

Machine Words in Isabelle/HOL

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Abstract

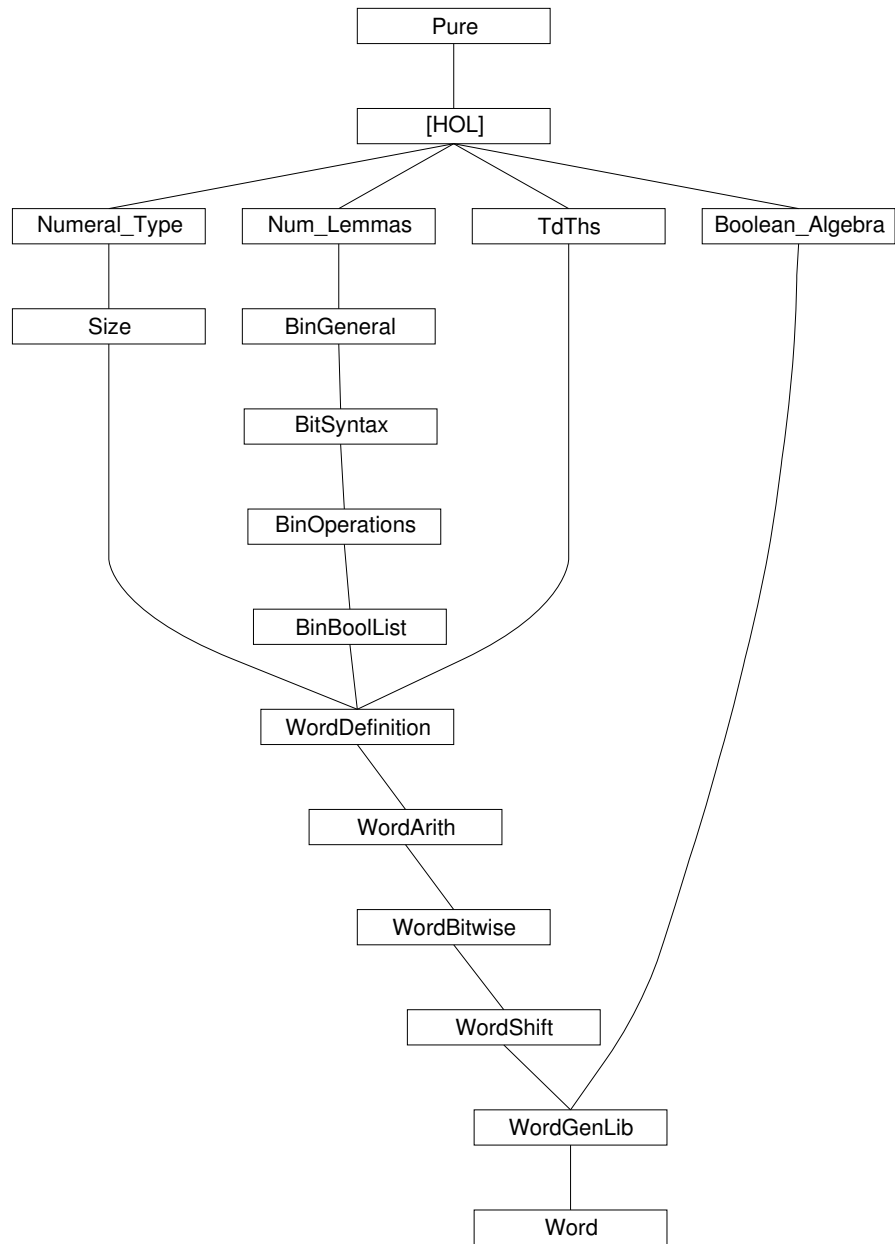
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
imports Main
begin
```

1.1 Preliminary lemmas

```
lemma (in type-definition) univ:
   $UNIV = Abs\ 'A$ 
 $\langle proof \rangle$ 
```

```
lemma (in type-definition) card:  $card\ (UNIV :: 'b\ set) = card\ A$ 
 $\langle proof \rangle$ 
```

1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations  $CARD('t) => CONST\ card\ (CONST\ UNIV :: 't\ set)$ 
```

$\langle ML \rangle$

```
lemma card-unit [simp]:  $CARD(unit) = 1$ 
 $\langle proof \rangle$ 
```

```
lemma card-bool [simp]:  $CARD(bool) = 2$ 
 $\langle proof \rangle$ 
```

```
lemma card-prod [simp]:  $CARD('a \times 'b) = CARD('a::finite) * CARD('b::finite)$ 
 $\langle proof \rangle$ 
```

```
lemma card-sum [simp]:  $CARD('a + 'b) = CARD('a::finite) + CARD('b::finite)$ 
 $\langle proof \rangle$ 
```

```
lemma card-option [simp]:  $CARD('a\ option) = Suc\ CARD('a::finite)$ 
 $\langle proof \rangle$ 
```

```
lemma card-set [simp]:  $CARD('a\ set) = 2 ^ CARD('a::finite)$ 
 $\langle proof \rangle$ 
```

```
lemma card-nat [simp]:  $CARD(nat) = 0$ 
 $\langle proof \rangle$ 
```

1.3 Classes with at least 1 and 2

Class *finite* already captures “at least 1”

```
lemma zero-less-card-finite [simp]:  $0 < CARD('a::finite)$ 
 $\langle proof \rangle$ 
```

lemma *one-le-card-finite* [simp]: $\text{Suc } 0 \leq \text{CARD}('a::\text{finite})$
 $\langle \text{proof} \rangle$

Class for cardinality ”at least 2”

class *card2* = *finite* +
assumes *two-le-card*: $2 \leq \text{CARD}('a)$

lemma *one-less-card*: $\text{Suc } 0 < \text{CARD}('a::\text{card2})$
 $\langle \text{proof} \rangle$

lemma *one-less-int-card*: $1 < \text{int } \text{CARD}('a::\text{card2})$
 $\langle \text{proof} \rangle$

