f = \%w. \ w \ \text{BIT } b$, standard]

lemmas *bintrunc-BIT-I* = *trans* [OF *bintrunc-BIT thobini1*]

lemmas *bintrunc-Min-I* = *trans* [OF *bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps*(1-3) [THEN *thobini1* [THEN [2] trans], standard]

lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)

lemmas *bintrunc-Min-minus-I* = *bmsts*(2)

lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-0-Min*: $\text{bintrunc } 0 \ \text{Int.Min} = \text{Int.Pl}$

by *auto*

lemma *bintrunc-0-BIT*: $\text{bintrunc } 0 \ (w \ \text{BIT } b) = \text{Int.Pl}$

by *auto*

lemma *bintrunc-Suc-lem*:

$\text{bintrunc } (\text{Suc } n) \ x = y \implies m = \text{Suc } n \implies \text{bintrunc } m \ x = y$

by *auto*

lemmas *bintrunc-Suc-Ialts* =

bintrunc-Min-I [THEN *bintrunc-Suc-lem*, standard]

bintrunc-BIT-I [THEN *bintrunc-Suc-lem*, standard]

lemmas *sbintrunc-BIT-I* = *trans* [OF *sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs(1-3) [THEN thobini1 [THEN [2] trans], standard]

lemmas *sbintrunc-Suc-minus-Is* =
sbintrunc-minus-simps(1-3) [THEN thobini1 [THEN [2] trans], standard]

lemma *sbintrunc-Suc-lem*:
 $sbintrunc (Suc\ n)\ x = y \implies m = Suc\ n \implies sbintrunc\ m\ x = y$
by *auto*

lemmas *sbintrunc-Suc-Ialts* =
sbintrunc-Suc-Is [THEN *sbintrunc-Suc-lem*, standard]

lemma *sbintrunc-bintrunc-lt*:
 $m > n \implies sbintrunc\ n\ (bintrunc\ m\ w) = sbintrunc\ n\ w$
by (rule *bin-eqI*) (auto simp: *nth-sbintr nth-bintr*)

lemma *bintrunc-sbintrunc-le*:
 $m \leq Suc\ n \implies bintrunc\ m\ (sbintrunc\ n\ w) = bintrunc\ m\ w$
apply (rule *bin-eqI*)
apply (auto simp: *nth-sbintr nth-bintr*)
apply (subgoal-tac $x=n$, safe, arith+)[1]
apply (subgoal-tac $x=n$, safe, arith+)[1]
done

lemmas *bintrunc-sbintrunc [simp]* = *order-refl* [THEN *bintrunc-sbintrunc-le*]
lemmas *sbintrunc-bintrunc [simp]* = *lessI* [THEN *sbintrunc-bintrunc-lt*]
lemmas *bintrunc-bintrunc [simp]* = *order-refl* [THEN *bintrunc-bintrunc-l*]
lemmas *sbintrunc-sbintrunc [simp]* = *order-refl* [THEN *sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc' [simp]*:
 $0 < n \implies bintrunc\ n\ (sbintrunc\ (n - 1)\ w) = bintrunc\ n\ w$
by (cases n) (auto simp del: *bintrunc.Suc*)

lemma *sbintrunc-bintrunc' [simp]*:
 $0 < n \implies sbintrunc\ (n - 1)\ (bintrunc\ n\ w) = sbintrunc\ (n - 1)\ w$
by (cases n) (auto simp del: *bintrunc.Suc*)

lemma *bin-sbin-eq-iff*:
 $bintrunc\ (Suc\ n)\ x = bintrunc\ (Suc\ n)\ y \iff$
 $sbintrunc\ n\ x = sbintrunc\ n\ y$
apply (rule *iffI*)
apply (rule *box-equals* [OF - *sbintrunc-bintrunc sbintrunc-bintrunc*])
apply *simp*
apply (rule *box-equals* [OF - *bintrunc-sbintrunc bintrunc-sbintrunc*])
apply *simp*
done

lemma *bin-sbin-eq-iff'*:
 $0 < n \implies bintrunc\ n\ x = bintrunc\ n\ y \iff$

$$sbintrunc\ (n - 1)\ x = sbintrunc\ (n - 1)\ y$$
by (cases n) (simp-all add: bin-sbin-eq-iff del: bintrunc.Suc)

lemmas bintrunc-sbintruncS0 [simp] = bintrunc-sbintrunc' [unfolded One-nat-def]
lemmas sbintrunc-bintruncS0 [simp] = sbintrunc-bintrunc' [unfolded One-nat-def]

lemmas bintrunc-bintrunc-l' = le-add1 [THEN bintrunc-bintrunc-l]
lemmas sbintrunc-sbintrunc-l' = le-add1 [THEN sbintrunc-sbintrunc-l]

lemmas nat-non0-gr =
 trans [OF iszero-def [THEN Not-eq-iff [THEN iffD2]] refl, standard]

lemmas bintrunc-pred-simps [simp] =
 bintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemmas sbintrunc-pred-simps [simp] =
 sbintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemma no-bintr-alt:
 number-of (bintrunc n w) = w mod 2 ^ n
by (simp add: number-of-eq bintrunc-mod2p)

lemma no-bintr-alt1: bintrunc n = (%w. w mod 2 ^ n :: int)
by (rule ext) (rule bintrunc-mod2p)

lemma range-bintrunc: range (bintrunc n) = {i. 0 <= i & i < 2 ^ n}
apply (unfold no-bintr-alt1)
apply (auto simp add: image-iff)
apply (rule exI)
apply (auto intro: int-mod-lem [THEN iffD1, symmetric])
done

lemma no-bintr:
 number-of (bintrunc n w) = (number-of w mod 2 ^ n :: int)
by (simp add : bintrunc-mod2p number-of-eq)

lemma no-sbintr-alt2:
 sbintrunc n = (%w. (w + 2 ^ n) mod 2 ^ Suc n - 2 ^ n :: int)
by (rule ext) (simp add : sbintrunc-mod2p)

lemma no-sbintr:
 number-of (sbintrunc n w) =
 ((number-of w + 2 ^ n) mod 2 ^ Suc n - 2 ^ n :: int)
by (simp add : no-sbintr-alt2 number-of-eq)

lemma range-sbintrunc:
 range (sbintrunc n) = {i. - (2 ^ n) <= i & i < 2 ^ n}

```

apply (unfold no-sbintr-alt2)
apply (auto simp add: image-iff eq-diff-eq)
apply (rule exI)
apply (auto intro: int-mod-lem [THEN iffD1, symmetric])
done

```

```

lemma sb-inc-lem:
   $(a::int) + 2^k < 0 \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$ 
  apply (erule int-mod-ge' [where n = 2 ^ (Suc k) and b = a + 2 ^ k, simplified
zless2p])
  apply (rule TrueI)
done

```

```

lemma sb-inc-lem':
   $(a::int) < - (2^k) \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$ 
  by (rule sb-inc-lem) simp

```

```

lemma sbintrunc-inc:
   $x < - (2^n) \implies x + 2^{(Suc\ n)} \leq sbintrunc\ n\ x$ 
  unfolding no-sbintr-alt2 by (drule sb-inc-lem') simp

```

```

lemma sb-dec-lem:
   $(0::int) \leq - (2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$ 
  by (rule int-mod-le' [where n = 2 ^ (Suc k) and b = a + 2 ^ k,
simplified zless2p, OF - TrueI, simplified])

```

```

lemma sb-dec-lem':
   $(2::int) ^ k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$ 
  by (rule iffD1 [OF diff-le-eq', THEN sb-dec-lem, simplified])

```

```

lemma sbintrunc-dec:
   $x \geq (2^n) \implies x - 2^{(Suc\ n)} \geq sbintrunc\ n\ x$ 
  unfolding no-sbintr-alt2 by (drule sb-dec-lem') simp

```

```

lemmas zmod-uminus' = zmod-uminus [where b=c, standard]
lemmas zpower-zmod' = zpower-zmod [where m=c and y=k, standard]

```

```

lemmas brdmod1s' [symmetric] =
  mod-add-left-eq mod-add-right-eq
  zmod-zsub-left-eq zmod-zsub-right-eq
  zmod-zmult1-eq zmod-zmult1-eq-rev

```

```

lemmas brdmods' [symmetric] =
  zpower-zmod' [symmetric]
  trans [OF mod-add-left-eq mod-add-right-eq]
  trans [OF zmod-zsub-left-eq zmod-zsub-right-eq]
  trans [OF zmod-zmult1-eq zmod-zmult1-eq-rev]
  zmod-uminus' [symmetric]
  mod-add-left-eq [where b = 1::int]

```

zmod-zsub-left-eq [where $b = 1$]

lemmas *bintr-arith1s* =
brdmod1s' [where $c=2^n::int$, folded pred-def succ-def bintrunc-mod2p, standard]
lemmas *bintr-ariths* =
brdmods' [where $c=2^n::int$, folded pred-def succ-def bintrunc-mod2p, standard]

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [OF *zless2p*, standard]

lemma *bintr-ge0*: $(0 :: int) \leq \text{number-of } (\text{bintrunc } n \ w)$
by (*simp add : no-bintr m2pths*)

lemma *bintr-lt2p*: $\text{number-of } (\text{bintrunc } n \ w) < (2^n :: int)$
by (*simp add : no-bintr m2pths*)

lemma *bintr-Min*:
 $\text{number-of } (\text{bintrunc } n \ \text{Int.Min}) = (2^n :: int) - 1$
by (*simp add : no-bintr m1mod2k*)

lemma *sbintr-ge*: $(-(2^n :: int) \leq \text{number-of } (\text{sbintrunc } n \ w))$
by (*simp add : no-sbintr m2pths*)

lemma *sbintr-lt*: $\text{number-of } (\text{sbintrunc } n \ w) < (2^n :: int)$
by (*simp add : no-sbintr m2pths*)

lemma *bintrunc-Suc*:
 $\text{bintrunc } (\text{Suc } n) \ bin = \text{bintrunc } n \ (\text{bin-rest } bin) \ \text{BIT } \text{bin-last } bin$
by (*case-tac bin rule: bin-exhaust*) *auto*

lemma *sign-Pls-ge-0*:
 $(\text{bin-sign } bin = \text{Int.Pl}) = (\text{number-of } bin \geq (0 :: int))$
by (*induct bin rule: numeral-induct*) *auto*

lemma *sign-Min-lt-0*:
 $(\text{bin-sign } bin = \text{Int.Min}) = (\text{number-of } bin < (0 :: int))$
by (*induct bin rule: numeral-induct*) *auto*

lemmas *sign-Min-neg* = *trans* [OF *sign-Min-lt-0 neg-def* [*symmetric*]]

lemma *bin-rest-trunc*:
 $!!bin. (\text{bin-rest } (\text{bintrunc } n \ bin)) = \text{bintrunc } (n - 1) \ (\text{bin-rest } bin)$
by (*induct n*) *auto*

lemma *bin-rest-power-trunc* [rule-format] :
 $(\text{bin-rest } ^k) (\text{bintrunc } n \ bin) =$
 $\text{bintrunc } (n - k) \ ((\text{bin-rest } ^k) \ bin)$
by (*induct k*) (*auto simp: bin-rest-trunc*)

lemma *bin-rest-trunc-i*:


```

bintrunc n (bin-rest bin) = bin-rest (bintrunc (Suc n) bin)
by auto

```

```

lemma bin-rest-strunc:
!!bin. bin-rest (sbintrunc (Suc n) bin) = sbintrunc n (bin-rest bin)
by (induct n) auto

```

```

lemma bintrunc-rest [simp]:
!!bin. bintrunc n (bin-rest (bintrunc n bin)) = bin-rest (bintrunc n bin)
apply (induct n, simp)
apply (case-tac bin rule: bin-exhaust)
apply (auto simp: bintrunc-bintrunc-l)
done

```

```

lemma sbintrunc-rest [simp]:
!!bin. sbintrunc n (bin-rest (sbintrunc n bin)) = bin-rest (sbintrunc n bin)
apply (induct n, simp)
apply (case-tac bin rule: bin-exhaust)
apply (auto simp: bintrunc-bintrunc-l split: bit.splits)
done

```

```

lemma bintrunc-rest':
bintrunc n o bin-rest o bintrunc n = bin-rest o bintrunc n
by (rule ext) auto

```

```

lemma sbintrunc-rest' :
sbintrunc n o bin-rest o sbintrunc n = bin-rest o sbintrunc n
by (rule ext) auto

```

```

lemma rco-lem:
f o g o f = g o f ==> f o (g o f) ^ n = g ^ n o f
apply (rule ext)
apply (induct-tac n)
apply (simp-all (no-asm))
apply (drule fun-cong)
apply (unfold o-def)
apply (erule trans)
apply simp
done

```

```

lemma rco-alt: (f o g) ^ n o f = f o (g o f) ^ n
apply (rule ext)
apply (induct n)
apply (simp-all add: o-def)
done

```

```

lemmas rco-bintr = bintrunc-rest'
[THEN rco-lem [THEN fun-cong], unfolded o-def]
lemmas rco-sbintr = sbintrunc-rest'

```

[*THEN* *rco-lem* [*THEN* *fun-cong*], *unfolded o-def*]

4.6 Splitting and concatenation

primrec *bin-split* :: *nat* \Rightarrow *int* \Rightarrow *int* \times *int* **where**
Z: *bin-split* 0 *w* = (*w*, *Int.Pls*)
| *Suc*: *bin-split* (*Suc* *n*) *w* = (*let* (*w1*, *w2*) = *bin-split* *n* (*bin-rest* *w*)
in (*w1*, *w2* BIT *bin-last* *w*))

primrec *bin-cat* :: *int* \Rightarrow *nat* \Rightarrow *int* \Rightarrow *int* **where**
Z: *bin-cat* *w* 0 *v* = *w*
| *Suc*: *bin-cat* *w* (*Suc* *n*) *v* = *bin-cat* *w* *n* (*bin-rest* *v*) BIT *bin-last* *v*

4.7 Miscellaneous lemmas

lemma *funpow-minus-simp*:
 $0 < n \implies f^{^^n} = f \circ f^{^^(n-1)}$
by (*cases* *n*) *simp-all*

lemmas *funpow-pred-simp* [*simp*] =
funpow-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *replicate-minus-simp* =
trans [*OF gen-minus* [**where** *f* = %*n*. *replicate* *n* *x*] *replicate.replicate-Suc*,
standard]

lemmas *replicate-pred-simp* [*simp*] =
replicate-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *power-Suc-no* [*simp*] = *power-Suc* [*of number-of a, standard*]

lemmas *power-minus-simp* =
trans [*OF gen-minus* [**where** *f* = *power* *f*] *power-Suc*, *standard*]

lemmas *power-pred-simp* =
power-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *power-pred-simp-no* [*simp*] = *power-pred-simp* [**where** *f* = *number-of* *f*,
standard]

lemma *list-exhaust-size-gt0*:
assumes *y*: $\bigwedge a \text{ list. } y = a \# \text{ list} \implies P$
shows $0 < \text{length } y \implies P$
apply (*cases* *y*, *simp*)
apply (*rule* *y*)
apply *fastsimp*
done

lemma *list-exhaust-size-eq0*:
assumes *y*: $y = [] \implies P$
shows $\text{length } y = 0 \implies P$

```

apply (cases y)
  apply (rule y, simp)
apply simp
done

lemma size-Cons-lem-eq:
   $y = xa \# list \implies size\ y = Suc\ k \implies size\ list = k$ 
by auto

lemma size-Cons-lem-eq-bin:
   $y = xa \# list \implies size\ y = number-of\ (Int.succ\ k) \implies$ 
   $size\ list = number-of\ k$ 
by (auto simp: pred-def succ-def split add : split-if-asm)

lemmas ls-splits =
  prod.split split-split prod.split-asm split-split-asm split-if-asm

lemma not-B1-is-B0:  $y \neq bit.B1 \implies y = bit.B0$ 
by (cases y) auto

lemma B1-ass-B0:
  assumes  $y: y = bit.B0 \implies y = bit.B1$ 
shows  $y = bit.B1$ 
apply (rule classical)
apply (drule not-B1-is-B0)
apply (erule y)
done

— simplifications for specific word lengths
lemmas n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'

lemmas s2n-ths = n2s-ths [symmetric]

end

```

5 BitSyntax: Syntactic classes for bitwise operations

```

theory BitSyntax
imports BinGeneral
begin

class bit =
  fixes bitNOT :: 'a  $\Rightarrow$  'a (NOT - [70] 71)
  and bitAND :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr AND 64)
  and bitOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr OR 59)
  and bitXOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr XOR 59)

```

We want the bitwise operations to bind slightly weaker than $+$ and $-$, but $\sim\sim$ to bind slightly stronger than $*$.

Testing and shifting operations.

```

class bits = bit +
  fixes test-bit :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool (infixl !! 100)
    and lsb      :: 'a  $\Rightarrow$  bool
    and set-bit  :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool  $\Rightarrow$  'a
    and set-bits :: (nat  $\Rightarrow$  bool)  $\Rightarrow$  'a (binder BITS 10)
    and shiftr  :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl << 55)
    and shiftr  :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl >> 55)

```

```

class bitss = bits +
  fixes msb      :: 'a  $\Rightarrow$  bool

```

5.1 Bitwise operations on *bit*

```

instantiation bit :: bit
begin

```

```

primrec bitNOT-bit where
  NOT bit.B0 = bit.B1
  | NOT bit.B1 = bit.B0

```

```

primrec bitAND-bit where
  bit.B0 AND y = bit.B0
  | bit.B1 AND y = y

```

```

primrec bitOR-bit where
  bit.B0 OR y = y
  | bit.B1 OR y = bit.B1

```

```

primrec bitXOR-bit where
  bit.B0 XOR y = y
  | bit.B1 XOR y = NOT y

```

```

instance ..

```

```

end

```

```

lemmas bit-simps =
  bitNOT-bit.simps bitAND-bit.simps bitOR-bit.simps bitXOR-bit.simps

```

```

lemma bit-extra-simps [simp]:
  x AND bit.B0 = bit.B0
  x AND bit.B1 = x
  x OR bit.B1 = bit.B1
  x OR bit.B0 = x
  x XOR bit.B1 = NOT x

```

```

 $x \text{ XOR } \text{bit.B0} = x$ 
by (cases  $x$ , auto)+

lemma bit-ops-comm:
  ( $x::\text{bit}$ )  $\text{AND } y = y \text{ AND } x$ 
  ( $x::\text{bit}$ )  $\text{OR } y = y \text{ OR } x$ 
  ( $x::\text{bit}$ )  $\text{XOR } y = y \text{ XOR } x$ 
  by (cases  $y$ , auto)+

lemma bit-ops-same [simp]:
  ( $x::\text{bit}$ )  $\text{AND } x = x$ 
  ( $x::\text{bit}$ )  $\text{OR } x = x$ 
  ( $x::\text{bit}$ )  $\text{XOR } x = \text{bit.B0}$ 
  by (cases  $x$ , auto)+

lemma bit-not-not [simp]:  $\text{NOT } (\text{NOT } (x::\text{bit})) = x$ 
  by (cases  $x$ ) auto

end

```

6 BinOperations: Bitwise Operations on Binary Integers

```

theory BinOperations
imports BinGeneral BitSyntax
begin

```

6.1 Logical operations

bit-wise logical operations on the int type

```

instantiation int :: bit
begin

```

```

definition
  int-not-def [code del]: bitNOT = bin-rec Int.Min Int.Pls
    ( $\lambda w \ b \ s. \ s \ \text{BIT } (\text{NOT } b)$ )

```

```

definition
  int-and-def [code del]: bitAND = bin-rec ( $\lambda x. \text{Int.Pls}$ ) ( $\lambda y. \ y$ )
    ( $\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT } (b \ \text{AND } \text{bin-last } y)$ )

```

```

definition
  int-or-def [code del]: bitOR = bin-rec ( $\lambda x. \ x$ ) ( $\lambda y. \ \text{Int.Min}$ )
    ( $\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT } (b \ \text{OR } \text{bin-last } y)$ )

```

```

definition
  int-xor-def [code del]: bitXOR = bin-rec ( $\lambda x. \ x$ ) bitNOT

```

$$(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT } (b \ \text{XOR } \text{bin-last } y))$$

instance ..

end

lemma *int-not-simps* [*simp*]:
 $\text{NOT } \text{Int.Pls} = \text{Int.Min}$
 $\text{NOT } \text{Int.Min} = \text{Int.Pls}$
 $\text{NOT } (\text{Int.Bit0 } w) = \text{Int.Bit1 } (\text{NOT } w)$
 $\text{NOT } (\text{Int.Bit1 } w) = \text{Int.Bit0 } (\text{NOT } w)$
 $\text{NOT } (w \ \text{BIT } b) = (\text{NOT } w) \ \text{BIT } (\text{NOT } b)$
unfolding *int-not-def* **by** (*simp-all add: bin-rec-simps*)

declare *int-not-simps*(1–4) [*code*]

lemma *int-xor-Pls* [*simp, code*]:
 $\text{Int.Pls XOR } x = x$
unfolding *int-xor-def* **by** (*simp add: bin-rec-PM*)

lemma *int-xor-Min* [*simp, code*]:
 $\text{Int.Min XOR } x = \text{NOT } x$
unfolding *int-xor-def* **by** (*simp add: bin-rec-PM*)

lemma *int-xor-Bits* [*simp*]:
 $(x \ \text{BIT } b) \ \text{XOR } (y \ \text{BIT } c) = (x \ \text{XOR } y) \ \text{BIT } (b \ \text{XOR } c)$
apply (*unfold int-xor-def*)
apply (*rule bin-rec-simps (1) [THEN fun-cong, THEN trans]*)
apply (*rule ext, simp*)
prefer 2
apply *simp*
apply (*rule ext*)
apply (*simp add: int-not-simps [symmetric]*)
done

lemma *int-xor-Bits2* [*simp, code*]:
 $(\text{Int.Bit0 } x) \ \text{XOR } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \ \text{XOR } y)$
 $(\text{Int.Bit0 } x) \ \text{XOR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \ \text{XOR } y)$
 $(\text{Int.Bit1 } x) \ \text{XOR } (\text{Int.Bit0 } y) = \text{Int.Bit1 } (x \ \text{XOR } y)$
 $(\text{Int.Bit1 } x) \ \text{XOR } (\text{Int.Bit1 } y) = \text{Int.Bit0 } (x \ \text{XOR } y)$
unfolding *BIT-simps [symmetric]* *int-xor-Bits* **by** *simp-all*

lemma *int-xor-x-simps'*:
 $w \ \text{XOR } (\text{Int.Pls BIT bit.B0}) = w$
 $w \ \text{XOR } (\text{Int.Min BIT bit.B1}) = \text{NOT } w$
apply (*induct w rule: bin-induct*)
apply *simp-all[4]*
apply (*unfold int-xor-Bits*)
apply *clarsimp+*

done

lemma *int-xor-extra-simps* [*simp*, *code*]:

$w \text{ XOR } \text{Int.Pls} = w$

$w \text{ XOR } \text{Int.Min} = \text{NOT } w$

using *int-xor-x-simps'* **by** *simp-all*

lemma *int-or-Pls* [*simp*, *code*]:

$\text{Int.Pls OR } x = x$

by (*unfold int-or-def*) (*simp add: bin-rec-PM*)

lemma *int-or-Min* [*simp*, *code*]:

$\text{Int.Min OR } x = \text{Int.Min}$

by (*unfold int-or-def*) (*simp add: bin-rec-PM*)

lemma *int-or-Bits* [*simp*]:

$(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \text{ OR } c)$

unfolding *int-or-def* **by** (*simp add: bin-rec-simps*)

lemma *int-or-Bits2* [*simp*, *code*]:

$(\text{Int.Bit0 } x) \text{ OR } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ OR } y)$

$(\text{Int.Bit0 } x) \text{ OR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ OR } y)$

$(\text{Int.Bit1 } x) \text{ OR } (\text{Int.Bit0 } y) = \text{Int.Bit1 } (x \text{ OR } y)$

$(\text{Int.Bit1 } x) \text{ OR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ OR } y)$

unfolding *BIT-simps* [*symmetric*] *int-or-Bits* **by** *simp-all*

lemma *int-or-x-simps'*:

$w \text{ OR } (\text{Int.Pls BIT bit.B0}) = w$

$w \text{ OR } (\text{Int.Min BIT bit.B1}) = \text{Int.Min}$

apply (*induct w rule: bin-induct*)

apply *simp-all*[4]

apply (*unfold int-or-Bits*)

apply *clarsimp+*

done

lemma *int-or-extra-simps* [*simp*, *code*]:

$w \text{ OR } \text{Int.Pls} = w$

$w \text{ OR } \text{Int.Min} = \text{Int.Min}$

using *int-or-x-simps'* **by** *simp-all*

lemma *int-and-Pls* [*simp*, *code*]:

$\text{Int.Pls AND } x = \text{Int.Pls}$

unfolding *int-and-def* **by** (*simp add: bin-rec-PM*)

lemma *int-and-Min* [*simp*, *code*]:

$\text{Int.Min AND } x = x$

unfolding *int-and-def* **by** (*simp add: bin-rec-PM*)

lemma *int-and-Bits* [*simp*]:

$(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \text{ AND } c)$
unfolding *int-and-def* **by** (*simp add: bin-rec-simps*)

lemma *int-and-Bits2* [*simp, code*]:
 $(\text{Int.Bit0 } x) \text{ AND } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit0 } x) \text{ AND } (\text{Int.Bit1 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit1 } x) \text{ AND } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit1 } x) \text{ AND } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ AND } y)$
unfolding *BIT-simps* [*symmetric*] *int-and-Bits* **by** *simp-all*

lemma *int-and-x-simps'*:
 $w \text{ AND } (\text{Int.Pls BIT bit.B0}) = \text{Int.Pls}$
 $w \text{ AND } (\text{Int.Min BIT bit.B1}) = w$
apply (*induct w rule: bin-induct*)
apply *simp-all*[4]
apply (*unfold int-and-Bits*)
apply *clarsimp+*
done

lemma *int-and-extra-simps* [*simp, code*]:
 $w \text{ AND } \text{Int.Pls} = \text{Int.Pls}$
 $w \text{ AND } \text{Int.Min} = w$
using *int-and-x-simps'* **by** *simp-all*

lemma *bin-ops-comm*:
shows
int-and-comm: $!!y::\text{int}. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $!!y::\text{int}. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $!!y::\text{int}. x \text{ XOR } y = y \text{ XOR } x$
apply (*induct x rule: bin-induct*)
apply *simp-all*[6]
apply (*case-tac y rule: bin-exhaust, simp add: bit-ops-comm*)
done

lemma *bin-ops-same* [*simp*]:
 $(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$
 $(x::\text{int}) \text{ XOR } x = \text{Int.Pls}$
by (*induct x rule: bin-induct*) *auto*

lemma *int-not-not* [*simp*]: $\text{NOT } (\text{NOT } (x::\text{int})) = x$
by (*induct x rule: bin-induct*) *auto*

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min

lemma *bbw-ao-absorb*:

```
!!y::int. x AND (y OR x) = x & x OR (y AND x) = x
apply (induct x rule: bin-induct)
  apply auto
  apply (case-tac [!] y rule: bin-exhaust)
  apply auto
  apply (case-tac [!] bit)
  apply auto
done
```

lemma *bbw-ao-absorbs-other*:

```
x AND (x OR y) = x ∧ (y AND x) OR x = (x::int)
(y OR x) AND x = x ∧ x OR (x AND y) = (x::int)
(x OR y) AND x = x ∧ (x AND y) OR x = (x::int)
apply (auto simp: bbw-ao-absorb int-or-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (simp add: bbw-ao-absorb)
done
```

lemmas *bbw-ao-absorbs [simp] = bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:

```
!!y::int. (NOT x) XOR y = NOT (x XOR y) &
  x XOR (NOT y) = NOT (x XOR y)
apply (induct x rule: bin-induct)
  apply auto
  apply (case-tac y rule: bin-exhaust, auto,
    case-tac b, auto)+
done
```

lemma *bbw-assocs'*:

```
!!y z::int. (x AND y) AND z = x AND (y AND z) &
  (x OR y) OR z = x OR (y OR z) &
  (x XOR y) XOR z = x XOR (y XOR z)
apply (induct x rule: bin-induct)
  apply (auto simp: int-xor-not)
  apply (case-tac [!] y rule: bin-exhaust)
  apply (case-tac [!] z rule: bin-exhaust)
  apply (case-tac [!] bit)
  apply (case-tac [!] b)
  apply (auto simp del: BIT-simps)
done
```

lemma *int-and-assoc*:

$(x \text{ AND } y) \text{ AND } (z::\text{int}) = x \text{ AND } (y \text{ AND } z)$

by (*simp add: bbw-assocs'*)

lemma *int-or-assoc*:

$(x \text{ OR } y) \text{ OR } (z::\text{int}) = x \text{ OR } (y \text{ OR } z)$

by (*simp add: bbw-assocs'*)

lemma *int-xor-assoc*:

$(x \text{ XOR } y) \text{ XOR } (z::\text{int}) = x \text{ XOR } (y \text{ XOR } z)$

by (*simp add: bbw-assocs'*)

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [*simp*]:

$(y::\text{int}) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$

$(y::\text{int}) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$

$(y::\text{int}) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$

apply (*auto simp: bbw-assocs [symmetric]*)

apply (*auto simp: bin-ops-comm*)

done

lemma *bbw-not-dist*:

$!!y::\text{int}. \text{NOT } (x \text{ OR } y) = (\text{NOT } x) \text{ AND } (\text{NOT } y)$

$!!y::\text{int}. \text{NOT } (x \text{ AND } y) = (\text{NOT } x) \text{ OR } (\text{NOT } y)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac [!] y rule: bin-exhaust*)

apply (*case-tac [!] bit, auto simp del: BIT-simps*)

done

lemma *bbw-oa-dist*:

$!!y \ z::\text{int}. (x \text{ AND } y) \text{ OR } z =$

$(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac y rule: bin-exhaust*)

apply (*case-tac z rule: bin-exhaust*)

apply (*case-tac ba, auto simp del: BIT-simps*)

done

lemma *bbw-ao-dist*:

$!!y \ z::\text{int}. (x \text{ OR } y) \text{ AND } z =$

$(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac y rule: bin-exhaust*)

apply (*case-tac z rule: bin-exhaust*)

apply (*case-tac ba, auto simp del: BIT-simps*)

done

```

lemma plus-and-or [rule-format]:
  ALL y::int. (x AND y) + (x OR y) = x + y
apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
  apply clarsimp
  apply (case-tac y rule: bin-exhaust)
  apply clarsimp
  apply (unfold Bit-def)
  apply clarsimp
  apply (erule-tac x = x in allE)
  apply (simp split: bit.split)
done

```

```

lemma le-int-or:
  !!x. bin-sign y = Int.Pls ==> x <= x OR y
apply (induct y rule: bin-induct)
  apply clarsimp
  apply clarsimp
  apply (case-tac x rule: bin-exhaust)
  apply (case-tac b)
  apply (case-tac [!] bit)
  apply (auto simp: less-eq-int-code)
done

```

```

lemmas int-and-le =
  xtr3 [OF bbw-ao-absorbs (2) [THEN conjunct2, symmetric] le-int-or]

```

```

lemma bin-nth-ops:
  !!x y. bin-nth (x AND y) n = (bin-nth x n & bin-nth y n)
  !!x y. bin-nth (x OR y) n = (bin-nth x n | bin-nth y n)
  !!x y. bin-nth (x XOR y) n = (bin-nth x n ~ bin-nth y n)
  !!x. bin-nth (NOT x) n = (~ bin-nth x n)
apply (induct n)
  apply safe
  apply (case-tac [!] x rule: bin-exhaust)
  apply (simp-all del: BIT-simps)
  apply (case-tac [!] y rule: bin-exhaust)
  apply (simp-all del: BIT-simps)
  apply (auto dest: not-B1-is-B0 intro: B1-ass-B0)
done

```

```

lemma bin-add-not: x + NOT x = Int.Min

```

```

apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
apply (case-tac bit, auto)
done

```

```

lemma bin-trunc-ao:
  !!x y. (bintrunc n x) AND (bintrunc n y) = bintrunc n (x AND y)
  !!x y. (bintrunc n x) OR (bintrunc n y) = bintrunc n (x OR y)
apply (induct n)
  apply auto
  apply (case-tac [!] x rule: bin-exhaust)
  apply (case-tac [!] y rule: bin-exhaust)
  apply auto
done

```

```

lemma bin-trunc-xor:
  !!x y. bintrunc n (bintrunc n x XOR bintrunc n y) =
    bintrunc n (x XOR y)
apply (induct n)
  apply auto
  apply (case-tac [!] x rule: bin-exhaust)
  apply (case-tac [!] y rule: bin-exhaust)
  apply auto
done

```

```

lemma bin-trunc-not:
  !!x. bintrunc n (NOT (bintrunc n x)) = bintrunc n (NOT x)
apply (induct n)
  apply auto
  apply (case-tac [!] x rule: bin-exhaust)
  apply auto
done

```

```

lemma bintr-bintr-i:
  x = bintrunc n y ==> bintrunc n x = bintrunc n y
by auto

```

```

lemmas bin-trunc-and = bin-trunc-ao(1) [THEN bintr-bintr-i]
lemmas bin-trunc-or = bin-trunc-ao(2) [THEN bintr-bintr-i]

```

6.2 Setting and clearing bits

```

primrec
  bin-sc :: nat => bit => int => int
where
  Z: bin-sc 0 b w = bin-rest w BIT b

```

| *Suc*: $\text{bin-sc } (\text{Suc } n) \ b \ w = \text{bin-sc } n \ b \ (\text{bin-rest } w) \ \text{BIT } \text{bin-last } w$

lemma *bin-nth-sc [simp]*:

!!*w*. $\text{bin-nth } (\text{bin-sc } n \ b \ w) \ n = (b = \text{bit.B1})$
by (*induct n*) *auto*

lemma *bin-sc-sc-same [simp]*:

!!*w*. $\text{bin-sc } n \ c \ (\text{bin-sc } n \ b \ w) = \text{bin-sc } n \ c \ w$
by (*induct n*) *auto*

lemma *bin-sc-sc-diff*:

!!*w m*. $m \sim n ==>$
 $\text{bin-sc } m \ c \ (\text{bin-sc } n \ b \ w) = \text{bin-sc } n \ b \ (\text{bin-sc } m \ c \ w)$
apply (*induct n*)
apply (*case-tac* [!] *m*)
apply *auto*
done

lemma *bin-nth-sc-gen*:

!!*w m*. $\text{bin-nth } (\text{bin-sc } n \ b \ w) \ m = (\text{if } m = n \text{ then } b = \text{bit.B1} \text{ else } \text{bin-nth } w \ m)$
by (*induct n*) (*case-tac* [!] *m*, *auto*)

lemma *bin-sc-nth [simp]*:

!!*w*. $(\text{bin-sc } n \ (\text{If } (\text{bin-nth } w \ n) \ \text{bit.B1} \ \text{bit.B0}) \ w) = w$
by (*induct n*) *auto*

lemma *bin-sign-sc [simp]*:

!!*w*. $\text{bin-sign } (\text{bin-sc } n \ b \ w) = \text{bin-sign } w$
by (*induct n*) *auto*

lemma *bin-sc-bintr [simp]*:

!!*w m*. $\text{bintrunc } m \ (\text{bin-sc } n \ x \ (\text{bintrunc } m \ (w))) = \text{bintrunc } m \ (\text{bin-sc } n \ x \ w)$
apply (*induct n*)
apply (*case-tac* [!] *w* *rule*: *bin-exhaust*)
apply (*case-tac* [!] *m*, *auto*)
done

lemma *bin-clr-le*:

!!*w*. $\text{bin-sc } n \ \text{bit.B0} \ w \leq w$
apply (*induct n*)
apply (*case-tac* [!] *w* *rule*: *bin-exhaust*)
apply (*auto simp del*: *BIT-simps*)
apply (*unfold Bit-def*)
apply (*simp-all split*: *bit.split*)
done

lemma *bin-set-ge*:

```

!!w. bin-sc n bit.B1 w >= w
apply (induct n)
apply (case-tac [!] w rule: bin-exhaust)
apply (auto simp del: BIT-simps)
apply (unfold Bit-def)
apply (simp-all split: bit.split)
done

```

```

lemma bintr-bin-clr-le:
!!w m. bintrunc n (bin-sc m bit.B0 w) <= bintrunc n w
apply (induct n)
apply simp
apply (case-tac w rule: bin-exhaust)
apply (case-tac m)
apply (auto simp del: BIT-simps)
apply (unfold Bit-def)
apply (simp-all split: bit.split)
done

```

```

lemma bintr-bin-set-ge:
!!w m. bintrunc n (bin-sc m bit.B1 w) >= bintrunc n w
apply (induct n)
apply simp
apply (case-tac w rule: bin-exhaust)
apply (case-tac m)
apply (auto simp del: BIT-simps)
apply (unfold Bit-def)
apply (simp-all split: bit.split)
done

```

```

lemma bin-sc-FP [simp]: bin-sc n bit.B0 Int.Pls = Int.Pls
by (induct n) auto

```

```

lemma bin-sc-TM [simp]: bin-sc n bit.B1 Int.Min = Int.Min
by (induct n) auto

```

```

lemmas bin-sc-simps = bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP

```

```

lemma bin-sc-minus:
0 < n ==> bin-sc (Suc (n - 1)) b w = bin-sc n b w
by auto

```

```

lemmas bin-sc-Suc-minus =
trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]

```

```

lemmas bin-sc-Suc-pred [simp] =
bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]

```

6.3 Operations on lists of booleans

primrec *bl-to-bin-aux* :: *bool list* \Rightarrow *int* \Rightarrow *int* **where**
Nil: *bl-to-bin-aux* [] *w* = *w*
| *Cons*: *bl-to-bin-aux* (*b* # *bs*) *w* =
bl-to-bin-aux *bs* (*w* BIT (if *b* then *bit.B1* else *bit.B0*))

definition *bl-to-bin* :: *bool list* \Rightarrow *int* **where**
bl-to-bin-def : *bl-to-bin* *bs* = *bl-to-bin-aux* *bs* *Int.Pls*

primrec *bin-to-bl-aux* :: *nat* \Rightarrow *int* \Rightarrow *bool list* \Rightarrow *bool list* **where**
Z: *bin-to-bl-aux* 0 *w* *bl* = *bl*
| *Suc*: *bin-to-bl-aux* (*Suc* *n*) *w* *bl* =
bin-to-bl-aux *n* (*bin-rest* *w*) ((*bin-last* *w* = *bit.B1*) # *bl*)

definition *bin-to-bl* :: *nat* \Rightarrow *int* \Rightarrow *bool list* **where**
bin-to-bl-def : *bin-to-bl* *n* *w* = *bin-to-bl-aux* *n* *w* []

primrec *bl-of-nth* :: *nat* \Rightarrow (*nat* \Rightarrow *bool*) \Rightarrow *bool list* **where**
Suc: *bl-of-nth* (*Suc* *n*) *f* = *f* *n* # *bl-of-nth* *n* *f*
| *Z*: *bl-of-nth* 0 *f* = []

primrec *takefill* :: '*a* \Rightarrow *nat* \Rightarrow '*a* list \Rightarrow '*a* list **where**
Z: *takefill* *fill* 0 *xs* = []
| *Suc*: *takefill* *fill* (*Suc* *n*) *xs* = (
case *xs* of [] => *fill* # *takefill* *fill* *n* *xs*
| *y* # *ys* => *y* # *takefill* *fill* *n* *ys*)

definition *map2* :: ('*a* \Rightarrow '*b* \Rightarrow '*c*) \Rightarrow '*a* list \Rightarrow '*b* list \Rightarrow '*c* list **where**
map2 *f* *as* *bs* = *map* (*split* *f*) (*zip* *as* *bs*)

6.4 Splitting and concatenation

definition *bin-rcat* :: *nat* \Rightarrow *int* list \Rightarrow *int* **where**
bin-rcat *n* = *foldl* (%*u* *v*. *bin-cat* *u* *n* *v*) *Int.Pls*

fun *bin-rsplit-aux* :: *nat* \Rightarrow *nat* \Rightarrow *int* \Rightarrow *int* list \Rightarrow *int* list **where**
bin-rsplit-aux *n* *m* *c* *bs* =
(if *m* = 0 | *n* = 0 then *bs* else
let (*a*, *b*) = *bin-split* *n* *c*
in *bin-rsplit-aux* *n* (*m* - *n*) *a* (*b* # *bs*))

definition *bin-rsplit* :: *nat* \Rightarrow *nat* \times *int* \Rightarrow *int* list **where**
bin-rsplit *n* *w* = *bin-rsplit-aux* *n* (*fst* *w*) (*snd* *w*) []

fun *bin-rsplittl-aux* :: *nat* \Rightarrow *nat* \Rightarrow *int* \Rightarrow *int* list \Rightarrow *int* list **where**
bin-rsplittl-aux *n* *m* *c* *bs* =
(if *m* = 0 | *n* = 0 then *bs* else
let (*a*, *b*) = *bin-split* (*min* *m* *n*) *c*
in *bin-rsplittl-aux* *n* (*m* - *n*) *a* (*b* # *bs*))

definition *bin-rsplittl* :: *nat* \Rightarrow *nat* \times *int* \Rightarrow *int list* **where**

bin-rsplittl *n w* = *bin-rsplittl-aux* *n* (*fst w*) (*snd w*) []

declare *bin-rsplit-aux.simps* [*simp del*]

declare *bin-rsplittl-aux.simps* [*simp del*]

lemma *bin-sign-cat*:

!!*y*. *bin-sign* (*bin-cat* *x n y*) = *bin-sign* *x*

by (*induct n*) *auto*

lemma *bin-cat-Suc-Bit*:

bin-cat *w* (*Suc n*) (*v BIT b*) = *bin-cat* *w n v BIT b*

by *auto*

lemma *bin-nth-cat*:

!!*n y*. *bin-nth* (*bin-cat* *x k y*) *n* =

(if *n* < *k* then *bin-nth* *y n* else *bin-nth* *x* (*n* - *k*))

apply (*induct k*)

apply *clarsimp*

apply (*case-tac n*, *auto*)

done

lemma *bin-nth-split*:

!!*b c*. *bin-split* *n c* = (*a*, *b*) ==>

(*ALL k*. *bin-nth* *a k* = *bin-nth* *c* (*n* + *k*)) &

(*ALL k*. *bin-nth* *b k* = (*k* < *n* & *bin-nth* *c k*))

apply (*induct n*)

apply *clarsimp*

apply (*clarsimp simp*: *Let-def split*: *ls-splits*)

apply (*case-tac k*)

apply *auto*

done

lemma *bin-cat-assoc*:

!!*z*. *bin-cat* (*bin-cat* *x m y*) *n z* = *bin-cat* *x* (*m* + *n*) (*bin-cat* *y n z*)

by (*induct n*) *auto*

lemma *bin-cat-assoc-sym*: !!*z m*.

bin-cat *x m* (*bin-cat* *y n z*) = *bin-cat* (*bin-cat* *x* (*m* - *n*) *y*) (*min m n*) *z*

apply (*induct n*, *clarsimp*)

apply (*case-tac m*, *auto*)

done

lemma *bin-cat-Pls* [*simp*]:

!!*w*. *bin-cat* *Int.Pls* *n w* = *bintrunc* *n w*

by (*induct n*) *auto*

lemma *bintr-cat1*:

!!b. $\text{bintrunc } (k + n) (\text{bin-cat } a \ n \ b) = \text{bin-cat } (\text{bintrunc } k \ a) \ n \ b$
 by (induct n) auto

lemma *bintr-cat*: $\text{bintrunc } m (\text{bin-cat } a \ n \ b) =$
 $\text{bin-cat } (\text{bintrunc } (m - n) \ a) \ n \ (\text{bintrunc } (\min \ m \ n) \ b)$
 by (rule bin-eqI) (auto simp: bin-nth-cat nth-bintr)

lemma *bintr-cat-same* [simp]:
 $\text{bintrunc } n (\text{bin-cat } a \ n \ b) = \text{bintrunc } n \ b$
 by (auto simp add : bintr-cat)

lemma *cat-bintr* [simp]:
 $\text{bin-cat } a \ n (\text{bintrunc } n \ b) = \text{bin-cat } a \ n \ b$
 by (induct n) auto

lemma *split-bintrunc*:
 $\text{bin-split } n \ c = (a, b) \implies b = \text{bintrunc } n \ c$
 by (induct n) (auto simp: Let-def split: ls-splits)

lemma *bin-cat-split*:
 $\text{bin-split } n \ w = (u, v) \implies w = \text{bin-cat } u \ n \ v$
 by (induct n) (auto simp: Let-def split: ls-splits)

lemma *bin-split-cat*:
 $\text{bin-split } n (\text{bin-cat } v \ n \ w) = (v, \text{bintrunc } n \ w)$
 by (induct n) auto

lemma *bin-split-Pls* [simp]:
 $\text{bin-split } n \ \text{Int.Pls} = (\text{Int.Pls}, \text{Int.Pls})$
 by (induct n) (auto simp: Let-def split: ls-splits)

lemma *bin-split-Min* [simp]:
 $\text{bin-split } n \ \text{Int.Min} = (\text{Int.Min}, \text{bintrunc } n \ \text{Int.Min})$
 by (induct n) (auto simp: Let-def split: ls-splits)

lemma *bin-split-trunc*:
 $\text{bin-split } (\min \ m \ n) \ c = (a, b) \implies$
 $\text{bin-split } n (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, b)$
 apply (induct n, clarsimp)
 apply (simp add: bin-rest-trunc Let-def split: ls-splits)
 apply (case-tac m)
 apply (auto simp: Let-def split: ls-splits)
 done

lemma *bin-split-trunc1*:
 $\text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-split } n (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, \text{bintrunc } m \ b)$
 apply (induct n, clarsimp)
 apply (simp add: bin-rest-trunc Let-def split: ls-splits)

```

apply (case-tac m)
apply (auto simp: Let-def split: ls-splits)
done

```

```

lemma bin-cat-num:
  !!b. bin-cat a n b = a * 2 ^ n + bintrunc n b
apply (induct n, clarsimp)
apply (simp add: Bit-def cong: number-of-False-cong)
done

```

```

lemma bin-split-num:
  !!b. bin-split n b = (b div 2 ^ n, b mod 2 ^ n)
apply (induct n, clarsimp)
apply (simp add: bin-rest-div zdiv-zmult2-eq)
apply (case-tac b rule: bin-exhaust)
apply simp
apply (simp add: Bit-def mod-mult-mult1 p1mod22k
  split: bit.split
  cong: number-of-False-cong)
done

```

6.5 Miscellaneous lemmas

```

lemma nth-2p-bin:
  !!m. bin-nth (2 ^ n) m = (m = n)
apply (induct n)
apply clarsimp
apply safe
apply (case-tac m)
apply (auto simp: trans [OF numeral-1-eq-1 [symmetric] number-of-eq])
apply (case-tac m)
apply (auto simp: Bit-B0-2t [symmetric])
done

```

```

lemma ex-eq-or:
  (EX m. n = Suc m & (m = k | P m)) = (n = Suc k | (EX m. n = Suc m & P
  m))
  by auto
end

```

7 BinBoolList: Bool lists and integers

```

theory BinBoolList
imports BinOperations

```

begin

7.1 Arithmetic in terms of bool lists

primrec *rbl-succ* :: *bool list* ==> *bool list* **where**

Nil: *rbl-succ Nil* = *Nil*

| *Cons*: *rbl-succ* (*x* # *xs*) = (if *x* then *False* # *rbl-succ xs* else *True* # *xs*)

primrec *rbl-pred* :: *bool list* ==> *bool list* **where**

Nil: *rbl-pred Nil* = *Nil*

| *Cons*: *rbl-pred* (*x* # *xs*) = (if *x* then *False* # *xs* else *True* # *rbl-pred xs*)

primrec *rbl-add* :: *bool list* ==> *bool list* ==> *bool list* **where**

Nil: *rbl-add Nil x* = *Nil*

| *Cons*: *rbl-add* (*y* # *ys*) *x* = (let *ws* = *rbl-add ys* (*tl x*) in
 (*y* ~ = *hd x*) # (if *hd x* & *y* then *rbl-succ ws* else *ws*))

primrec *rbl-mult* :: *bool list* ==> *bool list* ==> *bool list* **where**

Nil: *rbl-mult Nil x* = *Nil*

| *Cons*: *rbl-mult* (*y* # *ys*) *x* = (let *ws* = *False* # *rbl-mult ys x* in
 if *y* then *rbl-add ws x* else *ws*)

lemma *butlast-power*:

(*butlast* ^ *n*) *bl* = *take* (*length bl* - *n*) *bl*

by (*induct n*) (*auto simp: butlast-take*)

lemma *bin-to-bl-aux-Pls-minus-simp* [*simp*]:

0 < *n* ==> *bin-to-bl-aux n Int.Pls bl* =

bin-to-bl-aux (*n* - 1) *Int.Pls* (*False* # *bl*)

by (*cases n*) *auto*

lemma *bin-to-bl-aux-Min-minus-simp* [*simp*]:

0 < *n* ==> *bin-to-bl-aux n Int.Min bl* =

bin-to-bl-aux (*n* - 1) *Int.Min* (*True* # *bl*)

by (*cases n*) *auto*

lemma *bin-to-bl-aux-Bit-minus-simp* [*simp*]:

0 < *n* ==> *bin-to-bl-aux n* (*w BIT b*) *bl* =

bin-to-bl-aux (*n* - 1) *w* ((*b* = *bit.B1*) # *bl*)

by (*cases n*) *auto*

lemma *bin-to-bl-aux-Bit0-minus-simp* [*simp*]:

0 < *n* ==> *bin-to-bl-aux n* (*Int.Bit0 w*) *bl* =

bin-to-bl-aux (*n* - 1) *w* (*False* # *bl*)

by (*cases n*) *auto*

lemma *bin-to-bl-aux-Bit1-minus-simp* [*simp*]:

```

0 < n ==> bin-to-bl-aux n (Int.Bit1 w) bl =
  bin-to-bl-aux (n - 1) w (True # bl)
by (cases n) auto

```

```

lemma bl-to-bin-aux-append:
  bl-to-bin-aux (bs @ cs) w = bl-to-bin-aux cs (bl-to-bin-aux bs w)
by (induct bs arbitrary: w) auto

```

```

lemma bin-to-bl-aux-append:
  bin-to-bl-aux n w bs @ cs = bin-to-bl-aux n w (bs @ cs)
by (induct n arbitrary: w bs) auto

```

```

lemma bl-to-bin-append:
  bl-to-bin (bs @ cs) = bl-to-bin-aux cs (bl-to-bin bs)
unfolding bl-to-bin-def by (rule bl-to-bin-aux-append)

```

```

lemma bin-to-bl-aux-alt:
  bin-to-bl-aux n w bs = bin-to-bl n w @ bs
unfolding bin-to-bl-def by (simp add : bin-to-bl-aux-append)

```

```

lemma bin-to-bl-0: bin-to-bl 0 bs = []
unfolding bin-to-bl-def by auto

```

```

lemma size-bin-to-bl-aux:
  size (bin-to-bl-aux n w bs) = n + length bs
by (induct n arbitrary: w bs) auto

```

```

lemma size-bin-to-bl: size (bin-to-bl n w) = n
unfolding bin-to-bl-def by (simp add : size-bin-to-bl-aux)

```

```

lemma bin-bl-bin':
  bl-to-bin (bin-to-bl-aux n w bs) =
    bl-to-bin-aux bs (bintrunc n w)
by (induct n arbitrary: w bs) (auto simp add : bl-to-bin-def)

```

```

lemma bin-bl-bin: bl-to-bin (bin-to-bl n w) = bintrunc n w
unfolding bin-to-bl-def bin-bl-bin' by auto

```

```

lemma bl-bin-bl':
  bin-to-bl (n + length bs) (bl-to-bin-aux bs w) =
    bin-to-bl-aux n w bs
apply (induct bs arbitrary: w n)
apply auto
apply (simp-all only : add-Suc [symmetric])
apply (auto simp add : bin-to-bl-def)
done

```

```

lemma bl-bin-bl: bin-to-bl (length bs) (bl-to-bin bs) = bs
  unfolding bl-to-bin-def
  apply (rule box-equals)
  apply (rule bl-bin-bl')
  prefer 2
  apply (rule bin-to-bl-aux.Z)
  apply simp
  done

```

```

declare
  bin-to-bl-0 [simp]
  size-bin-to-bl [simp]
  bin-bl-bin [simp]
  bl-bin-bl [simp]

```

```

lemma bl-to-bin-inj:
  bl-to-bin bs = bl-to-bin cs ==> length bs = length cs ==> bs = cs
  apply (rule-tac box-equals)
  defer
  apply (rule bl-bin-bl)
  apply (rule bl-bin-bl)
  apply simp
  done

```

```

lemma bl-to-bin-False: bl-to-bin (False # bl) = bl-to-bin bl
  unfolding bl-to-bin-def by auto

```

```

lemma bl-to-bin-Nil: bl-to-bin [] = Int.Pls
  unfolding bl-to-bin-def by auto

```

```

lemma bin-to-bl-Pls-aux:
  bin-to-bl-aux n Int.Pls bl = replicate n False @ bl
  by (induct n arbitrary: bl) (auto simp: replicate-app-Cons-same)

```

```

lemma bin-to-bl-Pls: bin-to-bl n Int.Pls = replicate n False
  unfolding bin-to-bl-def by (simp add: bin-to-bl-Pls-aux)

```

```

lemma bin-to-bl-Min-aux [rule-format] :
  ALL bl. bin-to-bl-aux n Int.Min bl = replicate n True @ bl
  by (induct n) (auto simp: replicate-app-Cons-same)

```

```

lemma bin-to-bl-Min: bin-to-bl n Int.Min = replicate n True
  unfolding bin-to-bl-def by (simp add: bin-to-bl-Min-aux)

```

```

lemma bl-to-bin-rep-F:
  bl-to-bin (replicate n False @ bl) = bl-to-bin bl
  apply (simp add: bin-to-bl-Pls-aux [symmetric] bin-bl-bin')
  apply (simp add: bl-to-bin-def)
  done

```

lemma *bin-to-bl-trunc*:

$n \leq m \implies \text{bin-to-bl } n (\text{bintrunc } m w) = \text{bin-to-bl } n w$
by (*auto intro: bl-to-bin-inj*)

declare

bin-to-bl-trunc [*simp*]
bl-to-bin-False [*simp*]
bl-to-bin-Nil [*simp*]

lemma *bin-to-bl-aux-bintr* [*rule-format*] :

ALL m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =
replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl
apply (*induct n*)
apply *clarsimp*
apply *clarsimp*
apply (*case-tac m*)
apply (*clarsimp simp: bin-to-bl-Pls-aux*)
apply (*erule thin-rl*)
apply (*induct-tac n*)
apply *auto*
done

lemmas *bin-to-bl-bintr* =

bin-to-bl-aux-bintr [**where** *bl = []*, *folded bin-to-bl-def*]

lemma *bl-to-bin-rep-False*: *bl-to-bin (replicate n False) = Int.Pls*

by (*induct n*) *auto*

lemma *len-bin-to-bl-aux*:

length (bin-to-bl-aux n w bs) = n + length bs
by (*induct n arbitrary: w bs*) *auto*

lemma *len-bin-to-bl* [*simp*]: *length (bin-to-bl n w) = n*

unfolding *bin-to-bl-def len-bin-to-bl-aux* **by** *auto*

lemma *sign-bl-bin'*:

bin-sign (bl-to-bin-aux bs w) = bin-sign w
by (*induct bs arbitrary: w*) *auto*

lemma *sign-bl-bin*: *bin-sign (bl-to-bin bs) = Int.Pls*

unfolding *bl-to-bin-def* **by** (*simp add : sign-bl-bin'*)

lemma *bl-sbin-sign-aux*:

hd (bin-to-bl-aux (Suc n) w bs) =
(bin-sign (sbintrunc n w) = Int.Min)
apply (*induct n arbitrary: w bs*)
apply *clarsimp*
apply (*cases w rule: bin-exhaust*)

```

  apply (simp split add : bit.split)
  apply clarsimp
done

```

lemma *bl-sbin-sign*:

```

  hd (bin-to-bl (Suc n) w) = (bin-sign (sbintrunc n w) = Int.Min)
  unfolding bin-to-bl-def by (rule bl-sbin-sign-aux)

```

lemma *bin-nth-of-bl-aux* [rule-format]:

```

  ∀ w. bin-nth (bl-to-bin-aux bl w) n =
    (n < size bl & rev bl ! n | n >= length bl & bin-nth w (n - size bl))
  apply (induct-tac bl)
  apply clarsimp
  apply clarsimp
  apply (cut-tac x=n and y=size list in linorder-less-linear)
  apply (erule disjE, simp add: nth-append)+
  apply auto
done

```

lemma *bin-nth-of-bl*: $\text{bin-nth } (\text{bl-to-bin } bl) \ n = (n < \text{length } bl \ \& \ \text{rev } bl \ ! \ n)$

```

  unfolding bl-to-bin-def by (simp add : bin-nth-of-bl-aux)

```

lemma *bin-nth-bl* [rule-format] : $\text{ALL } m \ w. \ n < m \ \longrightarrow$

```

  bin-nth w n = nth (rev (bin-to-bl m w)) n
  apply (induct n)
  apply clarsimp
  apply (case-tac m, clarsimp)
  apply (clarsimp simp: bin-to-bl-def)
  apply (simp add: bin-to-bl-aux-alt)
  apply clarsimp
  apply (case-tac m, clarsimp)
  apply (clarsimp simp: bin-to-bl-def)
  apply (simp add: bin-to-bl-aux-alt)
done

```

lemma *nth-rev* [rule-format] :

```

  n < length xs  $\longrightarrow$  rev xs ! n = xs ! (length xs - 1 - n)
  apply (induct-tac xs)
  apply simp
  apply (clarsimp simp add : nth-append nth.simps split add : nat.split)
  apply (rule-tac f = %n. list ! n in arg-cong)
  apply arith
done

```

lemmas *nth-rev-alt* = *nth-rev* [where $xs = \text{rev } ys$, simplified, standard]

lemma *nth-bin-to-bl-aux* [rule-format] :

```

  ALL w n bl. n < m + length bl  $\longrightarrow$  (bin-to-bl-aux m w bl) ! n =
    (if n < m then bin-nth w (m - 1 - n) else bl ! (n - m))

```

```

apply (induct m)
  apply clarsimp
apply clarsimp
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac n - m)
  apply arith
apply simp
apply (rule-tac f = %n. bl ! n in arg-cong)
apply arith
done

lemma nth-bin-to-bl:  $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$ 
  unfolding bin-to-bl-def by (simp add : nth-bin-to-bl-aux)

lemma bl-to-bin-lt2p-aux [rule-format]:
   $\forall w. \text{bl-to-bin-aux } bs \ w < (w + 1) * (2 ^ \text{length } bs)$ 
  apply (induct bs)
  apply clarsimp
apply clarsimp
apply safe
apply (erule allE, erule xtr8 [rotated],
    simp add: numeral-simps algebra-simps cong add : number-of-False-cong)+
done

lemma bl-to-bin-lt2p:  $\text{bl-to-bin } bs < (2 ^ \text{length } bs)$ 
  apply (unfold bl-to-bin-def)
apply (rule xtr1)
prefer 2
apply (rule bl-to-bin-lt2p-aux)
apply simp
done

lemma bl-to-bin-ge2p-aux [rule-format] :
   $\forall w. \text{bl-to-bin-aux } bs \ w \geq w * (2 ^ \text{length } bs)$ 
  apply (induct bs)
  apply clarsimp
apply clarsimp
apply safe
apply (erule allE, erule preorder-class.order-trans [rotated],
    simp add: numeral-simps algebra-simps cong add : number-of-False-cong)+
done

lemma bl-to-bin-ge0:  $\text{bl-to-bin } bs \geq 0$ 
  apply (unfold bl-to-bin-def)
apply (rule xtr4)
apply (rule bl-to-bin-ge2p-aux)
apply simp
done

```


lemma *butlast-rest-bin*:

butlast (*bin-to-bl* *n w*) = *bin-to-bl* (*n* - 1) (*bin-rest w*)
apply (*unfold bin-to-bl-def*)
apply (*cases w rule: bin-exhaust*)
apply (*cases n, clarsimp*)
apply *clarsimp*
apply (*auto simp add: bin-to-bl-aux-alt*)
done

lemmas *butlast-bin-rest = butlast-rest-bin*

[**where** *w=bl-to-bin bl* **and** *n=length bl, simplified, standard*]

lemma *butlast-rest-bl2bin-aux*:

bl \sim [] \implies
bl-to-bin-aux (*butlast bl*) *w* = *bin-rest* (*bl-to-bin-aux bl w*)
by (*induct bl arbitrary: w*) *auto*

lemma *butlast-rest-bl2bin*:

bl-to-bin (*butlast bl*) = *bin-rest* (*bl-to-bin bl*)
apply (*unfold bl-to-bin-def*)
apply (*cases bl*)
apply (*auto simp add: butlast-rest-bl2bin-aux*)
done

lemma *trunc-bl2bin-aux* [*rule-format*]:

ALL w. bintrunc m (bl-to-bin-aux bl w) =
bl-to-bin-aux (*drop* (*length bl* - *m*) *bl*) (*bintrunc* (*m* - *length bl*) *w*)
apply (*induct-tac bl*)
apply *clarsimp*
apply *clarsimp*
apply *safe*
apply (*case-tac m - size list*)
apply (*simp add : diff-is-0-eq [THEN iffD1, THEN Suc-diff-le]*)
apply *simp*
apply (*rule-tac f = %nat. bl-to-bin-aux list (Int.Bit1 (bintrunc nat w))*)
in *arg-cong*)
apply *simp*
apply (*case-tac m - size list*)
apply (*simp add: diff-is-0-eq [THEN iffD1, THEN Suc-diff-le]*)
apply *simp*
apply (*rule-tac f = %nat. bl-to-bin-aux list (Int.Bit0 (bintrunc nat w))*)
in *arg-cong*)
apply *simp*
done

lemma *trunc-bl2bin*:

bintrunc m (bl-to-bin bl) = bl-to-bin (*drop* (*length bl* - *m*) *bl*)
unfolding *bl-to-bin-def* **by** (*simp add : trunc-bl2bin-aux*)

lemmas *trunc-bl2bin-len* [*simp*] =
trunc-bl2bin [*of length bl bl, simplified, standard*]

lemma *bl2bin-drop*:
 $bl\text{-}to\text{-}bin\ (drop\ k\ bl) = bintrunc\ (length\ bl - k)\ (bl\text{-}to\text{-}bin\ bl)$
apply (*rule trans*)
prefer 2
apply (*rule trunc-bl2bin [symmetric]*)
apply (*cases k <= length bl*)
apply *auto*
done

lemma *nth-rest-power-bin* [*rule-format*] :
 $ALL\ n.\ bin\text{-}nth\ ((bin\text{-}rest\ ^\wedge k)\ w)\ n = bin\text{-}nth\ w\ (n + k)$
apply (*induct k, clarsimp*)
apply *clarsimp*
apply (*simp only: bin-nth.Suc [symmetric] add-Suc*)
done

lemma *take-rest-power-bin*:
 $m \leq n \implies take\ m\ (bin\text{-}to\text{-}bl\ n\ w) = bin\text{-}to\text{-}bl\ m\ ((bin\text{-}rest\ ^\wedge (n - m))\ w)$
apply (*rule nth-equalityI*)
apply *simp*
apply (*clarsimp simp add: nth-bin-to-bl nth-rest-power-bin*)
done

lemma *hd-butlast*: $size\ xs > 1 \implies hd\ (butlast\ xs) = hd\ xs$
by (*cases xs*) *auto*

lemma *last-bin-last'*:
 $size\ xs > 0 \implies last\ xs = (bin\text{-}last\ (bl\text{-}to\text{-}bin\text{-}aux\ xs\ w) = bit.B1)$
by (*induct xs arbitrary: w*) *auto*

lemma *last-bin-last*:
 $size\ xs > 0 \implies last\ xs = (bin\text{-}last\ (bl\text{-}to\text{-}bin\ xs) = bit.B1)$
unfolding *bl-to-bin-def* **by** (*erule last-bin-last'*)

lemma *bin-last-last*:
 $bin\text{-}last\ w = (if\ last\ (bin\text{-}to\text{-}bl\ (Suc\ n)\ w)\ then\ bit.B1\ else\ bit.B0)$
apply (*unfold bin-to-bl-def*)
apply *simp*
apply (*auto simp add: bin-to-bl-aux-alt*)
done

lemma *map2-Nil* [*simp*]: $map2\ f\ []\ ys = []$
unfolding *map2-def* **by** *auto*

```

lemma map2-Cons [simp]:
  map2 f (x # xs) (y # ys) = f x y # map2 f xs ys
  unfolding map2-def by auto

lemma bl-xor-aux-bin [rule-format] : ALL v w bs cs.
  map2 (%x y. x ~ = y) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v XOR w) (map2 (%x y. x ~ = y) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

lemma bl-or-aux-bin [rule-format] : ALL v w bs cs.
  map2 (op | ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v OR w) (map2 (op | ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

lemma bl-and-aux-bin [rule-format] : ALL v w bs cs.
  map2 (op & ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v AND w) (map2 (op & ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

lemma bl-not-aux-bin [rule-format] :
  ALL w cs. map Not (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (NOT w) (map Not cs)
apply (induct n)
apply clarsimp

```

```

apply clarsimp
apply (case-tac w rule: bin-exhaust)
apply (case-tac b)
apply auto
done

lemmas bl-not-bin = bl-not-aux-bin
  [where cs = [], unfolded bin-to-bl-def [symmetric] map.simps]

lemmas bl-and-bin = bl-and-aux-bin [where bs=[] and cs=[],
  unfolded map2-Nil, folded bin-to-bl-def]

lemmas bl-or-bin = bl-or-aux-bin [where bs=[] and cs=[],
  unfolded map2-Nil, folded bin-to-bl-def]

lemmas bl-xor-bin = bl-xor-aux-bin [where bs=[] and cs=[],
  unfolded map2-Nil, folded bin-to-bl-def]

lemma drop-bin2bl-aux [rule-format] :
  ALL m bin bs. drop m (bin-to-bl-aux n bin bs) =
    bin-to-bl-aux (n - m) bin (drop (m - n) bs)
apply (induct n, clarsimp)
apply clarsimp
apply (case-tac bin rule: bin-exhaust)
apply (case-tac m <= n, simp)
apply (case-tac m - n, simp)
apply simp
apply (rule-tac f = %nat. drop nat bs in arg-cong)
apply simp
done

lemma drop-bin2bl: drop m (bin-to-bl n bin) = bin-to-bl (n - m) bin
  unfolding bin-to-bl-def by (simp add : drop-bin2bl-aux)

lemma take-bin2bl-lem1 [rule-format] :
  ALL w bs. take m (bin-to-bl-aux m w bs) = bin-to-bl m w
apply (induct m, clarsimp)
apply clarsimp
apply (simp add: bin-to-bl-aux-alt)
apply (simp add: bin-to-bl-def)
apply (simp add: bin-to-bl-aux-alt)
done

lemma take-bin2bl-lem [rule-format] :
  ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =
    take m (bin-to-bl (m + n) w)
apply (induct n)
apply clarify
apply (simp-all (no-asm) add: bin-to-bl-def take-bin2bl-lem1)

```

apply *simp*
done

lemma *bin-split-take* [rule-format] :
 $ALL\ b\ c.\ bin_split\ n\ c = (a, b) \dashrightarrow$
 $bin_to_bl\ m\ a = take\ m\ (bin_to_bl\ (m + n)\ c)$
apply (induct *n*)
apply *clarsimp*
apply (*clarsimp simp: Let-def split: ls-splits*)
apply (*simp add: bin-to-bl-def*)
apply (*simp add: take-bin2bl-lem*)
done

lemma *bin-split-take1*:
 $k = m + n \implies bin_split\ n\ c = (a, b) \implies$
 $bin_to_bl\ m\ a = take\ m\ (bin_to_bl\ k\ c)$
by (*auto elim: bin-split-take*)

lemma *nth-takefill* [rule-format] : $ALL\ m\ l.\ m < n \dashrightarrow$
 $takefill\ fill\ n\ l!\ m = (if\ m < length\ l\ then\ l!\ m\ else\ fill)$
apply (induct *n*, *clarsimp*)
apply *clarsimp*
apply (*case-tac m*)
apply (*simp split: list.split*)
apply *clarsimp*
apply (*erule allE*)
apply (*erule (1) impE*)
apply (*simp split: list.split*)
done

lemma *takefill-alt* [rule-format] :
 $ALL\ l.\ takefill\ fill\ n\ l = take\ n\ l\ @\ replicate\ (n - length\ l)\ fill$
by (induct *n*) (*auto split: list.split*)

lemma *takefill-replicate* [*simp*]:
 $takefill\ fill\ n\ (replicate\ m\ fill) = replicate\ n\ fill$
by (*simp add : takefill-alt replicate-add [symmetric]*)

lemma *takefill-le'* [rule-format] :
 $ALL\ l\ n.\ n = m + k \dashrightarrow takefill\ x\ m\ (takefill\ x\ n\ l) = takefill\ x\ m\ l$
by (induct *m*) (*auto split: list.split*)

lemma *length-takefill* [*simp*]: $length\ (takefill\ fill\ n\ l) = n$
by (*simp add : takefill-alt*)

lemma *take-takefill'*:
 $!!w\ n.\ n = k + m \implies take\ k\ (takefill\ fill\ n\ w) = takefill\ fill\ k\ w$
by (induct *k*) (*auto split add : list.split*)

lemma *drop-takefill*:

!!w. $\text{drop } k \ (\text{takefill fill } (m + k) \ w) = \text{takefill fill } m \ (\text{drop } k \ w)$
by (induct k) (auto split add : list.split)

lemma *takefill-le* [simp]:

$m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
by (auto simp: le-iff-add takefill-le')

lemma *take-takefill* [simp]:

$m \leq n \implies \text{take } m \ (\text{takefill fill } n \ w) = \text{takefill fill } m \ w$
by (auto simp: le-iff-add take-takefill')

lemma *takefill-append*:

$\text{takefill fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill fill } m \ w)$
by (induct xs) auto

lemma *takefill-same'*:

$l = \text{length } xs \implies \text{takefill fill } l \ xs = xs$
by clarify (induct xs, auto)

lemmas *takefill-same* [simp] = *takefill-same'* [OF refl]

lemma *takefill-bintrunc*:

$\text{takefill False } n \ bl = \text{rev } (\text{bin-to-bl } n \ (\text{bl-to-bin } (\text{rev } bl)))$
apply (rule nth-equalityI)
apply simp
apply (clarsimp simp: nth-takefill nth-rev nth-bin-to-bl bin-nth-of-bl)
done

lemma *bl-bin-bl-rtf*:

$\text{bin-to-bl } n \ (\text{bl-to-bin } bl) = \text{rev } (\text{takefill False } n \ (\text{rev } bl))$
by (simp add : takefill-bintrunc)

lemmas *bl-bin-bl-rep-drop* =

bl-bin-bl-rtf [simplified takefill-alt,
simplified, simplified rev-take, simplified]

lemma *tf-rev*:

$n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$
apply (rule nth-equalityI)
apply (auto simp add: nth-takefill nth-rev)
apply (rule-tac f = %n. bl ! n in arg-cong)
apply arith
done

lemma *takefill-minus*:

$0 < n \implies \text{takefill fill } (\text{Suc } (n - 1)) \ w = \text{takefill fill } n \ w$
by auto

```

lemmas takefill-Suc-cases =
  list.cases [THEN takefill.Suc [THEN trans], standard]

lemmas takefill-Suc-Nil = takefill-Suc-cases (1)
lemmas takefill-Suc-Cons = takefill-Suc-cases (2)

lemmas takefill-minus-simps = takefill-Suc-cases [THEN [2]
  takefill-minus [symmetric, THEN trans], standard]

lemmas takefill-pred-simps [simp] =
  takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]

lemma bl-to-bin-aux-cat:
  !!nv v. bl-to-bin-aux bs (bin-cat w nv v) =
    bin-cat w (nv + length bs) (bl-to-bin-aux bs v)
  apply (induct bs)
  apply simp
  apply (simp add: bin-cat-Suc-Bit [symmetric] del: bin-cat.simps)
  done

lemma bin-to-bl-aux-cat:
  !!w bs. bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs =
    bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)
  by (induct nw) auto

lemmas bl-to-bin-aux-alt =
  bl-to-bin-aux-cat [where nv = 0 and v = Int.Pls,
    simplified bl-to-bin-def [symmetric], simplified]

lemmas bin-to-bl-cat =
  bin-to-bl-aux-cat [where bs = [], folded bin-to-bl-def]

lemmas bl-to-bin-aux-app-cat =
  trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]

lemmas bin-to-bl-aux-cat-app =
  trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]

lemmas bl-to-bin-app-cat = bl-to-bin-aux-app-cat
  [where w = Int.Pls, folded bl-to-bin-def]

lemmas bin-to-bl-cat-app = bin-to-bl-aux-cat-app
  [where bs = [], folded bin-to-bl-def]

lemma bl-to-bin-app-cat-alt:

```

$bin-cat (bl-to-bin cs) n w = bl-to-bin (cs @ bin-to-bl n w)$
by (*simp add : bl-to-bin-app-cat*)

lemma *mask-lem*: $(bl-to-bin (True \# replicate n False)) =$
 $Int.succ (bl-to-bin (replicate n True))$
apply (*unfold bl-to-bin-def*)
apply (*induct n*)
apply *simp*
apply (*simp only: Suc-eq-plus1 replicate-add*
 $append-Cons [symmetric] bl-to-bin-aux-append$)
apply *simp*
done

lemma *length-bl-of-nth* [*simp*]: $length (bl-of-nth n f) = n$
by (*induct n*) *auto*

lemma *nth-bl-of-nth* [*simp*]:
 $m < n \implies rev (bl-of-nth n f) ! m = f m$
apply (*induct n*)
apply *simp*
apply (*clarsimp simp add : nth-append*)
apply (*rule-tac f = f in arg-cong*)
apply *simp*
done

lemma *bl-of-nth-inj*:
 $(!!k. k < n \implies f k = g k) \implies bl-of-nth n f = bl-of-nth n g$
by (*induct n*) *auto*

lemma *bl-of-nth-nth-le* [*rule-format*] : *ALL xs.*
 $length xs \geq n \implies bl-of-nth n (nth (rev xs)) = drop (length xs - n) xs$
apply (*induct n, clarsimp*)
apply *clarsimp*
apply (*rule trans [OF - hd-Cons-tl]*)
apply (*frule Suc-le-lessD*)
apply (*simp add: nth-rev trans [OF drop-Suc drop-tl, symmetric]*)
apply (*subst hd-drop-conv-nth*)
apply *force*
apply *simp-all*
apply (*rule-tac f = %n. drop n xs in arg-cong*)
apply *simp*
done

lemmas *bl-of-nth-nth* [*simp*] = *order-refl [THEN bl-of-nth-nth-le, simplified]*

lemma *size-rbl-pred*: $length (rbl-pred bl) = length bl$
by (*induct bl*) *auto*

lemma *size-rbl-succ*: $\text{length } (\text{rbl-succ } bl) = \text{length } bl$
by (*induct* *bl*) *auto*

lemma *size-rbl-add*:
 $!!cl. \text{length } (\text{rbl-add } bl \ cl) = \text{length } bl$
by (*induct* *bl*) (*auto simp: Let-def size-rbl-succ*)

lemma *size-rbl-mult*:
 $!!cl. \text{length } (\text{rbl-mult } bl \ cl) = \text{length } bl$
by (*induct* *bl*) (*auto simp add: Let-def size-rbl-add*)

lemmas *rbl-sizes* [*simp*] =
size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult

lemmas *rbl-Nils* =
rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil

lemma *rbl-pred*:
 $!!bin. \text{rbl-pred } (\text{rev } (\text{bin-to-bl } n \ bin)) = \text{rev } (\text{bin-to-bl } n \ (\text{Int.pred } bin))$
apply (*induct* *n*, *simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bin rule: bin-exhaust*)
apply (*case-tac b*)
apply (*clarsimp simp: bin-to-bl-aux-alt*) +
done

lemma *rbl-succ*:
 $!!bin. \text{rbl-succ } (\text{rev } (\text{bin-to-bl } n \ bin)) = \text{rev } (\text{bin-to-bl } n \ (\text{Int.succ } bin))$
apply (*induct* *n*, *simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bin rule: bin-exhaust*)
apply (*case-tac b*)
apply (*clarsimp simp: bin-to-bl-aux-alt*) +
done

lemma *rbl-add*:
 $!!bina \ binb. \text{rbl-add } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } n \ binb)) =$
 $\text{rev } (\text{bin-to-bl } n \ (bina + binb))$
apply (*induct* *n*, *simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bina rule: bin-exhaust*)
apply (*case-tac binb rule: bin-exhaust*)
apply (*case-tac b*)
apply (*case-tac [!] ba*)
apply (*auto simp: rbl-succ succ-def bin-to-bl-aux-alt Let-def add-ac*)
done

lemma *rbl-add-app2*:

```
!!blb. length blb >= length bla ==>
  rbl-add bla (blb @ blc) = rbl-add bla blb
apply (induct bla, simp)
apply clarsimp
apply (case-tac blb, clarsimp)
apply (clarsimp simp: Let-def)
done
```

lemma *rbl-add-take2*:

```
!!blb. length blb >= length bla ==>
  rbl-add bla (take (length bla) blb) = rbl-add bla blb
apply (induct bla, simp)
apply clarsimp
apply (case-tac blb, clarsimp)
apply (clarsimp simp: Let-def)
done
```

lemma *rbl-add-long*:

```
m >= n ==> rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
  rev (bin-to-bl n (bina + binb))
apply (rule box-equals [OF - rbl-add-take2 rbl-add])
apply (rule-tac f = rbl-add (rev (bin-to-bl n bina)) in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
apply simp
done
```

lemma *rbl-mult-app2*:

```
!!blb. length blb >= length bla ==>
  rbl-mult bla (blb @ blc) = rbl-mult bla blb
apply (induct bla, simp)
apply clarsimp
apply (case-tac blb, clarsimp)
apply (clarsimp simp: Let-def rbl-add-app2)
done
```

lemma *rbl-mult-take2*:

```
length blb >= length bla ==>
  rbl-mult bla (take (length bla) blb) = rbl-mult bla blb
apply (rule trans)
apply (rule rbl-mult-app2 [symmetric])
apply simp
apply (rule-tac f = rbl-mult bla in arg-cong)
apply (rule append-take-drop-id)
done
```

lemma *rbl-mult-gt1*:

```

m >= length bl ==> rbl-mult bl (rev (bin-to-bl m binb)) =
  rbl-mult bl (rev (bin-to-bl (length bl) binb))
apply (rule trans)
apply (rule rbl-mult-take2 [symmetric])
apply simp-all
apply (rule-tac f = rbl-mult bl in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
done

```

lemma *rbl-mult-gt*:

```

m > n ==> rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
  rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb))
by (auto intro: trans [OF rbl-mult-gt1])

```

lemmas *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

lemma *rbbl-Cons*:

```

b # rev (bin-to-bl n x) = rev (bin-to-bl (Suc n) (x BIT If b bit.B1 bit.B0))
apply (unfold bin-to-bl-def)
apply simp
apply (simp add: bin-to-bl-aux-alt)
done

```

lemma *rbl-mult: !!bina binb.*

```

rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
  rev (bin-to-bl n (bina * binb))
apply (induct n)
apply simp
apply (unfold bin-to-bl-def)
apply clarsimp
apply (case-tac bina rule: bin-exhaust)
apply (case-tac binb rule: bin-exhaust)
apply (case-tac b)
apply (case-tac [!] ba)
apply (auto simp: bin-to-bl-aux-alt Let-def)
apply (auto simp: rbbl-Cons rbl-mult-Suc rbl-add)
done

```

lemma *rbl-add-split*:

```

P (rbl-add (y # ys) (x # xs)) =
  (ALL ws. length ws = length ys --> ws = rbl-add ys xs -->
    (y --> ((x --> P (False # rbl-succ ws)) & (~ x --> P (True # ws)))) &
    (~ y --> P (x # ws)))
apply (auto simp add: Let-def)
apply (case-tac [!] y)
apply auto
done

```

lemma *rbl-mult-split*:

$P \text{ (rbl-mult } (y \# ys) xs) =$
 $(ALL \text{ ws. length ws} = \text{Suc (length ys)} \longrightarrow \text{ws} = \text{False} \# \text{rbl-mult ys xs} \longrightarrow$
 $(y \longrightarrow P \text{ (rbl-add ws xs)}) \& (\sim y \longrightarrow P \text{ ws}))$
by (*clarsimp simp add : Let-def*)

lemma *and-len*: $xs = ys \implies xs = ys \& \text{length xs} = \text{length ys}$
by *auto*

lemma *size-if*: $\text{size (if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then size } xs \text{ else size } ys)$
by *auto*

lemma *tl-if*: $\text{tl (if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then tl } xs \text{ else tl } ys)$
by *auto*

lemma *hd-if*: $\text{hd (if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then hd } xs \text{ else hd } ys)$
by *auto*

lemma *if-Not-x*: $(\text{if } p \text{ then } \sim x \text{ else } x) = (p = (\sim x))$
by *auto*

lemma *if-x-Not*: $(\text{if } p \text{ then } x \text{ else } \sim x) = (p = x)$
by *auto*

lemma *if-same-and*: $(\text{If } p \text{ } x \text{ } y \& \text{ If } p \text{ } u \text{ } v) = (\text{if } p \text{ then } x \& u \text{ else } y \& v)$
by *auto*

lemma *if-same-eq*: $(\text{If } p \text{ } x \text{ } y = (\text{If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (u) \text{ else } y = (v))$
by *auto*

lemma *if-same-eq-not*:
 $(\text{If } p \text{ } x \text{ } y = (\sim \text{ If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (\sim u) \text{ else } y = (\sim v))$
by *auto*

lemma *if-Cons*: $(\text{if } p \text{ then } x \# xs \text{ else } y \# ys) = \text{If } p \text{ } x \text{ } y \# \text{ If } p \text{ } xs \text{ } ys$
by *auto*

lemma *if-single*:
 $(\text{if } xc \text{ then } [xab] \text{ else } [an]) = [\text{if } xc \text{ then } xab \text{ else } an]$
by *auto*

lemma *if-bool-simps*:
 $\text{If } p \text{ True } y = (p \mid y) \& \text{ If } p \text{ False } y = (\sim p \& y) \&$
 $\text{If } p \text{ } y \text{ True} = (p \longrightarrow y) \& \text{ If } p \text{ } y \text{ False} = (p \& y)$
by *auto*

lemmas *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *seqr* = *eq-reflection* [**where** *x* = *size w*, *standard*]

lemmas *tl-Nil* = *tl.simps* (1)
lemmas *tl-Cons* = *tl.simps* (2)

7.2 Repeated splitting or concatenation

lemma *sclem*:

size (*concat* (*map* (*bin-to-bl* *n*) *xs*)) = *length xs* * *n*
by (*induct xs*) *auto*

lemma *bin-cat-foldl-lem* [*rule-format*] :

ALL x. foldl (%*u. bin-cat u n*) *x xs* =
bin-cat x (*size xs* * *n*) (*foldl* (%*u. bin-cat u n*) *y xs*)
apply (*induct xs*)
apply *simp*
apply *clarify*
apply (*simp* (*no-asm*))
apply (*frule asm-rl*)
apply (*drule spec*)
apply (*erule trans*)
apply (*drule-tac x* = *bin-cat y n a in spec*)
apply (*simp add* : *bin-cat-assoc-sym min-max.inf-absorb2*)
done

lemma *bin-rcat-bl*:

(*bin-rcat n wl*) = *bl-to-bin* (*concat* (*map* (*bin-to-bl n*) *wl*))
apply (*unfold bin-rcat-def*)
apply (*rule sym*)
apply (*induct wl*)
apply (*auto simp add* : *bl-to-bin-append*)
apply (*simp add* : *bl-to-bin-aux-alt sclem*)
apply (*simp add* : *bin-cat-foldl-lem [symmetric]*)
done

lemmas *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps bin-rsplitl-aux.simps*

lemmas *rsplit-aux-simps* = *bin-rsplit-aux-simps*

lemmas *th-if-simp1* = *split-if* [**where** *P* = *op* = *l*,
THEN iffD1, *THEN conjunct1*, *THEN mp*, *standard*]

lemmas *th-if-simp2* = *split-if* [**where** *P* = *op* = *l*,
THEN iffD1, *THEN conjunct2*, *THEN mp*, *standard*]

lemmas *rsplit-aux-simp1s* = *rsplit-aux-simps* [*THEN th-if-simp1*]

lemmas *rsplit-aux-simp2ls* = *rsplit-aux-simps* [*THEN th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [*simp*] = *rsplit-aux-simp2ls* [*unfolded Let-def*]

lemmas *rscl* = *bin-rsplit-aux-simp2s* (2)

lemmas *rsplit-aux-0-simps* [*simp*] =
rsplit-aux-simp1s [*OF disjI1*] *rsplit-aux-simp1s* [*OF disjI2*]

lemma *bin-rsplit-aux-append*:
bin-rsplit-aux *n m c* (*bs @ cs*) = *bin-rsplit-aux* *n m c* *bs @ cs*
apply (*induct* *n m c* *bs* *rule*: *bin-rsplit-aux.induct*)
apply (*subst* *bin-rsplit-aux.simps*)
apply (*subst* *bin-rsplit-aux.simps*)
apply (*clarsimp* *split*: *ls-splits*)
apply *auto*
done

lemma *bin-rsplittl-aux-append*:
bin-rsplittl-aux *n m c* (*bs @ cs*) = *bin-rsplittl-aux* *n m c* *bs @ cs*
apply (*induct* *n m c* *bs* *rule*: *bin-rsplittl-aux.induct*)
apply (*subst* *bin-rsplittl-aux.simps*)
apply (*subst* *bin-rsplittl-aux.simps*)
apply (*clarsimp* *split*: *ls-splits*)
apply *auto*
done

lemmas *rsplit-aux-apps* [**where** *bs* = []] =
bin-rsplit-aux-append *bin-rsplittl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

lemmas *rsplit-aux-alt*s = *rsplit-aux-apps*
[*unfolded append-Nil* *rsplit-def-auxs* [*symmetric*]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) \ w = \text{bin-split } n \ w$
by *auto*

lemmas *bin-split-minus-simp* =
bin-split.Suc [*THEN* [2] *bin-split-minus* [*symmetric*, *THEN trans*], *standard*]

lemma *bin-split-pred-simp* [*simp*]:
 $(0::\text{nat}) < \text{number-of } \text{bin} \implies$
bin-split (*number-of* *bin*) *w* =
(let (*w1*, *w2*) = *bin-split* (*number-of* (*Int.pred* *bin*)) (*bin-rest* *w*)
in (*w1*, *w2* *BIT* *bin-last* *w*))
by (*simp* *only*: *nobm1* *bin-split-minus-simp*)

declare *bin-split-pred-simp* [*simp*]

lemma *bin-rsplit-aux-simp-alt*:
bin-rsplit-aux *n m c* *bs* =
(*if* *m* = 0 \vee *n* = 0

```

    then bs
  else let (a, b) = bin-split n c in bin-rsplit n (m - n, a) @ b # bs)
unfolding bin-rsplit-aux.simps [of n m c bs]
apply simp
apply (subst rsplit-aux-alts)
apply (simp add: bin-rsplit-def)
done

lemmas bin-rsplit-simp-alt =
  trans [OF bin-rsplit-def
    bin-rsplit-aux-simp-alt, standard]

lemmas bthrs = bin-rsplit-simp-alt [THEN [2] trans]

lemma bin-rsplit-size-sign' [rule-format] :
  n > 0 ==> (ALL nw w. rev sw = bin-rsplit n (nw, w) -->
    (ALL v: set sw. bintrunc n v = v))
apply (induct sw)
apply clarsimp
apply clarsimp
apply (drule bthrs)
apply (simp (no-asm-use) add: Let-def split: ls-splits)
apply clarify
apply (erule impE, rule exI, erule exI)
apply (drule split-bintrunc)
apply simp
done

lemmas bin-rsplit-size-sign = bin-rsplit-size-sign' [OF asm-rl
  rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]],
  standard]

lemma bin-nth-rsplit [rule-format] :
  n > 0 ==> m < n ==> (ALL w k nw. rev sw = bin-rsplit n (nw, w) -->
    k < size sw --> bin-nth (sw ! k) m = bin-nth w (k * n + m))
apply (induct sw)
apply clarsimp
apply clarsimp
apply (drule bthrs)
apply (simp (no-asm-use) add: Let-def split: ls-splits)
apply clarify
apply (erule allE, erule impE, erule exI)
apply (case-tac k)
apply clarsimp
prefer 2
apply clarsimp
apply (erule allE)
apply (erule (1) impE)
apply (drule bin-nth-split, erule conjE, erule allE,

```

```

      erule trans, simp add : add-ac)+
done

```

lemma *bin-rsplit-all*:

```

  0 < nw ==> nw <= n ==> bin-rsplit n (nw, w) = [bintrunc n w]
unfolding bin-rsplit-def
by (clarsimp dest!: split-bintrunc simp: rsplit-aux-simp2ls split: ls-splits)

```

lemma *bin-rsplit-l* [rule-format] :

```

  ALL bin. bin-rsplittl n (m, bin) = bin-rsplit n (m, bintrunc m bin)
apply (rule-tac a = m in wf-less-than [THEN wf-induct])
apply (simp (no-asm) add : bin-rsplittl-def bin-rsplit-def)
apply (rule allI)
apply (subst bin-rsplittl-aux.simps)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: Let-def split: ls-splits)
apply (drule bin-split-trunc)
apply (drule sym [THEN trans], assumption)
apply (subst rsplit-aux-alts(1))
apply (subst rsplit-aux-alts(2))
apply clarsimp
unfolding bin-rsplit-def bin-rsplittl-def
apply simp
done

```

lemma *bin-rsplit-rcat* [rule-format] :

```

  n > 0 --> bin-rsplit n (n * size ws, bin-rcat n ws) = map (bintrunc n) ws
apply (unfold bin-rsplit-def bin-rcat-def)
apply (rule-tac xs = ws in rev-induct)
apply clarsimp
apply clarsimp
apply (subst rsplit-aux-alts)
unfolding bin-split-cat
apply simp
done

```

lemma *bin-rsplit-aux-len-le* [rule-format] :

```

  ∀ ws m. n ≠ 0 → ws = bin-rsplit-aux n nw w bs →
    length ws ≤ m ↔ nw + length bs * n ≤ m * n
apply (induct n nw w bs rule: bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (simp add: lrlem Let-def split: ls-splits)
done

```

lemma *bin-rsplit-len-le*:

```

  n ≠ 0 --> ws = bin-rsplit n (nw, w) --> (length ws <= m) = (nw <= m *
n)
unfolding bin-rsplit-def by (clarsimp simp add : bin-rsplit-aux-len-le)

```



```

lemma bin-rsplit-aux-len [rule-format] :
  n ≠ 0 --> length (bin-rsplit-aux n nw w cs) =
    (nw + n - 1) div n + length cs
apply (induct n nw w cs rule: bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: Let-def split: ls-splits)
apply (erule thin-rl)
apply (case-tac m)
apply simp
apply (case-tac m <= n)
apply auto
done

lemma bin-rsplit-len:
  n ≠ 0 ==> length (bin-rsplit n (nw, w)) = (nw + n - 1) div n
unfolding bin-rsplit-def by (clarsimp simp add: bin-rsplit-aux-len)

lemma bin-rsplit-aux-len-indep:
  n ≠ 0 ==> length bs = length cs ==>
    length (bin-rsplit-aux n nw v bs) =
    length (bin-rsplit-aux n nw w cs)
proof (induct n nw w cs arbitrary: v bs rule: bin-rsplit-aux.induct)
  case (1 n m w cs v bs) show ?case
  proof (cases m = 0)
    case True then show ?thesis using ⟨length bs = length cs⟩ by simp
  next
    case False
    from 1.hyps ⟨m ≠ 0⟩ ⟨n ≠ 0⟩ have hyp:  $\bigwedge v \text{ bs. } \text{length } bs = \text{Suc } (\text{length } cs)$ 
  ==>
    length (bin-rsplit-aux n (m - n) v bs) =
    length (bin-rsplit-aux n (m - n) (fst (bin-split n w)) (snd (bin-split n w) #
cs))
    by auto
    show ?thesis using ⟨length bs = length cs⟩ ⟨n ≠ 0⟩
    by (auto simp add: bin-rsplit-aux-simp-alt Let-def bin-rsplit-len
split: ls-splits)
  qed
qed

lemma bin-rsplit-len-indep:
  n ≠ 0 ==> length (bin-rsplit n (nw, v)) = length (bin-rsplit n (nw, w))
apply (unfold bin-rsplit-def)
apply (simp (no-asm))
apply (erule bin-rsplit-aux-len-indep)
apply (rule refl)
done

end

```

8 TdThs: Type Definition Theorems

```
theory TdThs
imports Main
begin
```

9 More lemmas about normal type definitions

lemma

```
tdD1: type-definition Rep Abs A ==> ∀x. Rep x ∈ A and
tdD2: type-definition Rep Abs A ==> ∀x. Abs (Rep x) = x and
tdD3: type-definition Rep Abs A ==> ∀y. y ∈ A ⟶ Rep (Abs y) = y
by (auto simp: type-definition-def)
```

lemma *td-nat-int*:

```
type-definition int nat (Collect (op <= 0))
unfolding type-definition-def by auto
```

```
context type-definition
begin
```

lemmas *Rep'* [iff] = *Rep* [simplified]

declare *Rep-inverse* [simp] *Rep-inject* [simp]

lemma *Abs-eqD*: *Abs* *x* = *Abs* *y* ==> *x* ∈ *A* ==> *y* ∈ *A* ==> *x* = *y*
by (*simp add: Abs-inject*)

lemma *Abs-inverse'*:

```
r : A ==> Abs r = a ==> Rep a = r
by (safe elim!: Abs-inverse)
```

lemma *Rep-comp-inverse*:

```
Rep o f = g ==> Abs o g = f
using Rep-inverse by (auto intro: ext)
```

lemma *Rep-eqD* [elim!]: *Rep* *x* = *Rep* *y* ==> *x* = *y*
by *simp*

lemma *Rep-inverse'*: *Rep* *a* = *r* ==> *Abs* *r* = *a*
by (*safe intro!: Rep-inverse*)

lemma *comp-Abs-inverse*:

```
f o Abs = g ==> g o Rep = f
using Rep-inverse by (auto intro: ext)
```

lemma *set-Rep*:

```
A = range Rep
proof (rule set-ext)
```

```

fix x
show  $(x \in A) = (x \in \text{range } \text{Rep})$ 
  by (auto dest: Abs-inverse [of x, symmetric])
qed

```

```

lemma set-Rep-Abs:  $A = \text{range } (\text{Rep } o \text{ Abs})$ 
proof (rule set-ext)
  fix x
  show  $(x \in A) = (x \in \text{range } (\text{Rep } o \text{ Abs}))$ 
    by (auto dest: Abs-inverse [of x, symmetric])
qed

```

```

lemma Abs-inj-on: inj-on Abs A
  unfolding inj-on-def
  by (auto dest: Abs-inject [THEN iffD1])

```

```

lemma image:  $\text{Abs } ` A = \text{UNIV}$ 
  by (auto intro!: image-eqI)

```

```

lemmas td-thm = type-definition-axioms

```

```

lemma fns1:
   $\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep} \mid \text{fa } o \text{ Abs} = \text{Abs } o \text{ fr} ==> \text{Abs } o \text{ fr } o \text{ Rep} = \text{fa}$ 
  by (auto dest: Rep-comp-inverse elim: comp-Abs-inverse simp: o-assoc)

```

```

lemmas fns1a = disjI1 [THEN fns1]
lemmas fns1b = disjI2 [THEN fns1]

```

```

lemma fns4:
   $\text{Rep } o \text{ fa } o \text{ Abs} = \text{fr} ==>$ 
   $\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep} \ \& \ \text{fa } o \text{ Abs} = \text{Abs } o \text{ fr}$ 
  by (auto intro!: ext)

```

```

end

```

```

interpretation nat-int: type-definition int nat Collect (op <= 0)
  by (rule td-nat-int)

```

```

declare
  nat-int.Rep-cases [cases del]
  nat-int.Abs-cases [cases del]
  nat-int.Rep-induct [induct del]
  nat-int.Abs-induct [induct del]

```

9.1 Extended form of type definition predicate

```

lemma td-conds:
   $\text{norm } o \text{ norm} = \text{norm} ==> (\text{fr } o \text{ norm} = \text{norm } o \text{ fr}) =$ 
   $(\text{norm } o \text{ fr } o \text{ norm} = \text{fr } o \text{ norm} \ \& \ \text{norm } o \text{ fr } o \text{ norm} = \text{norm } o \text{ fr})$ 

```

```

apply safe
  apply (simp-all add: o-assoc [symmetric])
  apply (simp-all add: o-assoc)
done

lemma fn-comm-power:
   $fa \circ tr = tr \circ fr \implies fa^{\wedge n} \circ tr = tr \circ fr^{\wedge n}$ 
  apply (rule ext)
  apply (induct n)
  apply (auto dest: fun-cong)
done

lemmas fn-comm-power' =
  ext [THEN fn-comm-power, THEN fun-cong, unfolded o-def, standard]

locale td-ext = type-definition +
  fixes norm
  assumes eq-norm:  $\bigwedge x. Rep\ (Abs\ x) = norm\ x$ 
begin

lemma Abs-norm [simp]:
   $Abs\ (norm\ x) = Abs\ x$ 
  using eq-norm [of x] by (auto elim: Rep-inverse')

lemma td-th:
   $g \circ Abs = f \implies f\ (Rep\ x) = g\ x$ 
  by (drule comp-Abs-inverse [symmetric]) simp

lemma eq-norm':  $Rep \circ Abs = norm$ 
  by (auto simp: eq-norm intro!: ext)

lemma norm-Rep [simp]:  $norm\ (Rep\ x) = Rep\ x$ 
  by (auto simp: eq-norm' intro: td-th)

lemmas td = td-thm

lemma set-iff-norm:  $w : A \iff w = norm\ w$ 
  by (auto simp: set-Rep-Abs eq-norm' eq-norm [symmetric])

lemma inverse-norm:
   $(Abs\ n = w) = (Rep\ w = norm\ n)$ 
  apply (rule iffI)
  apply (clarsimp simp add: eq-norm)
  apply (simp add: eq-norm' [symmetric])
done

lemma norm-eq-iff:
   $(norm\ x = norm\ y) = (Abs\ x = Abs\ y)$ 

```

by (*simp add: eq-norm' [symmetric]*)

lemma *norm-comps*:

Abs o norm = Abs

norm o Rep = Rep

norm o norm = norm

by (*auto simp: eq-norm' [symmetric] o-def*)

lemmas *norm-norm [simp] = norm-comps*

lemma *fns5*:

Rep o fa o Abs = fr ==>

fr o norm = fr & norm o fr = fr

by (*fold eq-norm'*) (*auto intro!: ext*)

lemma *fns2*:

Abs o fr o Rep = fa ==>

(norm o fr o norm = fr o norm) = (Rep o fa = fr o Rep)

apply (*fold eq-norm'*)

apply *safe*

prefer 2

apply (*simp add: o-assoc*)

apply (*rule ext*)

apply (*drule-tac x=Rep x in fun-cong*)

apply *auto*

done

lemma *fns3*:

Abs o fr o Rep = fa ==>

(norm o fr o norm = norm o fr) = (fa o Abs = Abs o fr)

apply (*fold eq-norm'*)

apply *safe*

prefer 2

apply (*simp add: o-assoc [symmetric]*)

apply (*rule ext*)

apply (*drule fun-cong*)

apply *simp*

done

lemma *fns*:

fr o norm = norm o fr ==>

(fa o Abs = Abs o fr) = (Rep o fa = fr o Rep)

apply *safe*

apply (*frule fns1b*)

prefer 2

apply (*frule fns1a*)

apply (*rule fns3 [THEN iffD1]*)

prefer 3

```

    apply (rule fns2 [THEN iffD1])
    apply (simp-all add: o-assoc [symmetric])
    apply (simp-all add: o-assoc)
  done

lemma range-norm:
  range (Rep o Abs) = A
  by (simp add: set-Rep-Abs)

end

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

end

```

10 WordDefinition: Definition of Word Type

```

theory WordDefinition
imports Size BinBoolList TdThs
begin

```

10.1 Type definition

```

typedef (open word) 'a word = {(0::int) ..< 2^len-of TYPE('a::len0)}
  morphisms uint Abs-word by auto

```

definition *word-of-int* :: *int* \Rightarrow *'a::len0 word* **where**
 — representation of words using unsigned or signed bins, only difference in these
 is the type class
word-of-int *w* = *Abs-word* (*bintrunc* (*len-of TYPE ('a)*) *w*)

lemma *uint-word-of-int* [*code*]: *uint* (*word-of-int* *w* :: *'a::len0 word*) = *w mod 2 ^*
len-of TYPE ('a)
by (*auto simp add: word-of-int-def bintrunc-mod2p intro: Abs-word-inverse*)

```

code-datatype word-of-int

```

```

notation fcomp (infixl o> 60)
notation scomp (infixl o→ 60)

```

```

instantiation word :: ({len0, typerep}) random
begin

```

definition
random-word *i* = *Random.range* (*max i* (*2 ^ len-of TYPE ('a)*)) *o→* (λk . *Pair* (
let j = word-of-int (*Code-Numeral.int-of k*) :: *'a word*

in (j, λ::unit. Code-Evaluation.term-of j)))

instance ..

end

no-notation *fcomp* (**infixl** *o* > 60)

no-notation *scomp* (**infixl** *o* → 60)

10.2 Type conversions and casting

definition *sint* :: 'a :: len word => int **where**

— treats the most-significant-bit as a sign bit

sint-uint: *sint* *w* = *sbintrunc* (*len-of TYPE* ('a) - 1) (*uint* *w*)

definition *unat* :: 'a :: len0 word => nat **where**

unat *w* = *nat* (*uint* *w*)

definition *uints* :: nat => int set **where**

— the sets of integers representing the words

uints *n* = *range* (*bintrunc* *n*)

definition *sints* :: nat => int set **where**

sints *n* = *range* (*sbintrunc* (*n* - 1))

definition *unats* :: nat => nat set **where**

unats *n* = {*i*. *i* < 2 ^ *n*}

definition *norm-sint* :: nat => int => int **where**

norm-sint *n* *w* = (*w* + 2 ^ (*n* - 1)) mod 2 ^ *n* - 2 ^ (*n* - 1)

definition *scast* :: 'a :: len word => 'b :: len word **where**

— cast a word to a different length

scast *w* = *word-of-int* (*sint* *w*)

definition *ucast* :: 'a :: len0 word => 'b :: len0 word **where**

ucast *w* = *word-of-int* (*uint* *w*)

instantiation *word* :: (len0) size

begin

definition

word-size: *size* (*w* :: 'a word) = *len-of TYPE*('a)

instance ..

end

definition *source-size* :: ('a :: len0 word => 'b) => nat **where**

— whether a cast (or other) function is to a longer or shorter length
source-size $c = (\text{let } arb = \text{undefined} ; x = c \text{ } arb \text{ in } \text{size } arb)$

definition *target-size* :: ('a => 'b :: len0 word) => nat **where**
target-size $c = \text{size } (c \text{ undefined})$

definition *is-up* :: ('a :: len0 word => 'b :: len0 word) => bool **where**
is-up $c \iff \text{source-size } c \leq \text{target-size } c$

definition *is-down* :: ('a :: len0 word => 'b :: len0 word) => bool **where**
is-down $c \iff \text{target-size } c \leq \text{source-size } c$

definition *of-bl* :: bool list => 'a :: len0 word **where**
of-bl $bl = \text{word-of-int } (bl\text{-to-bin } bl)$

definition *to-bl* :: 'a :: len0 word => bool list **where**
to-bl $w = \text{bin-to-bl } (\text{len-of TYPE } ('a)) \text{ (uint } w)$

definition *word-reverse* :: 'a :: len0 word => 'a word **where**
word-reverse $w = \text{of-bl } (\text{rev } (\text{to-bl } w))$

definition *word-int-case* :: (int => 'b) => ('a :: len0 word) => 'b **where**
word-int-case $f \text{ } w = f \text{ (uint } w)$

syntax

of-int :: int => 'a

translations

case $x \text{ of } \text{CONST } of\text{-int } y \Rightarrow b == \text{CONST } \text{word-int-case } (\%y. b) \text{ } x$

10.3 Arithmetic operations

instantiation *word* :: (len0) {number, uminus, minus, plus, one, zero, times, Divides.div, ord, bit}

begin

definition

word-0-wi: $0 = \text{word-of-int } 0$

definition

word-1-wi: $1 = \text{word-of-int } 1$

definition

word-add-def: $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$

definition

word-sub-wi: $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$

definition

word-minus-def: $- a = \text{word-of-int } (- \text{uint } a)$

definition

word-mult-def: $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$

definition

word-div-def: $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div } \text{uint } b)$

definition

word-mod-def: $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod } \text{uint } b)$

definition

word-number-of-def: $\text{number-of } w = \text{word-of-int } w$

definition

word-le-def: $a \leq b \longleftrightarrow \text{uint } a \leq \text{uint } b$

definition

word-less-def: $x < y \longleftrightarrow x \leq y \wedge x \neq (y :: 'a \text{ word})$

definition

word-and-def:
 $(a :: 'a \text{ word}) \text{ AND } b = \text{word-of-int } (\text{uint } a \text{ AND } \text{uint } b)$

definition

word-or-def:
 $(a :: 'a \text{ word}) \text{ OR } b = \text{word-of-int } (\text{uint } a \text{ OR } \text{uint } b)$

definition

word-xor-def:
 $(a :: 'a \text{ word}) \text{ XOR } b = \text{word-of-int } (\text{uint } a \text{ XOR } \text{uint } b)$

definition

word-not-def:
 $\text{NOT } (a :: 'a \text{ word}) = \text{word-of-int } (\text{NOT } (\text{uint } a))$

instance ..**end****definition**

word-succ :: $'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

where

word-succ $a = \text{word-of-int } (\text{Int.succ } (\text{uint } a))$

definition

word-pred :: $'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

where

word-pred $a = \text{word-of-int } (\text{Int.pred } (\text{uint } a))$

definition *udvd* :: 'a::len word => 'a::len word => bool (**infixl** *udvd* 50) **where**
a udvd b == *EX n>=0. uint b = n * uint a*

definition *word-sle* :: 'a :: len word => 'a word => bool ((-/ <=s -) [50, 51] 50)
where
a <=s b == *sint a <= sint b*

definition *word-sless* :: 'a :: len word => 'a word => bool ((-/ <s -) [50, 51] 50)
where
(x <s y) == *(x <=s y & x ~ = y)*

10.4 Bit-wise operations

instantiation *word* :: (len0) bits
begin

definition
word-test-bit-def: *test-bit a = bin-nth (uint a)*

definition
word-set-bit-def: *set-bit a n x =*
word-of-int (bin-sc n (If x bit.B1 bit.B0) (uint a))

definition
word-set-bits-def: *(BITS n. f n) = of-bl (bl-of-nth (len-of TYPE ('a)) f)*

definition
word-lsb-def: *lsb a ⟷ bin-last (uint a) = bit.B1*

definition *shiffl1* :: 'a word ⇒ 'a word **where**
shiffl1 w = word-of-int (uint w BIT bit.B0)

definition *shiftr1* :: 'a word ⇒ 'a word **where**
 — shift right as unsigned or as signed, ie logical or arithmetic
shiftr1 w = word-of-int (bin-rest (uint w))

definition
shiffl-def: *w << n = (shiffl1 ^^ n) w*

definition
shiftr-def: *w >> n = (shiftr1 ^^ n) w*

instance ..

end

instantiation *word* :: (len) bitss
begin

definition*word-msb-def:* $msb\ a \longleftrightarrow bin\text{-}sign\ (sint\ a) = Int.Min$ **instance ..****end****definition** *setBit* :: 'a :: len0 word => nat => 'a word **where***setBit* w n == *set-bit* w n True**definition** *clearBit* :: 'a :: len0 word => nat => 'a word **where***clearBit* w n == *set-bit* w n False**10.5 Shift operations****definition** *sshiftr1* :: 'a :: len word => 'a word **where***sshiftr1* w == *word-of-int* (bin-rest (sint w))**definition** *bshiftr1* :: bool => 'a :: len word => 'a word **where***bshiftr1* b w == *of-bl* (b # butlast (to-bl w))**definition** *sshiftr* :: 'a :: len word => nat => 'a word (**infixl** >>> 55) **where***w* >>> n == (*sshiftr1* ^ n) w**definition** *mask* :: nat => 'a::len word **where***mask* n == (1 << n) - 1**definition** *revcast* :: 'a :: len0 word => 'b :: len0 word **where***revcast* w == *of-bl* (*takefill* False (len-of TYPE('b)) (to-bl w))**definition** *slice1* :: nat => 'a :: len0 word => 'b :: len0 word **where***slice1* n w == *of-bl* (*takefill* False n (to-bl w))**definition** *slice* :: nat => 'a :: len0 word => 'b :: len0 word **where***slice* n w == *slice1* (size w - n) w**10.6 Rotation****definition** *rotater1* :: 'a list => 'a list **where***rotater1* ys ==

case ys of [] => [] | x # xs => last xs # butlast xs

definition *rotater* :: nat => 'a list => 'a list **where***rotater* n == *rotater1* ^ n**definition** *word-rotr* :: nat => 'a :: len0 word => 'a :: len0 word **where***word-rotr* n w == *of-bl* (*rotater* n (to-bl w))**definition** *word-rotl* :: nat => 'a :: len0 word => 'a :: len0 word **where**

$word\text{-}rotl\ n\ w == of\text{-}bl\ (rotate\ n\ (to\text{-}bl\ w))$

definition $word\text{-}roti :: int \Rightarrow 'a :: len0\ word \Rightarrow 'a :: len0\ word$ **where**
 $word\text{-}roti\ i\ w == if\ i \geq 0\ then\ word\text{-}rotr\ (nat\ i)\ w$
 $else\ word\text{-}rotl\ (nat\ (-\ i))\ w$