1.4 Numeral Types

typedef (open) *num0* = *UNIV* :: *nat set* $\langle \text{proof} \rangle$
typedef (open) *num1* = *UNIV* :: *unit set* $\langle \text{proof} \rangle$

typedef (open) *'a bit0* = $\{0 \dots 2 * \text{int } \text{CARD}('a::\text{finite})\}$
 $\langle \text{proof} \rangle$

typedef (open) *'a bit1* = $\{0 \dots 1 + 2 * \text{int } \text{CARD}('a::\text{finite})\}$
 $\langle \text{proof} \rangle$

lemma *card-num0* [simp]: $\text{CARD}(\text{num0}) = 0$
 $\langle \text{proof} \rangle$

lemma *card-num1* [simp]: $\text{CARD}(\text{num1}) = 1$
 $\langle \text{proof} \rangle$

lemma *card-bit0* [simp]: $\text{CARD}('a \text{ bit0}) = 2 * \text{CARD}('a::\text{finite})$
 $\langle \text{proof} \rangle$

lemma *card-bit1* [simp]: $\text{CARD}('a \text{ bit1}) = \text{Suc } (2 * \text{CARD}('a::\text{finite}))$
 $\langle \text{proof} \rangle$

instance *num1* :: *finite*
 $\langle \text{proof} \rangle$

instance *bit0* :: (*finite*) *card2*
 $\langle \text{proof} \rangle$

instance *bit1* :: (*finite*) *card2*
 $\langle \text{proof} \rangle$

1.5 Locale for modular arithmetic subtypes

locale *mod-type* =

```

fixes  $n :: int$ 
and  $Rep :: 'a :: \{zero, one, plus, times, uminus, minus\} \Rightarrow int$ 
and  $Abs :: int \Rightarrow 'a :: \{zero, one, plus, times, uminus, minus\}$ 
assumes  $type: type-definition\ Rep\ Abs\ \{0..<n\}$ 
and  $size1: 1 < n$ 
and  $zero-def: 0 = Abs\ 0$ 
and  $one-def: 1 = Abs\ 1$ 
and  $add-def: x + y = Abs\ ((Rep\ x + Rep\ y)\ mod\ n)$ 
and  $mult-def: x * y = Abs\ ((Rep\ x * Rep\ y)\ mod\ n)$ 
and  $diff-def: x - y = Abs\ ((Rep\ x - Rep\ y)\ mod\ n)$ 
and  $minus-def: -x = Abs\ ((- Rep\ x)\ mod\ n)$ 
begin

```

```

lemma  $size0: 0 < n$ 
 $\langle proof \rangle$ 

```

```

lemmas  $definitions =$ 
 $zero-def\ one-def\ add-def\ mult-def\ minus-def\ diff-def$ 

```

```

lemma  $Rep-less-n: Rep\ x < n$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-le-n: Rep\ x \leq n$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-inject-sym: x = y \longleftrightarrow Rep\ x = Rep\ y$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-inverse: Abs\ (Rep\ x) = x$ 
 $\langle proof \rangle$ 

```

```

lemma  $Abs-inverse: m \in \{0..<n\} \Longrightarrow Rep\ (Abs\ m) = m$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-Abs-mod: Rep\ (Abs\ (m\ mod\ n)) = m\ mod\ n$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-Abs-0: Rep\ (Abs\ 0) = 0$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-0: Rep\ 0 = 0$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-Abs-1: Rep\ (Abs\ 1) = 1$ 
 $\langle proof \rangle$ 

```

```

lemma  $Rep-1: Rep\ 1 = 1$ 
 $\langle proof \rangle$ 

```

lemma *Rep-mod*: $\text{Rep } x \text{ mod } n = \text{Rep } x$
 $\langle \text{proof} \rangle$

lemmas *Rep-simps* =
Rep-inject-sym Rep-inverse Rep-Abs-mod Rep-mod Rep-Abs-0 Rep-Abs-1

lemma *comm-ring-1*: $\text{OFCLASS}('a, \text{comm-ring-1-class})$
 $\langle \text{proof} \rangle$

end

locale *mod-ring* = *mod-type* +
constrains $n :: \text{int}$
and $\text{Rep} :: 'a :: \{\text{number-ring}\} \Rightarrow \text{int}$
and $\text{Abs} :: \text{int} \Rightarrow 'a :: \{\text{number-ring}\}$
begin

lemma *of-nat-eq*: $\text{of-nat } k = \text{Abs } (\text{int } k \text{ mod } n)$
 $\langle \text{proof} \rangle$

lemma *of-int-eq*: $\text{of-int } z = \text{Abs } (z \text{ mod } n)$
 $\langle \text{proof} \rangle$

lemma *Rep-number-of*:
 $\text{Rep } (\text{number-of } w) = \text{number-of } w \text{ mod } n$
 $\langle \text{proof} \rangle$

lemma *iszero-number-of*:
 $\text{iszero } (\text{number-of } w :: 'a) \longleftrightarrow \text{number-of } w \text{ mod } n = 0$
 $\langle \text{proof} \rangle$

lemma *cases*:
assumes $1: \bigwedge z. \llbracket (x :: 'a) = \text{of-int } z; 0 \leq z; z < n \rrbracket \Longrightarrow P$
shows P
 $\langle \text{proof} \rangle$

lemma *induct*:
 $(\bigwedge z. \llbracket 0 \leq z; z < n \rrbracket \Longrightarrow P (\text{of-int } z)) \Longrightarrow P (x :: 'a)$
 $\langle \text{proof} \rangle$

end

1.6 Number ring instances

Unfortunately a number ring instance is not possible for *num1*, since 0 and 1 are not distinct.

instantiation *num1* :: $\{\text{comm-ring}, \text{comm-monoid-mult}, \text{number}\}$
begin


```

lemma num1-eq-iff:  $(x::\text{num1}) = (y::\text{num1}) \longleftrightarrow \text{True}$ 
   $\langle \text{proof} \rangle$ 

instance  $\langle \text{proof} \rangle$ 

end

instantiation
  bit0 and bit1 :: (finite) {zero,one,plus,times,uminus,minus}
begin

definition Abs-bit0' :: int  $\Rightarrow$  'a bit0 where
  Abs-bit0' x = Abs-bit0 (x mod int CARD('a bit0))

definition Abs-bit1' :: int  $\Rightarrow$  'a bit1 where
  Abs-bit1' x = Abs-bit1 (x mod int CARD('a bit1))

definition 0 = Abs-bit0 0
definition 1 = Abs-bit0 1
definition x + y = Abs-bit0' (Rep-bit0 x + Rep-bit0 y)
definition x * y = Abs-bit0' (Rep-bit0 x * Rep-bit0 y)
definition x - y = Abs-bit0' (Rep-bit0 x - Rep-bit0 y)
definition  $- x$  = Abs-bit0' ( $- \text{Rep-bit0 } x$ )

definition 0 = Abs-bit1 0
definition 1 = Abs-bit1 1
definition x + y = Abs-bit1' (Rep-bit1 x + Rep-bit1 y)
definition x * y = Abs-bit1' (Rep-bit1 x * Rep-bit1 y)
definition x - y = Abs-bit1' (Rep-bit1 x - Rep-bit1 y)
definition  $- x$  = Abs-bit1' ( $- \text{Rep-bit1 } x$ )

instance  $\langle \text{proof} \rangle$ 

end

interpretation bit0:
  mod-type int CARD('a::finite bit0)
  Rep-bit0 :: 'a::finite bit0  $\Rightarrow$  int
  Abs-bit0 :: int  $\Rightarrow$  'a::finite bit0
   $\langle \text{proof} \rangle$ 

interpretation bit1:
  mod-type int CARD('a::finite bit1)
  Rep-bit1 :: 'a::finite bit1  $\Rightarrow$  int
  Abs-bit1 :: int  $\Rightarrow$  'a::finite bit1
   $\langle \text{proof} \rangle$ 

instance bit0 :: (finite) comm-ring-1
   $\langle \text{proof} \rangle$ 