10.7 Split and cat operations

definition $word\text{-}cat :: 'a :: len0\ word \Rightarrow 'b :: len0\ word \Rightarrow 'c :: len0\ word$ **where**
 $word\text{-}cat\ a\ b == word\text{-}of\text{-}int\ (bin\text{-}cat\ (uint\ a)\ (len\text{-}of\ TYPE\ ('b))\ (uint\ b))$

definition $word\text{-}split :: 'a :: len0\ word \Rightarrow ('b :: len0\ word) * ('c :: len0\ word)$
where
 $word\text{-}split\ a ==$
 $case\ bin\text{-}split\ (len\text{-}of\ TYPE\ ('c))\ (uint\ a)\ of$
 $(u, v) \Rightarrow (word\text{-}of\text{-}int\ u, word\text{-}of\text{-}int\ v)$

definition $word\text{-}rcat :: 'a :: len0\ word\ list \Rightarrow 'b :: len0\ word$ **where**
 $word\text{-}rcat\ ws ==$
 $word\text{-}of\text{-}int\ (bin\text{-}rcat\ (len\text{-}of\ TYPE\ ('a))\ (map\ uint\ ws))$

definition $word\text{-}rsplit :: 'a :: len0\ word \Rightarrow 'b :: len\ word\ list$ **where**
 $word\text{-}rsplit\ w ==$
 $map\ word\text{-}of\text{-}int\ (bin\text{-}rsplit\ (len\text{-}of\ TYPE\ ('b))\ (len\text{-}of\ TYPE\ ('a),\ uint\ w))$

definition $max\text{-}word :: 'a :: len\ word$ — Largest representable machine integer.
where
 $max\text{-}word \equiv word\text{-}of\text{-}int\ (2 ^ len\text{-}of\ TYPE\ ('a) - 1)$

primrec $of\text{-}bool :: bool \Rightarrow 'a :: len\ word$ **where**
 $of\text{-}bool\ False = 0$
 $| of\text{-}bool\ True = 1$

lemmas $of\text{-}nth\text{-}def = word\text{-}set\text{-}bits\text{-}def$

lemmas $word\text{-}size\text{-}gt\text{-}0\ [iff] =$

$xtr1\ [OF\ word\text{-}size\ len\text{-}gt\text{-}0,\ standard]$

lemmas $lens\text{-}gt\text{-}0 = word\text{-}size\text{-}gt\text{-}0\ len\text{-}gt\text{-}0$

lemmas $lens\text{-}not\text{-}0\ [iff] = lens\text{-}gt\text{-}0\ [THEN\ gr\text{-}implies\text{-}not0,\ standard]$

lemma $uints\text{-}num: uints\ n = \{i. 0 \leq i \wedge i < 2 ^ n\}$
by $(simp\ add: uints\text{-}def\ range\text{-}bintrunc)$

lemma $sints\text{-}num: sints\ n = \{i. -(2 ^ (n - 1)) \leq i \wedge i < 2 ^ (n - 1)\}$
by $(simp\ add: sints\text{-}def\ range\text{-}sbintrunc)$

lemmas $atLeastLessThan\text{-}alt = atLeastLessThan\text{-}def\ [unfolded\ atLeast\text{-}def\ lessThan\text{-}def\ Collect\text{-}conj\text{-}eq\ [symmetric]]$

lemma *mod-in-reps*: $m > 0 \implies y \bmod m : \{0::\text{int} \dots m\}$
unfolding *atLeastLessThan-alt* **by** *auto*

lemma
uint-0:0 \leq *uint* *x* **and**
uint-lt: *uint* (*x*::*'a*::*len0* *word*) $<$ $2^{\text{len-of TYPE('a)}}$
by (*auto simp: uint [simplified]*)

lemma *uint-mod-same*:
uint *x* $\bmod 2^{\text{len-of TYPE('a)}}$ = *uint* (*x*::*'a*::*len0* *word*)
by (*simp add: int-mod-eq uint-lt uint-0*)

lemma *td-ext-uint*:
td-ext (*uint* :: *'a* *word* \Rightarrow *int*) *word-of-int* (*uints* (*len-of TYPE('a::len0)*))
 ($\%w::\text{int}. w \bmod 2^{\text{len-of TYPE('a)}}$)
apply (*unfold td-ext-def'*)
apply (*simp add: uints-num word-of-int-def bintrunc-mod2p*)
apply (*simp add: uint-mod-same uint-0 uint-lt*
 word.uint-inverse word.Abs-word-inverse int-mod-lem)
done

lemmas *int-word-uint* = *td-ext-uint* [*THEN td-ext.eq-norm, standard*]

interpretation *word-uint*:
td-ext *uint*::*'a*::*len0* *word* \Rightarrow *int*
 word-of-int
 uints (*len-of TYPE('a::len0)*)
 $\lambda w. w \bmod 2^{\text{len-of TYPE('a::len0)}}$
by (*rule td-ext-uint*)

lemmas *td-uint* = *word-uint.td-thm*

lemmas *td-ext-ubin* = *td-ext-uint*
 [*simplified len-gt-0 no-bintr-alt1 [symmetric]*]

interpretation *word-ubin*:
td-ext *uint*::*'a*::*len0* *word* \Rightarrow *int*
 word-of-int
 uints (*len-of TYPE('a::len0)*)
 bintrunc (*len-of TYPE('a::len0)*)
by (*rule td-ext-ubin*)

lemma *sint-sbintrunc'*:
sint (*word-of-int* *bin* :: *'a* *word*) =
 (*sbintrunc* (*len-of TYPE('a :: len) - 1*) *bin*)
unfolding *sint-uint*
by (*auto simp: word-ubin.eq-norm sbintrunc-bintrunc-lt*)

lemma *uint-sint*:

uint w = bintrunc (len-of TYPE('a)) (sint (w :: 'a :: len word))
unfolding *sint-uint* **by** (*auto simp: bintrunc-sbintrunc-le*)

lemma *bintr-uint'*:

n >= size w ==> bintrunc n (uint w) = uint w
apply (*unfold word-size*)
apply (*subst word-ubin.norm-Rep [symmetric]*)
apply (*simp only: bintrunc-bintrunc-min word-size*)
apply (*simp add: min-max.inf-absorb2*)
done

lemma *wi-bintr'*:

wb = word-of-int bin ==> n >= size wb ==>
word-of-int (bintrunc n bin) = wb
unfolding *word-size*
by (*clarsimp simp add: word-ubin.norm-eq-iff [symmetric] min-max.inf-absorb1*)

lemmas *bintr-uint = bintr-uint' [unfolded word-size]*

lemmas *wi-bintr = wi-bintr' [unfolded word-size]*

lemma *td-ext-sbin*:

td-ext (sint :: 'a word => int) word-of-int (sints (len-of TYPE('a::len)))
(sbintrunc (len-of TYPE('a) - 1))
apply (*unfold td-ext-def' sint-uint*)
apply (*simp add : word-ubin.eq-norm*)
apply (*cases len-of TYPE('a)*)
apply (*auto simp add : sints-def*)
apply (*rule sym [THEN trans]*)
apply (*rule word-ubin.Abs-norm*)
apply (*simp only: bintrunc-sbintrunc*)
apply (*drule sym*)
apply *simp*
done

lemmas *td-ext-sint = td-ext-sbin*

[*simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]*]

interpretation *word-sint*:

td-ext sint :: 'a::len word => int
word-of-int
sints (len-of TYPE('a::len))
 $\%w. (w + 2^{(len-of TYPE('a::len) - 1)}) \bmod 2^{len-of TYPE('a::len) - 2^{(len-of TYPE('a::len) - 1)}}$
by (*rule td-ext-sint*)

interpretation *word-sbin*:

td-ext sint :: 'a::len word => int

```

      word-of-int
      sints (len-of TYPE('a::len))
      sbintrunc (len-of TYPE('a::len) - 1)
    by (rule td-ext-sbin)

lemmas int-word-sint = td-ext-sint [THEN td-ext.eq-norm, standard]

lemmas td-sint = word-sint.td

lemma word-number-of-alt: number-of b == word-of-int (number-of b)
  unfolding word-number-of-def by (simp add: number-of-eq)

lemma word-no-wi: number-of = word-of-int
  by (auto simp: word-number-of-def intro: ext)

lemma to-bl-def':
  (to-bl :: 'a :: len0 word => bool list) =
    bin-to-bl (len-of TYPE('a)) o uint
  by (auto simp: to-bl-def intro: ext)

lemmas word-reverse-no-def [simp] = word-reverse-def [of number-of w, standard]

lemmas uints-mod = uints-def [unfolded no-bintr-alt1]

lemma uint-bintrunc: uint (number-of bin :: 'a word) =
  number-of (bintrunc (len-of TYPE ('a :: len0)) bin)
  unfolding word-number-of-def number-of-eq
  by (auto intro: word-ubin.eq-norm)

lemma sint-sbintrunc: sint (number-of bin :: 'a word) =
  number-of (sbintrunc (len-of TYPE ('a :: len) - 1) bin)
  unfolding word-number-of-def number-of-eq
  by (subst word-sbin.eq-norm) simp

lemma unat-bintrunc:
  unat (number-of bin :: 'a :: len0 word) =
  number-of (bintrunc (len-of TYPE('a)) bin)
  unfolding unat-def nat-number-of-def
  by (simp only: uint-bintrunc)

declare
  uint-bintrunc [simp]
  sint-sbintrunc [simp]
  unat-bintrunc [simp]

lemma size-0-eq: size (w :: 'a :: len0 word) = 0 ==> v = w
  apply (unfold word-size)
  apply (rule word-uint.Rep-eqD)

```

```

apply (rule box-equals)
  defer
    apply (rule word-ubin.norm-Rep)+
apply simp
done

lemmas uint-lem = word-uint.Rep [unfolded uints-num mem-Collect-eq]
lemmas sint-lem = word-sint.Rep [unfolded sints-num mem-Collect-eq]
lemmas uint-ge-0 [iff] = uint-lem [THEN conjunct1, standard]
lemmas uint-lt2p [iff] = uint-lem [THEN conjunct2, standard]
lemmas sint-ge = sint-lem [THEN conjunct1, standard]
lemmas sint-lt = sint-lem [THEN conjunct2, standard]

lemma sign-uint-Pls [simp]:
  bin-sign (uint x) = Int.Pl
  by (simp add: sign-Pls-ge-0 number-of-eq)

lemmas uint-m2p-neg = iffD2 [OF diff-less-0-iff-less uint-lt2p, standard]
lemmas uint-m2p-not-non-neg =
  iffD2 [OF linorder-not-le uint-m2p-neg, standard]

lemma lt2p-lem:
  len-of TYPE('a) <= n ==> uint (w :: 'a :: len0 word) < 2 ^ n
  by (rule xtr8 [OF - uint-lt2p]) simp

lemmas uint-le-0-iff [simp] =
  uint-ge-0 [THEN leD, THEN linorder-antisym-conv1, standard]

lemma uint-nat: uint w == int (unat w)
  unfolding unat-def by auto

lemma uint-number-of:
  uint (number-of b :: 'a :: len0 word) = number-of b mod 2 ^ len-of TYPE('a)
  unfolding word-number-of-alt
  by (simp only: int-word-uint)

lemma unat-number-of:
  bin-sign b = Int.Pl ==>
  unat (number-of b :: 'a :: len0 word) = number-of b mod 2 ^ len-of TYPE ('a)
  apply (unfold unat-def)
  apply (clarsimp simp only: uint-number-of)
  apply (rule nat-mod-distrib [THEN trans])
  apply (erule sign-Pls-ge-0 [THEN iffD1])
  apply (simp-all add: nat-power-eq)
done

lemma sint-number-of: sint (number-of b :: 'a :: len word) = (number-of b +
  2 ^ (len-of TYPE('a) - 1)) mod 2 ^ len-of TYPE('a) -
  2 ^ (len-of TYPE('a) - 1)

```


unfolding *word-number-of-alt* **by** (*rule int-word-sint*)

lemma *word-of-int-bin* [*simp*] :
 (*word-of-int* (*number-of bin*) :: 'a :: len0 word) = (*number-of bin*)
unfolding *word-number-of-alt* **by** *auto*

lemma *word-int-case-wi*:
word-int-case *f* (*word-of-int* *i* :: 'b word) =
f (*i mod 2* ^ *len-of TYPE('b::len0)*)
unfolding *word-int-case-def* **by** (*simp add: word-uint.eq-norm*)

lemma *word-int-split*:
P (*word-int-case* *f* *x*) =
 (*ALL i. x* = (*word-of-int* *i* :: 'b :: len0 word) &
0 <= *i* & *i* < 2 ^ *len-of TYPE('b)* \longrightarrow *P* (*f i*))
unfolding *word-int-case-def*
by (*auto simp: word-uint.eq-norm int-mod-eq'*)

lemma *word-int-split-asm*:
P (*word-int-case* *f* *x*) =
 (\sim (*EX n. x* = (*word-of-int* *n* :: 'b::len0 word) &
0 <= *n* & *n* < 2 ^ *len-of TYPE('b::len0)* & \sim *P* (*f n*)))
unfolding *word-int-case-def*
by (*auto simp: word-uint.eq-norm int-mod-eq'*)

lemmas *uint-range'* =
word-uint.Rep [*unfolded uints-num mem-Collect-eq, standard*]
lemmas *sint-range'* = *word-sint.Rep* [*unfolded One-nat-def*
sints-num mem-Collect-eq, standard]

lemma *uint-range-size*: *0* <= *uint w* & *uint w* < 2 ^ *size w*
unfolding *word-size* **by** (*rule uint-range'*)

lemma *sint-range-size*:
 (2 ^ (*size w* - *Suc 0*)) <= *sint w* & *sint w* < 2 ^ (*size w* - *Suc 0*)
unfolding *word-size* **by** (*rule sint-range'*)

lemmas *sint-above-size* = *sint-range-size*
 [*THEN conjunct2, THEN* [2] *xtr8, folded One-nat-def, standard*]

lemmas *sint-below-size* = *sint-range-size*
 [*THEN conjunct1, THEN* [2] *order-trans, folded One-nat-def, standard*]

lemma *test-bit-eq-iff*: (*test-bit* (*u::'a::len0 word*) = *test-bit* *v*) = (*u* = *v*)
unfolding *word-test-bit-def* **by** (*simp add: bin-nth-eq-iff*)

lemma *test-bit-size* [*rule-format*] : (*w::'a::len0 word*) !! *n* \longrightarrow *n* < *size w*
apply (*unfold word-test-bit-def*)
apply (*subst word-ubin.norm-Rep* [*symmetric*])

```

apply (simp only: nth-bintr word-size)
apply fast
done

```

```

lemma word-eqI [rule-format] :
  fixes u :: 'a::len0 word
  shows (ALL n. n < size u --> u !! n = v !! n) ==> u = v
  apply (rule test-bit-eq-iff [THEN iffD1])
  apply (rule ext)
  apply (erule allE)
  apply (erule impCE)
  prefer 2
  apply assumption
  apply (auto dest!: test-bit-size simp add: word-size)
done

```

```

lemmas word-eqD = test-bit-eq-iff [THEN iffD2, THEN fun-cong, standard]

```

```

lemma test-bit-bin': w !! n = (n < size w & bin-nth (uint w) n)
  unfolding word-test-bit-def word-size
  by (simp add: nth-bintr [symmetric])

```

```

lemmas test-bit-bin = test-bit-bin' [unfolded word-size]

```

```

lemma bin-nth-uint-imp': bin-nth (uint w) n --> n < size w
  apply (unfold word-size)
  apply (rule impI)
  apply (rule nth-bintr [THEN iffD1, THEN conjunct1])
  apply (subst word-ubin.norm-Rep)
  apply assumption
done

```

```

lemma bin-nth-sint':
  n >= size w --> bin-nth (sint w) n = bin-nth (sint w) (size w - 1)
  apply (rule impI)
  apply (subst word-sbin.norm-Rep [symmetric])
  apply (simp add: nth-sbintr word-size)
  apply auto
done

```

```

lemmas bin-nth-uint-imp = bin-nth-uint-imp' [rule-format, unfolded word-size]

```

```

lemmas bin-nth-sint = bin-nth-sint' [rule-format, unfolded word-size]

```

```

lemma td-bl:
  type-definition (to-bl :: 'a::len0 word => bool list)
    of-bl
    {bl. length bl = len-of TYPE('a)}
  apply (unfold type-definition-def of-bl-def to-bl-def)

```

```

apply (simp add: word-ubin.eq-norm)
apply safe
apply (drule sym)
apply simp
done

interpretation word-bl:
  type-definition to-bl :: 'a::len0 word => bool list
    of-bl
    {bl. length bl = len-of TYPE('a::len0)}
  by (rule td-bl)

lemma word-size-bl: size w == size (to-bl w)
  unfolding word-size by auto

lemma to-bl-use-of-bl:
  (to-bl w = bl) = (w = of-bl bl ∧ length bl = length (to-bl w))
  by (fastsimp elim!: word-bl.Abs-inverse [simplified])

lemma to-bl-word-rev: to-bl (word-reverse w) = rev (to-bl w)
  unfolding word-reverse-def by (simp add: word-bl.Abs-inverse)

lemma word-rev-rev [simp] : word-reverse (word-reverse w) = w
  unfolding word-reverse-def by (simp add: word-bl.Abs-inverse)

lemma word-rev-gal: word-reverse w = u ==> word-reverse u = w
  by auto

lemmas word-rev-gal' = sym [THEN word-rev-gal, symmetric, standard]

lemmas length-bl-gt-0 [iff] = xtr1 [OF word-bl.Rep' len-gt-0, standard]
lemmas bl-not-Nil [iff] =
  length-bl-gt-0 [THEN length-greater-0-conv [THEN iffD1], standard]
lemmas length-bl-neq-0 [iff] = length-bl-gt-0 [THEN gr-implies-not0]

lemma hd-bl-sign-sint: hd (to-bl w) = (bin-sign (sint w) = Int.Min)
  apply (unfold to-bl-def sint-uint)
  apply (rule trans [OF - bl-sbin-sign])
  apply simp
  done

lemma of-bl-drop':
  lend = length bl - len-of TYPE ('a :: len0) ==>
    of-bl (drop lend bl) = (of-bl bl :: 'a word)
  apply (unfold of-bl-def)
  apply (clarsimp simp add: trunc-bl2bin [symmetric])
  done

lemmas of-bl-no = of-bl-def [folded word-number-of-def]

```

```

lemma test-bit-of-bl:
  (of-bl bl::'a::len0 word) !! n = (rev bl ! n ∧ n < len-of TYPE('a) ∧ n < length
  bl)
  apply (unfold of-bl-def word-test-bit-def)
  apply (auto simp add: word-size word-ubin.eq-norm nth-bintr bin-nth-of-bl)
  done

lemma no-of-bl:
  (number-of bin ::'a::len0 word) = of-bl (bin-to-bl (len-of TYPE('a)) bin)
  unfolding word-size of-bl-no by (simp add : word-number-of-def)

lemma uint-bl: to-bl w == bin-to-bl (size w) (uint w)
  unfolding word-size to-bl-def by auto

lemma to-bl-bin: bl-to-bin (to-bl w) = uint w
  unfolding uint-bl by (simp add : word-size)

lemma to-bl-of-bin:
  to-bl (word-of-int bin::'a::len0 word) = bin-to-bl (len-of TYPE('a)) bin
  unfolding uint-bl by (clarsimp simp add: word-ubin.eq-norm word-size)

lemmas to-bl-no-bin [simp] = to-bl-of-bin [folded word-number-of-def]

lemma to-bl-to-bin [simp] : bl-to-bin (to-bl w) = uint w
  unfolding uint-bl by (simp add : word-size)

lemmas uint-bl-bin [simp] = trans [OF bin-bl-bin word-ubin.norm-Rep, standard]

lemmas num-AB-u [simp] = word-uint.Rep-inverse
  [unfolded o-def word-number-of-def [symmetric], standard]
lemmas num-AB-s [simp] = word-sint.Rep-inverse
  [unfolded o-def word-number-of-def [symmetric], standard]

lemma uints-unats: uints n = int ‘ unats n
  apply (unfold unats-def uints-num)
  apply safe
  apply (rule-tac image-eqI)
  apply (erule-tac nat-0-le [symmetric])
  apply auto
  apply (erule-tac nat-less-iff [THEN iffD2])
  apply (rule-tac [2] zless-nat-eq-int-zless [THEN iffD1])
  apply (auto simp add : nat-power-eq int-power)
  done

lemma unats-uints: unats n = nat ‘ uints n
  by (auto simp add : uints-unats image-iff)

```

lemmas *bintr-num* = *word-ubin.norm-eq-iff*
 [*symmetric*, *folded word-number-of-def*, *standard*]

lemmas *sbintr-num* = *word-sbin.norm-eq-iff*
 [*symmetric*, *folded word-number-of-def*, *standard*]

lemmas *num-of-bintr* = *word-ubin.Abs-norm* [*folded word-number-of-def*, *standard*]

lemmas *num-of-sbintr* = *word-sbin.Abs-norm* [*folded word-number-of-def*, *standard*]

lemma *num-of-bintr'*:
bintrunc (*len-of TYPE*('a :: len0)) *a* = *b* ==>
number-of a = (*number-of b* :: 'a word)
apply *safe*
apply (*rule-tac num-of-bintr* [*symmetric*])
done

lemma *num-of-sbintr'*:
sbintrunc (*len-of TYPE*('a :: len) - 1) *a* = *b* ==>
number-of a = (*number-of b* :: 'a word)
apply *safe*
apply (*rule-tac num-of-sbintr* [*symmetric*])
done

lemmas *num-abs-bintr* = *sym* [*THEN trans*,
OF num-of-bintr word-number-of-def, *standard*]

lemmas *num-abs-sbintr* = *sym* [*THEN trans*,
OF num-of-sbintr word-number-of-def, *standard*]

lemma *ucast-id*: *ucast w* = *w*
unfolding *ucast-def* **by** *auto*

lemma *scast-id*: *scast w* = *w*
unfolding *scast-def* **by** *auto*

lemma *ucast-bl*: *ucast w* == *of-bl* (*to-bl w*)
unfolding *ucast-def of-bl-def uint-bl*
by (*auto simp add* : *word-size*)

lemma *nth-ucast*:
(*ucast w*::'a::len0 word) !! *n* = (*w* !! *n* & *n* < *len-of TYPE*('a))
apply (*unfold ucast-def test-bit-bin*)
apply (*simp add*: *word-ubin.eq-norm nth-bintr word-size*)
apply (*fast elim*!: *bin-nth-uint-imp*)
done

```

lemma ucast-bintr [simp]:
  ucast (number-of w :: 'a::len0 word) =
    number-of (bintrunc (len-of TYPE('a)) w)
unfolding ucast-def by simp

lemma scast-sbintr [simp]:
  scast (number-of w :: 'a::len word) =
    number-of (sbintrunc (len-of TYPE('a) - Suc 0) w)
unfolding scast-def by simp

lemmas source-size = source-size-def [unfolded Let-def word-size]
lemmas target-size = target-size-def [unfolded Let-def word-size]
lemmas is-down = is-down-def [unfolded source-size target-size]
lemmas is-up = is-up-def [unfolded source-size target-size]

lemmas is-up-down = trans [OF is-up is-down [symmetric], standard]

lemma down-cast-same': uc = ucast ==> is-down uc ==> uc = scast
apply (unfold is-down)
apply safe
apply (rule ext)
apply (unfold ucast-def scast-def uint-sint)
apply (rule word-ubin.norm-eq-iff [THEN iffD1])
apply simp
done

lemma word-rev-tf':
  r = to-bl (of-bl bl) ==> r = rev (takefill False (length r) (rev bl))
unfolding of-bl-def uint-bl
by (clarsimp simp add: bl-bin-bl-rtf word-ubin.eq-norm word-size)

lemmas word-rev-tf = refl [THEN word-rev-tf', unfolded word-bl.Rep', standard]

lemmas word-rep-drop = word-rev-tf [simplified takefill-alt,
  simplified, simplified rev-take, simplified]

lemma to-bl-ucast:
  to-bl (ucast (w::'b::len0 word) :: 'a::len0 word) =
    replicate (len-of TYPE('a) - len-of TYPE('b)) False @
    drop (len-of TYPE('b) - len-of TYPE('a)) (to-bl w)
apply (unfold ucast-bl)
apply (rule trans)
apply (rule word-rep-drop)
apply simp
done

```

lemma *ucast-up-app'*:

uc = *ucast* ==> *source-size uc* + *n* = *target-size uc* ==>
to-bl (uc w) = *replicate n False @ (to-bl w)*
by (*auto simp add : source-size target-size to-bl-ucast*)

lemma *ucast-down-drop'*:

uc = *ucast* ==> *source-size uc* = *target-size uc* + *n* ==>
to-bl (uc w) = *drop n (to-bl w)*
by (*auto simp add : source-size target-size to-bl-ucast*)

lemma *scast-down-drop'*:

sc = *scast* ==> *source-size sc* = *target-size sc* + *n* ==>
to-bl (sc w) = *drop n (to-bl w)*
apply (*subgoal-tac sc = ucast*)
apply *safe*
apply *simp*
apply (*erule refl [THEN ucast-down-drop']*)
apply (*rule refl [THEN down-cast-same', symmetric]*)
apply (*simp add : source-size target-size is-down*)
done

lemma *sint-up-scast'*:

sc = *scast* ==> *is-up sc* ==> *sint (sc w)* = *sint w*
apply (*unfold is-up*)
apply *safe*
apply (*simp add: scast-def word-sbin.eq-norm*)
apply (*rule box-equals*)
prefer 3
apply (*rule word-sbin.norm-Rep*)
apply (*rule sbintrunc-sbintrunc-l*)
defer
apply (*subst word-sbin.norm-Rep*)
apply (*rule refl*)
apply *simp*
done

lemma *uint-up-ucast'*:

uc = *ucast* ==> *is-up uc* ==> *uint (uc w)* = *uint w*
apply (*unfold is-up*)
apply *safe*
apply (*rule bin-eqI*)
apply (*fold word-test-bit-def*)
apply (*auto simp add: nth-ucast*)
apply (*auto simp add: test-bit-bin*)
done

lemmas *down-cast-same* = *refl [THEN down-cast-same']*

lemmas *ucast-up-app* = *refl [THEN ucast-up-app']*

lemmas *ucast-down-drop* = *refl [THEN ucast-down-drop']*

lemmas *scast-down-drop* = refl [THEN *scast-down-drop*']

lemmas *uint-up-ucast* = refl [THEN *uint-up-ucast*']

lemmas *sint-up-scast* = refl [THEN *sint-up-scast*']

lemma *ucast-up-ucast'*: *uc* = *ucast* ==> *is-up uc* ==> *ucast (uc w)* = *ucast w*
apply (*simp* (*no-asm*) *add*: *ucast-def*)
apply (*clarsimp simp add*: *uint-up-ucast*)
done

lemma *scast-up-scast'*: *sc* = *scast* ==> *is-up sc* ==> *scast (sc w)* = *scast w*
apply (*simp* (*no-asm*) *add*: *scast-def*)
apply (*clarsimp simp add*: *sint-up-scast*)
done

lemma *ucast-of-bl-up'*:
w = *of-bl bl* ==> *size bl* <= *size w* ==> *ucast w* = *of-bl bl*
by (*auto simp add* : *nth-ucast word-size test-bit-of-bl intro!*: *word-eqI*)

lemmas *ucast-up-ucast* = refl [THEN *ucast-up-ucast*']

lemmas *scast-up-scast* = refl [THEN *scast-up-scast*']

lemmas *ucast-of-bl-up* = refl [THEN *ucast-of-bl-up*']

lemmas *ucast-up-ucast-id* = trans [OF *ucast-up-ucast ucast-id*]

lemmas *scast-up-scast-id* = trans [OF *scast-up-scast scast-id*]

lemmas *isduu* = *is-up-down* [where *c* = *ucast*, THEN *iffD2*]

lemmas *isdus* = *is-up-down* [where *c* = *scast*, THEN *iffD2*]

lemmas *ucast-down-ucast-id* = *isduu* [THEN *ucast-up-ucast-id*]

lemmas *scast-down-scast-id* = *isdus* [THEN *ucast-up-ucast-id*]

lemma *up-ucast-surj*:

is-up (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

surj (*ucast* :: 'a word => 'b word)

by (*rule surjI*, *erule ucast-up-ucast-id*)

lemma *up-scast-surj*:

is-up (*scast* :: 'b::len word => 'a::len word) ==>

surj (*scast* :: 'a word => 'b word)

by (*rule surjI*, *erule scast-up-scast-id*)

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

by (*rule inj-on-inverseI*, *erule scast-down-scast-id*)

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

by (*rule inj-on-inverseI*, *erule ucast-down-ucast-id*)


```

lemma of-bl-append-same: of-bl (X @ to-bl w) = w
  by (rule word-bl.Rep-eqD) (simp add: word-rep-drop)

lemma ucast-down-no':
  uc = ucast ==> is-down uc ==> uc (number-of bin) = number-of bin
  apply (unfold word-number-of-def is-down)
  apply (clarsimp simp add: ucast-def word-ubin.eq-norm)
  apply (rule word-ubin.norm-eq-iff [THEN iffD1])
  apply (erule bintrunc-bintrunc-ge)
  done

lemmas ucast-down-no = ucast-down-no' [OF refl]

lemma ucast-down-bl': uc = ucast ==> is-down uc ==> uc (of-bl bl) = of-bl bl
  unfolding of-bl-no by clarify (erule ucast-down-no)

lemmas ucast-down-bl = ucast-down-bl' [OF refl]

lemmas slice-def' = slice-def [unfolded word-size]
lemmas test-bit-def' = word-test-bit-def [THEN fun-cong]

lemmas word-log-defs = word-and-def word-or-def word-xor-def word-not-def
lemmas word-log-bin-defs = word-log-defs

Executable equality
instantiation word :: ({len0}) eq
begin

definition eq-word :: 'a word  $\Rightarrow$  'a word  $\Rightarrow$  bool where
  eq-word k l  $\longleftrightarrow$  HOL.eq (uint k) (uint l)

instance proof
qed (simp add: eq eq-word-def)

end

end

```

11 WordArith: Word Arithmetic

```

theory WordArith
imports WordDefinition
begin

lemma word-less-alt: (a < b) = (uint a < uint b)
  unfolding word-less-def word-le-def
  by (auto simp del: word-uint.Rep-inject)

```

simp: *word-uint.Rep-inject* [*symmetric*])

lemma *signed-linorder*: *class.linorder word-sle word-sless*
proof
qed (*unfold word-sle-def word-sless-def, auto*)

interpretation *signed*: *linorder word-sle word-sless*
by (*rule signed-linorder*)

lemmas *word-arith-wis* =
word-add-def word-mult-def word-minus-def
word-succ-def word-pred-def word-0-wi word-1-wi

lemma *udvdI*:
 $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$
by (*auto simp: udvd-def*)

lemmas *word-div-no* [*simp*] =
word-div-def [*of number-of a number-of b, standard*]

lemmas *word-mod-no* [*simp*] =
word-mod-def [*of number-of a number-of b, standard*]

lemmas *word-less-no* [*simp*] =
word-less-def [*of number-of a number-of b, standard*]

lemmas *word-le-no* [*simp*] =
word-le-def [*of number-of a number-of b, standard*]

lemmas *word-sless-no* [*simp*] =
word-sless-def [*of number-of a number-of b, standard*]

lemmas *word-sle-no* [*simp*] =
word-sle-def [*of number-of a number-of b, standard*]

lemmas *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]
lemmas *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

lemma *int-one-bin*: $(1 :: \text{int}) == (\text{Int.Pls BIT bit.B1})$
unfolding *Pls-def Bit-def* **by** *auto*

lemma *word-1-no*:
 $(1 :: 'a :: \text{len0 word}) == \text{number-of } (\text{Int.Pls BIT bit.B1})$
unfolding *word-1-wi word-number-of-def int-one-bin* **by** *auto*

lemma *word-m1-wi*: $-1 == \text{word-of-int } -1$
by (*rule word-number-of-alt*)

```

lemma word-m1-wi-Min:  $-1 = \text{word-of-int Int.Min}$ 
  by (simp add: word-m1-wi number-of-eq)

lemma word-0-bl: of-bl [] = 0
  unfolding word-0-wi of-bl-def by (simp add : Pls-def)

lemma word-1-bl: of-bl [True] = 1
  unfolding word-1-wi of-bl-def
  by (simp add : bl-to-bin-def Bit-def Pls-def)

lemma uint-0 [simp] : (uint 0 = 0)
  unfolding word-0-wi
  by (simp add: word-ubin.eq-norm Pls-def [symmetric])

lemma of-bl-0 [simp] : of-bl (replicate n False) = 0
  by (simp add : word-0-wi of-bl-def bl-to-bin-rep-False Pls-def)

lemma to-bl-0:
  to-bl (0::'a::len0 word) = replicate (len-of TYPE('a)) False
  unfolding uint-bl
  by (simp add : word-size bin-to-bl-Pls Pls-def [symmetric])

lemma uint-0-iff: (uint x = 0) = (x = 0)
  by (auto intro!: word-uint.Rep-eqD)

lemma unat-0-iff: (unat x = 0) = (x = 0)
  unfolding unat-def by (auto simp add : nat-eq-iff uint-0-iff)

lemma unat-0 [simp]: unat 0 = 0
  unfolding unat-def by auto

lemma size-0-same': size w = 0 ==> w = (v :: 'a :: len0 word)
  apply (unfold word-size)
  apply (rule box-equals)
  defer
  apply (rule word-uint.Rep-inverse)+
  apply (rule word-ubin.norm-eq-iff [THEN iffD1])
  apply simp
  done

lemmas size-0-same = size-0-same' [folded word-size]

lemmas unat-eq-0 = unat-0-iff
lemmas unat-eq-zero = unat-0-iff

lemma unat-gt-0: (0 < unat x) = (x ~ = 0)
by (auto simp: unat-0-iff [symmetric])

lemma ucast-0 [simp] : ucast 0 = 0

```

unfolding *ucast-def*
by *simp (simp add: word-0-wi)*

lemma *sint-0 [simp] : sint 0 = 0*
unfolding *sint-uint*
by *(simp add: Pls-def [symmetric])*

lemma *scast-0 [simp] : scast 0 = 0*
apply *(unfold scast-def)*
apply *simp*
apply *(simp add: word-0-wi)*
done

lemma *sint-n1 [simp] : sint -1 = -1*
apply *(unfold word-m1-wi-Min)*
apply *(simp add: word-sbin.eq-norm)*
apply *(unfold Min-def number-of-eq)*
apply *simp*
done

lemma *scast-n1 [simp] : scast -1 = -1*
apply *(unfold scast-def sint-n1)*
apply *(unfold word-number-of-alt)*
apply *(rule refl)*
done

lemma *uint-1 [simp] : uint (1 :: 'a :: len word) = 1*
unfolding *word-1-wi*
by *(simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)*

lemma *unat-1 [simp] : unat (1 :: 'a :: len word) = 1*
by *(unfold unat-def uint-1) auto*

lemma *ucast-1 [simp] : ucast (1 :: 'a :: len word) = 1*
unfolding *ucast-def word-1-wi*
by *(simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)*

lemmas *ariths =*
bintr-ariths [THEN word-ubin.norm-eq-iff [THEN iffD1],
folded word-ubin.eq-norm, standard]

lemma *wi-homs:*
shows
wi-hom-add: word-of-int a + word-of-int b = word-of-int (a + b) and
*wi-hom-mult: word-of-int a * word-of-int b = word-of-int (a * b) and*
wi-hom-neg: - word-of-int a = word-of-int (- a) and
wi-hom-succ: word-succ (word-of-int a) = word-of-int (Int.succ a) and

wi-hom-pred: *word-pred* (*word-of-int* *a*) = *word-of-int* (*Int.pred* *a*)
by (*auto simp*: *word-arith-wis arths*)

lemmas *wi-hom-syms* = *wi-homs* [*symmetric*]

lemma *word-sub-def*: $a - b == a + - (b :: 'a :: \text{len0 word})$
unfolding *word-sub-wi diff-def*
by (*simp only* : *word-uint.Rep-inverse wi-hom-syms*)

lemmas *word-diff-minus* = *word-sub-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-of-int-sub-hom*:
 $(\text{word-of-int } a) - \text{word-of-int } b = \text{word-of-int } (a - b)$
unfolding *word-sub-def diff-def* **by** (*simp only* : *wi-homs*)

lemmas *new-word-of-int-homs* =
word-of-int-sub-hom wi-homs word-0-wi word-1-wi

lemmas *new-word-of-int-hom-syms* = *new-word-of-int-homs* [*symmetric, standard*]

lemmas *word-of-int-hom-syms* =
new-word-of-int-hom-syms [*unfolded succ-def pred-def*]

lemmas *word-of-int-homs* =
new-word-of-int-homs [*unfolded succ-def pred-def*]

lemmas *word-of-int-add-hom* = *word-of-int-homs* (2)
lemmas *word-of-int-mult-hom* = *word-of-int-homs* (3)
lemmas *word-of-int-minus-hom* = *word-of-int-homs* (4)
lemmas *word-of-int-succ-hom* = *word-of-int-homs* (5)
lemmas *word-of-int-pred-hom* = *word-of-int-homs* (6)
lemmas *word-of-int-0-hom* = *word-of-int-homs* (7)
lemmas *word-of-int-1-hom* = *word-of-int-homs* (8)

lemmas *word-arith-alts* =
word-sub-wi [*unfolded succ-def pred-def, standard*]
word-arith-wis [*unfolded succ-def pred-def, standard*]

lemmas *word-sub-alt* = *word-arith-alts* (1)
lemmas *word-add-alt* = *word-arith-alts* (2)
lemmas *word-mult-alt* = *word-arith-alts* (3)
lemmas *word-minus-alt* = *word-arith-alts* (4)
lemmas *word-succ-alt* = *word-arith-alts* (5)
lemmas *word-pred-alt* = *word-arith-alts* (6)
lemmas *word-0-alt* = *word-arith-alts* (7)
lemmas *word-1-alt* = *word-arith-alts* (8)

11.1 Transferring goals from words to ints

lemma *word-ths*:

shows

word-succ-p1: $\text{word-succ } a = a + 1$ **and**

word-pred-m1: $\text{word-pred } a = a - 1$ **and**

word-pred-succ: $\text{word-pred } (\text{word-succ } a) = a$ **and**

word-succ-pred: $\text{word-succ } (\text{word-pred } a) = a$ **and**

word-mult-succ: $\text{word-succ } a * b = b + a * b$

by (*rule* *word-uint.Abs-cases* [of *b*],

rule *word-uint.Abs-cases* [of *a*],

simp add: *pred-def succ-def add-commute mult-commute*

ring-distrib new-word-of-int-homs) +

lemmas *uint-cong* = *arg-cng* [**where** *f* = *uint*]

lemmas *uint-word-ariths* =

word-arith-alt [*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

lemmas *uint-word-arith-bintr* = *uint-word-ariths* [*folded bintrunc-mod2p*]

lemmas *sint-word-ariths* = *uint-word-arith-bintr*

[*THEN uint-sint* [*symmetric*, *THEN trans*],

unfolded uint-sint bintr-arith1s bintr-ariths

len-gt-0 [*THEN bin-sbin-eq-iff*] *word-sbin.norm-Rep*, *standard*]

lemmas *uint-div-alt* = *word-div-def*

[*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

lemmas *uint-mod-alt* = *word-mod-def*

[*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

lemma *word-pred-0-n1*: $\text{word-pred } 0 = \text{word-of-int } -1$

unfolding *word-pred-def number-of-eq*

by (*simp add* : *pred-def word-no-wi*)

lemma *word-pred-0-Min*: $\text{word-pred } 0 = \text{word-of-int } \text{Int.Min}$

by (*simp add*: *word-pred-0-n1 number-of-eq*)

lemma *word-m1-Min*: $-1 = \text{word-of-int } \text{Int.Min}$

unfolding *Min-def* **by** (*simp only*: *word-of-int-hom-syms*)

lemma *succ-pred-no* [*simp*]:

$\text{word-succ } (\text{number-of } \text{bin}) = \text{number-of } (\text{Int.succ } \text{bin}) \ \&$

$\text{word-pred } (\text{number-of } \text{bin}) = \text{number-of } (\text{Int.pred } \text{bin})$

unfolding *word-number-of-def* **by** (*simp add* : *new-word-of-int-homs*)

lemma *word-sp-01* [*simp*] :

$\text{word-succ } -1 = 0 \ \& \ \text{word-succ } 0 = 1 \ \& \ \text{word-pred } 0 = -1 \ \& \ \text{word-pred } 1 = 0$

by (*unfold word-0-no word-1-no*) *auto*

lemma *word-of-int-Ex*:

$\exists y. x = \text{word-of-int } y$
by (*rule-tac* $x = \text{uint } x$ **in** *exI*) *simp*

lemma *word-arith-egs*:

fixes $a :: 'a::\text{len0 word}$
fixes $b :: 'a::\text{len0 word}$
shows
word-add-0: $0 + a = a$ **and**
word-add-0-right: $a + 0 = a$ **and**
word-mult-1: $1 * a = a$ **and**
word-mult-1-right: $a * 1 = a$ **and**
word-add-commute: $a + b = b + a$ **and**
word-add-assoc: $a + b + c = a + (b + c)$ **and**
word-add-left-commute: $a + (b + c) = b + (a + c)$ **and**
word-mult-commute: $a * b = b * a$ **and**
word-mult-assoc: $a * b * c = a * (b * c)$ **and**
word-mult-left-commute: $a * (b * c) = b * (a * c)$ **and**
word-left-distrib: $(a + b) * c = a * c + b * c$ **and**
word-right-distrib: $a * (b + c) = a * b + a * c$ **and**
word-left-minus: $- a + a = 0$ **and**
word-diff-0-right: $a - 0 = a$ **and**
word-diff-self: $a - a = 0$
using *word-of-int-Ex* [*of a*]
 word-of-int-Ex [*of b*]
 word-of-int-Ex [*of c*]
by (*auto simp: word-of-int-hom-syms* [*symmetric*]
 zadd-0-right add-commute add-assoc add-left-commute
 mult-commute mult-assoc mult-left-commute
 left-distrib right-distrib)

lemmas *word-add-ac* = *word-add-commute word-add-assoc word-add-left-commute*

lemmas *word-mult-ac* = *word-mult-commute word-mult-assoc word-mult-left-commute*

lemmas *word-plus-ac0* = *word-add-0 word-add-0-right word-add-ac*

lemmas *word-times-ac1* = *word-mult-1 word-mult-1-right word-mult-ac*

11.2 Order on fixed-length words

lemma *word-order-trans*: $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

lemma *word-order-refl*: $z \leq (z :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

lemma *word-order-antisym*: $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** (*auto intro!: word-uint.Rep-eqD*)

```

lemma word-order-linear:
   $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$ 
  unfolding word-le-def by auto

lemma word-zero-le [simp] :
   $0 \leq (y :: 'a :: \text{len0 word})$ 
  unfolding word-le-def by auto

instance word :: (len0) semigroup-add
  by intro-classes (simp add: word-add-assoc)

instance word :: (len0) linorder
  by intro-classes (auto simp: word-less-def word-le-def)

instance word :: (len0) ring
  by intro-classes
    (auto simp: word-arith-eqs word-diff-minus
      word-diff-self [unfolded word-diff-minus])

lemma word-m1-ge [simp] : word-pred 0  $\geq$  y
  unfolding word-le-def
  by (simp only : word-pred-0-n1 word-uint.eq-norm m1mod2k) auto

lemmas word-n1-ge [simp] = word-m1-ge [simplified word-sp-01]

lemmas word-not-simps [simp] =
  word-zero-le [THEN leD] word-m1-ge [THEN leD] word-n1-ge [THEN leD]

lemma word-gt-0:  $0 < y = (0 \sim (y :: 'a :: \text{len0 word}))$ 
  unfolding word-less-def by auto

lemmas word-gt-0-no [simp] = word-gt-0 [of number-of y, standard]

lemma word-sless-alt:  $(a <_s b) == (\text{sint } a < \text{sint } b)$ 
  unfolding word-sle-def word-sless-def
  by (auto simp add: less-le)

lemma word-le-nat-alt:  $(a \leq b) = (\text{unat } a \leq \text{unat } b)$ 
  unfolding unat-def word-le-def
  by (rule nat-le-eq-zle [symmetric]) simp

lemma word-less-nat-alt:  $(a < b) = (\text{unat } a < \text{unat } b)$ 
  unfolding unat-def word-less-alt
  by (rule nat-less-eq-zless [symmetric]) simp

lemma wi-less:
   $(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) =$ 
   $(n \bmod 2^{\text{len-of TYPE('a)}} < m \bmod 2^{\text{len-of TYPE('a)}})$ 

```



```

unfolding word-less-alt by (simp add: word-uint.eq-norm)

lemma wi-le:
  (word-of-int n <= (word-of-int m :: 'a :: len0 word)) =
    (n mod 2 ^ len-of TYPE('a) <= m mod 2 ^ len-of TYPE('a))
unfolding word-le-def by (simp add: word-uint.eq-norm)

lemma udvd-nat-alt: a udvd b = (EX n>=0. unat b = n * unat a)
apply (unfold udvd-def)
apply safe
apply (simp add: unat-def nat-mult-distrib)
apply (simp add: uint-nat int-mult)
apply (rule exI)
apply safe
prefer 2
apply (erule notE)
apply (rule refl)
apply force
done

lemma udvd-iff-dvd: x udvd y <-> unat x dvd unat y
unfolding dvd-def udvd-nat-alt by force

lemmas unat-mono = word-less-nat-alt [THEN iffD1, standard]

lemma word-zero-neq-one: 0 < len-of TYPE ('a :: len0) ==> (0 :: 'a word) ~=
1
unfolding word-arith-wis
by (auto simp add: word-ubin.norm-eq-iff [symmetric] gr0-conv-Suc)

lemmas lenw1-zero-neq-one = len-gt-0 [THEN word-zero-neq-one]

lemma no-no [simp] : number-of (number-of b) = number-of b
by (simp add: number-of-eq)

lemma unat-minus-one: x ~= 0 ==> unat (x - 1) = unat x - 1
apply (unfold unat-def)
apply (simp only: int-word-uint word-arith-alts rdmods)
apply (subgoal-tac uint x >= 1)
prefer 2
apply (drule contrapos-nn)
apply (erule word-uint.Rep-inverse' [symmetric])
apply (insert uint-ge-0 [of x])[1]
apply arith
apply (rule box-equals)
apply (rule nat-diff-distrib)
prefer 2
apply assumption
apply simp

```

```

apply (subst mod-pos-pos-trivial)
  apply arith
  apply (insert uint-lt2p [of x])[1]
  apply arith
  apply (rule refl)
apply simp
done

```

```

lemma measure-unat:  $p \sim = 0 \implies \text{unat } (p - 1) < \text{unat } p$ 
  by (simp add: unat-minus-one) (simp add: unat-0-iff [symmetric])

```

```

lemmas uint-add-ge0 [simp] =
  add-nonneg-nonneg [OF uint-ge-0 uint-ge-0, standard]
lemmas uint-mult-ge0 [simp] =
  mult-nonneg-nonneg [OF uint-ge-0 uint-ge-0, standard]

```

```

lemma uint-sub-lt2p [simp]:
  uint (x :: 'a :: len0 word) - uint (y :: 'b :: len0 word) <
    2 ^ len-of TYPE('a)
  using uint-ge-0 [of y] uint-lt2p [of x] by arith

```

11.3 Conditions for the addition (etc) of two words to overflow

```

lemma uint-add-lem:
  (uint x + uint y < 2 ^ len-of TYPE('a)) =
  (uint (x + y :: 'a :: len0 word) = uint x + uint y)
  by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-mult-lem:
  (uint x * uint y < 2 ^ len-of TYPE('a)) =
  (uint (x * y :: 'a :: len0 word) = uint x * uint y)
  by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-sub-lem:
  (uint x >= uint y) = (uint (x - y) = uint x - uint y)
  by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-add-le: uint (x + y) <= uint x + uint y
  unfolding uint-word-ariths by (auto simp: mod-add-if-z)

```

```

lemma uint-sub-ge: uint (x - y) >= uint x - uint y
  unfolding uint-word-ariths by (auto simp: mod-sub-if-z)

```

```

lemmas uint-sub-if' =
  trans [OF uint-word-ariths(1) mod-sub-if-z, simplified, standard]
lemmas uint-plus-if' =
  trans [OF uint-word-ariths(2) mod-add-if-z, simplified, standard]

```

11.4 Definition of uint_arith

lemma *word-of-int-inverse*:

```
word-of-int r = a ==> 0 <= r ==> r < 2 ^ len-of TYPE('a) ==>
  uint (a::'a::len0 word) = r
apply (erule word-uint.Abs-inverse' [rotated])
apply (simp add: uints-num)
done
```

lemma *uint-split*:

```
fixes x::'a::len0 word
shows P (uint x) =
  (ALL i. word-of-int i = x & 0 <= i & i < 2^len-of TYPE('a) --> P i)
apply (fold word-int-case-def)
apply (auto dest!: word-of-int-inverse simp: int-word-uint int-mod-eq'
  split: word-int-split)
done
```

lemma *uint-split-asm*:

```
fixes x::'a::len0 word
shows P (uint x) =
  (~ (EX i. word-of-int i = x & 0 <= i & i < 2^len-of TYPE('a) & ~ P i))
by (auto dest!: word-of-int-inverse
  simp: int-word-uint int-mod-eq'
  split: uint-split)
```

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =

```
word-le-def word-less-alt
word-uint.Rep-inject [symmetric]
uint-sub-if' uint-plus-if'
```

lemma *power-False-cong*: *False ==> a ^ b = c ^ d*

by *auto*

ML <<

```
fun uint-arith-ss-of ss =
  ss addsimps @{thms uint-arith-simps}
  delsimps @{thms word-uint.Rep-inject}
  addsplits @{thms split-if-asm}
  addcongs @{thms power-False-cong}
```

fun *uint-arith-tacs* *ctxt* =

```
let
fun arith-tac' n t =
  Arith-Data.verbose-arith-tac ctxt n t
  handle Cooper.COOPER - => Seq.empty;
```

```

    val cs = claset-of ctxt;
    val ss = simpset-of ctxt;
  in
    [ clarify-tac cs 1,
      full-simp-tac (uint-arith-ss-of ss) 1,
      ALLGOALS (full-simp-tac (HOL-ss addsplits @ {thms uint-splits}
                               addcongs @ {thms power-False-cong})),
      rewrite-goals-tac @ {thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
                           REPEAT (etac conjE n) THEN
                           REPEAT (dtac @ {thm word-of-int-inverse} n
                                   THEN atac n
                                   THEN atac n)),
      TRYALL arith-tac' ]
  end

fun uint-arith-tac ctxt = SELECT-GOAL (EVERY (uint-arith-tacs ctxt))
>>

method-setup uint-arith =
  << Scan.succeed (SIMPLE-METHOD' o uint-arith-tac) >>
  solving word arithmetic via integers and arith

```

11.5 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*:

$((x :: 'a :: \text{len0 word}) \leq x + y) = (\text{uint } x + \text{uint } y < 2^{\text{size } x})$
unfolding *word-size* **by** *uint-arith*

lemmas *no-olen-add* = *no-plus-overflow-uint-size* [unfolded *word-size*]

lemma *no-ulen-sub*: $((x :: 'a :: \text{len0 word}) \geq x - y) = (\text{uint } y \leq \text{uint } x)$
by *uint-arith*

lemma *no-olen-add'*:

fixes $x :: 'a :: \text{len0 word}$
shows $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)})}$
by (*simp add: word-add-ac add-ac no-olen-add*)

lemmas *olen-add-equiv* = *trans* [*OF no-olen-add no-olen-add'* [*symmetric*], *standard*]

lemmas *uint-plus-simple-iff* = *trans* [*OF no-olen-add uint-add-lem, standard*]

lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1, standard*]

lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem, standard*]

lemmas *uint-minus-simple-alt* = *uint-sub-lem* [folded *word-le-def*]

lemmas *word-sub-le-iff* = *no-ulen-sub* [folded *word-le-def*]

lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2, standard*]

lemma *word-less-sub1*:

$$(x :: 'a :: \text{len0 word}) \sim = 0 ==> (1 < x) = (0 < x - 1)$$

by *uint-arith*

lemma *word-le-sub1*:

$$(x :: 'a :: \text{len0 word}) \sim = 0 ==> (1 <= x) = (0 <= x - 1)$$

by *uint-arith*

lemma *sub-wrap-lt*:

$$((x :: 'a :: \text{len0 word}) < x - z) = (x < z)$$

by *uint-arith*

lemma *sub-wrap*:

$$((x :: 'a :: \text{len0 word}) <= x - z) = (z = 0 \mid x < z)$$

by *uint-arith*

lemma *plus-minus-not-NULL-ab*:

$$(x :: 'a :: \text{len0 word}) <= ab - c ==> c <= ab ==> c \sim = 0 ==> x + c \sim = 0$$

by *uint-arith*

lemma *plus-minus-no-overflow-ab*:

$$(x :: 'a :: \text{len0 word}) <= ab - c ==> c <= ab ==> x <= x + c$$

by *uint-arith*

lemma *le-minus'*:

$$(a :: 'a :: \text{len0 word}) + c <= b ==> a <= a + c ==> c <= b - a$$

by *uint-arith*

lemma *le-plus'*:

$$(a :: 'a :: \text{len0 word}) <= b ==> c <= b - a ==> a + c <= b$$

by *uint-arith*

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*, standard]

lemma *word-plus-mono-right*:

$$(y :: 'a :: \text{len0 word}) <= z ==> x <= x + z ==> x + y <= x + z$$

by *uint-arith*

lemma *word-less-minus-cancel*:

$$y - x < z - x ==> x <= z ==> (y :: 'a :: \text{len0 word}) < z$$

by *uint-arith*

lemma *word-less-minus-mono-left*:

$$(y :: 'a :: \text{len0 word}) < z ==> x <= y ==> y - x < z - x$$

by *uint-arith*

lemma *word-less-minus-mono*:

$a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - (d :: 'a :: \text{len word})$
by *uint-arith*

lemma *word-le-minus-cancel*:

$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) \leq z$
by *uint-arith*

lemma *word-le-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq y \implies y - x \leq z - x$
by *uint-arith*

lemma *word-le-minus-mono*:

$a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c$
 $\implies a - b \leq c - (d :: 'a :: \text{len word})$
by *uint-arith*

lemma *plus-le-left-cancel-wrap*:

$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$
by *uint-arith*

lemma *plus-le-left-cancel-nowrap*:

$(x :: 'a :: \text{len0 word}) \leq x + y' \implies x \leq x + y \implies$
 $(x + y' < x + y) = (y' < y)$
by *uint-arith*

lemma *word-plus-mono-right2*:

$(a :: 'a :: \text{len0 word}) \leq a + b \implies c \leq b \implies a \leq a + c$
by *uint-arith*

lemma *word-less-add-right*:

$(x :: 'a :: \text{len0 word}) < y - z \implies z \leq y \implies x + z < y$
by *uint-arith*

lemma *word-less-sub-right*:

$(x :: 'a :: \text{len0 word}) < y + z \implies y \leq x \implies x - y < z$
by *uint-arith*

lemma *word-le-plus-either*:

$(x :: 'a :: \text{len0 word}) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$
by *uint-arith*

lemma *word-less-nowrapI*:

$(x :: 'a :: \text{len0 word}) < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
by *uint-arith*

lemma *inc-le*: $(i :: 'a :: \text{len word}) < m \implies i + 1 \leq m$

by *uint-arith*

lemma *inc-i*:

$(1 :: 'a :: \text{len word}) \leq i \implies i < m \implies 1 \leq (i + 1) \ \& \ i + 1 \leq m$
by *uint-arith*

lemma *udvd-incr-lem*:

$up < uq \implies up = ua + n * \text{uint } K \implies$
 $uq = ua + n' * \text{uint } K \implies up + \text{uint } K \leq uq$
apply *clarsimp*
apply (*drule less-le-mult*)
apply *safe*
done

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K \leq q$
apply (*unfold word-less-alt word-le-def*)
apply (*drule* (2) *udvd-incr-lem*)
apply (*erule uint-add-le [THEN order-trans]*)
done

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p \leq q - K$
apply (*unfold word-less-alt word-le-def*)
apply (*drule* (2) *udvd-incr-lem*)
apply (*drule le-diff-eq [THEN iffD2]*)
apply (*erule order-trans*)
apply (*rule uint-sub-ge*)
done

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [**where** *ua=0, simplified*]

lemmas *udvd-incr0* = *udvd-incr'* [**where** *ua=0, simplified*]

lemmas *udvd-decr0* = *udvd-decr'* [**where** *ua=0, simplified*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy \leq k - z$
apply (*unfold udvd-def*)
apply *clarify*
apply (*erule* (2) *udvd-decr0*)
done

ML $\ll \text{Delsimprocs Numeral-Simprocs.cancel-factors} \gg$

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq$
 $p \implies$
 $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$
apply (*unfold udvd-def*)

```

apply clarify
apply (simp add: uint-arith-simps split: split-if-asm)
prefer 2
apply (insert uint-range' [of s][1])
apply arith
apply (drule add-commute [THEN xtr1])
apply (simp add: diff-less-eq [symmetric])
apply (drule less-le-mult)
apply arith
apply simp
done

```

ML \ll *Addsimprocs Numeral-Simprocs.cancel-factors* \gg

```

lemma word-succ-rbl:
  to-bl w = bl ==> to-bl (word-succ w) = (rev (rbl-succ (rev bl)))
apply (unfold word-succ-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-succ)
done

```

```

lemma word-pred-rbl:
  to-bl w = bl ==> to-bl (word-pred w) = (rev (rbl-pred (rev bl)))
apply (unfold word-pred-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-pred)
done

```

```

lemma word-add-rbl:
  to-bl v = vbl ==> to-bl w = wbl ==>
    to-bl (v + w) = (rev (rbl-add (rev vbl) (rev wbl)))
apply (unfold word-add-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-add)
done

```

```

lemma word-mult-rbl:
  to-bl v = vbl ==> to-bl w = wbl ==>
    to-bl (v * w) = (rev (rbl-mult (rev vbl) (rev wbl)))
apply (unfold word-mult-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-mult)
done

```


lemma *rtb-rbl-ariths*:

rev (to-bl w) = ys \implies rev (to-bl (word-succ w)) = rbl-succ ys

rev (to-bl w) = ys \implies rev (to-bl (word-pred w)) = rbl-pred ys

*[| rev (to-bl v) = ys; rev (to-bl w) = xs |]
 \implies rev (to-bl (v * w)) = rbl-mult ys xs*

*[| rev (to-bl v) = ys; rev (to-bl w) = xs |]
 \implies rev (to-bl (v + w)) = rbl-add ys xs*

by (*auto simp: rev-swap [symmetric] word-succ-rbl
word-pred-rbl word-mult-rbl word-add-rbl*)

11.6 Arithmetic type class instantiations

instance *word* :: (len0) comm-monoid-add ..

instance *word* :: (len0) comm-monoid-mult
apply (*intro-classes*)
apply (*simp add: word-mult-commute*)
apply (*simp add: word-mult-1*)
done

instance *word* :: (len0) comm-semiring
by (*intro-classes*) (*simp add : word-left-distrib*)

instance *word* :: (len0) ab-group-add ..

instance *word* :: (len0) comm-ring ..

instance *word* :: (len) comm-semiring-1
by (*intro-classes*) (*simp add: lenw1-zero-neq-one*)

instance *word* :: (len) comm-ring-1 ..

instance *word* :: (len0) comm-semiring-0 ..

instance *word* :: (len0) order ..

lemma *zero-bintrunc*:

*iszero (number-of x :: 'a :: len word) =
(bintrunc (len-of TYPE('a)) x = Int.Pls)*
apply (*unfold iszero-def word-0-wi word-no-wi*)
apply (*rule word-ubin.norm-eq-iff [symmetric, THEN trans]*)
apply (*simp add : Pls-def [symmetric]*)
done

lemmas *word-le-0-iff* [*simp*] =

word-zero-le [*THEN leD*, *THEN linorder-antisym-conv1*]

lemma *word-of-nat*: *of-nat n = word-of-int (int n)*
by (*induct n*) (*auto simp add : word-of-int-hom-syms*)

lemma *word-of-int*: *of-int = word-of-int*
apply (*rule ext*)
apply (*unfold of-int-def*)
apply (*rule contentsI*)
apply *safe*
apply (*simp-all add: word-of-nat word-of-int-homs*)
defer
apply (*rule Rep-Integ-ne [THEN nonemptyE]*)
apply (*rule bexI*)
prefer 2
apply *assumption*
apply (*auto simp add: RI-eq-diff*)
done

lemma *word-of-int-nat*:
 $0 \leq x \implies \text{word-of-int } x = \text{of-nat } (\text{nat } x)$
by (*simp add: of-nat-nat word-of-int*)

lemma *word-number-of-eq*:
 $\text{number-of } w = (\text{of-int } w :: 'a :: \text{len word})$
unfolding *word-number-of-def word-of-int* **by** *auto*

instance *word* :: (*len*) *number-ring*
by (*intro-classes*) (*simp add : word-number-of-eq*)

lemma *iszero-word-no* [*simp*] :
 $\text{iszero } (\text{number-of } \text{bin} :: 'a :: \text{len word}) =$
 $\text{iszero } (\text{number-of } (\text{bintrunc } (\text{len-of TYPE('a)}) \text{bin}) :: \text{int})$
apply (*simp add: zero-bintrunc number-of-is-id*)
apply (*unfold iszero-def Pls-def*)
apply (*rule refl*)
done

11.7 Word and nat

lemma *td-ext-unat'*:
 $n = \text{len-of TYPE } ('a :: \text{len}) \implies$
 $\text{td-ext } (\text{unat} :: 'a \text{ word} \implies \text{nat}) \text{ of-nat}$
 $(\text{unats } n) (\%i. i \bmod 2^n)$
apply (*unfold td-ext-def' unat-def word-of-nat unats-uints*)
apply (*auto intro!: imageI simp add : word-of-int-hom-syms*)
apply (*erule word-uint.Abs-inverse [THEN arg-cong]*)
apply (*simp add: int-word-uint nat-mod-distrib nat-power-eq*)
done

```

lemmas td-ext-unat = refl [THEN td-ext-unat]
lemmas unat-of-nat = td-ext-unat [THEN td-ext.eq-norm, standard]

interpretation word-unat:
  td-ext unat::'a::len word ==> nat
    of-nat
    unats (len-of TYPE('a::len))
    %i. i mod 2 ^ len-of TYPE('a::len)
  by (rule td-ext-unat)

lemmas td-unat = word-unat.td-thm

lemmas unat-lt2p [iff] = word-unat.Rep [unfolded unats-def mem-Collect-eq]

lemma unat-le: y <= unat (z :: 'a :: len word) ==> y : unats (len-of TYPE
('a))
  apply (unfold unats-def)
  apply clarsimp
  apply (rule xtrans, rule unat-lt2p, assumption)
  done

lemma word-nchotomy:
  ALL w. EX n. (w :: 'a :: len word) = of-nat n & n < 2 ^ len-of TYPE ('a)
  apply (rule allI)
  apply (rule word-unat.Abs-cases)
  apply (unfold unats-def)
  apply auto
  done

lemma of-nat-eq:
  fixes w :: 'a::len word
  shows (of-nat n = w) = ( $\exists q. n = \text{unat } w + q * 2 ^ \text{len-of TYPE('a)}$ )
  apply (rule trans)
  apply (rule word-unat.inverse-norm)
  apply (rule iffI)
  apply (rule mod-eqD)
  apply simp
  apply clarsimp
  done

lemma of-nat-eq-size:
  (of-nat n = w) = (EX q. n = unat w + q * 2 ^ size w)
  unfolding word-size by (rule of-nat-eq)

lemma of-nat-0:
  (of-nat m = (0::'a::len word)) = ( $\exists q. m = q * 2 ^ \text{len-of TYPE('a)}$ )
  by (simp add: of-nat-eq)

```

lemmas *of-nat-2p = mult-1 [symmetric, THEN iffD2 [OF of-nat-0 exI]]*

lemma *of-nat-gt-0: of-nat k ~ = 0 ==> 0 < k*
by (*cases k*) *auto*

lemma *of-nat-neq-0:*
 $0 < k ==> k < 2 \wedge \text{len-of TYPE } ('a :: \text{len}) ==> \text{of-nat } k \sim = (0 :: 'a \text{ word})$
by (*clarsimp simp add : of-nat-0*)

lemma *Abs-fnat-hom-add:*
 $\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$
by *simp*

lemma *Abs-fnat-hom-mult:*
 $\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$
by (*simp add: word-of-nat word-of-int-mult-hom zmult-int*)

lemma *Abs-fnat-hom-Suc:*
 $\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$
by (*simp add: word-of-nat word-of-int-succ-hom add-ac*)

lemma *Abs-fnat-hom-0: (0::'a::len word) = of-nat 0*
by (*simp add: word-of-nat word-0-wi*)

lemma *Abs-fnat-hom-1: (1::'a::len word) = of-nat (Suc 0)*
by (*simp add: word-of-nat word-1-wi*)

lemmas *Abs-fnat-homs =*
Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc
Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add:*
 $a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$
by *simp*

lemma *word-arith-nat-mult:*
 $a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$
by (*simp add: Abs-fnat-hom-mult [symmetric]*)

lemma *word-arith-nat-Suc:*
 $\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$
by (*subst Abs-fnat-hom-Suc [symmetric]*) *simp*

lemma *word-arith-nat-div:*
 $a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$
by (*simp add: word-div-def word-of-nat zdiv-int uint-nat*)

lemma *word-arith-nat-mod:*
 $a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$

by (*simp add: word-mod-def word-of-nat zmod-int uint-nat*)

lemmas *word-arith-nat-defs* =
word-arith-nat-add word-arith-nat-mult
word-arith-nat-Suc Abs-fnat-hom-0
Abs-fnat-hom-1 word-arith-nat-div
word-arith-nat-mod

lemmas *unat-cong* = *arg-cong* [**where** *f* = *unat*]

lemmas *unat-word-ariths* = *word-arith-nat-defs*
 [*THEN trans* [*OF unat-cong unat-of-nat*], *standard*]

lemmas *word-sub-less-iff* = *word-sub-le-iff*
 [*simplified linorder-not-less* [*symmetric*], *simplified*]

lemma *unat-add-lem*:
 (*unat x + unat y < 2 ^ len-of TYPE('a)*) =
 (*unat (x + y :: 'a :: len word) = unat x + unat y*)
unfolding *unat-word-ariths*
by (*auto intro!: trans* [*OF - nat-mod-lem*])

lemma *unat-mult-lem*:
 (*unat x * unat y < 2 ^ len-of TYPE('a)*) =
 (*unat (x * y :: 'a :: len word) = unat x * unat y*)
unfolding *unat-word-ariths*
by (*auto intro!: trans* [*OF - nat-mod-lem*])

lemmas *unat-plus-if'* =
trans [*OF unat-word-ariths(1) mod-nat-add, simplified, standard*]

lemma *le-no-overflow*:
x <= b ==> a <= a + b ==> x <= a + (b :: 'a :: len0 word)
apply (*erule order-trans*)
apply (*erule olen-add-egv* [*THEN iffD1*])
done

lemmas *un-ui-le* = *trans*
 [*OF word-le-nat-alt* [*symmetric*]
word-le-def,
standard]

lemma *unat-sub-if-size*:
unat (x - y) = (if unat y <= unat x
then unat x - unat y
else unat x + 2 ^ size x - unat y)
apply (*unfold word-size*)
apply (*simp add: un-ui-le*)
apply (*auto simp add: unat-def uint-sub-if'*)

```

apply (rule nat-diff-distrib)
prefer 3
apply (simp add: algebra-simps)
apply (rule nat-diff-distrib [THEN trans])
prefer 3
apply (subst nat-add-distrib)
prefer 3
apply (simp add: nat-power-eq)
apply auto
apply uint-arith
done

```

lemmas *unat-sub-if' = unat-sub-if-size [unfolded word-size]*

```

lemma unat-div: unat ((x :: 'a :: len word) div y) = unat x div unat y
apply (simp add : unat-word-ariths)
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
apply (rule div-le-dividend)
done

```

```

lemma unat-mod: unat ((x :: 'a :: len word) mod y) = unat x mod unat y
apply (clarsimp simp add : unat-word-ariths)
apply (cases unat y)
prefer 2
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
apply (rule mod-le-divisor)
apply auto
done

```

```

lemma uint-div: uint ((x :: 'a :: len word) div y) = uint x div uint y
unfolding uint-nat by (simp add : unat-div zdiv-int)

```

```

lemma uint-mod: uint ((x :: 'a :: len word) mod y) = uint x mod uint y
unfolding uint-nat by (simp add : unat-mod zmod-int)

```

11.8 Definition of unat_arith tactic

```

lemma unat-split:
  fixes x::'a::len word
  shows P (unat x) =
    (ALL n. of-nat n = x & n < 2^len-of TYPE('a) --> P n)
  by (auto simp: unat-of-nat)

```

```

lemma unat-split-asm:
  fixes x::'a::len word
  shows P (unat x) =
    (~ (EX n. of-nat n = x & n < 2^len-of TYPE('a) & ~ P n))
  by (auto simp: unat-of-nat)

```

lemmas *of-nat-inverse* =
word-unat.Abs-inverse' [*rotated, unfolded unats-def, simplified*]

lemmas *unat-splits* = *unat-split unat-split-asm*

lemmas *unat-arith-simps* =
word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [*symmetric*]
unat-sub-if' unat-plus-if' unat-div unat-mod

ML \ll
fun unat-arith-ss-of ss =
 ss addsimps @{thms unat-arith-simps}
 delsimps @{thms word-unat.Rep-inject}
 addsplits @{thms split-if-asm}
 addcongs @{thms power-False-cong}

fun unat-arith-tacs ctxt =
 let
 fun arith-tac' n t =
 Arith-Data.verbose-arith-tac ctxt n t
 handle Cooper.COOPER - => Seq.empty;
 val cs = claset-of ctxt;
 val ss = simpset-of ctxt;
 in
 [*clarify-tac cs 1,*
 full-simp-tac (unat-arith-ss-of ss) 1,
 ALLGOALS (full-simp-tac (HOL-ss addsplits @{thms unat-splits}
 addcongs @{thms power-False-cong})),
 rewrite-goals-tac @{thms word-size},
 ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
 REPEAT (etac conjE n) THEN
 REPEAT (dtac @{thm of-nat-inverse} n THEN atac n)),
 TRYALL arith-tac']
 end

fun unat-arith-tac ctxt = SELECT-GOAL (EVERY (unat-arith-tacs ctxt))
 \gg

method-setup *unat-arith* =
 \ll *Scan.succeed (SIMPLE-METHOD' o unat-arith-tac)* \gg
solving word arithmetic via natural numbers and arith

lemma *no-plus-overflow-unat-size*:
 $((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$
unfolding *word-size* **by** *unat-arith*

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$

by *unat-arith*

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

lemmas *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem, standard*]

lemma *word-div-mult*:

($0 :: 'a :: \text{len word}$) $< y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$
 $x * y \text{ div } y = x$

apply *unat-arith*

apply *clarsimp*

apply (*subst unat-mult-lem* [*THEN iffD1*])

apply *auto*

done

lemma *div-lt'*: ($i :: 'a :: \text{len word}$) $\leq k \text{ div } x \implies$

$\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$

apply *unat-arith*

apply *clarsimp*

apply (*drule mult-le-mono1*)

apply (*erule order-le-less-trans*)

apply (*rule xtr7* [*OF unat-lt2p div-mult-le*])

done

lemmas *div-lt''* = *order-less-imp-le* [*THEN div-lt'*]

lemma *div-lt-mult*: ($i :: 'a :: \text{len word}$) $< k \text{ div } x \implies 0 < x \implies i * x < k$

apply (*frule div-lt''* [*THEN unat-mult-lem* [*THEN iffD1*]])

apply (*simp add: unat-arith-simps*)

apply (*drule* (1) *mult-less-mono1*)

apply (*erule order-less-le-trans*)

apply (*rule div-mult-le*)

done

lemma *div-le-mult*:

($i :: 'a :: \text{len word}$) $\leq k \text{ div } x \implies 0 < x \implies i * x \leq k$

apply (*frule div-lt'* [*THEN unat-mult-lem* [*THEN iffD1*]])

apply (*simp add: unat-arith-simps*)

apply (*drule mult-le-mono1*)

apply (*erule order-trans*)

apply (*rule div-mult-le*)

done

lemma *div-lt-uint'*:

($i :: 'a :: \text{len word}$) $\leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$

apply (*unfold uint-nat*)

apply (*drule div-lt'*)

apply (*simp add: zmult-int zless-nat-eq-int-zless* [*symmetric*]
nat-power-eq)

done

lemmas *div-lt-uint'' = order-less-imp-le [THEN div-lt-uint']*

lemma *word-le-exists'*:

(*x :: 'a :: len0 word*) $\leq y \implies$
 (*EX z. y = x + z & uint x + uint z < 2 ^ len-of TYPE('a)*)
 apply (rule *exI*)
 apply (rule *conjI*)
 apply (rule *zadd-diff-inverse*)
 apply *uint-arith*
 done

lemmas *plus-minus-not-NULL = order-less-imp-le [THEN plus-minus-not-NULL-ab]*

lemmas *plus-minus-no-overflow =*
order-less-imp-le [THEN plus-minus-no-overflow-ab]

lemmas *mcs = word-less-minus-cancel word-less-minus-mono-left*
word-le-minus-cancel word-le-minus-mono-left

lemmas *word-l-diffs = mcs [where y = w + x, unfolded add-diff-cancel, standard]*
 lemmas *word-diff-ls = mcs [where z = w + x, unfolded add-diff-cancel, standard]*
 lemmas *word-plus-mcs = word-diff-ls*
[where y = v + x, unfolded add-diff-cancel, standard]

lemmas *le-unat-uoI = unat-le [THEN word-unat.Abs-inverse]*

lemmas *thd = refl [THEN [2] split-div-lemma [THEN iffD2], THEN conjunct1]*

lemma *thd1*:
*a div b * b ≤ (a::nat)*
 using *gt-or-eq-0 [of b]*
 apply (rule *disjE*)
 apply (erule *xtr4 [OF thd mult-commute]*)
 apply *clarsimp*
 done

lemmas *uno-simps [THEN le-unat-uoI, standard] =*
mod-le-divisor div-le-dividend thd1

lemma *word-mod-div-equality*:
*(n div b) * b + (n mod b) = (n :: 'a :: len word)*
 apply (unfold *word-less-nat-alt word-arith-nat-defs*)
 apply (cut-tac *y=unat b in gt-or-eq-0*)
 apply (erule *disjE*)
 apply (simp *add: mod-div-equality uno-simps*)
 apply *simp*
 done

```

lemma word-div-mult-le:  $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$ 
  apply (unfold word-le-nat-alt word-arith-nat-defs)
  apply (cut-tac  $y = \text{unat } b$  in gt-or-eq-0)
  apply (erule disjE)
  apply (simp add: div-mult-le uno-simps)
  apply simp
done

```

```

lemma word-mod-less-divisor:  $0 < n \implies m \bmod n < (n :: 'a :: \text{len word})$ 
  apply (simp only: word-less-nat-alt word-arith-nat-defs)
  apply (clarsimp simp add: uno-simps)
done

```

```

lemma word-of-int-power-hom:
   $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a :: \text{len word})$ 
  by (induct n) (simp-all add: word-of-int-hom-syms power-Suc)

```

```

lemma word-arith-power-alt:
   $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a :: \text{len word})$ 
  by (simp add: word-of-int-power-hom [symmetric])

```

```

lemma of-bl-length-less:
   $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 ^ k$ 
  apply (unfold of-bl-no [unfolded word-number-of-def]
    word-less-alt word-number-of-alt)
  apply safe
  apply (simp (no-asm) add: word-of-int-power-hom word-uint.eq-norm
    del: word-of-int-bin)
  apply (simp add: mod-pos-pos-trivial)
  apply (subst mod-pos-pos-trivial)
  apply (rule bl-to-bin-ge0)
  apply (rule order-less-trans)
  apply (rule bl-to-bin-lt2p)
  apply simp
  apply (rule bl-to-bin-lt2p)
done

```

11.9 Cardinality, finiteness of set of words

```

lemmas card-lessThan' = card-lessThan [unfolded lessThan-def]

```

```

lemmas card-eq = word-unat.Abs-inj-on [THEN card-image,
  unfolded word-unat.image, unfolded unats-def, standard]

```

```

lemmas card-word = trans [OF card-eq card-lessThan', standard]

```

```

lemma finite-word-UNIV: finite (UNIV :: 'a :: len word set)
apply (rule contrapos-np)

```

```

prefer 2
apply (erule card-infinite)
apply (simp add: card-word)
done

```

```

lemma card-word-size:
  card (UNIV :: 'a :: len word set) = (2 ^ size (x :: 'a word))
unfolding word-size by (rule card-word)

end

```

12 WordBitwise: Bitwise Operations on Words

```

theory WordBitwise
imports WordArith
begin

```

```

lemmas bin-log-bintrs = bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or

```

```

lemmas wils1 = bin-log-bintrs [THEN word-ubin.norm-eq-iff [THEN iffD1],
  folded word-ubin.eq-norm, THEN eq-reflection, standard]

```

```

lemmas word-log-binary-defs =
  word-and-def word-or-def word-xor-def

```

```

lemmas word-no-log-defs [simp] =
  word-not-def [where a=number-of a,
    unfolded word-no-wi wils1, folded word-no-wi, standard]
  word-log-binary-defs [where a=number-of a and b=number-of b,
    unfolded word-no-wi wils1, folded word-no-wi, standard]

```

```

lemmas word-wi-log-defs = word-no-log-defs [unfolded word-no-wi]

```

```

lemma uint-or: uint (x OR y) = (uint x) OR (uint y)
by (simp add: word-or-def word-no-wi [symmetric] number-of-is-id
  bin-trunc-ao(2) [symmetric])

```

```

lemma uint-and: uint (x AND y) = (uint x) AND (uint y)
by (simp add: word-and-def number-of-is-id word-no-wi [symmetric]
  bin-trunc-ao(1) [symmetric])

```

```

lemma word-ops-nth-size:
  n < size (x::'a::len0 word) ==>
    (x OR y) !! n = (x !! n | y !! n) &

```

$$\begin{aligned} (x \text{ AND } y) !! n &= (x !! n \& y !! n) \& \\ (x \text{ XOR } y) !! n &= (x !! n \sim = y !! n) \& \\ (\text{NOT } x) !! n &= (\sim x !! n) \end{aligned}$$

unfolding *word-size word-no-wi word-test-bit-def word-log-defs*
by (*clarsimp simp add : word-ubin.eq-norm nth-bintr bin-nth-ops*)

lemma *word-ao-nth*:
fixes *x :: 'a::len0 word*
shows $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \&$
 $(x \text{ AND } y) !! n = (x !! n \& y !! n)$
apply (*cases n < size x*)
apply (*drule-tac y = y in word-ops-nth-size*)
apply *simp*
apply (*simp add : test-bit-bin word-size*)
done

lemmas *bwsimps =*
word-of-int-homs(2)
word-0-wi-Pls
word-m1-wi-Min
word-wi-log-defs

lemma *word-bw-assocs*:
fixes *x :: 'a::len0 word*
shows
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
using *word-of-int-Ex [where x=x]*
word-of-int-Ex [where x=y]
word-of-int-Ex [where x=z]
by (*auto simp: bwsimps bbw-assocs*)

lemma *word-bw-comms*:
fixes *x :: 'a::len0 word*
shows
 $x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
using *word-of-int-Ex [where x=x]*
word-of-int-Ex [where x=y]
by (*auto simp: bwsimps bin-ops-comm*)

lemma *word-bw-lcs*:
fixes *x :: 'a::len0 word*
shows
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$

```

y OR x OR z = x OR y OR z
y XOR x XOR z = x XOR y XOR z
using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
      word-of-int-Ex [where x=z]
by (auto simp: bwsimps)

```

```

lemma word-log-esimps [simp]:
  fixes x :: 'a::len0 word
  shows
    x AND 0 = 0
    x AND -1 = x
    x OR 0 = x
    x OR -1 = -1
    x XOR 0 = x
    x XOR -1 = NOT x
    0 AND x = 0
    -1 AND x = x
    0 OR x = x
    -1 OR x = -1
    0 XOR x = x
    -1 XOR x = NOT x
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps)

```

```

lemma word-not-dist:
  fixes x :: 'a::len0 word
  shows
    NOT (x OR y) = NOT x AND NOT y
    NOT (x AND y) = NOT x OR NOT y
  using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
  by (auto simp: bwsimps bbw-not-dist)

```

```

lemma word-bw-same:
  fixes x :: 'a::len0 word
  shows
    x AND x = x
    x OR x = x
    x XOR x = 0
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps)

```

```

lemma word-ao-absorbs [simp]:
  fixes x :: 'a::len0 word
  shows
    x AND (y OR x) = x
    x OR y AND x = x
    x AND (x OR y) = x

```

```

y AND x OR x = x
(y OR x) AND x = x
x OR x AND y = x
(x OR y) AND x = x
x AND y OR x = x
using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
by (auto simp: bwsimps)

```

```

lemma word-not-not [simp]:
  NOT NOT (x::'a::len0 word) = x
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps)

```

```

lemma word-ao-dist:
  fixes x :: 'a::len0 word
  shows (x OR y) AND z = x AND z OR y AND z
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
        word-of-int-Ex [where x=z]
  by (auto simp: bwsimps bbw-ao-dist simp del: bin-ops-comm)

```

```

lemma word-oa-dist:
  fixes x :: 'a::len0 word
  shows x AND y OR z = (x OR z) AND (y OR z)
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
        word-of-int-Ex [where x=z]
  by (auto simp: bwsimps bbw-oa-dist simp del: bin-ops-comm)

```

```

lemma word-add-not [simp]:
  fixes x :: 'a::len0 word
  shows x + NOT x = -1
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps bin-add-not)

```

```

lemma word-plus-and-or [simp]:
  fixes x :: 'a::len0 word
  shows (x AND y) + (x OR y) = x + y
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
  by (auto simp: bwsimps plus-and-or)

```

```

lemma leoa:
  fixes x :: 'a::len0 word
  shows (w = (x OR y)) ==> (y = (w AND y)) by auto
lemma leao:
  fixes x' :: 'a::len0 word
  shows (w' = (x' AND y')) ==> (x' = (x' OR w')) by auto

```

```

lemmas word-ao-equiv = leao [COMP leoa [COMP iffI]]

lemma le-word-or2: x <= x OR (y::'a::len0 word)
  unfolding word-le-def uint-or
  by (auto intro: le-int-or)

lemmas le-word-or1 = xtr3 [OF word-bw-comms (2) le-word-or2, standard]
lemmas word-and-le1 =
  xtr3 [OF word-ao-absorbs (4) [symmetric] le-word-or2, standard]
lemmas word-and-le2 =
  xtr3 [OF word-ao-absorbs (8) [symmetric] le-word-or2, standard]

lemma bl-word-not: to-bl (NOT w) = map Not (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-not-bin number-of-is-id word-no-wi [symmetric] bin-to-bl-def[symmetric])

lemma bl-word-xor: to-bl (v XOR w) = map2 op ~ = (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-xor-bin
  by (simp add: number-of-is-id word-no-wi [symmetric])

lemma bl-word-or: to-bl (v OR w) = map2 op | (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-or-bin number-of-is-id word-no-wi [symmetric])

lemma bl-word-and: to-bl (v AND w) = map2 op & (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-and-bin number-of-is-id word-no-wi [symmetric])

lemma word-lsb-alt: lsb (w::'a::len0 word) = test-bit w 0
  by (auto simp: word-test-bit-def word-lsb-def)

lemma word-lsb-1-0: lsb (1::'a::len word) & ~ lsb (0::'b::len0 word)
  unfolding word-lsb-def word-1-no word-0-no by auto

lemma word-lsb-last: lsb (w::'a::len word) = last (to-bl w)
  apply (unfold word-lsb-def uint-bl bin-to-bl-def)
  apply (rule-tac bin=uint w in bin-exhaust)
  apply (cases size w)
  apply auto
  apply (auto simp add: bin-to-bl-aux-alt)
  done

lemma word-lsb-int: lsb w = (uint w mod 2 = 1)
  unfolding word-lsb-def bin-last-mod by auto

lemma word-msb-sint: msb w = (sint w < 0)
  unfolding word-msb-def
  by (simp add : sign-Min-lt-0 number-of-is-id)

```

```

lemma word-msb-no':
  w = number-of bin ==> msb (w::'a::len word) = bin-nth bin (size w - 1)
  unfolding word-msb-def word-number-of-def
  by (clarsimp simp add: word-sbin.eq-norm word-size bin-sign-lem)

lemmas word-msb-no = refl [THEN word-msb-no', unfolded word-size]

lemma word-msb-nth': msb (w::'a::len word) = bin-nth (uint w) (size w - 1)
  apply (unfold word-size)
  apply (rule trans [OF - word-msb-no])
  apply (simp add : word-number-of-def)
  done

lemmas word-msb-nth = word-msb-nth' [unfolded word-size]

lemma word-msb-alt: msb (w::'a::len word) = hd (to-bl w)
  apply (unfold word-msb-nth uint-bl)
  apply (subst hd-conv-nth)
  apply (rule length-greater-0-conv [THEN iffD1])
  apply simp
  apply (simp add : nth-bin-to-bl word-size)
  done

lemma word-set-nth:
  set-bit w n (test-bit w n) = (w::'a::len0 word)
  unfolding word-test-bit-def word-set-bit-def by auto

lemma bin-nth-uint':
  bin-nth (uint w) n = (rev (bin-to-bl (size w) (uint w)) ! n & n < size w)
  apply (unfold word-size)
  apply (safe elim!: bin-nth-uint-imp)
  apply (frule bin-nth-uint-imp)
  apply (fast dest!: bin-nth-bl)+
  done

lemmas bin-nth-uint = bin-nth-uint' [unfolded word-size]

lemma test-bit-bl: w !! n = (rev (to-bl w) ! n & n < size w)
  unfolding to-bl-def word-test-bit-def word-size
  by (rule bin-nth-uint)

lemma to-bl-nth: n < size w ==> to-bl w ! n = w !! (size w - Suc n)
  apply (unfold test-bit-bl)
  apply clarsimp
  apply (rule trans)
  apply (rule nth-rev-alt)
  apply (auto simp add: word-size)
  done

```


lemma *test-bit-set*:

fixes $w :: 'a::len0$ word
shows $(set-bit\ w\ n\ x) \iff n = (n < size\ w \ \&\ x)$
unfolding *word-size word-test-bit-def word-set-bit-def*
by (*clarsimp simp add : word-ubin.eq-norm nth-bintr*)

lemma *test-bit-set-gen*:

fixes $w :: 'a::len0$ word
shows $test-bit\ (set-bit\ w\ n\ x)\ m =$
 $(if\ m = n\ then\ n < size\ w \ \&\ x\ else\ test-bit\ w\ m)$
apply (*unfold word-size word-test-bit-def word-set-bit-def*)
apply (*clarsimp simp add: word-ubin.eq-norm nth-bintr bin-nth-sc-gen*)
apply (*auto elim!: test-bit-size [unfolded word-size]*
 $simp\ add: word-test-bit-def [symmetric]$)
done

lemma *of-bl-rep-False*: $of-bl\ (replicate\ n\ False\ @\ bs) = of-bl\ bs$

unfolding *of-bl-def bl-to-bin-rep-F* **by** *auto*

lemma *msb-nth'*:

fixes $w :: 'a::len$ word
shows $msb\ w = w \iff (size\ w - 1)$
unfolding *word-msb-nth' word-test-bit-def* **by** *simp*

lemmas $msb-nth = msb-nth' [unfolded\ word-size]$

lemmas $msb0 = len-gt-0 [THEN\ diff-Suc-less, THEN$

$word-ops-nth-size [unfolded\ word-size], standard]$

lemmas $msb1 = msb0 [where\ i = 0]$

lemmas $word-ops-msb = msb1 [unfolded\ msb-nth [symmetric, unfolded\ One-nat-def]]$

lemmas $lsb0 = len-gt-0 [THEN\ word-ops-nth-size [unfolded\ word-size], standard]$

lemmas $word-ops-lsb = lsb0 [unfolded\ word-lsb-alt]$

lemma *td-ext-nth'*:

$n = size\ (w::'a::len0\ word) \implies ofn = set-bits \implies [w, ofn\ g] = l \implies$

$td-ext\ test-bit\ ofn\ \{f.\ ALL\ i.\ f\ i \implies i < n\} \ (\%h\ i.\ h\ i \ \&\ i < n)$

apply (*unfold word-size td-ext-def'*)

apply (*safe del: subset-antisym*)

apply (*rule-tac [3] ext*)

apply (*rule-tac [4] ext*)

apply (*unfold word-size of-nth-def test-bit-bl*)

apply *safe*

defer

apply (*clarsimp simp: word-bl.Abs-inverse*)**+**

apply (*rule word-bl.Rep-inverse'*)

apply (*rule sym [THEN trans]*)

apply (*rule bl-of-nth-nth*)

```

apply simp
apply (rule bl-of-nth-inj)
apply (clarsimp simp add : test-bit-bl word-size)
done

lemmas td-ext-nth = td-ext-nth' [OF refl refl refl, unfolded word-size]

interpretation test-bit:
  td-ext op !! :: 'a::len0 word => nat => bool
    set-bits
    {f. ∀ i. f i ⟶ i < len-of TYPE('a::len0)}
    ( $\lambda h i. h i \wedge i < \text{len-of } \text{TYPE}('a::\text{len0})$ )
  by (rule td-ext-nth)

declare test-bit.Rep' [simp del]
declare test-bit.Rep' [rule del]

lemmas td-nth = test-bit.td-thm

lemma word-set-set-same:
  fixes w :: 'a::len0 word
  shows set-bit (set-bit w n x) n y = set-bit w n y
  by (rule word-eqI) (simp add : test-bit-set-gen word-size)

lemma word-set-set-diff:
  fixes w :: 'a::len0 word
  assumes m ~ n
  shows set-bit (set-bit w m x) n y = set-bit (set-bit w n y) m x
  by (rule word-eqI) (clarsimp simp add : test-bit-set-gen word-size prems)

lemma test-bit-no':
  fixes w :: 'a::len0 word
  shows w = number-of bin ==> test-bit w n = (n < size w & bin-nth bin n)
  unfolding word-test-bit-def word-number-of-def word-size
  by (simp add : nth-bintr [symmetric] word-ubin.eq-norm)

lemmas test-bit-no =
  refl [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

lemma nth-0: ~ (0::'a::len0 word) !! n
  unfolding test-bit-no word-0-no by auto

lemma nth-sint:
  fixes w :: 'a::len word
  defines l ≡ len-of TYPE ('a)
  shows bin-nth (sint w) n = (if n < l - 1 then w !! n else w !! (l - 1))
  unfolding sint-uint l-def
  by (clarsimp simp add: nth-sbintr word-test-bit-def [symmetric])

```

lemma *word-lsb-no*:

lsb (number-of bin :: 'a :: len word) = (bin-last bin = bit.B1)

unfolding *word-lsb-alt test-bit-no* **by** *auto*

lemma *word-set-no*:

set-bit (number-of bin :: 'a :: len0 word) n b =

number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)

apply (*unfold word-set-bit-def word-number-of-def [symmetric]*)

apply (*rule word-eqI*)

apply (*clarsimp simp: word-size bin-nth-sc-gen number-of-is-id*
test-bit-no nth-bintr)

done

lemmas *setBit-no = setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemmas *clearBit-no = clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemma *to-bl-n1*:

to-bl (-1 :: 'a :: len0 word) = replicate (len-of TYPE ('a)) True

apply (*rule word-bl.Abs-inverse'*)

apply *simp*

apply (*rule word-eqI*)

apply (*clarsimp simp add: word-size test-bit-no*)

apply (*auto simp add: word-bl.Abs-inverse test-bit-bl word-size*)

done

lemma *word-msb-n1*: *msb (-1 :: 'a :: len word)*

unfolding *word-msb-alt word-msb-alt to-bl-n1* **by** *simp*

declare *word-set-set-same [simp] word-set-nth [simp]*

test-bit-no [simp] word-set-no [simp] nth-0 [simp]

setBit-no [simp] clearBit-no [simp]

word-lsb-no [simp] word-msb-no [simp] word-msb-n1 [simp] word-lsb-1-0 [simp]

lemma *word-set-nth-iff*:

(set-bit w n b = w) = (w !! n = b | n >= size (w :: 'a :: len0 word))

apply (*rule iffI*)

apply (*rule disjCI*)

apply (*drule word-eqD*)

apply (*erule sym [THEN trans]*)

apply (*simp add: test-bit-set*)

apply (*erule disjE*)

apply *clarsimp*

apply (*rule word-eqI*)

apply (*clarsimp simp add: test-bit-set-gen*)

apply (*drule test-bit-size*)

apply *force*

done

lemma *test-bit-2p'*:

$w = \text{word-of-int } (2 \wedge n) \implies$
 $w !! m = (m = n \ \& \ m < \text{size } (w :: 'a :: \text{len word}))$
unfolding *word-test-bit-def word-size*
by (*auto simp add: word-ubin.eq-norm nth-bintr nth-2p-bin*)

lemmas *test-bit-2p = refl [THEN test-bit-2p', unfolded word-size]*

lemma *nth-w2p*:

$((2 :: 'a :: \text{len word}) \wedge n) !! m \longleftrightarrow m = n \wedge m < \text{len-of TYPE}('a :: \text{len})$
unfolding *test-bit-2p [symmetric] word-of-int [symmetric]*
by (*simp add: of-int-power*)

lemma *uint-2p*:

$(0 :: 'a :: \text{len word}) < 2 \wedge n \implies \text{uint } (2 \wedge n :: 'a :: \text{len word}) = 2 \wedge n$
apply (*unfold word-arith-power-alt*)
apply (*case-tac len-of TYPE ('a)*)
apply *clarsimp*
apply (*case-tac nat*)
apply *clarsimp*
apply (*case-tac n*)
apply (*clarsimp simp add : word-1-wi [symmetric]*)
apply (*clarsimp simp add : word-0-wi [symmetric]*)
apply (*drule word-gt-0 [THEN iffD1]*)
apply (*safe intro!: word-eqI bin-nth-lem ext*)
apply (*auto simp add: test-bit-2p nth-2p-bin word-test-bit-def [symmetric]*)
done

lemma *word-of-int-2p*: $(\text{word-of-int } (2 \wedge n) :: 'a :: \text{len word}) = 2 \wedge n$

apply (*unfold word-arith-power-alt*)
apply (*case-tac len-of TYPE ('a)*)
apply *clarsimp*
apply (*case-tac nat*)
apply (*rule word-ubin.norm-eq-iff [THEN iffD1]*)
apply (*rule box-equals*)
apply (*rule-tac [2] bintr-ariths (1))+*
apply (*clarsimp simp add : number-of-is-id*)
apply *simp*
done

lemma *bang-is-le*: $x !! m \implies 2 \wedge m \leq (x :: 'a :: \text{len word})$

apply (*rule xtr3*)
apply (*rule-tac [2] y = x in le-word-or2*)
apply (*rule word-eqI*)
apply (*auto simp add: word-ao-nth nth-w2p word-size*)
done

lemma *word-clr-le*:

```

fixes  $w :: 'a::len0$  word
shows  $w \geq \text{set-bit } w \ n \ \text{False}$ 
apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
apply simp
apply (rule order-trans)
  apply (rule bintr-bin-clr-le)
apply simp
done

lemma word-set-ge:
  fixes  $w :: 'a::len$  word
  shows  $w \leq \text{set-bit } w \ n \ \text{True}$ 
  apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
  apply simp
  apply (rule order-trans [OF - bintr-bin-set-ge])
  apply simp
  done

end

```

13 WordShift: Shifting, Rotating, and Splitting Words

```

theory WordShift
imports WordBitwise
begin

```

13.1 Bit shifting

```

lemma shiffl1-number [simp] :
  shiffl1 (number-of  $w$ ) = number-of ( $w \ \text{BIT} \ \text{bit.B0}$ )
  apply (unfold shiffl1-def word-number-of-def)
  apply (simp add: word-ubin.norm-eq-iff [symmetric] word-ubin.eq-norm
    del: BIT-simps)
  apply (subst refl [THEN bintrunc-BIT-I, symmetric])
  apply (subst bintrunc-bintrunc-min)
  apply simp
  done

```

```

lemma shiffl1-0 [simp] : shiffl1 0 = 0
  unfolding word-0-no shiffl1-number by auto

```

```

lemmas shiffl1-def-u = shiffl1-def [folded word-number-of-def]

```

```

lemma shiffl1-def-s: shiffl1  $w$  = number-of ( $\text{sint } w \ \text{BIT} \ \text{bit.B0}$ )
  by (rule trans [OF - shiffl1-number]) simp

```

```

lemma shiftr1-0 [simp] : shiftr1 0 = 0
  unfolding shiftr1-def
  by simp (simp add: word-0-wi)

lemma sshiftr1-0 [simp] : sshiftr1 0 = 0
  apply (unfold sshiftr1-def)
  apply simp
  apply (simp add : word-0-wi)
  done

lemma sshiftr1-n1 [simp] : sshiftr1 -1 = -1
  unfolding sshiftr1-def by auto

lemma shiftl-0 [simp] : (0::'a::len0 word) << n = 0
  unfolding shiftl-def by (induct n) auto

lemma shiftr-0 [simp] : (0::'a::len0 word) >> n = 0
  unfolding shiftr-def by (induct n) auto

lemma sshiftr-0 [simp] : 0 >>> n = 0
  unfolding sshiftr-def by (induct n) auto

lemma sshiftr-n1 [simp] : -1 >>> n = -1
  unfolding sshiftr-def by (induct n) auto

lemma nth-shiftl1: shiftl1 w !! n = (n < size w & n > 0 & w !! (n - 1))
  apply (unfold shiftl1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm word-size)
  apply (cases n)
  apply auto
  done

lemma nth-shiftl' [rule-format]:
  ALL n. ((w::'a::len0 word) << m) !! n = (n < size w & n >= m & w !! (n -
m))
  apply (unfold shiftl-def)
  apply (induct m)
  apply (force elim!: test-bit-size)
  apply (clarsimp simp add : nth-shiftl1 word-size)
  apply arith
  done

lemmas nth-shiftl = nth-shiftl' [unfolded word-size]

lemma nth-shiftr1: shiftr1 w !! n = w !! Suc n
  apply (unfold shiftr1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm)
  apply safe

```

```

apply (drule bin-nth.Suc [THEN iffD2, THEN bin-nth-uint-imp])
apply simp
done

```

```

lemma nth-shiftr:
   $\bigwedge n. ((w::'a::len0 \text{ word}) >> m) !! n = w !! (n + m)$ 
apply (unfold shiftr-def)
apply (induct m)
apply (auto simp add : nth-shiftr1)
done

```

```

lemma uint-shiftr1: uint (shiftr1 w) = bin-rest (uint w)
apply (unfold shiftr1-def word-ubin.eq-norm bin-rest-trunc-i)
apply (subst bintr-uint [symmetric, OF order-refl])
apply (simp only : bintrunc-bintrunc-l)
apply simp
done

```

```

lemma nth-sshiftr1:
  sshiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
apply (unfold sshiftr1-def word-test-bit-def)
apply (simp add: nth-bintr word-ubin.eq-norm
               bin-nth.Suc [symmetric] word-size
               del: bin-nth.simps)
apply (simp add: nth-bintr uint-sint del : bin-nth.simps)
apply (auto simp add: bin-nth-sint)
done

```

```

lemma nth-sshiftr [rule-format] :
  ALL n. sshiftr w m !! n = (n < size w &
    (if n + m >= size w then w !! (size w - 1) else w !! (n + m)))
apply (unfold sshiftr-def)
apply (induct-tac m)
apply (simp add: test-bit-bl)
apply (clarsimp simp add: nth-sshiftr1 word-size)
apply safe
  apply arith
  apply arith
apply (erule thin-rl)
apply (case-tac n)
apply safe
apply simp
apply simp
apply (erule thin-rl)
apply (case-tac n)
apply safe
apply simp

```

```

    apply simp
    apply arith+
done

lemma shiftr1-div-2: uint (shiftr1 w) = uint w div 2
  apply (unfold shiftr1-def bin-rest-div)
  apply (rule word-uint.Abs-inverse)
  apply (simp add: uints-num pos-imp-zdiv-nonneg-iff)
  apply (rule xtr7)
  prefer 2
  apply (rule zdiv-le-dividend)
  apply auto
done

lemma sshiftr1-div-2: sint (sshiftr1 w) = sint w div 2
  apply (unfold sshiftr1-def bin-rest-div [symmetric])
  apply (simp add: word-sbin.eq-norm)
  apply (rule trans)
  defer
  apply (subst word-sbin.norm-Rep [symmetric])
  apply (rule refl)
  apply (subst word-sbin.norm-Rep [symmetric])
  apply (unfold One-nat-def)
  apply (rule sbintrunc-rest)
done

lemma shiftr-div-2n: uint (shiftr w n) = uint w div 2 ^ n
  apply (unfold shiftr-def)
  apply (induct n)
  apply simp
  apply (simp add: shiftr1-div-2 mult-commute
    zdiv-zmult2-eq [symmetric])
done

lemma sshiftr-div-2n: sint (sshiftr w n) = sint w div 2 ^ n
  apply (unfold sshiftr-def)
  apply (induct n)
  apply simp
  apply (simp add: sshiftr1-div-2 mult-commute
    zdiv-zmult2-eq [symmetric])
done

```

13.1.1 shift functions in terms of lists of bools

```

lemmas bshiftr1-no-bin [simp] =
  bshiftr1-def [where w=number-of w, unfolded to-bl-no-bin, standard]

lemma bshiftr1-bl: to-bl (bshiftr1 b w) = b # butlast (to-bl w)
  unfolding bshiftr1-def by (rule word-bl.Abs-inverse) simp

```


lemma *shifftl1-of-bl*: *shifftl1* (*of-bl* *bl*) = *of-bl* (*bl* @ [*False*])
unfolding *uint-bl of-bl-no*
by (*simp add: bl-to-bin-aux-append bl-to-bin-def*)

lemma *shifftl1-bl*: *shifftl1* (*w*::'*a*::len0 *word*) = *of-bl* (*to-bl* *w* @ [*False*])
proof –
have *shifftl1* *w* = *shifftl1* (*of-bl* (*to-bl* *w*)) **by** *simp*
also have ... = *of-bl* (*to-bl* *w* @ [*False*]) **by** (*rule shifftl1-of-bl*)
finally show ?*thesis* .
qed

lemma *bl-shifftl1*:
to-bl (*shifftl1* (*w*::'*a*::len *word*)) = *tl* (*to-bl* *w*) @ [*False*]
apply (*simp add: shifftl1-bl word-rep-drop drop-Suc drop-Cons'*)
apply (*fast intro!: Suc-leI*)
done

lemma *shiftr1-bl*: *shiftr1* *w* = *of-bl* (*butlast* (*to-bl* *w*))
apply (*unfold shiftr1-def uint-bl of-bl-def*)
apply (*simp add: butlast-rest-bin word-size*)
apply (*simp add: bin-rest-trunc [symmetric, unfolded One-nat-def]*)
done

lemma *bl-shiftr1*:
to-bl (*shiftr1* (*w*::'*a*::len *word*)) = *False* # *butlast* (*to-bl* *w*)
unfolding *shiftr1-bl*
by (*simp add : word-rep-drop len-gt-0 [THEN Suc-leI]*)

lemma *shifftl1-rev*:
shifftl1 (*w*::'*a*::len *word*) = *word-reverse* (*shiftr1* (*word-reverse* *w*))
apply (*unfold word-reverse-def*)
apply (*rule word-bl.Rep-inverse' [symmetric]*)
apply (*simp add: bl-shifftl1 bl-shiftr1 word-bl.Abs-inverse*)
apply (*cases to-bl w*)
apply *auto*
done

lemma *shiftr1-rev*:
shiftr1 (*w*::'*a*::len *word*) *n* = *word-reverse* (*shiftr* (*word-reverse* *w*) *n*)
apply (*unfold shiftr1-def shiftr-def*)
apply (*induct n*)
apply (*auto simp add : shifftl1-rev*)
done

lemmas *rev-shifftl* =
shiftr1-rev [**where** *w* = *word-reverse* *w*, *simplified*, *standard*]

lemmas *shiftr-rev* = *rev-shiftr* [*THEN* *word-rev-gal'*, *standard*]

lemmas *rev-shiftr* = *shiftr-rev* [*THEN* *word-rev-gal'*, *standard*]

lemma *bl-sshiftr1*:

to-bl (*sshiftr1* (*w* :: '*a* :: *len word*')) = *hd* (*to-bl w*) # *butlast* (*to-bl w*)

apply (*unfold sshiftr1-def uint-bl word-size*)

apply (*simp add: butlast-rest-bin word-ubin.eq-norm*)

apply (*simp add: sint-uint*)

apply (*rule nth-equalityI*)

apply *clarsimp*

apply *clarsimp*

apply (*case-tac i*)

apply (*simp-all add: hd-conv-nth length-0-conv [symmetric]*)

nth-bin-to-bl bin-nth.Suc [symmetric]

nth-sbintr

del: bin-nth.Suc)

apply *force*

apply (*rule impI*)

apply (*rule-tac f = bin-nth (uint w) in arg-cong*)

apply *simp*

done

lemma *drop-shiftr*:

drop n (*to-bl* ((*w* :: '*a* :: *len word*) >> *n*)) = *take* (*size w* - *n*) (*to-bl w*)

apply (*unfold shiftr-def*)

apply (*induct n*)

prefer 2

apply (*simp add: drop-Suc bl-shiftr1 butlast-drop [symmetric]*)

apply (*rule butlast-take [THEN trans]*)

apply (*auto simp: word-size*)

done

lemma *drop-sshiftr*:

drop n (*to-bl* ((*w* :: '*a* :: *len word*) >>> *n*)) = *take* (*size w* - *n*) (*to-bl w*)

apply (*unfold sshiftr-def*)

apply (*induct n*)

prefer 2

apply (*simp add: drop-Suc bl-sshiftr1 butlast-drop [symmetric]*)

apply (*rule butlast-take [THEN trans]*)

apply (*auto simp: word-size*)

done

lemma *take-shiftr* [*rule-format*] :

n <= *size* (*w* :: '*a* :: *len word*) --> *take n* (*to-bl* (*w* >> *n*)) =

replicate n False

apply (*unfold shiftr-def*)

apply (*induct n*)

prefer 2

```

apply (simp add: bl-shiftr1)
apply (rule impI)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

lemma take-sshiftr' [rule-format] :
  n <= size (w :: 'a :: len word) --> hd (to-bl (w >>> n)) = hd (to-bl w) &
  take n (to-bl (w >>> n)) = replicate n (hd (to-bl w))
apply (unfold sshiftr-def)
apply (induct n)
prefer 2
apply (simp add: bl-sshiftr1)
apply (rule impI)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

lemmas hd-sshiftr = take-sshiftr' [THEN conjunct1, standard]
lemmas take-sshiftr = take-sshiftr' [THEN conjunct2, standard]

lemma atd-lem: take n xs = t ==> drop n xs = d ==> xs = t @ d
by (auto intro: append-take-drop-id [symmetric])

lemmas bl-shiftr = atd-lem [OF take-shiftr drop-shiftr]
lemmas bl-sshiftr = atd-lem [OF take-sshiftr drop-sshiftr]

lemma shiftl-of-bl: of-bl bl << n = of-bl (bl @ replicate n False)
unfolding shiftl-def
by (induct n) (auto simp: shiftl1-of-bl replicate-app-Cons-same)

lemma shiftl-bl:
  (w :: 'a :: len0 word) << (n :: nat) = of-bl (to-bl w @ replicate n False)
proof -
  have w << n = of-bl (to-bl w) << n by simp
  also have ... = of-bl (to-bl w @ replicate n False) by (rule shiftl-of-bl)
  finally show ?thesis .
qed

lemmas shiftl-number [simp] = shiftl-def [where w=number-of w, standard]

lemma bl-shiftl:
  to-bl (w << n) = drop n (to-bl w) @ replicate (min (size w) n) False
by (simp add: shiftl-bl word-rep-drop word-size)

lemma shiftl-zero-size:
  fixes x :: 'a :: len0 word
  shows size x <= n ==> x << n = 0
  apply (unfold word-size)

```

```

apply (rule word-eqI)
apply (clarsimp simp add: shiftl-bl word-size test-bit-of-bl nth-append)
done

```

```

lemma shiftl1-2t: shiftl1 (w :: 'a :: len word) = 2 * w
apply (simp add: shiftl1-def-u)
apply (simp only: double-number-of-Bit0 [symmetric])
apply simp
done

```

```

lemma shiftl1-p: shiftl1 (w :: 'a :: len word) = w + w
apply (simp add: shiftl1-def-u)
apply (simp only: double-number-of-Bit0 [symmetric])
apply simp
done

```

```

lemma shiftl-t2n: shiftl (w :: 'a :: len word) n = 2 ^ n * w
unfolding shiftl-def
by (induct n) (auto simp: shiftl1-2t power-Suc)

```

```

lemma shiftr1-bintr [simp]:
  (shiftr1 (number-of w) :: 'a :: len0 word) =
    number-of (bin-rest (bintrunc (len-of TYPE ('a)) w))
unfolding shiftr1-def word-number-of-def
by (simp add : word-ubin.eq-norm)

```

```

lemma sshiftr1-sbintr [simp] :
  (sshiftr1 (number-of w) :: 'a :: len word) =
    number-of (bin-rest (sbintrunc (len-of TYPE ('a) - 1) w))
unfolding sshiftr1-def word-number-of-def
by (simp add : word-sbin.eq-norm)

```

```

lemma shiftr-no':
  w = number-of bin ==>
  (w::'a::len0 word) >> n = number-of ((bin-rest ^ n) (bintrunc (size w) bin))
applyclarsimp
apply (rule word-eqI)
apply (auto simp: nth-shiftr nth-rest-power-bin nth-bintr word-size)
done

```

```

lemma sshiftr-no':
  w = number-of bin ==> w >>> n = number-of ((bin-rest ^ n)
    (sbintrunc (size w - 1) bin))
applyclarsimp
apply (rule word-eqI)
apply (auto simp: nth-sshiftr nth-rest-power-bin nth-sbintr word-size)
apply (subgoal-tac na + n = len-of TYPE('a) - Suc 0, simp, simp)+

```

done

lemmas *sshiftr-no* [*simp*] =
sshiftr-no' [**where** *w* = number-of *w*, OF refl, unfolded word-size, standard]

lemmas *shiftr-no* [*simp*] =
shiftr-no' [**where** *w* = number-of *w*, OF refl, unfolded word-size, standard]

lemma *shiftr1-bl-of'*:
 $us = shiftr1\ (of\text{-}bl\ bl) \implies length\ bl \leq size\ us \implies$
 $us = of\text{-}bl\ (butlast\ bl)$
by (*clarsimp simp: shiftr1-def of-bl-def word-size butlast-rest-bl2bin*
word-ubin.eq-norm trunc-bl2bin)

lemmas *shiftr1-bl-of* = refl [*THEN shiftr1-bl-of'*, unfolded word-size]

lemma *shiftr-bl-of'* [*rule-format*]:
 $us = of\text{-}bl\ bl \gg n \implies length\ bl \leq size\ us \dashv\vdash$
 $us = of\text{-}bl\ (take\ (length\ bl - n)\ bl)$
apply (*unfold shiftr-def*)
apply *hypsubst*
apply (*unfold word-size*)
apply (*induct n*)
apply *clarsimp*
apply *clarsimp*
apply (*subst shiftr1-bl-of*)
apply *simp*
apply (*simp add: butlast-take*)
done

lemmas *shiftr-bl-of* = refl [*THEN shiftr-bl-of'*, unfolded word-size]

lemmas *shiftr-bl* = *word-bl.Rep'* [*THEN eq-imp-le*, *THEN shiftr-bl-of*,
simplified word-size, simplified, THEN eq-reflection, standard]

lemma *msb-shift'*: $msb\ (w::'a::len\ word) \longleftrightarrow (w \gg (size\ w - 1)) \approx = 0$
apply (*unfold shiftr-bl word-msb-alt*)
apply (*simp add: word-size Suc-le-eq take-Suc*)
apply (*cases hd (to-bl w)*)
apply (*auto simp: word-1-bl word-0-bl*
of-bl-rep-False [where n=1 and bs=[], simplified])
done

lemmas *msb-shift* = *msb-shift'* [*unfolded word-size*]

lemma *align-lem-or* [*rule-format*] :
ALL x m. length x = n + m $\dashv\vdash$ length y = n + m $\dashv\vdash$
drop m x = replicate n False $\dashv\vdash$ take m y = replicate m False $\dashv\vdash$
map2 op | x y = take m x @ drop m y

```

apply (induct-tac y)
  apply force
apply clarsimp
apply (case-tac x, force)
apply (case-tac m, auto)
apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma align-lem-and [rule-format] :
  ALL x m. length x = n + m --> length y = n + m -->
    drop m x = replicate n False --> take m y = replicate m False -->
      map2 op & x y = replicate (n + m) False
apply (induct-tac y)
  apply force
apply clarsimp
apply (case-tac x, force)
apply (case-tac m, auto)
apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma aligned-bl-add-size':
  size x - n = m ==> n <= size x ==> drop m (to-bl x) = replicate n False
==>
  take m (to-bl y) = replicate m False ==>
    to-bl (x + y) = take m (to-bl x) @ drop m (to-bl y)
apply (subgoal-tac x AND y = 0)
prefer 2
apply (rule word-bl.Rep-eqD)
apply (simp add: bl-word-and to-bl-0)
apply (rule align-lem-and [THEN trans])
  apply (simp-all add: word-size)[5]
apply simp
apply (subst word-plus-and-or [symmetric])
apply (simp add : bl-word-or)
apply (rule align-lem-or)
  apply (simp-all add: word-size)
done