```

```

instance bit1 :: (finite) comm-ring-1
  ⟨proof⟩

instantiation bit0 and bit1 :: (finite) number-ring
begin

definition (number-of w :: - bit0) = of-int w

definition (number-of w :: - bit1) = of-int w

instance ⟨proof⟩

end

interpretation bit0:
  mod-ring int CARD('a::finite bit0)
    Rep-bit0 :: 'a::finite bit0 ⇒ int
    Abs-bit0 :: int ⇒ 'a::finite bit0
  ⟨proof⟩

interpretation bit1:
  mod-ring int CARD('a::finite bit1)
    Rep-bit1 :: 'a::finite bit1 ⇒ int
    Abs-bit1 :: int ⇒ 'a::finite bit1
  ⟨proof⟩

Set up cases, induction, and arithmetic

lemmas bit0-cases [case-names of-int, cases type: bit0] = bit0.cases
lemmas bit1-cases [case-names of-int, cases type: bit1] = bit1.cases

lemmas bit0-induct [case-names of-int, induct type: bit0] = bit0.induct
lemmas bit1-induct [case-names of-int, induct type: bit1] = bit1.induct

lemmas bit0-iszero-number-of [simp] = bit0.iszero-number-of
lemmas bit1-iszero-number-of [simp] = bit1.iszero-number-of

```

1.7 Syntax

```

syntax
  -NumeralType :: num-const => type (-)
  -NumeralType0 :: type (0)
  -NumeralType1 :: type (1)

```

```

translations
  (type) 1 == (type) num1
  (type) 0 == (type) num0

```

⟨*ML*⟩

1.8 Examples

```

lemma  $CARD(0) = 0$   $\langle proof \rangle$ 
lemma  $CARD(17) = 17$   $\langle proof \rangle$ 
lemma  $8 * 11 \wedge 3 - 6 = (2::5)$   $\langle proof \rangle$ 

end

```

2 Size: The len classes

```

theory Size
imports Natural-Type
begin

```

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in *Natural-Type*.

```

class len0 =
  fixes len-of :: 'a itself  $\Rightarrow$  nat

```

Some theorems are only true on words with length greater 0.

```

class len = len0 +
  assumes len-gt-0 [iff]:  $0 < \text{len-of } TYPE \ ('a)$ 

```

```

instantiation num0 and num1 :: len0
begin

```

```

definition
  len-num0:  $\text{len-of } (x::\text{num0 } \textit{itself}) = 0$ 

```

```

definition
  len-num1:  $\text{len-of } (x::\text{num1 } \textit{itself}) = 1$ 

```

```

instance  $\langle proof \rangle$ 

```

```

end

```

```

instantiation bit0 and bit1 :: (len0) len0
begin

```

```

definition
  len-bit0:  $\text{len-of } (x::'a::\text{len0 } \textit{bit0 } \textit{itself}) = 2 * \text{len-of } TYPE \ ('a)$ 

```

```

definition
  len-bit1:  $\text{len-of } (x::'a::\text{len0 } \textit{bit1 } \textit{itself}) = 2 * \text{len-of } TYPE \ ('a) + 1$ 

```

```

instance  $\langle proof \rangle$ 

```

end

lemmas *len-of-numeral-defs* [*simp*] = *len-num0 len-num1 len-bit0 len-bit1*

instance *num1* :: *len* <*proof*>

instance *bit0* :: (*len*) *len* <*proof*>

instance *bit1* :: (*len0*) *len* <*proof*>

lemma *len-of TYPE(17) = 17* <*proof*>

lemma *len-of TYPE(0) = 0* <*proof*>

lemma *len-of TYPE('a::len0) = x*
<*proof*>

end

3 Num-Lemmas: Useful Numerical Lemmas

theory *Num-Lemmas*

imports *Main Parity*

begin

lemma *contentsI*: $y = \{x\} \implies \text{contents } y = x$
<*proof*>

lemmas *split-split* = *prod.split* [*unfolded prod-case-split*]

lemmas *split-split-asm* = *prod.split-asm* [*unfolded prod-case-split*]

lemmas *split-splits* = *split-split split-split-asm*

lemmas *funpow-0* = *funpow.simps(1)*

lemmas *funpow-Suc* = *funpow.simps(2)*

lemma *nonemptyE*: $S \sim \{\} \implies (!x. x : S \implies R) \implies R$ <*proof*>

lemma *gt-or-eq-0*: $0 < y \vee 0 = (y::nat)$ <*proof*>

declare *iszero-0* [*iff*]

lemmas *xtr1* = *xtrans(1)*

lemmas *xtr2* = *xtrans(2)*

lemmas *xtr3* = *xtrans(3)*

lemmas *xtr4* = *xtrans(4)*

lemmas *xtr5* = *xtrans(5)*

lemmas *xtr6* = *xtrans(6)*

lemmas *xtr7* = *xtrans(7)*

lemmas *xtr8* = *xtrans(8)*

lemmas *nat-simps* = *diff-add-inverse2 diff-add-inverse*

lemmas *nat-iffs* = *le-add1 le-add2*

lemma *sum-imp-diff*: $j = k + i \implies j - i = (k :: \text{nat})$ $\langle \text{proof} \rangle$

lemma *nobm1*:

$0 < (\text{number-of } w :: \text{nat}) \implies$
 $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Int.pred } w)$
 $\langle \text{proof} \rangle$

lemma *zless2*: $0 < (2 :: \text{int})$ $\langle \text{proof} \rangle$

lemmas *zless2p* [*simp*] = *zless2* [*THEN zero-less-power*]

lemmas *zle2p* [*simp*] = *zless2p* [*THEN order-less-imp-le*]

lemmas *pos-mod-sign2* = *zless2* [*THEN pos-mod-sign* [**where** $b = 2 :: \text{int}$]]

lemmas *pos-mod-bound2* = *zless2* [*THEN pos-mod-bound* [**where** $b = 2 :: \text{int}$]]

— the inverse(s) of *number-of*

lemma *nmod2*: $n \bmod (2 :: \text{int}) = 0 \mid n \bmod 2 = 1$ $\langle \text{proof} \rangle$

lemma *emep1*:

$\text{even } n \implies \text{even } d \implies 0 \leq d \implies (n + 1) \bmod (d :: \text{int}) = (n \bmod d) + 1$
 $\langle \text{proof} \rangle$

lemmas *eme1p* = *emep1* [*simplified add-commute*]

lemma *le-diff-eq'*: $(a \leq c - b) = (b + a \leq (c :: \text{int}))$ $\langle \text{proof} \rangle$

lemma *less-diff-eq'*: $(a < c - b) = (b + a < (c :: \text{int}))$ $\langle \text{proof} \rangle$

lemma *diff-le-eq'*: $(a - b \leq c) = (a \leq b + (c :: \text{int}))$ $\langle \text{proof} \rangle$

lemma *diff-less-eq'*: $(a - b < c) = (a < b + (c :: \text{int}))$ $\langle \text{proof} \rangle$

lemmas *m1mod2k* = *zless2p* [*THEN zmod-minus1*]

lemmas *m1mod22k* = *mult-pos-pos* [*OF zless2 zless2p, THEN zmod-minus1*]

lemmas *p1mod22k'* = *zless2p* [*THEN order-less-imp-le, THEN pos-zmod-mult-2*]

lemmas *z1pmod2'* = *zero-le-one* [*THEN pos-zmod-mult-2, simplified*]

lemmas *z1pdiv2'* = *zero-le-one* [*THEN pos-zdiv-mult-2, simplified*]

lemma *p1mod22k*:

$(2 * b + 1) \bmod (2 * 2 ^ n) = 2 * (b \bmod 2 ^ n) + (1 :: \text{int})$
 $\langle \text{proof} \rangle$

lemma *z1pmod2*:

$(2 * b + 1) \bmod 2 = (1 :: \text{int})$ $\langle \text{proof} \rangle$

lemma *z1pdiv2*:

$(2 * b + 1) \text{div } 2 = (b :: \text{int})$ $\langle \text{proof} \rangle$

lemmas *zdiv-le-dividend* = *xtr3* [*OF div-by-1* [*symmetric*] *zdiv-mono2*,
simplified int-one-le-iff-zero-less, *simplified*, *standard*]

lemma *axxbyy*:

$$a + m + m = b + n + n ==> (a = 0 \mid a = 1) ==> (b = 0 \mid b = 1) ==> \\ a = b \ \& \ m = (n :: \text{int}) \ \langle \text{proof} \rangle$$

lemma *axxmod2*:

$$(1 + x + x) \bmod 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \bmod 2 = (0 :: \text{int}) \ \langle \text{proof} \rangle$$

lemma *axxdiv2*:

$$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int}) \ \langle \text{proof} \rangle$$

lemmas *iszero-minus* = *trans* [*THEN trans*,

OF iszero-def neg-equal-0-iff-equal iszero-def [*symmetric*], *standard*]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel* [*symmetric*] *add-commute*,
standard]

lemmas *add-diff-cancel2* = *add-commute* [*THEN diff-eq-eq* [*THEN iffD2*], *standard*]

lemma *zmod-uminus*: $-(a :: \text{int}) \bmod b \bmod b = -a \bmod b$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-distrib*: $((a :: \text{int}) - b) \bmod c = (a \bmod c - b \bmod c) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-right-eq*: $((a :: \text{int}) - b) \bmod c = (a - b \bmod c) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-left-eq*: $((a :: \text{int}) - b) \bmod c = (a \bmod c - b) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-self* [*simp*]:

$$((b :: \text{int}) - a) \bmod a = b \bmod a \\ \langle \text{proof} \rangle$$

lemma *zmod-zmult1-eq-rev*:

$$b * a \bmod c = b \bmod c * a \bmod (c :: \text{int}) \\ \langle \text{proof} \rangle$$

lemmas *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]

zmod-zsub-left-eq *zmod-zsub-right-eq* *mod-add-left-eq*
mod-add-right-eq *zmod-zmult1-eq* *zmod-zmult1-eq-rev*

lemma *mod-plus-right*:

$$((a + x) \bmod m = (b + x) \bmod m) = (a \bmod m = b \bmod (m :: \text{nat}))$$

$\langle \text{proof} \rangle$

lemma *nat-minus-mod*: $(n - n \text{ mod } m) \text{ mod } m = (0 :: \text{nat})$
 $\langle \text{proof} \rangle$

lemmas *nat-minus-mod-plus-right* = *trans* [*OF nat-minus-mod mod-0* [*symmetric*],
THEN mod-plus-right [*THEN iffD2*], *standard*, *simplified*]

lemmas *push-mods'* = *mod-add-eq* [*standard*]
mod-mult-eq [*standard*] *zmod-zsub-distrib* [*standard*]
zmod-uminus [*symmetric*, *standard*]

lemmas *push-mods* = *push-mods'* [*THEN eq-reflection*, *standard*]

lemmas *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN eq-reflection*, *standard*]

lemmas *mod-simps* =
mod-mult-self2-is-0 [*THEN eq-reflection*]
mod-mult-self1-is-0 [*THEN eq-reflection*]
mod-mod-trivial [*THEN eq-reflection*]

lemma *nat-mod-eq*:
 $!!b. b < n ==> a \text{ mod } n = b \text{ mod } n ==> a \text{ mod } n = (b :: \text{nat})$
 $\langle \text{proof} \rangle$

lemmas *nat-mod-eq'* = *refl* [*THEN* [2] *nat-mod-eq*]

lemma *nat-mod-lem*:
 $(0 :: \text{nat}) < n ==> b < n = (b \text{ mod } n = b)$
 $\langle \text{proof} \rangle$

lemma *mod-nat-add*:
 $(x :: \text{nat}) < z ==> y < z ==>$
 $(x + y) \text{ mod } z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 $\langle \text{proof} \rangle$

lemma *mod-nat-sub*:
 $(x :: \text{nat}) < z ==> (x - y) \text{ mod } z = x - y$
 $\langle \text{proof} \rangle$

lemma *int-mod-lem*:
 $(0 :: \text{int}) < n ==> (0 \leq b \ \& \ b < n) = (b \text{ mod } n = b)$
 $\langle \text{proof} \rangle$

lemma *int-mod-eq*:
 $(0 :: \text{int}) \leq b ==> b < n ==> a \text{ mod } n = b \text{ mod } n ==> a \text{ mod } n = b$
 $\langle \text{proof} \rangle$

lemmas *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

lemma *int-mod-le*: $0 \leq a \implies 0 < (n :: \text{int}) \implies a \bmod n \leq a$
 ⟨proof⟩

lemma *int-mod-le'*: $0 \leq b - n \implies 0 < (n :: \text{int}) \implies b \bmod n \leq b - n$
 ⟨proof⟩

lemma *int-mod-ge*: $a < n \implies 0 < (n :: \text{int}) \implies a \leq a \bmod n$
 ⟨proof⟩

lemma *int-mod-ge'*: $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$
 ⟨proof⟩

lemma *mod-add-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 ⟨proof⟩

lemma *mod-sub-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
 ⟨proof⟩

lemmas *zmde* = *zmod-zdiv-equality* [THEN *diff-eq-eq* [THEN *iffD2*], *symmetric*]
lemmas *mcl* = *mult-cancel-left* [THEN *iffD1*, THEN *make-pos-rule*]

lemma *zdiv-mult-self*: $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$
 ⟨proof⟩

lemma *eqne*: $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$
 ⟨proof⟩

lemmas *Rep-Integ-ne* = *Integ.Rep-Integ*
 [THEN *equiv-intrel* [THEN *eqne*, *simplified Integ-def* [symmetric]], *standard*]

lemmas *riq* = *Integ.Rep-Integ* [*simplified Integ-def*]
lemmas *intrel-refl* = *refl* [THEN *equiv-intrel-iff* [THEN *iffD1*], *standard*]
lemmas *Rep-Integ-equiv* = *quotient-eq-iff*
 [OF *equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]
lemmas *Rep-Integ-same* =
Rep-Integ-equiv [THEN *intrel-refl* [THEN *rev-iffD2*], *standard*]

lemma *RI-int*: $(a, 0) : \text{Rep-Integ } (\text{int } a)$
 ⟨proof⟩

lemmas *RI-intrel* [*simp*] = *UNIV-I* [THEN *quotientI*,
 THEN *Integ.Abs-Integ-inverse* [*simplified Integ-def*], *standard*]

lemma *RI-minus*: $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (-x)$
 $\langle \text{proof} \rangle$

lemma *RI-add*:
 $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$
 $(a + c, b + d) : \text{Rep-Integ } (x + y)$
 $\langle \text{proof} \rangle$

lemma *mem-same*: $a : S \implies a = b \implies b : S$
 $\langle \text{proof} \rangle$

lemma *RI-eq-diff'*: $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$
 $\langle \text{proof} \rangle$

lemma *RI-eq-diff*: $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$
 $\langle \text{proof} \rangle$

lemma *mod-power-lem*:
 $a > 1 \implies a^n \bmod a^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int})^n)$
 $\langle \text{proof} \rangle$

lemma *min-pm* [simp]: $\min a b + (a - b) = (a :: \text{nat})$ $\langle \text{proof} \rangle$

lemmas *min-pm1* [simp] = trans [OF add-commute min-pm]

lemma *rev-min-pm* [simp]: $\min b a + (a - b) = (a :: \text{nat})$ $\langle \text{proof} \rangle$

lemmas *rev-min-pm1* [simp] = trans [OF add-commute rev-min-pm]

lemma *pl-pl-rels*:
 $a + b = c + d \implies$
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$ $\langle \text{proof} \rangle$

lemmas *pl-pl-rels'* = add-commute [THEN [2] trans, THEN pl-pl-rels]

lemma *minus-eq*: $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$ $\langle \text{proof} \rangle$

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$ $\langle \text{proof} \rangle$

lemmas *pl-pl-mm'* = add-commute [THEN [2] trans, THEN pl-pl-mm]

lemma *min-minus* [simp]: $\min m (m - k) = (m - k :: \text{nat})$ $\langle \text{proof} \rangle$

lemmas *min-minus'* [simp] = trans [OF min-max.inf-commute min-minus]

lemma *nat-no-eq-iff*:
 $(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$
 $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$

$\langle \text{proof} \rangle$

lemmas $dme = \text{box-equals}$ [*OF div-mod-equality add-0-right add-0-right*]

lemmas $dtle = \text{xtr3}$ [*OF dme [symmetric] le-add1*]

lemmas $th2 = \text{order-trans}$ [*OF order-refl [THEN [2] mult-le-mono] dtle*]

lemma $td\text{-gal}$:

$0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$

$\langle \text{proof} \rangle$

lemmas $td\text{-gal-lt} = td\text{-gal}$ [*simplified not-less [symmetric], simplified*]

lemma div-mult-le : $(a :: \text{nat}) \text{ div } b * b \leq a$

$\langle \text{proof} \rangle$

lemmas $sdl = \text{split-div-lemma}$ [*THEN iffD1, symmetric*]

lemma given-quot : $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$

$\langle \text{proof} \rangle$

lemma given-quot-alt : $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$

$\langle \text{proof} \rangle$

lemma diff-mod-le : $(a :: \text{nat}) < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$

$\langle \text{proof} \rangle$

lemma less-le-mult' :

$w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: \text{int})$

$\langle \text{proof} \rangle$

lemmas $\text{less-le-mult} = \text{less-le-mult'}$ [*simplified left-distrib, simplified*]

lemmas $\text{less-le-mult-minus} = \text{iffD2}$ [*OF le-diff-eq less-le-mult, simplified left-diff-distrib, standard*]

lemma lrlem' :

assumes $d: (i :: \text{nat}) \leq j \vee m < j'$

assumes $R1: i * k \leq j * k \implies R$

assumes $R2: \text{Suc } m * k' \leq j' * k' \implies R$

shows R $\langle \text{proof} \rangle$

lemma lrlem : $(0 :: \text{nat}) < sc \implies$

$(sc - n + (n + lb * n) \leq m * n) = (sc + lb * n \leq m * n)$

$\langle \text{proof} \rangle$

lemma gen-minus : $0 < n \implies f \ n = f \ (\text{Suc } (n - 1))$

$\langle \text{proof} \rangle$

lemma mpl-lem : $j \leq (i :: \text{nat}) \implies k < j \implies i - j + k < i$ $\langle \text{proof} \rangle$

lemma *nonneg-mod-div*:

$0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b :: \text{int}) \ \& \ 0 \leq a \operatorname{div} b$
 $\langle \text{proof} \rangle$

end

4 BinGeneral: Basic Definitions for Binary Integers

theory *BinGeneral*

imports *Num-Lemmas*

begin

4.1 Further properties of numerals

datatype *bit* = *B0* | *B1*

definition

$\text{Bit} :: \text{int} \Rightarrow \text{bit} \Rightarrow \text{int}$ (**infixl** *BIT* 90) **where**
 $k \text{ BIT } b = (\text{case } b \text{ of } B0 \Rightarrow 0 \mid B1 \Rightarrow 1) + k + k$

lemma *BIT-B0-eq-Bit0* [*simp*]: $w \text{ BIT } B0 = \text{Int.Bit0 } w$
 $\langle \text{proof} \rangle$

lemma *BIT-B1-eq-Bit1* [*simp*]: $w \text{ BIT } B1 = \text{Int.Bit1 } w$
 $\langle \text{proof} \rangle$

lemmas *BIT-simps* = *BIT-B0-eq-Bit0 BIT-B1-eq-Bit1*

hide-const (**open**) *B0 B1*

lemma *Min-ne-Pls* [*iff*]:

$\text{Int.Min} \sim = \text{Int.PlS}$
 $\langle \text{proof} \rangle$

lemmas *Pls-ne-Min* [*iff*] = *Min-ne-Pls* [*symmetric*]

lemmas *PlsMin-defs* [*intro!*] =

Pls-def Min-def Pls-def [*symmetric*] *Min-def* [*symmetric*]

lemmas *PlsMin-simps* [*simp*] = *PlsMin-defs* [*THEN Eq-TrueI*]

lemma *number-of-False-cong*:

$\text{False} \implies \text{number-of } x = \text{number-of } y$
 $\langle \text{proof} \rangle$

lemma *BIT-eq*: $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$
 $\langle \text{proof} \rangle$

lemmas *BIT-eqE* [*elim!*] = *BIT-eq* [*THEN conjE, standard*]

lemma *BIT-eq-iff* [*simp*]:
 $(u \text{ BIT } b = v \text{ BIT } c) = (u = v \wedge b = c)$
 $\langle \text{proof} \rangle$

lemmas *BIT-eqI* [*intro!*] = *conjI* [*THEN BIT-eq-iff* [*THEN iffD2*]]

lemma *less-Bits*:
 $(v \text{ BIT } b < w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *le-Bits*:
 $(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim = \text{bit.B1} \mid c \sim = \text{bit.B0}))$
 $\langle \text{proof} \rangle$

lemma *no-no* [*simp*]: *number-of* (*number-of* i) = i
 $\langle \text{proof} \rangle$

lemma *Bit-B0*:
 $k \text{ BIT } \text{bit.B0} = k + k$
 $\langle \text{proof} \rangle$

lemma *Bit-B1*:
 $k \text{ BIT } \text{bit.B1} = k + k + 1$
 $\langle \text{proof} \rangle$

lemma *Bit-B0-2t*: $k \text{ BIT } \text{bit.B0} = 2 * k$
 $\langle \text{proof} \rangle$

lemma *Bit-B1-2t*: $k \text{ BIT } \text{bit.B1} = 2 * k + 1$
 $\langle \text{proof} \rangle$

lemma *B-mod-2'*:
 $X = 2 \implies (w \text{ BIT } \text{bit.B1}) \bmod X = 1 \ \& \ (w \text{ BIT } \text{bit.B0}) \bmod X = 0$
 $\langle \text{proof} \rangle$

lemma *B1-mod-2* [*simp*]: $(\text{Int.Bit1 } w) \bmod 2 = 1$
 $\langle \text{proof} \rangle$

lemma *B0-mod-2* [*simp*]: $(\text{Int.Bit0 } w) \bmod 2 = 0$
 $\langle \text{proof} \rangle$

lemma *neB1E* [*elim!*]:

assumes $ne: y \neq bit.B1$
assumes $y: y = bit.B0 \implies P$
shows P
 $\langle proof \rangle$

lemma *bin-ex-rl*: $EX\ w\ b.\ w\ BIT\ b = bin$
 $\langle proof \rangle$

lemma *bin-exhaust*:
assumes $Q: \bigwedge x\ b.\ bin = x\ BIT\ b \implies Q$
shows Q
 $\langle proof \rangle$

4.2 Destructors for binary integers

definition *bin-rl* :: $int \Rightarrow int \times bit$ **where**
 $[code\ del]: bin-rl\ w = (THE\ (r,\ l).\ w = r\ BIT\ l)$

lemma *bin-rl-char*: $(bin-rl\ w = (r,\ l)) = (r\ BIT\ l = w)$
 $\langle proof \rangle$

definition
bin-rest-def $[code\ del]: bin-rest\ w = fst\ (bin-rl\ w)$

definition
bin-last-def $[code\ del]: bin-last\ w = snd\ (bin-rl\ w)$

primrec *bin-nth* **where**
 $Z: bin-nth\ w\ 0 = (bin-last\ w = bit.B1)$
 $| Suc: bin-nth\ w\ (Suc\ n) = bin-nth\ (bin-rest\ w)\ n$

lemma *bin-rl*: $bin-rl\ w = (bin-rest\ w,\ bin-last\ w)$
 $\langle proof \rangle$

lemma *bin-rl-simps* $[simp]$:
 $bin-rl\ Int.Pls = (Int.Pls,\ bit.B0)$
 $bin-rl\ Int.Min = (Int.Min,\ bit.B1)$
 $bin-rl\ (Int.Bit0\ r) = (r,\ bit.B0)$
 $bin-rl\ (Int.Bit1\ r) = (r,\ bit.B1)$
 $bin-rl\ (r\ BIT\ b) = (r,\ b)$
 $\langle proof \rangle$

declare *bin-rl-simps*(1–4) $[code]$

lemmas *bin-rl-simp* $[simp] = iffD1\ [OF\ bin-rl-char\ bin-rl]$

lemma *bin-abs-lem*:
 $bin = (w\ BIT\ b) \implies \sim\ bin = Int.Min \dashv\dashv \sim\ bin = Int.Pls \dashv\dashv$
 $nat\ (abs\ w) < nat\ (abs\ bin)$

$\langle \text{proof} \rangle$

lemma *bin-induct*:
 assumes *PPls*: $P \text{ Int.Pls}$
 and *PMin*: $P \text{ Int.Min}$
 and *PBit*: $\forall \text{bin bit}. P \text{ bin} \implies P (\text{bin BIT bit})$
 shows $P \text{ bin}$
 $\langle \text{proof} \rangle$

lemma *numeral-induct*:
 assumes *Pls*: $P \text{ Int.Pls}$
 assumes *Min*: $P \text{ Int.Min}$
 assumes *Bit0*: $\bigwedge w. \llbracket P w; w \neq \text{Int.Pls} \rrbracket \implies P (\text{Int.Bit0 } w)$
 assumes *Bit1*: $\bigwedge w. \llbracket P w; w \neq \text{Int.Min} \rrbracket \implies P (\text{Int.Bit1 } w)$
 shows $P x$
 $\langle \text{proof} \rangle$

lemma *bin-rest-simps* [*simp*]:
 $\text{bin-rest Int.Pls} = \text{Int.Pls}$
 $\text{bin-rest Int.Min} = \text{Int.Min}$
 $\text{bin-rest} (\text{Int.Bit0 } w) = w$
 $\text{bin-rest} (\text{Int.Bit1 } w) = w$
 $\text{bin-rest} (w \text{ BIT } b) = w$
 $\langle \text{proof} \rangle$

declare *bin-rest-simps*(1–4) [*code*]

lemma *bin-last-simps* [*simp*]:
 $\text{bin-last Int.Pls} = \text{bit.B0}$
 $\text{bin-last Int.Min} = \text{bit.B1}$
 $\text{bin-last} (\text{Int.Bit0 } w) = \text{bit.B0}$
 $\text{bin-last} (\text{Int.Bit1 } w) = \text{bit.B1}$
 $\text{bin-last} (w \text{ BIT } b) = b$
 $\langle \text{proof} \rangle$

declare *bin-last-simps*(1–4) [*code*]

lemma *bin-r-l-extras* [*simp*]:
 $\text{bin-last } 0 = \text{bit.B0}$
 $\text{bin-last } (-1) = \text{bit.B1}$
 $\text{bin-last } -1 = \text{bit.B1}$
 $\text{bin-last } 1 = \text{bit.B1}$
 $\text{bin-rest } 1 = 0$
 $\text{bin-rest } 0 = 0$
 $\text{bin-rest } (-1) = -1$
 $\text{bin-rest } -1 = -1$
 $\langle \text{proof} \rangle$

lemma *bin-last-mod*:

$\text{bin-last } w = (\text{if } w \bmod 2 = 0 \text{ then bit.B0 else bit.B1})$
 $\langle \text{proof} \rangle$

lemma *bin-rest-div*:
 $\text{bin-rest } w = w \text{ div } 2$
 $\langle \text{proof} \rangle$

lemma *Bit-div2 [simp]*: $(w \text{ BIT } b) \text{ div } 2 = w$
 $\langle \text{proof} \rangle$

lemma *Bit0-div2 [simp]*: $(\text{Int.Bit0 } w) \text{ div } 2 = w$
 $\langle \text{proof} \rangle$

lemma *Bit1-div2 [simp]*: $(\text{Int.Bit1 } w) \text{ div } 2 = w$
 $\langle \text{proof} \rangle$

lemma *bin-nth-lem [rule-format]*:
 $\text{ALL } y. \text{bin-nth } x = \text{bin-nth } y \implies x = y$
 $\langle \text{proof} \rangle$

lemma *bin-nth-eq-iff*: $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$
 $\langle \text{proof} \rangle$

lemmas *bin-eqI* = *ext [THEN bin-nth-eq-iff [THEN iffD1], standard]*

lemma *bin-nth-Pls [simp]*: $\sim \text{bin-nth Int.Pls } n$
 $\langle \text{proof} \rangle$

lemma *bin-nth-Min [simp]*: $\text{bin-nth Int.Min } n$
 $\langle \text{proof} \rangle$

lemma *bin-nth-0-BIT*: $\text{bin-nth } (w \text{ BIT } b) 0 = (b = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *bin-nth-Suc-BIT*: $\text{bin-nth } (w \text{ BIT } b) (\text{Suc } n) = \text{bin-nth } w n$
 $\langle \text{proof} \rangle$

lemma *bin-nth-minus [simp]*: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) n = \text{bin-nth } w (n - 1)$
 $\langle \text{proof} \rangle$

lemma *bin-nth-minus-Bit0 [simp]*:
 $0 < n \implies \text{bin-nth } (\text{Int.Bit0 } w) n = \text{bin-nth } w (n - 1)$
 $\langle \text{proof} \rangle$

lemma *bin-nth-minus-Bit1 [simp]*:
 $0 < n \implies \text{bin-nth } (\text{Int.Bit1 } w) n = \text{bin-nth } w (n - 1)$
 $\langle \text{proof} \rangle$

lemmas $\text{bin-nth-0} = \text{bin-nth.simps}(1)$
lemmas $\text{bin-nth-Suc} = \text{bin-nth.simps}(2)$

lemmas $\text{bin-nth-simps} =$
 $\text{bin-nth-0} \text{ bin-nth-Suc } \text{bin-nth-Pls} \text{ bin-nth-Min } \text{bin-nth-minus}$
 $\text{bin-nth-minus-Bit0} \text{ bin-nth-minus-Bit1}$

4.3 Recursion combinator for binary integers

lemma $\text{brlem}: (\text{bin} = \text{Int.Min}) = (- \text{bin} + \text{Int.pred } 0 = 0)$
 $\langle \text{proof} \rangle$

function

$\text{bin-rec} :: 'a \Rightarrow 'a \Rightarrow (\text{int} \Rightarrow \text{bit} \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{int} \Rightarrow 'a$

where

$\text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ } \text{bin} = (\text{if } \text{bin} = \text{Int.Pls} \text{ then } f1$
 $\text{else if } \text{bin} = \text{Int.Min} \text{ then } f2$
 $\text{else case bin-rl bin of } (w, b) \Rightarrow f3 \text{ } w \text{ } b \text{ } (\text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ } w))$
 $\langle \text{proof} \rangle$

termination

$\langle \text{proof} \rangle$

declare $\text{bin-rec.simps} [\text{simp del}]$

lemma bin-rec-PM :

$f = \text{bin-rec } f1 \text{ } f2 \text{ } f3 \implies f \text{ Int.Pls} = f1 \text{ \& } f \text{ Int.Min} = f2$
 $\langle \text{proof} \rangle$

lemma $\text{bin-rec-Pls}: \text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ Int.Pls} = f1$
 $\langle \text{proof} \rangle$

lemma $\text{bin-rec-Min}: \text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ Int.Min} = f2$
 $\langle \text{proof} \rangle$

lemma bin-rec-Bit0 :

$f3 \text{ Int.Pls } \text{bit.B0 } f1 = f1 \implies$
 $\text{bin-rec } f1 \text{ } f2 \text{ } f3 (\text{Int.Bit0 } w) = f3 \text{ } w \text{ } \text{bit.B0 } (\text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ } w)$
 $\langle \text{proof} \rangle$

lemma bin-rec-Bit1 :

$f3 \text{ Int.Min } \text{bit.B1 } f2 = f2 \implies$
 $\text{bin-rec } f1 \text{ } f2 \text{ } f3 (\text{Int.Bit1 } w) = f3 \text{ } w \text{ } \text{bit.B1 } (\text{bin-rec } f1 \text{ } f2 \text{ } f3 \text{ } w)$
 $\langle \text{proof} \rangle$

lemma bin-rec-Bit :

$f = \text{bin-rec } f1 \text{ } f2 \text{ } f3 \implies f3 \text{ Int.Pls } \text{bit.B0 } f1 = f1 \implies$
 $f3 \text{ Int.Min } \text{bit.B1 } f2 = f2 \implies f (w \text{ BIT } b) = f3 \text{ } w \text{ } b (f \text{ } w)$
 $\langle \text{proof} \rangle$

lemmas *bin-rec-simps* = refl [THEN *bin-rec-Bit*] *bin-rec-Pls bin-rec-Min bin-rec-Bit0 bin-rec-Bit1*