```

```

lemmas aligned-bl-add-size = refl [THEN aligned-bl-add-size']

```

13.1.2 Mask

```

lemma nth-mask': m = mask n ==> test-bit m i = (i < n & i < size m)
  apply (unfold mask-def test-bit-bl)
  apply (simp only: word-1-bl [symmetric] shiffl-of-bl)

```

```

apply (clarsimp simp add: word-size)
apply (simp only: of-bl-no mask-lem number-of-succ add-diff-cancel2)
apply (fold of-bl-no)
apply (simp add: word-1-bl)
apply (rule test-bit-of-bl [THEN trans, unfolded test-bit-bl word-size])
apply auto
done

lemmas nth-mask [simp] = refl [THEN nth-mask']

lemma mask-bl: mask n = of-bl (replicate n True)
  by (auto simp add : test-bit-of-bl word-size intro: word-eqI)

lemma mask-bin: mask n = number-of (bintrunc n Int.Min)
  by (auto simp add: nth-bintr word-size intro: word-eqI)

lemma and-mask-bintr: w AND mask n = number-of (bintrunc n (uint w))
  apply (rule word-eqI)
  apply (simp add: nth-bintr word-size word-ops-nth-size)
  apply (auto simp add: test-bit-bin)
  done

lemma and-mask-no: number-of i AND mask n = number-of (bintrunc n i)
  by (auto simp add : nth-bintr word-size word-ops-nth-size intro: word-eqI)

lemmas and-mask-wi = and-mask-no [unfolded word-number-of-def]

lemma bl-and-mask:
  to-bl (w AND mask n :: 'a :: len word) =
    replicate (len-of TYPE('a) - n) False @
    drop (len-of TYPE('a) - n) (to-bl w)
  apply (rule nth-equalityI)
  apply simp
  apply (clarsimp simp add: to-bl-nth word-size)
  apply (simp add: word-size word-ops-nth-size)
  apply (auto simp add: word-size test-bit-bl nth-append nth-rev)
  done

lemmas and-mask-mod-2p =
  and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma and-mask-lt-2p: uint (w AND mask n) < 2 ^ n
  apply (simp add : and-mask-bintr no-bintr-alt)
  apply (rule xtr8)
  prefer 2
  apply (rule pos-mod-bound)
  apply auto
  done

```

lemmas $eq\text{-}mod\text{-}iff = trans [symmetric, OF\ int\text{-}mod\text{-}lem\ eq\text{-}sym\text{-}conv]$

lemma $mask\text{-}eq\text{-}iff: (w\ AND\ mask\ n) = w <-> uint\ w < 2^{\wedge} n$
apply (*simp add: and-mask-bintr word-number-of-def*)
apply (*simp add: word-ubin.inverse-norm*)
apply (*simp add: eq-mod-iff bintrunc-mod2p min-def*)
apply (*fast intro!: lt2p-lem*)
done

lemma $and\text{-}mask\text{-}dvd: 2^{\wedge} n\ dvd\ uint\ w = (w\ AND\ mask\ n = 0)$
apply (*simp add: dvd-eq-mod-eq-0 and-mask-mod-2p*)
apply (*simp add: word-uint.norm-eq-iff [symmetric] word-of-int-homs*)
apply (*subst word-uint.norm-Rep [symmetric]*)
apply (*simp only: bintrunc-bintrunc-min bintrunc-mod2p [symmetric] min-def*)
apply *auto*
done

lemma $and\text{-}mask\text{-}dvd\text{-}nat: 2^{\wedge} n\ dvd\ unat\ w = (w\ AND\ mask\ n = 0)$
apply (*unfold unat-def*)
apply (*rule trans [OF - and-mask-dvd]*)
apply (*unfold dvd-def*)
apply *auto*
apply (*drule uint-ge-0 [THEN nat-int.Abs-inverse' [simplified], symmetric]*)
apply (*simp add : int-mult int-power*)
apply (*simp add : nat-mult-distrib nat-power-eq*)
done

lemma $word\text{-}2p\text{-}lem:$
 $n < size\ w ==> w < 2^{\wedge} n = (uint\ (w :: 'a :: len\ word) < 2^{\wedge} n)$
apply (*unfold word-size word-less-alt word-number-of-alt*)
apply (*clarsimp simp add: word-of-int-power-hom word-uint.eq-norm*
 $int\text{-}mod\text{-}eq'$
 $simp\ del: word\text{-}of\text{-}int\text{-}bin$)
done

lemma $less\text{-}mask\text{-}eq: x < 2^{\wedge} n ==> x\ AND\ mask\ n = (x :: 'a :: len\ word)$
apply (*unfold word-less-alt word-number-of-alt*)
apply (*clarsimp simp add: and-mask-mod-2p word-of-int-power-hom*
 $word\text{-}uint.eq\text{-}norm$
 $simp\ del: word\text{-}of\text{-}int\text{-}bin$)
apply (*drule xtr8 [rotated]*)
apply (*rule int-mod-le*)
apply (*auto simp add : mod-pos-pos-trivial*)
done

lemmas $mask\text{-}eq\text{-}iff\text{-}w2p =$
 $trans [OF\ mask\text{-}eq\text{-}iff\ word\text{-}2p\text{-}lem [symmetric], standard]$

lemmas $and\text{-}mask\text{-}less' =$

iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND mask } n < 2^n$
unfolding *word-size* **by** (*erule and-mask-less'*)

lemma *word-mod-2p-is-mask'*:
 $c = 2^n \implies c > 0 \implies x \bmod c = (x :: 'a :: \text{len word}) \text{ AND mask } n$
by (*clarsimp simp add: word-mod-def uint-2p and-mask-mod-2p*)

lemmas *word-mod-2p-is-mask* = *refl* [*THEN word-mod-2p-is-mask'*]

lemma *mask-egs*:
 $(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $(a \text{ AND mask } n) * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $-(a \text{ AND mask } n) \text{ AND mask } n = -a \text{ AND mask } n$
 $\text{word-succ } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-succ } a \text{ AND mask } n$
 $\text{word-pred } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-pred } a \text{ AND mask } n$
using *word-of-int-Ex* [**where** $x=a$] *word-of-int-Ex* [**where** $x=b$]
by (*auto simp: and-mask-wi bintr-ariths bintr-arith1s new-word-of-int-homs*)

lemma *mask-power-eq*:
 $(x \text{ AND mask } n)^k \text{ AND mask } n = x^k \text{ AND mask } n$
using *word-of-int-Ex* [**where** $x=x$]
by (*clarsimp simp: and-mask-wi word-of-int-power-hom bintr-ariths*)

13.1.3 Recast

lemmas *revcast-def'* = *revcast-def* [*simplified*]
lemmas *revcast-def''* = *revcast-def'* [*simplified word-size*]
lemmas *revcast-no-def* [*simp*] =
revcast-def' [**where** $w=\text{number-of } w$, *unfolded word-size, standard*]

lemma *to-bl-revcast*:
 $\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
apply (*unfold revcast-def' word-size*)
apply (*rule word-bl.Abs-inverse*)
apply *simp*
done

lemma *revcast-rev-ucast'*:
 $cs = [rc, uc] \implies rc = \text{revcast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$

```

    rc = word-reverse uc
  apply (unfold ucast-def revcast-def' Let-def word-reverse-def)
  apply (clarsimp simp add : to-bl-of-bin takefill-bintrunc)
  apply (simp add : word-bl.Abs-inverse word-size)
done

lemmas revcast-rev-ucast = revcast-rev-ucast' [OF refl refl refl]

lemmas revcast-ucast = revcast-rev-ucast
  [where w = word-reverse w, simplified word-rev-rev, standard]

lemmas ucast-revcast = revcast-rev-ucast [THEN word-rev-gal', standard]
lemmas ucast-rev-revcast = revcast-ucast [THEN word-rev-gal', standard]

```

— linking revcast and cast via shift

```

lemmas wsst-TYs = source-size target-size word-size

lemma revcast-down-uu':
  rc = revcast ==> source-size rc = target-size rc + n ==>
    rc (w :: 'a :: len word) = ucast (w >> n)
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule ucast-down-drop)
  prefer 2
  apply (rule trans, rule drop-shiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

lemma revcast-down-us':
  rc = revcast ==> source-size rc = target-size rc + n ==>
    rc (w :: 'a :: len word) = ucast (w >>> n)
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule ucast-down-drop)
  prefer 2
  apply (rule trans, rule drop-sshiftr)
  apply (auto simp: takefill-alt wsst-TYs)
done

lemma revcast-down-su':
  rc = revcast ==> source-size rc = target-size rc + n ==>
    rc (w :: 'a :: len word) = scast (w >> n)
  apply (simp add: revcast-def')
  apply (rule word-bl.Rep-inverse')
  apply (rule trans, rule scast-down-drop)
  prefer 2
  apply (rule trans, rule drop-shiftr)

```

```

apply (auto simp: takefill-alt wsst-TYs)
done

```

lemma *revcast-down-ss'*:

```

rc = revcast ==> source-size rc = target-size rc + n ==>
  rc (w :: 'a :: len word) = scast (w >>> n)
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (rule trans, rule scast-down-drop)
prefer 2
apply (rule trans, rule drop-sshiftr)
apply (auto simp: takefill-alt wsst-TYs)
done

```

lemmas *revcast-down-uu* = refl [THEN *revcast-down-uu*']

lemmas *revcast-down-us* = refl [THEN *revcast-down-us*']

lemmas *revcast-down-su* = refl [THEN *revcast-down-su*']

lemmas *revcast-down-ss* = refl [THEN *revcast-down-ss*']

lemma *cast-down-rev*:

```

uc = ucast ==> source-size uc = target-size uc + n ==>
  uc w = revcast ((w :: 'a :: len word) <<< n)
apply (unfold shiftl-rev)
apply clarify
apply (simp add: revcast-rev-ucast)
apply (rule word-rev-gal')
apply (rule trans [OF - revcast-rev-ucast])
apply (rule revcast-down-uu [symmetric])
apply (auto simp add: wsst-TYs)
done

```

lemma *revcast-up'*:

```

rc = revcast ==> source-size rc + n = target-size rc ==>
  rc w = (ucast w :: 'a :: len word) <<< n
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (simp add: takefill-alt)
apply (rule bl-shiftr [THEN trans])
apply (subst ucast-up-app)
apply (auto simp add: wsst-TYs)
done

```

lemmas *revcast-up* = refl [THEN *revcast-up*']

lemmas *rc1* = *revcast-up* [THEN

revcast-rev-ucast [symmetric, THEN *trans*, THEN *word-rev-gal*, symmetric]]

lemmas *rc2* = *revcast-down-uu* [THEN

revcast-rev-ucast [symmetric, THEN *trans*, THEN *word-rev-gal*, symmetric]]

```

lemmas ucast-up =
  rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]
lemmas ucast-down =
  rc2 [simplified rev-shiftr revcast-ucast [symmetric]]

```

13.1.4 Slices

```

lemmas slice1-no-bin [simp] =
  slice1-def [where w=number-of w, unfolded to-bl-no-bin, standard]

```

```

lemmas slice-no-bin [simp] =
  trans [OF slice-def [THEN meta-eq-to-obj-eq]
    slice1-no-bin [THEN meta-eq-to-obj-eq],
    unfolded word-size, standard]

```

```

lemma slice1-0 [simp] : slice1 n 0 = 0
  unfolding slice1-def by (simp add : to-bl-0)

```

```

lemma slice-0 [simp] : slice n 0 = 0
  unfolding slice-def by auto

```

```

lemma slice-take': slice n w = of-bl (take (size w - n) (to-bl w))
  unfolding slice-def' slice1-def
  by (simp add : takefill-alt word-size)

```

```

lemmas slice-take = slice-take' [unfolded word-size]

```

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

```

lemmas shiftr-slice = trans
  [OF shiftr-bl [THEN meta-eq-to-obj-eq] slice-take [symmetric], standard]

```

```

lemma slice-shiftr: slice n w = ucast (w >> n)
  apply (unfold slice-take shiftr-bl)
  apply (rule ucast-of-bl-up [symmetric])
  apply (simp add: word-size)
  done

```

```

lemma nth-slice:
  (slice n w :: 'a :: len0 word) !! m =
  (w !! (m + n) & m < len-of TYPE ('a))
  unfolding slice-shiftr
  by (simp add : nth-ucast nth-shiftr)

```

```

lemma slice1-down-alt':
  sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>
  to-bl sl = takefill False fs (drop k (to-bl w))
  unfolding slice1-def word-size of-bl-def uint-bl
  by (clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop drop-takefill)

```

```

lemma slice1-up-alt':
  sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>
    to-bl sl = takefill False fs (replicate k False @ (to-bl w))
apply (unfold slice1-def word-size of-bl-def uint-bl)
apply (clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop
    takefill-append [symmetric])
apply (rule-tac f = %k. takefill False (len-of TYPE('a))
    (replicate k False @ bin-to-bl (len-of TYPE('b)) (uint w)) in arg-cong)
apply arith
done

lemmas sd1 = slice1-down-alt' [OF refl refl, unfolded word-size]
lemmas su1 = slice1-up-alt' [OF refl refl, unfolded word-size]
lemmas slice1-down-alt = le-add-diff-inverse [THEN sd1]
lemmas slice1-up-alt =
  le-add-diff-inverse [symmetric, THEN su1]
  le-add-diff-inverse2 [symmetric, THEN su1]

lemma ucast-slice1: ucast w = slice1 (size w) w
  unfolding slice1-def ucast-bl
  by (simp add : takefill-same' word-size)

lemma ucast-slice: ucast w = slice 0 w
  unfolding slice-def by (simp add : ucast-slice1)

lemmas slice-id = trans [OF ucast-slice [symmetric] ucast-id]

lemma revcast-slice1':
  rc = revcast w ==> slice1 (size rc) w = rc
  unfolding slice1-def revcast-def' by (simp add : word-size)

lemmas revcast-slice1 = refl [THEN revcast-slice1']

lemma slice1-tf-tf':
  to-bl (slice1 n w :: 'a :: len0 word) =
    rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))
  unfolding slice1-def by (rule word-rev-tf)

lemmas slice1-tf-tf = slice1-tf-tf'
  [THEN word-bl.Rep-inverse', symmetric, standard]

lemma rev-slice1:
  n + k = len-of TYPE('a) + len-of TYPE('b) ==>
    slice1 n (word-reverse w :: 'b :: len0 word) =
    word-reverse (slice1 k w :: 'a :: len0 word)
  apply (unfold word-reverse-def slice1-tf-tf)
  apply (rule word-bl.Rep-inverse')
  apply (rule rev-swap [THEN iffD1])
  apply (rule trans [symmetric])

```

```

apply (rule tf-rev)
apply (simp add: word-bl.Abs-inverse)
apply (simp add: word-bl.Abs-inverse)
done

```

lemma *rev-slice'*:

```

res = slice n (word-reverse w) ==> n + k + size res = size w ==>
  res = word-reverse (slice k w)
apply (unfold slice-def word-size)
apply clarify
apply (rule rev-slice1)
apply arith
done

```

lemmas *rev-slice* = *refl* [*THEN rev-slice'*, *unfolded word-size*]

lemmas *sym-notr* =

```

not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

```

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:

```

(sint (x :: 'a :: len word) + sint y = sint (x + y)) =
  (((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)
apply (unfold word-size)
apply (cases len-of TYPE('a'), simp)
apply (subst msb-shift [THEN sym-notr])
apply (simp add: word-ops-msb)
apply (simp add: word-msb-sint)
apply safe
  apply simp-all
  apply (unfold sint-word-ariths)
  apply (unfold word-sbin.set-iff-norm [symmetric] sints-num)
  apply safe
    apply (insert sint-range' [where x=x])
    apply (insert sint-range' [where x=y])
    defer
    apply (simp (no-asm), arith)
    apply (simp (no-asm), arith)
    defer
    defer
    apply (simp (no-asm), arith)
    apply (simp (no-asm), arith)
  apply (rule notI [THEN notnotD],
    drule leI not-leE,
    drule sbintrunc-inc sbintrunc-dec,
    simp) +
done

```

13.2 Split and cat

lemmas *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

lemmas *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-rsplit-no*:

(*word-rsplit* (*number-of bin* :: '*b* :: len0 word) :: '*a* word list) =
 map *number-of* (*bin-rsplit* (*len-of TYPE*('a :: len))
 (*len-of TYPE*('b), *bintrunc* (*len-of TYPE*('b)) bin))

apply (*unfold word-rsplit-def word-no-wi*)

apply (*simp add: word-ubin.eq-norm*)

done

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*

[*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

lemma *test-bit-cat*:

wc = *word-cat a b* ==> *wc* !! *n* = (*n* < *size wc* &
 (*if n* < *size b* then *b* !! *n* else *a* !! (*n* - *size b*)))

apply (*unfold word-cat-bin' test-bit-bin*)

apply (*auto simp add : word-ubin.eq-norm nth-bintr bin-nth-cat word-size*)

apply (*erule bin-nth-uint-imp*)

done

lemma *word-cat-bl*: *word-cat a b* = *of-bl* (*to-bl a* @ *to-bl b*)

apply (*unfold of-bl-def to-bl-def word-cat-bin'*)

apply (*simp add: bl-to-bin-app-cat*)

done

lemma *of-bl-append*:

(*of-bl* (*xs* @ *ys*) :: '*a* :: len word) = *of-bl xs* * $2^{length\ ys}$ + *of-bl ys*

apply (*unfold of-bl-def*)

apply (*simp add: bl-to-bin-app-cat bin-cat-num*)

apply (*simp add: word-of-int-power-hom [symmetric] new-word-of-int-hom-syms*)

done

lemma *of-bl-False* [*simp*]:

of-bl (*False#xs*) = *of-bl xs*

by (*rule word-eqI*)

(*auto simp add: test-bit-of-bl nth-append*)

lemma *of-bl-True*:

(*of-bl* (*True#xs*)::'*a*::len word) = $2^{length\ xs}$ + *of-bl xs*

by (*subst of-bl-append [where xs=[True], simplified]*)

(*simp add: word-1-bl*)

lemma *of-bl-Cons*:

of-bl (*x#xs*) = *of-bool x* * $2^{length\ xs}$ + *of-bl xs*

by (*cases x*) (*simp-all add: of-bl-True*)

```

lemma split-uint-lem: bin-split n (uint (w :: 'a :: len0 word)) = (a, b) ==>
  a = bintrunc (len-of TYPE('a) - n) a & b = bintrunc (len-of TYPE('a)) b
apply (frule word-ubin.norm-Rep [THEN ssubst])
apply (drule bin-split-trunc1)
apply (drule sym [THEN trans])
apply assumption
apply safe
done

```

```

lemma word-split-bl':
  std = size c - size b ==> (word-split c = (a, b)) ==>
    (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c)))
apply (unfold word-split-bin')
apply safe
defer
apply (clarsimp split: prod.splits)
apply (drule word-ubin.norm-Rep [THEN ssubst])
apply (drule split-bintrunc)
apply (simp add : of-bl-def bl2bin-drop word-size
  word-ubin.norm-eq-iff [symmetric] min-def del : word-ubin.norm-Rep)
apply (clarsimp split: prod.splits)
apply (frule split-uint-lem [THEN conjunct1])
apply (unfold word-size)
apply (cases len-of TYPE('a) >= len-of TYPE('b))
defer
apply (simp add: word-0-bl word-0-wi-Pls)
apply (simp add : of-bl-def to-bl-def)
apply (subst bin-split-take1 [symmetric])
prefer 2
apply assumption
apply simp
apply (erule thin-rl)
apply (erule arg-cong [THEN trans])
apply (simp add : word-ubin.norm-eq-iff [symmetric])
done

```

```

lemma word-split-bl: std = size c - size b ==>
  (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c))) <->
  word-split c = (a, b)
apply (rule iffI)
defer
apply (erule (1) word-split-bl')
apply (case-tac word-split c)
apply (auto simp add : word-size)
apply (frule word-split-bl' [rotated])
apply (auto simp add : word-size)
done

```

```

lemma word-split-bl-eq:

```



```

(word-split (c::'a::len word) :: ('c :: len0 word * 'd :: len0 word)) =
  (of-bl (take (len-of TYPE('a::len) - len-of TYPE('d::len0)) (to-bl c)),
   of-bl (drop (len-of TYPE('a) - len-of TYPE('d)) (to-bl c)))
apply (rule word-split-bl [THEN iffD1])
apply (unfold word-size)
apply (rule refl conjI)+
done

```

— keep quantifiers for use in simplification

```

lemma test-bit-split':
  word-split c = (a, b) --> (ALL n m. b !! n = (n < size b & c !! n) &
    a !! m = (m < size a & c !! (m + size b)))
apply (unfold word-split-bin' test-bit-bin)
apply (clarify)
apply (clarsimp simp: word-ubin.eq-norm nth-bintr word-size split: prod.splits)
apply (drule bin-nth-split)
apply safe
  apply (simp-all add: add-commute)
apply (erule bin-nth-uint-imp)+
done

```

```

lemmas test-bit-split =
  test-bit-split' [THEN mp, simplified all-simps, standard]

```

```

lemma test-bit-split-eq: word-split c = (a, b) <->
  ((ALL n::nat. b !! n = (n < size b & c !! n)) &
   (ALL m::nat. a !! m = (m < size a & c !! (m + size b))))
apply (rule-tac iffI)
apply (rule-tac conjI)
  apply (erule test-bit-split [THEN conjunct1])
  apply (erule test-bit-split [THEN conjunct2])
apply (case-tac word-split c)
apply (frule test-bit-split)
apply (erule trans)
apply (fastsimp intro ! : word-eqI simp add : word-size)
done

```

— this odd result is analogous to *ucast-id*, result to the length given by the result type

```

lemma word-cat-id: word-cat a b = b
  unfolding word-cat-bin' by (simp add: word-ubin.inverse-norm)

```

— limited hom result

```

lemma word-cat-hom:
  len-of TYPE('a::len0) <= len-of TYPE('b::len0) + len-of TYPE('c::len0)
  ==>
  (word-cat (word-of-int w :: 'b word) (b :: 'c word) :: 'a word) =
  word-of-int (bin-cat w (size b) (uint b))

```

```

apply (unfold word-cat-def word-size)
apply (clarsimp simp add: word-ubin.norm-eq-iff [symmetric]
      word-ubin.eq-norm bintr-cat min-max.inf-absorb1)
done

```

```

lemma word-cat-split-alt:
  size w <= size u + size v ==> word-split w = (u, v) ==> word-cat u v = w
apply (rule word-eqI)
apply (drule test-bit-split)
apply (clarsimp simp add: test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemmas word-cat-split-size =
  sym [THEN [2] word-cat-split-alt [symmetric], standard]

```

13.2.1 Split and slice

```

lemma split-slices:
  word-split w = (u, v) ==> u = slice (size v) w & v = slice 0 w
apply (drule test-bit-split)
apply (rule conjI)
apply (rule word-eqI, clarsimp simp: nth-slice word-size)+
done

```

```

lemma slice-cat1':
  wc = word-cat a b ==> size wc >= size a + size b ==> slice (size b) wc = a
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
done

```

```

lemmas slice-cat1 = refl [THEN slice-cat1']
lemmas slice-cat2 = trans [OF slice-id word-cat-id]

```

```

lemma cat-slices:
  a = slice n c ==> b = slice 0 c ==> n = size b ==>
    size a + size b >= size c ==> word-cat a b = c
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemma word-split-cat-alt:
  w = word-cat u v ==> size u + size v <= size w ==> word-split w = (u, v)
apply (case-tac word-split ?w)

```

```

apply (rule trans, assumption)
apply (drule test-bit-split)
apply safe
apply (rule word-eqI, clarsimp simp: test-bit-cat word-size)+
done

```

```

lemmas word-cat-bl-no-bin [simp] =
  word-cat-bl [where a=number-of a
    and b=number-of b,
    unfolded to-bl-no-bin, standard]

```

```

lemmas word-split-bl-no-bin [simp] =
  word-split-bl-eq [where c=number-of c, unfolded to-bl-no-bin, standard]

```

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

```

lemma word-rsplit-same: word-rsplit w = [w]
  unfolding word-rsplit-def by (simp add : bin-rsplit-all)

```

```

lemma word-rsplit-empty-iff-size:
  (word-rsplit w = []) = (size w = 0)
  unfolding word-rsplit-def bin-rsplit-def word-size
  by (simp add: bin-rsplit-aux-simp-alt Let-def split: Product-Type.split-split)

```

```

lemma test-bit-rsplit:
  sw = word-rsplit w ==> m < size (hd sw :: 'a :: len word) ==>
    k < length sw ==> (rev sw ! k) !! m = (w !! (k * size (hd sw) + m))
  apply (unfold word-rsplit-def word-test-bit-def)
  apply (rule trans)
  apply (rule-tac f = %x. bin-nth x m in arg-cong)
  apply (rule nth-map [symmetric])
  apply simp
  apply (rule bin-nth-rsplit)
  apply simp-all
  apply (simp add : word-size rev-map)
  apply (rule trans)
  defer
  apply (rule map-ident [THEN fun-cong])
  apply (rule refl [THEN map-cong])
  apply (simp add : word-ubin.eq-norm)
  apply (erule bin-rsplit-size-sign [OF len-gt-0 refl])
  done

```

```

lemma word-rcat-bl: word-rcat wl == of-bl (concat (map to-bl wl))
  unfolding word-rcat-def to-bl-def' of-bl-def
  by (clarsimp simp add : bin-rcat-bl)

```

```

lemma size-rcat-lem':
  size (concat (map to-bl wl)) = length wl * size (hd wl)
  unfolding word-size by (induct wl) auto

lemmas size-rcat-lem = size-rcat-lem' [unfolded word-size]

lemmas td-gal-lt-len = len-gt-0 [THEN td-gal-lt, standard]

lemma nth-rcat-lem' [rule-format] :
  sw = size (hd wl :: 'a :: len word) ==> (ALL n. n < size wl * sw -->
    rev (concat (map to-bl wl)) ! n =
    rev (to-bl (rev wl ! (n div sw))) ! (n mod sw))
  apply (unfold word-size)
  apply (induct wl)
  apply clarsimp
  apply (clarsimp simp add : nth-append size-rcat-lem)
  apply (simp (no-asm-use) only: mult-Suc [symmetric]
    td-gal-lt-len less-Suc-eq-le mod-div-equality')
  apply clarsimp
  done

lemmas nth-rcat-lem = refl [THEN nth-rcat-lem', unfolded word-size]

lemma test-bit-rcat:
  sw = size (hd wl :: 'a :: len word) ==> rc = word-rcat wl ==> rc !! n =
    (n < size rc & n div sw < size wl & (rev wl) ! (n div sw) !! (n mod sw))
  apply (unfold word-rcat-bl word-size)
  apply (clarsimp simp add :
    test-bit-of-bl size-rcat-lem word-size td-gal-lt-len)
  apply safe
  apply (auto simp add :
    test-bit-bl word-size td-gal-lt-len [THEN iffD2, THEN nth-rcat-lem])
  done

lemma foldl-eq-foldr [rule-format] :
  ALL x. foldl op + x xs = foldr op + (x # xs) (0 :: 'a :: comm-monoid-add)
  by (induct xs) (auto simp add : add-assoc)

lemmas test-bit-cong = arg-cong [where f = test-bit, THEN fun-cong]

lemmas test-bit-rsplit-alt =
  trans [OF nth-rev-alt [THEN test-bit-cong]
  test-bit-rsplit [OF refl asm-rl diff-Suc-less]]

— lazy way of expressing that u and v, and su and sv, have same types
lemma word-rsplit-len-indep':
  [u,v] = p ==> [su,sv] = q ==> word-rsplit u = su ==>
    word-rsplit v = sv ==> length su = length sv
  apply (unfold word-rsplit-def)

```

```

apply (auto simp add : bin-rsplit-len-indep)
done

lemmas word-rsplit-len-indep = word-rsplit-len-indep' [OF refl refl refl refl]

lemma length-word-rsplit-size:
  n = len-of TYPE ('a :: len) ==>
    (length (word-rsplit w :: 'a word list) <= m) = (size w <= m * n)
apply (unfold word-rsplit-def word-size)
apply (clarsimp simp add : bin-rsplit-len-le)
done

lemmas length-word-rsplit-lt-size =
  length-word-rsplit-size [unfolded Not-eq-iff linorder-not-less [symmetric]]

lemma length-word-rsplit-exp-size:
  n = len-of TYPE ('a :: len) ==>
    length (word-rsplit w :: 'a word list) = (size w + n - 1) div n
unfolding word-rsplit-def by (clarsimp simp add : word-size bin-rsplit-len)

lemma length-word-rsplit-even-size:
  n = len-of TYPE ('a :: len) ==> size w = m * n ==>
    length (word-rsplit w :: 'a word list) = m
by (clarsimp simp add : length-word-rsplit-exp-size given-quot-alt)

lemmas length-word-rsplit-exp-size' = refl [THEN length-word-rsplit-exp-size]

lemmas tdle = iffD2 [OF split-div-lemma refl, THEN conjunct1]
lemmas dtle = xtr4 [OF tdle mult-commute]

lemma word-rcat-rsplit: word-rcat (word-rsplit w) = w
apply (rule word-eqI)
apply (clarsimp simp add : test-bit-rcat word-size)
apply (subst refl [THEN test-bit-rsplit])
apply (simp-all add: word-size
  refl [THEN length-word-rsplit-size [simplified not-less [symmetric], simplified]])
apply safe
apply (erule xtr7, rule len-gt-0 [THEN dtle]) +
done

lemma size-word-rsplit-rcat-size':
  word-rcat (ws :: 'a :: len word list) = frcw ==>
    size frcw = length ws * len-of TYPE ('a) ==>
    size (hd [word-rsplit frcw, ws]) = size ws
apply (clarsimp simp add : word-size length-word-rsplit-exp-size')
apply (fast intro: given-quot-alt)
done

```

lemmas *size-word-rsplit-rcat-size* =
size-word-rsplit-rcat-size' [*simplified*]

lemma *msreus*:
fixes *n::nat*
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
by (*auto simp: add-commute*)

lemma *word-rsplit-rcat-size'*:
word-rcat (*ws* :: 'a :: len word list) = *frcw* ==>
size frcw = *length ws* * *len-of TYPE ('a)* ==> *word-rsplit frcw* = *ws*
apply (*frule size-word-rsplit-rcat-size, assumption*)
apply (*clarsimp simp add : word-size*)
apply (*rule nth-equalityI, assumption*)
apply *clarsimp*
apply (*rule word-eqI*)
apply (*rule trans*)
apply (*rule test-bit-rsplit-alt*)
apply (*clarsimp simp: word-size*)+
apply (*rule trans*)
apply (*rule test-bit-rcat [OF refl refl]*)
apply (*simp add : word-size msreus*)
apply (*subst nth-rev*)
apply *arith*
apply (*simp add : le0 [THEN [2] xtr7, THEN diff-Suc-less]*)
apply *safe*
apply (*simp add : diff-mult-distrib*)
apply (*rule mpl-lem*)
apply (*cases size ws*)
apply *simp-all*
done

lemmas *word-rsplit-rcat-size* = *refl [THEN word-rsplit-rcat-size']*

13.3 Rotation

lemmas *rotater-0'* [*simp*] = *rotater-def* [**where** *n* = 0, *simplified*]

lemmas *word-rot-defs* = *word-roti-def word-rotr-def word-rotl-def*

lemma *rotate-eq-mod*:
 $m \text{ mod } \text{length } xs = n \text{ mod } \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
apply (*rule box-equals*)
defer
apply (*rule rotate-conv-mod [symmetric]*)+
apply *simp*
done

```

lemmas rotate-eqs [standard] =
  trans [OF rotate0 [THEN fun-cong] id-apply]
  rotate-rotate [symmetric]
  rotate-id
  rotate-conv-mod
  rotate-eq-mod

```

13.3.1 Rotation of list to right

```

lemma rotate1-rl': rotater1 (l @ [a]) = a # l
  unfolding rotater1-def by (cases l) auto

```

```

lemma rotate1-rl [simp] : rotater1 (rotate1 l) = l
  apply (unfold rotater1-def)
  apply (cases l)
  apply (case-tac [2] list)
  apply auto
  done

```

```

lemma rotate1-lr [simp] : rotate1 (rotater1 l) = l
  unfolding rotater1-def by (cases l) auto

```

```

lemma rotater1-rev': rotater1 (rev xs) = rev (rotate1 xs)
  apply (cases xs)
  apply (simp add : rotater1-def)
  apply (simp add : rotate1-rl')
  done

```

```

lemma rotater-rev': rotater n (rev xs) = rev (rotate n xs)
  unfolding rotater-def by (induct n) (auto intro: rotater1-rev')

```

```

lemmas rotater-rev = rotater-rev' [where xs = rev ys, simplified, standard]

```

```

lemma rotater-drop-take:
  rotater n xs =
    drop (length xs - n mod length xs) xs @
    take (length xs - n mod length xs) xs
  by (clarsimp simp add : rotater-rev rotate-drop-take rev-take rev-drop)

```

```

lemma rotater-Suc [simp] :
  rotater (Suc n) xs = rotater1 (rotater n xs)
  unfolding rotater-def by auto

```

```

lemma rotate-inv-plus [rule-format] :
  ALL k. k = m + n --> rotater k (rotate n xs) = rotater m xs &
  rotate k (rotater n xs) = rotate m xs &
  rotater n (rotate k xs) = rotate m xs &
  rotate n (rotater k xs) = rotater m xs
  unfolding rotater-def rotate-def

```

```

by (induct n) (auto intro: funpow-swap1 [THEN trans])

lemmas rotate-inv-rel = le-add-diff-inverse2 [symmetric, THEN rotate-inv-plus]

lemmas rotate-inv-eq = order-refl [THEN rotate-inv-rel, simplified]

lemmas rotate-lr [simp] = rotate-inv-eq [THEN conjunct1, standard]
lemmas rotate-rl [simp] =
  rotate-inv-eq [THEN conjunct2, THEN conjunct1, standard]

lemma rotate-gal: (rotater n xs = ys) = (rotate n ys = xs)
by auto

lemma rotate-gal': (ys = rotater n xs) = (xs = rotate n ys)
by auto

lemma length-rotater [simp]:
  length (rotater n xs) = length xs
by (simp add : rotater-rev)

lemmas rrs0 = rotate-eqs [THEN restrict-to-left,
  simplified rotate-gal [symmetric] rotate-gal' [symmetric], standard]
lemmas rrs1 = rrs0 [THEN refl [THEN rev-iffD1]]
lemmas rotater-eqs = rrs1 [simplified length-rotater, standard]
lemmas rotater-0 = rotater-eqs (1)
lemmas rotater-add = rotater-eqs (2)

```

13.3.2 map, map2, commuting with rotate(r)

```

lemma last-map: xs ~[] ==> last (map f xs) = f (last xs)
by (induct xs) auto

lemma butlast-map:
  xs ~[] ==> butlast (map f xs) = map f (butlast xs)
by (induct xs) auto

lemma rotater1-map: rotater1 (map f xs) = map f (rotater1 xs)
unfolding rotater1-def
by (cases xs) (auto simp add: last-map butlast-map)

lemma rotater-map:
  rotater n (map f xs) = map f (rotater n xs)
unfolding rotater-def
by (induct n) (auto simp add : rotater1-map)

lemma but-last-zip [rule-format] :
  ALL ys. length xs = length ys --> xs ~[] -->
    last (zip xs ys) = (last xs, last ys) &
    butlast (zip xs ys) = zip (butlast xs) (butlast ys)

```



```

apply (induct xs)
apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

lemma but-last-map2 [rule-format] :
  ALL ys. length xs = length ys ==> xs ~ = [] ==>
    last (map2 f xs ys) = f (last xs) (last ys) &
    butlast (map2 f xs ys) = map2 f (butlast xs) (butlast ys)
apply (induct xs)
apply auto
  apply (unfold map2-def)
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

lemma rotater1-zip:
  length xs = length ys ==>
    rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)
apply (unfold rotater1-def)
apply (cases xs)
apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv but-last-zip)[1])+
done

lemma rotater1-map2:
  length xs = length ys ==>
    rotater1 (map2 f xs ys) = map2 f (rotater1 xs) (rotater1 ys)
unfolding map2-def by (simp add: rotater1-map rotater1-zip)

lemmas lrth =
  box-equals [OF asm-rl length-rotater [symmetric]
    length-rotater [symmetric],
    THEN rotater1-map2]

lemma rotater-map2:
  length xs = length ys ==>
    rotater n (map2 f xs ys) = map2 f (rotater n xs) (rotater n ys)
by (induct n) (auto intro!: lrth)

lemma rotate1-map2:
  length xs = length ys ==>
    rotate1 (map2 f xs ys) = map2 f (rotate1 xs) (rotate1 ys)
apply (unfold map2-def)
apply (cases xs)
  apply (cases ys, auto simp add : rotate1-def)+
done

lemmas lth = box-equals [OF asm-rl length-rotate [symmetric]
  length-rotate [symmetric], THEN rotate1-map2]

```

lemma *rotate-map2*:

length xs = length ys ==>
rotate n (map2 f xs ys) = map2 f (rotate n xs) (rotate n ys)
by (*induct n*) (*auto intro! lth*)

— corresponding equalities for word rotation

lemma *to-bl-rotl*:

to-bl (word-rotl n w) = rotate n (to-bl w)
by (*simp add: word-bl.Abs-inverse' word-rotl-def*)

lemmas *blrs0 = rotate-egs [THEN to-bl-rotl [THEN trans]]*

lemmas *word-rotl-egs =*

blrs0 [simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotl [symmetric]]

lemma *to-bl-rotr*:

to-bl (word-rotr n w) = rotater n (to-bl w)
by (*simp add: word-bl.Abs-inverse' word-rotr-def*)

lemmas *brrs0 = rotater-egs [THEN to-bl-rotr [THEN trans]]*

lemmas *word-rotr-egs =*

brrs0 [simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotr [symmetric]]

declare *word-rotr-egs (1) [simp]*

declare *word-rotl-egs (1) [simp]*

lemma

word-rot-rl [simp]:
word-rotl k (word-rotr k v) = v and
word-rot-lr [simp]:
word-rotr k (word-rotl k v) = v
by (*auto simp add: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric]*)

lemma

word-rot-gal:
(word-rotr n v = w) = (word-rotl n w = v) and
word-rot-gal':
(w = word-rotr n v) = (v = word-rotl n w)
by (*auto simp: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric]*
dest: sym)

lemma *word-rotr-rev:*

word-rotr n w = word-reverse (word-rotl n (word-reverse w))
by (*simp add: word-bl.Rep-inject [symmetric] to-bl-word-rev*
to-bl-rotr to-bl-rotl rotater-rev)

lemma *word-roti-0* [*simp*]: *word-roti 0 w = w*
by (*unfold word-rot-defs*) *auto*

lemmas *abl-cong = arg-cong* [**where** *f = of-bl*]

lemma *word-roti-add*:
word-roti (m + n) w = word-roti m (word-roti n w)

proof –

have *rotater-eq-lem*:

$\bigwedge m\ n\ xs.\ m = n \implies \text{rotater } m\ xs = \text{rotater } n\ xs$
by *auto*

have *rotate-eq-lem*:

$\bigwedge m\ n\ xs.\ m = n \implies \text{rotate } m\ xs = \text{rotate } n\ xs$
by *auto*

note *rpts* [*symmetric, standard*] =
rotate-inv-plus [*THEN conjunct1*]
rotate-inv-plus [*THEN conjunct2, THEN conjunct1*]
rotate-inv-plus [*THEN conjunct2, THEN conjunct2, THEN conjunct1*]
rotate-inv-plus [*THEN conjunct2, THEN conjunct2, THEN conjunct2*]

note *rrp = trans* [*symmetric, OF rotate-rotate rotate-eq-lem*]

note *rrrp = trans* [*symmetric, OF rotater-add* [*symmetric*] *rotater-eq-lem*]

show *?thesis*

apply (*unfold word-rot-defs*)

apply (*simp only: split: split-if*)

apply (*safe intro!: abl-cong*)

apply (*simp-all only: to-bl-rotl* [*THEN word-bl.Rep-inverse*]
to-bl-rotl
to-bl-rotr [*THEN word-bl.Rep-inverse*]
to-bl-rotr)

apply (*rule rrp rrrp rpts,*
simp add: nat-add-distrib [*symmetric*]
nat-diff-distrib [*symmetric*])+

done

qed

lemma *word-roti-conv-mod'*: *word-roti n w = word-roti (n mod int (size w)) w*

apply (*unfold word-rot-defs*)

apply (*cut-tac y=size w in gt-or-eq-0*)

apply (*erule disjE*)

apply *simp-all*

apply (*safe intro!: abl-cong*)

apply (*rule rotater-eqs*)

apply (*simp add: word-size nat-mod-distrib*)

apply (*simp add: rotater-add* [*symmetric*] *rotate-gal* [*symmetric*])

```

apply (rule rotater-eqs)
apply (simp add: word-size nat-mod-distrib)
apply (rule int-eq-0-conv [THEN iffD1])
apply (simp only: zmod-int zadd-int [symmetric])
apply (simp add: rdmodes)
done

```

lemmas *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded word-size]

13.3.3 Word rotation commutes with bit-wise operations

locale *word-rotate*

context *word-rotate*
begin

lemmas *word-rot-defs'* = *to-bl-rotl to-bl-rotr*

lemmas *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

lemmas *ths-map2* [OF *lbl-lbl*] = *rotate-map2 rotater-map2*

lemmas *ths-map* [where *xs* = *to-bl v*, *standard*] = *rotate-map rotater-map*

lemmas *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-map2 ths-map*

lemma *word-rot-logs*:

```

word-rotl n (NOT v) = NOT word-rotl n v
word-rotr n (NOT v) = NOT word-rotr n v
word-rotl n (x AND y) = word-rotl n x AND word-rotl n y
word-rotr n (x AND y) = word-rotr n x AND word-rotr n y
word-rotl n (x OR y) = word-rotl n x OR word-rotl n y
word-rotr n (x OR y) = word-rotr n x OR word-rotr n y
word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y
word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y
by (rule word-bl.Rep-eqD,
      rule word-rot-defs' [THEN trans],
      simp only: blwl-syms [symmetric],
      rule th1s [THEN trans],
      rule refl) +

```

end

lemmas *word-rot-logs* = *word-rotate.word-rot-logs*

lemmas *bl-word-rotl-dt* = *trans* [OF *to-bl-rotl rotate-drop-take*,
simplified word-bl.Rep', *standard*]

lemmas *bl-word-rotr-dt* = *trans* [*OF to-bl-rotr rotater-drop-take*,
simplified word-bl.Rep', *standard*]

lemma *bl-word-roti-dt'*:
 $n = \text{nat } ((- i) \bmod \text{int } (\text{size } (w :: 'a :: \text{len word}))) \implies$
 $\text{to-bl } (\text{word-roti } i \ w) = \text{drop } n \ (\text{to-bl } w) \ @ \ \text{take } n \ (\text{to-bl } w)$
apply (*unfold word-roti-def*)
apply (*simp add: bl-word-rotl-dt bl-word-rotr-dt word-size*)
apply *safe*
apply (*simp add: zmod-zminus1-eq-if*)
apply *safe*
apply (*simp add: nat-mult-distrib*)
apply (*simp add: nat-diff-distrib* [*OF pos-mod-sign pos-mod-conj*
 $[\text{THEN } \text{conjunct2}, \text{THEN } \text{order-less-imp-le}]$
 nat-mod-distrib])
apply (*simp add: nat-mod-distrib*)
done

lemmas *bl-word-roti-dt* = *bl-word-roti-dt'* [*unfolded word-size*]

lemmas *word-rotl-dt* = *bl-word-rotl-dt*
 $[\text{THEN } \text{word-bl.Rep-inverse' } [\text{symmetric}], \text{ standard}]$

lemmas *word-rotr-dt* = *bl-word-rotr-dt*
 $[\text{THEN } \text{word-bl.Rep-inverse' } [\text{symmetric}], \text{ standard}]$

lemmas *word-roti-dt* = *bl-word-roti-dt*
 $[\text{THEN } \text{word-bl.Rep-inverse' } [\text{symmetric}], \text{ standard}]$

lemma *word-rotx-0* [*simp*] : *word-rotr i 0 = 0 & word-rotl i 0 = 0*
by (*simp add : word-rotr-dt word-rotl-dt to-bl-0 replicate-add* [*symmetric*])

lemma *word-roti-0'* [*simp*] : *word-roti n 0 = 0*
unfolding *word-roti-def* **by** *auto*

lemmas *word-rotr-dt-no-bin'* [*simp*] =
word-rotr-dt [**where** *w=number-of w*, *unfolded to-bl-no-bin*, *standard*]

lemmas *word-rotl-dt-no-bin'* [*simp*] =
word-rotl-dt [**where** *w=number-of w*, *unfolded to-bl-no-bin*, *standard*]

declare *word-roti-def* [*simp*]

end

14 Boolean-Algebra: Boolean Algebras

theory *Boolean-Algebra*
imports *Main*

begin

locale *boolean* =
fixes *conj* :: 'a \Rightarrow 'a \Rightarrow 'a (**infixr** \sqcap 70)
fixes *disj* :: 'a \Rightarrow 'a \Rightarrow 'a (**infixr** \sqcup 65)
fixes *compl* :: 'a \Rightarrow 'a (\sim - [81] 80)
fixes *zero* :: 'a (**0**)
fixes *one* :: 'a (**1**)
assumes *conj-assoc*: $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$
assumes *disj-assoc*: $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$
assumes *conj-commute*: $x \sqcap y = y \sqcap x$
assumes *disj-commute*: $x \sqcup y = y \sqcup x$
assumes *conj-disj-distrib*: $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$
assumes *disj-conj-distrib*: $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$
assumes *conj-one-right* [*simp*]: $x \sqcap \mathbf{1} = x$
assumes *disj-zero-right* [*simp*]: $x \sqcup \mathbf{0} = x$
assumes *conj-cancel-right* [*simp*]: $x \sqcap \sim x = \mathbf{0}$
assumes *disj-cancel-right* [*simp*]: $x \sqcup \sim x = \mathbf{1}$

sublocale *boolean* < *conj*!: *abel-semigroup conj* **proof**
qed (*fact conj-assoc conj-commute*)+

sublocale *boolean* < *disj*!: *abel-semigroup disj* **proof**
qed (*fact disj-assoc disj-commute*)+

context *boolean*
begin

lemmas *conj-left-commute* = *conj.left-commute*

lemmas *disj-left-commute* = *disj.left-commute*

lemmas *conj-ac* = *conj.assoc conj.commute conj.left-commute*
lemmas *disj-ac* = *disj.assoc disj.commute disj.left-commute*

lemma *dual*: *boolean disj conj compl one zero*
apply (*rule boolean.intro*)
apply (*rule disj-assoc*)
apply (*rule conj-assoc*)
apply (*rule disj-commute*)
apply (*rule conj-commute*)
apply (*rule disj-conj-distrib*)
apply (*rule conj-disj-distrib*)
apply (*rule disj-zero-right*)
apply (*rule conj-one-right*)
apply (*rule disj-cancel-right*)
apply (*rule conj-cancel-right*)
done

14.1 Complement

lemma *complement-unique*:

assumes 1: $a \sqcap x = 0$

assumes 2: $a \sqcup x = 1$

assumes 3: $a \sqcap y = 0$

assumes 4: $a \sqcup y = 1$

shows $x = y$

proof –

have $(a \sqcap x) \sqcup (x \sqcap y) = (a \sqcap y) \sqcup (x \sqcap y)$ **using** 1 3 **by** *simp*

hence $(x \sqcap a) \sqcup (x \sqcap y) = (y \sqcap a) \sqcup (y \sqcap x)$ **using** *conj-commute* **by** *simp*

hence $x \sqcap (a \sqcup y) = y \sqcap (a \sqcup x)$ **using** *conj-disj-distrib* **by** *simp*

hence $x \sqcap 1 = y \sqcap 1$ **using** 2 4 **by** *simp*

thus $x = y$ **using** *conj-one-right* **by** *simp*

qed

lemma *compl-unique*: $\llbracket x \sqcap y = 0; x \sqcup y = 1 \rrbracket \implies \sim x = y$

by (*rule complement-unique* [*OF conj-cancel-right disj-cancel-right*])

lemma *double-compl* [*simp*]: $\sim(\sim x) = x$

proof (*rule compl-unique*)

from *conj-cancel-right* **show** $\sim x \sqcap x = 0$ **by** (*simp only: conj-commute*)

from *disj-cancel-right* **show** $\sim x \sqcup x = 1$ **by** (*simp only: disj-commute*)

qed

lemma *compl-eq-compl-iff* [*simp*]: $(\sim x = \sim y) = (x = y)$

by (*rule inj-eq* [*OF inj-on-inverseI*], *rule double-compl*)

14.2 Conjunction

lemma *conj-absorb* [*simp*]: $x \sqcap x = x$

proof –

have $x \sqcap x = (x \sqcap x) \sqcup 0$ **using** *disj-zero-right* **by** *simp*

also have $\dots = (x \sqcap x) \sqcup (x \sqcap \sim x)$ **using** *conj-cancel-right* **by** *simp*

also have $\dots = x \sqcap (x \sqcup \sim x)$ **using** *conj-disj-distrib* **by** (*simp only:*)

also have $\dots = x \sqcap 1$ **using** *disj-cancel-right* **by** *simp*

also have $\dots = x$ **using** *conj-one-right* **by** *simp*

finally show *?thesis* .

qed

lemma *conj-zero-right* [*simp*]: $x \sqcap 0 = 0$

proof –

have $x \sqcap 0 = x \sqcap (x \sqcap \sim x)$ **using** *conj-cancel-right* **by** *simp*

also have $\dots = (x \sqcap x) \sqcap \sim x$ **using** *conj-assoc* **by** (*simp only:*)

also have $\dots = x \sqcap \sim x$ **using** *conj-absorb* **by** *simp*

also have $\dots = 0$ **using** *conj-cancel-right* **by** *simp*

finally show *?thesis* .

qed

lemma *compl-one* [*simp*]: $\sim 1 = 0$

by (rule compl-unique [OF conj-zero-right disj-zero-right])

lemma conj-zero-left [simp]: $\mathbf{0} \sqcap x = \mathbf{0}$
by (subst conj-commute) (rule conj-zero-right)

lemma conj-one-left [simp]: $\mathbf{1} \sqcap x = x$
by (subst conj-commute) (rule conj-one-right)

lemma conj-cancel-left [simp]: $\sim x \sqcap x = \mathbf{0}$
by (subst conj-commute) (rule conj-cancel-right)

lemma conj-left-absorb [simp]: $x \sqcap (x \sqcap y) = x \sqcap y$
by (simp only: conj-assoc [symmetric] conj-absorb)

lemma conj-disj-distrib2:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
by (simp only: conj-commute conj-disj-distrib)

lemmas conj-disj-distribs =
 conj-disj-distrib conj-disj-distrib2

14.3 Disjunction

lemma disj-absorb [simp]: $x \sqcup x = x$
by (rule boolean.conj-absorb [OF dual])

lemma disj-one-right [simp]: $x \sqcup \mathbf{1} = \mathbf{1}$
by (rule boolean.conj-zero-right [OF dual])

lemma compl-zero [simp]: $\sim \mathbf{0} = \mathbf{1}$
by (rule boolean.compl-one [OF dual])

lemma disj-zero-left [simp]: $\mathbf{0} \sqcup x = x$
by (rule boolean.conj-one-left [OF dual])

lemma disj-one-left [simp]: $\mathbf{1} \sqcup x = \mathbf{1}$
by (rule boolean.conj-zero-left [OF dual])

lemma disj-cancel-left [simp]: $\sim x \sqcup x = \mathbf{1}$
by (rule boolean.conj-cancel-left [OF dual])

lemma disj-left-absorb [simp]: $x \sqcup (x \sqcup y) = x \sqcup y$
by (rule boolean.conj-left-absorb [OF dual])

lemma disj-conj-distrib2:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
by (rule boolean.conj-disj-distrib2 [OF dual])

lemmas disj-conj-distribs =

disj-conj-distrib disj-conj-distrib2

14.4 De Morgan’s Laws

lemma *de-Morgan-conj* [*simp*]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
proof (*rule compl-unique*)
have $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = ((x \sqcap y) \sqcap \sim x) \sqcup ((x \sqcap y) \sqcap \sim y)$
by (*rule conj-disj-distrib*)
also have $\dots = (y \sqcap (x \sqcap \sim x)) \sqcup (x \sqcap (y \sqcap \sim y))$
by (*simp only: conj-ac*)
finally show $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = \mathbf{0}$
by (*simp only: conj-cancel-right conj-zero-right disj-zero-right*)
next
have $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = (x \sqcup (\sim x \sqcup \sim y)) \sqcap (y \sqcup (\sim x \sqcup \sim y))$
by (*rule disj-conj-distrib2*)
also have $\dots = (\sim y \sqcup (x \sqcup \sim x)) \sqcap (\sim x \sqcup (y \sqcup \sim y))$
by (*simp only: disj-ac*)
finally show $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = \mathbf{1}$
by (*simp only: disj-cancel-right disj-one-right conj-one-right*)
qed

lemma *de-Morgan-disj* [*simp*]: $\sim (x \sqcup y) = \sim x \sqcap \sim y$
by (*rule boolean.de-Morgan-conj [OF dual]*)

end

14.5 Symmetric Difference

locale *boolean-xor* = *boolean* +
fixes *xor* :: $'a \Rightarrow 'a \Rightarrow 'a$ (**infixr** \oplus 65)
assumes *xor-def*: $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$
sublocale *boolean-xor* < *xor!*: *abel-semigroup xor* **proof**
fix $x\ y\ z :: 'a$
let $?t = (x \sqcap y \sqcap z) \sqcup (x \sqcap \sim y \sqcap \sim z) \sqcup$
 $(\sim x \sqcap y \sqcap \sim z) \sqcup (\sim x \sqcap \sim y \sqcap z)$
have $?t \sqcup (z \sqcap x \sqcap \sim x) \sqcup (z \sqcap y \sqcap \sim y) =$
 $?t \sqcup (x \sqcap y \sqcap \sim y) \sqcup (x \sqcap z \sqcap \sim z)$
by (*simp only: conj-cancel-right conj-zero-right*)
thus $(x \oplus y) \oplus z = x \oplus (y \oplus z)$
apply (*simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl*)
apply (*simp only: conj-disj-distrib conj-ac disj-ac*)
done
show $x \oplus y = y \oplus x$
by (*simp only: xor-def conj-commute disj-commute*)
qed

context *boolean-xor*
begin

lemmas *xor-assoc* = *xor.assoc*

lemmas *xor-commute* = *xor.commute*

lemmas *xor-left-commute* = *xor.left-commute*

lemmas *xor-ac* = *xor.assoc xor.commute xor.left-commute*

lemma *xor-def2*:

$$x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$$

by (*simp only: xor-def conj-disj-distrib*
disj-ac conj-ac conj-cancel-right disj-zero-left)

lemma *xor-zero-right* [*simp*]: $x \oplus \mathbf{0} = x$

by (*simp only: xor-def compl-zero conj-one-right conj-zero-right disj-zero-right*)

lemma *xor-zero-left* [*simp*]: $\mathbf{0} \oplus x = x$

by (*subst xor-commute*) (*rule xor-zero-right*)

lemma *xor-one-right* [*simp*]: $x \oplus \mathbf{1} = \sim x$

by (*simp only: xor-def compl-one conj-zero-right conj-one-right disj-zero-left*)

lemma *xor-one-left* [*simp*]: $\mathbf{1} \oplus x = \sim x$

by (*subst xor-commute*) (*rule xor-one-right*)

lemma *xor-self* [*simp*]: $x \oplus x = \mathbf{0}$

by (*simp only: xor-def conj-cancel-right conj-cancel-left disj-zero-right*)

lemma *xor-left-self* [*simp*]: $x \oplus (x \oplus y) = y$

by (*simp only: xor-assoc [symmetric] xor-self xor-zero-left*)

lemma *xor-compl-left* [*simp*]: $\sim x \oplus y = \sim (x \oplus y)$

apply (*simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl*)

apply (*simp only: conj-disj-distrib*)

apply (*simp only: conj-cancel-right conj-cancel-left*)

apply (*simp only: disj-zero-left disj-zero-right*)

apply (*simp only: disj-ac conj-ac*)

done

lemma *xor-compl-right* [*simp*]: $x \oplus \sim y = \sim (x \oplus y)$

apply (*simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl*)

apply (*simp only: conj-disj-distrib*)

apply (*simp only: conj-cancel-right conj-cancel-left*)

apply (*simp only: disj-zero-left disj-zero-right*)

apply (*simp only: disj-ac conj-ac*)

done

lemma *xor-cancel-right*: $x \oplus \sim x = \mathbf{1}$

by (*simp only: xor-compl-right xor-self compl-zero*)

lemma *xor-cancel-left*: $\sim x \oplus x = \mathbf{1}$

by (simp only: xor-compl-left xor-self compl-zero)

lemma conj-xor-distrib: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

proof –

have $(x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z) =$
 $(y \sqcap x \sqcap \sim x) \sqcup (z \sqcap x \sqcap \sim x) \sqcup (x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z)$

by (simp only: conj-cancel-right conj-zero-right disj-zero-left)

thus $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

by (simp (no-asm-use) only:

xor-def de-Morgan-disj de-Morgan-conj double-compl

conj-disj-distrib conj-ac disj-ac)

qed

lemma conj-xor-distrib2:

$(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$

proof –

have $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

by (rule conj-xor-distrib)

thus $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$

by (simp only: conj-commute)

qed

lemmas conj-xor-distribs =

conj-xor-distrib conj-xor-distrib2

end

end

15 WordGenLib: Miscellaneous Library for Words

theory WordGenLib

imports WordShift Boolean-Algebra

begin

declare of-nat-2p [simp]

lemma word-int-cases:

$\llbracket \bigwedge n. \llbracket (x :: 'a::len0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$

$\implies P$

by (cases x rule: word-uint.Abs-cases) (simp add: uints-num)

lemma word-nat-cases [cases type: word]:

$\llbracket \bigwedge n. \llbracket (x :: 'a::len \text{ word}) = \text{of-nat } n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$

$\implies P$

by (cases x rule: word-unat.Abs-cases) (simp add: unats-def)

```

lemma max-word-eq:
  (max-word::'a::len word) = 2 ^ len-of TYPE('a) - 1
  by (simp add: max-word-def word-of-int-hom-syms word-of-int-2p)

lemma max-word-max [simp,intro!]:
  n ≤ max-word
  by (cases n rule: word-int-cases)
      (simp add: max-word-def word-le-def int-word-uint int-mod-eq')

lemma word-of-int-2p-len:
  word-of-int (2 ^ len-of TYPE('a)) = (0::'a::len0 word)
  by (subst word-uint.Abs-norm [symmetric])
      (simp add: word-of-int-hom-syms)

lemma word-pow-0:
  (2::'a::len word) ^ len-of TYPE('a) = 0
proof -
  have word-of-int (2 ^ len-of TYPE('a)) = (0::'a word)
    by (rule word-of-int-2p-len)
  thus ?thesis by (simp add: word-of-int-2p)
qed

lemma max-word-wrap: x + 1 = 0 ⇒ x = max-word
  apply (simp add: max-word-eq)
  apply uint-arith
  apply auto
  apply (simp add: word-pow-0)
  done

lemma max-word-minus:
  max-word = (-1::'a::len word)
proof -
  have -1 + 1 = (0::'a word) by simp
  thus ?thesis by (rule max-word-wrap [symmetric])
qed

lemma max-word-bl [simp]:
  to-bl (max-word::'a::len word) = replicate (len-of TYPE('a)) True
  by (subst max-word-minus to-bl-n1) + simp

lemma max-test-bit [simp]:
  (max-word::'a::len word) !! n = (n < len-of TYPE('a))
  by (auto simp add: test-bit-bl word-size)

lemma word-and-max [simp]:
  x AND max-word = x
  by (rule word-eqI) (simp add: word-ops-nth-size word-size)

lemma word-or-max [simp]:

```

$x \text{ OR } \text{max-word} = \text{max-word}$
by (rule word-eqI) (simp add: word-ops-nth-size word-size)

lemma word-ao-dist2:
 $x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z::'a::\text{len0 word})$
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-oa-dist2:
 $x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::\text{len0 word}))$
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-and-not [simp]:
 $x \text{ AND NOT } x = (0::'a::\text{len0 word})$
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-or-not [simp]:
 $x \text{ OR NOT } x = \text{max-word}$
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-boolean:
 boolean (op AND) (op OR) bitNOT 0 max-word
apply (rule boolean.intro)
 apply (rule word-bw-assocs)
 apply (rule word-bw-assocs)
 apply (rule word-bw-comms)
 apply (rule word-bw-comms)
 apply (rule word-ao-dist2)
 apply (rule word-oa-dist2)
 apply (rule word-and-max)
 apply (rule word-log-esimps)
 apply (rule word-and-not)
 apply (rule word-or-not)
done

interpretation word-bool-alg:
 boolean op AND op OR bitNOT 0 max-word
by (rule word-boolean)

lemma word-xor-and-or:
 $x \text{ XOR } y = x \text{ AND NOT } y \text{ OR NOT } x \text{ AND } (y::'a::\text{len0 word})$
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

interpretation word-bool-alg:
 boolean-xor op AND op OR bitNOT 0 max-word op XOR
apply (rule boolean-xor.intro)
 apply (rule word-boolean)
 apply (rule boolean-xor-axioms.intro)
 apply (rule word-xor-and-or)
done

```

lemma shiftr-0 [iff]:
  (x::'a::len0 word) >> 0 = x
  by (simp add: shiftr-bl)

lemma shiftr-0 [simp]:
  (x :: 'a :: len word) << 0 = x
  by (simp add: shiftr-t2n)

lemma shiftr-1 [simp]:
  (1::'a::len word) << n = 2^n
  by (simp add: shiftr-t2n)

lemma uint-lt-0 [simp]:
  uint x < 0 = False
  by (simp add: linorder-not-less)

lemma shiftr1-1 [simp]:
  shiftr1 (1::'a::len word) = 0
  by (simp add: shiftr1-def word-0-alt)

lemma shiftr-1 [simp]:
  (1::'a::len word) >> n = (if n = 0 then 1 else 0)
  by (induct n) (auto simp: shiftr-def)

lemma word-less-1 [simp]:
  ((x::'a::len word) < 1) = (x = 0)
  by (simp add: word-less-nat-alt unat-0-iff)

lemma to-bl-mask:
  to-bl (mask n :: 'a::len word) =
    replicate (len-of TYPE('a) - n) False @
    replicate (min (len-of TYPE('a)) n) True
  by (simp add: mask-bl word-rep-drop min-def)

lemma map-replicate-True:
  n = length xs ==>
    map (λ(x,y). x & y) (zip xs (replicate n True)) = xs
  by (induct xs arbitrary: n) auto

lemma map-replicate-False:
  n = length xs ==> map (λ(x,y). x & y)
    (zip xs (replicate n False)) = replicate n False
  by (induct xs arbitrary: n) auto

lemma bl-and-mask:
  fixes w :: 'a::len word
  fixes n
  defines n' ≡ len-of TYPE('a) - n

```

```

shows to-bl (w AND mask n) = replicate n' False @ drop n' (to-bl w)
proof -
  note [simp] = map-replicate-True map-replicate-False
  have to-bl (w AND mask n) =
    map2 op & (to-bl w) (to-bl (mask n::'a::len word))
    by (simp add: bl-word-and)
  also
  have to-bl w = take n' (to-bl w) @ drop n' (to-bl w) by simp
  also
  have map2 op & ... (to-bl (mask n::'a::len word)) =
    replicate n' False @ drop n' (to-bl w)
    unfolding to-bl-mask n'-def map2-def
    by (subst zip-append) auto
  finally
  show ?thesis .
qed

```

```

lemma drop-rev-takefill:
  length xs ≤ n ==>
    drop (n - length xs) (rev (takefill False n (rev xs))) = xs
  by (simp add: takefill-alt rev-take)

```

```

lemma map-nth-0 [simp]:
  map (op !! (0::'a::len0 word)) xs = replicate (length xs) False
  by (induct xs) auto

```

```

lemma uint-plus-if-size:
  uint (x + y) =
    (if uint x + uint y < 2^size x then
      uint x + uint y
    else
      uint x + uint y - 2^size x)
  by (simp add: word-arith-alts int-word-uint mod-add-if-z
    word-size)

```

```

lemma unat-plus-if-size:
  unat (x + (y::'a::len word)) =
    (if unat x + unat y < 2^size x then
      unat x + unat y
    else
      unat x + unat y - 2^size x)
  apply (subst word-arith-nat-defs)
  apply (subst unat-of-nat)
  apply (simp add: mod-nat-add word-size)
  done

```

```

lemma word-neq-0-conv [simp]:
  fixes w :: 'a :: len word
  shows (w ≠ 0) = (0 < w)

```

proof –

have $0 \leq w$ **by** (*rule word-zero-le*)
 thus *?thesis* **by** (*auto simp add: word-less-def*)
qed

lemma *max-lt*:

$\text{unat } (\max a b \text{ div } c) = \text{unat } (\max a b) \text{ div } \text{unat } (c :: 'a :: \text{len word})$
 apply (*subst word-arith-nat-defs*)
 apply (*subst word-unat.eq-norm*)
 apply (*subst mod-if*)
 apply *clarsimp*
 apply (*erule notE*)
 apply (*insert div-le-dividend [of unat (max a b) unat c]*)
 apply (*erule order-le-less-trans*)
 apply (*insert unat-lt2p [of max a b]*)
 apply *simp*
done

lemma *uint-sub-if-size*:

$\text{uint } (x - y) =$
 (*if* $\text{uint } y \leq \text{uint } x$ *then*
 $\text{uint } x - \text{uint } y$
 else
 $\text{uint } x - \text{uint } y + 2^{\text{size } x}$)
 by (*simp add: word-arith-alts int-word-uint mod-sub-if-z word-size*)

lemma *unat-sub-simple*:

$x \leq y \implies \text{unat } (y - x) = \text{unat } y - \text{unat } x$
 by (*simp add: unat-def uint-sub-if-size word-le-def nat-diff-distrib*)

lemmas *unat-sub = unat-sub-simple*

lemma *word-less-sub1*:

fixes $x :: 'a :: \text{len word}$
 shows $x \neq 0 \implies 1 < x = (0 < x - 1)$
 by (*simp add: unat-sub-if-size word-less-nat-alt*)

lemma *word-le-sub1*:

fixes $x :: 'a :: \text{len word}$
 shows $x \neq 0 \implies 1 \leq x = (0 \leq x - 1)$
 by (*simp add: unat-sub-if-size order-le-less word-less-nat-alt*)

lemmas *word-less-sub1-numberof [simp] =*
 word-less-sub1 [of number-of w, standard]

lemmas *word-le-sub1-numberof [simp] =*
 word-le-sub1 [of number-of w, standard]

lemma *word-of-int-minus*:


```

word-of-int (2^len-of TYPE('a) - i) = (word-of-int (-i)::'a::len word)
proof -
  have x: 2^len-of TYPE('a) - i = -i + 2^len-of TYPE('a) by simp
  show ?thesis
    apply (subst x)
    apply (subst word-uint.Abs-norm [symmetric], subst mod-add-self2)
    apply simp
    done
qed

```

```

lemmas word-of-int-inj =
  word-uint.Abs-inject [unfolded uints-num, simplified]

```

```

lemma word-le-less-eq:
  (x :: 'z::len word) ≤ y = (x = y ∨ x < y)
  by (auto simp add: word-less-def)

```

```

lemma mod-plus-cong:
  assumes 1: (b::int) = b'
    and 2: x mod b' = x' mod b'
    and 3: y mod b' = y' mod b'
    and 4: x' + y' = z'
  shows (x + y) mod b = z' mod b'
proof -
  from 1 2[symmetric] 3[symmetric] have (x + y) mod b = (x' mod b' + y' mod
b') mod b'
    by (simp add: mod-add-eq[symmetric])
  also have ... = (x' + y') mod b'
    by (simp add: mod-add-eq[symmetric])
  finally show ?thesis by (simp add: 4)
qed

```

```

lemma mod-minus-cong:
  assumes 1: (b::int) = b'
    and 2: x mod b' = x' mod b'
    and 3: y mod b' = y' mod b'
    and 4: x' - y' = z'
  shows (x - y) mod b = z' mod b'
  using assms
  apply (subst zmod-zsub-left-eq)
  apply (subst zmod-zsub-right-eq)
  apply (simp add: zmod-zsub-left-eq [symmetric] zmod-zsub-right-eq [symmetric])
  done

```

```

lemma word-induct-less:
  [P (0::'a::len word); ∧n. [n < m; P n] ⇒ P (1 + n)] ⇒ P m
  apply (cases m)
  apply atomize
  apply (erule rev-mp)+

```

```

apply (rule-tac x=m in spec)
apply (induct-tac n)
  apply simp
apply clarsimp
apply (erule impE)
  apply clarsimp
apply (erule-tac x=n in allE)
apply (erule impE)
  apply (simp add: unat-arith-simps)
  apply (clarsimp simp: unat-of-nat)
apply simp
apply (erule-tac x=of-nat na in allE)
apply (erule impE)
  apply (simp add: unat-arith-simps)
  apply (clarsimp simp: unat-of-nat)
apply simp
done

```

lemma word-induct:

```


$$\llbracket P (0::'a::\text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$$

by (erule word-induct-less, simp)

```

lemma word-induct2 [induct type]:

```


$$\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n::'b::\text{len word})$$

apply (rule word-induct, simp)
apply (case-tac 1+n = 0, auto)
done

```

definition word-rec :: 'a \Rightarrow ('b::len word \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b word \Rightarrow 'a **where**
 word-rec forZero forSuc n \equiv nat-rec forZero (forSuc \circ of-nat) (unat n)

lemma word-rec-0: word-rec z s 0 = z

```

by (simp add: word-rec-def)

```

lemma word-rec-Suc:

```


$$1 + n \neq (0::'a::\text{len word}) \implies \text{word-rec } z s (1 + n) = s n (\text{word-rec } z s n)$$

apply (simp add: word-rec-def unat-word-ariths)
apply (subst nat-mod-eq')
apply (cut-tac x=n in unat-lt2p)
apply (drule Suc-mono)
apply (simp add: less-Suc-eq-le)
apply (simp only: order-less-le, simp)
apply (erule contrapos-pn, simp)
apply (drule arg-cong[where f=of-nat])
apply simp
apply (subst (asm) word-unat.Rep-inverse[of n])
apply simp
apply simp
done

```

lemma *word-rec-Pred*:

```

 $n \neq 0 \implies \text{word-rec } z \ s \ n = s \ (n - 1) \ (\text{word-rec } z \ s \ (n - 1))$ 
apply (rule subst[where  $t=n$  and  $s=1 + (n - 1)$ ])
apply simp
apply (subst word-rec-Suc)
apply simp
apply simp
done

```

lemma *word-rec-in*:

```

 $f \ n \ (\text{word-rec } z \ (\lambda \cdot. f) \ n) = \text{word-rec } (f \ z) \ (\lambda \cdot. f) \ n$ 
by (induct n) (simp-all add: word-rec-0 word-rec-Suc)

```

lemma *word-rec-in2*:

```

 $f \ n \ (\text{word-rec } z \ f \ n) = \text{word-rec } (f \ 0 \ z) \ (f \circ op + 1) \ n$ 
by (induct n) (simp-all add: word-rec-0 word-rec-Suc)

```

lemma *word-rec-twice*:

```

 $m \leq n \implies \text{word-rec } z \ f \ n = \text{word-rec } (\text{word-rec } z \ f \ (n - m)) \ (f \circ op + (n - m)) \ m$ 
apply (erule rev-mp)
apply (rule-tac  $x=z$  in spec)
apply (rule-tac  $x=f$  in spec)
apply (induct n)
apply (simp add: word-rec-0)
apply clarsimp
apply (rule-tac  $t=1 + n - m$  and  $s=1 + (n - m)$  in subst)
apply simp
apply (case-tac  $1 + (n - m) = 0$ )
apply (simp add: word-rec-0)
apply (rule-tac  $f = \text{word-rec } ?a \ ?b$  in arg-cong)
apply (rule-tac  $t=m$  and  $s=m + (1 + (n - m))$  in subst)
apply simp
apply (simp (no-asm-use))
apply (simp add: word-rec-Suc word-rec-in2)
apply (erule impE)
apply uint-arith
apply (drule-tac  $x=x \circ op + 1$  in spec)
apply (drule-tac  $x=x \ 0 \ xa$  in spec)
apply simp
apply (rule-tac  $t=\lambda a. x \ (1 + (n - m + a))$  and  $s=\lambda a. x \ (1 + (n - m) + a)$  in subst)
apply (clarsimp simp add: expand-fun-eq)
apply (rule-tac  $t=(1 + (n - m + xb))$  and  $s=1 + (n - m) + xb$  in subst)
apply simp
apply (rule refl)
apply (rule refl)
done

```

```

lemma word-rec-id: word-rec z ( $\lambda$ -. id) n = z
  by (induct n) (auto simp add: word-rec-0 word-rec-Suc)

lemma word-rec-id-eq:  $\forall m < n. f\ m = id \implies word-rec\ z\ f\ n = z$ 
apply (erule rev-mp)
apply (induct n)
  apply (auto simp add: word-rec-0 word-rec-Suc)
  apply (drule spec, erule mp)
  apply uint-arith
apply (drule-tac x=n in spec, erule impE)
  apply uint-arith
apply simp
done

lemma word-rec-max:
   $\forall m \geq n. m \neq -1 \longrightarrow f\ m = id \implies word-rec\ z\ f\ -1 = word-rec\ z\ f\ n$ 
apply (subst word-rec-twice[where n=-1 and m=-1 - n])
  apply simp
apply simp
apply (rule word-rec-id-eq)
apply clarsimp
apply (drule spec, rule mp, erule mp)
  apply (rule word-plus-mono-right2[OF - order-less-imp-le])
  prefer 2
  apply assumption
  apply simp
apply (erule contrapos-pn)
apply simp
apply (drule arg-cong[where f= $\lambda x. x - n$ ])
apply simp
done

lemma unatSuc:
   $1 + n \neq (0::'a::len\ word) \implies unat\ (1 + n) = Suc\ (unat\ n)$ 
  by unat-arith

lemmas word-no-1 [simp] = word-1-no [symmetric, unfolded BIT-simps]
lemmas word-no-0 [simp] = word-0-no [symmetric]

declare word-0-bl [simp]
declare bin-to-bl-def [simp]
declare to-bl-0 [simp]
declare of-bl-True [simp]

end

```

16 Word: Word Library interface

```
theory Word
imports WordGenLib
uses ~~/src/HOL/Tools/SMT/smt-word.ML
begin

setup ⟨⟨ SMT-Word.setup ⟩⟩

see Examples/WordExamples.thy for examples
end
```

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.