4.4 Truncating binary integers

definition

bin-sign-def [code del] : *bin-sign* = *bin-rec Int.Pls Int.Min (%w b s. s)*

lemma *bin-sign-simps* [simp]:

bin-sign Int.Pls = *Int.Pls*
bin-sign Int.Min = *Int.Min*
bin-sign (Int.Bit0 w) = *bin-sign w*
bin-sign (Int.Bit1 w) = *bin-sign w*
bin-sign (w BIT b) = *bin-sign w*
 ⟨proof⟩

declare *bin-sign-simps*(1–4) [code]

lemma *bin-sign-rest* [simp]:

bin-sign (bin-rest w) = (*bin-sign w*)
 ⟨proof⟩

consts

bintrunc :: nat => int => int

primrec

Z : *bintrunc 0 bin* = *Int.Pls*
Suc : *bintrunc (Suc n) bin* = *bintrunc n (bin-rest bin) BIT (bin-last bin)*

consts

sbintrunc :: nat => int => int

primrec

Z : *sbintrunc 0 bin* =
 (case *bin-last bin* of *bit.B1* => *Int.Min* | *bit.B0* => *Int.Pls*)
Suc : *sbintrunc (Suc n) bin* = *sbintrunc n (bin-rest bin) BIT (bin-last bin)*

lemma *sign-bintr*:

!!*w. bin-sign (bintrunc n w)* = *Int.Pls*
 ⟨proof⟩

lemma *bintrunc-mod2p*:

!!*w. bintrunc n w* = (*w mod 2 ^ n* :: int)
 ⟨proof⟩

lemma *sbintrunc-mod2p*:

!!*w. sbintrunc n w* = ((*w + 2 ^ n*) mod 2 ^ (*Suc n*) - 2 ^ *n* :: int)
 ⟨proof⟩

4.5 Simplifications for (s)bintrunc

lemma *bit-bool*:

$(b = (b' = \text{bit}.B1)) = (b' = (\text{if } b \text{ then } \text{bit}.B1 \text{ else } \text{bit}.B0))$
 $\langle \text{proof} \rangle$

lemmas *bit-bool1* [simp] = refl [THEN bit-bool [THEN iffD1], symmetric]

lemma *bin-sign-lem*:

$!!\text{bin. } (\text{bin-sign } (\text{sbintrunc } n \text{ bin}) = \text{Int.Min}) = \text{bin-nth bin } n$
 $\langle \text{proof} \rangle$

lemma *nth-bintr*:

$!!w \text{ m. } \text{bin-nth } (\text{bintrunc } m \text{ w}) \text{ n} = (n < m \ \& \ \text{bin-nth } w \text{ n})$
 $\langle \text{proof} \rangle$

lemma *nth-sbintr*:

$!!w \text{ m. } \text{bin-nth } (\text{sbintrunc } m \text{ w}) \text{ n} =$
 $(\text{if } n < m \text{ then } \text{bin-nth } w \text{ n else } \text{bin-nth } w \text{ m})$
 $\langle \text{proof} \rangle$

lemma *bin-nth-Bit*:

$\text{bin-nth } (w \text{ BIT } b) \text{ n} = (n = 0 \ \& \ b = \text{bit}.B1 \mid (\text{EX } m. \text{ n} = \text{Suc } m \ \& \ \text{bin-nth } w \text{ m}))$
 $\langle \text{proof} \rangle$

lemma *bin-nth-Bit0*:

$\text{bin-nth } (\text{Int.Bit0 } w) \text{ n} = (\text{EX } m. \text{ n} = \text{Suc } m \ \& \ \text{bin-nth } w \text{ m})$
 $\langle \text{proof} \rangle$

lemma *bin-nth-Bit1*:

$\text{bin-nth } (\text{Int.Bit1 } w) \text{ n} = (n = 0 \mid (\text{EX } m. \text{ n} = \text{Suc } m \ \& \ \text{bin-nth } w \text{ m}))$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-l*:

$n \leq m \implies (\text{bintrunc } m (\text{bintrunc } n \text{ w}) = \text{bintrunc } n \text{ w})$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-l*:

$n \leq m \implies (\text{sbintrunc } m (\text{sbintrunc } n \text{ w}) = \text{sbintrunc } n \text{ w})$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-ge*:

$n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m \text{ w}) = \text{bintrunc } n \text{ w})$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-min* [simp]:

$\text{bintrunc } m (\text{bintrunc } n \text{ w}) = \text{bintrunc } (\min m \text{ n}) \text{ w}$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-min* [simp]:

$$\text{sbintrunc } m \ (\text{sbintrunc } n \ w) = \text{sbintrunc } (\min m \ n) \ w$$

<proof>

lemmas *bintrunc-Pls* =

bintrunc.Suc [where *bin=Int.Pls*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *bintrunc-Min* [simp] =

bintrunc.Suc [where *bin=Int.Min*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *bintrunc-BIT* [simp] =

bintrunc.Suc [where *bin=w BIT b*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemma *bintrunc-Bit0* [simp]:

$$\text{bintrunc } (\text{Suc } n) \ (\text{Int.Bit0 } w) = \text{Int.Bit0 } (\text{bintrunc } n \ w)$$

<proof>

lemma *bintrunc-Bit1* [simp]:

$$\text{bintrunc } (\text{Suc } n) \ (\text{Int.Bit1 } w) = \text{Int.Bit1 } (\text{bintrunc } n \ w)$$

<proof>

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

bintrunc-Bit0 bintrunc-Bit1

lemmas *sbintrunc-Suc-Pls* =

sbintrunc.Suc [where *bin=Int.Pls*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Suc-Min* =

sbintrunc.Suc [where *bin=Int.Min*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Suc-BIT* [simp] =

sbintrunc.Suc [where *bin=w BIT b*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemma *sbintrunc-Suc-Bit0* [simp]:

$$\text{sbintrunc } (\text{Suc } n) \ (\text{Int.Bit0 } w) = \text{Int.Bit0 } (\text{sbintrunc } n \ w)$$

<proof>

lemma *sbintrunc-Suc-Bit1* [simp]:

$$\text{sbintrunc } (\text{Suc } n) \ (\text{Int.Bit1 } w) = \text{Int.Bit1 } (\text{sbintrunc } n \ w)$$

<proof>

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls sbintrunc-Suc-Min sbintrunc-Suc-BIT*

sbintrunc-Suc-Bit0 sbintrunc-Suc-Bit1

lemmas *sbintrunc-Pls* =
 sbintrunc.Z [**where** *bin*=*Int.Pls*,
 simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-Min* =
 sbintrunc.Z [**where** *bin*=*Int.Min*,
 simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B0* [*simp*] =
 sbintrunc.Z [**where** *bin*=*w BIT bit.B0*,
 simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
 sbintrunc.Z [**where** *bin*=*w BIT bit.B1*,
 simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemma *sbintrunc-0-Bit0* [*simp*]: *sbintrunc 0 (Int.Bit0 w) = Int.Pls*
 ⟨*proof*⟩

lemma *sbintrunc-0-Bit1* [*simp*]: *sbintrunc 0 (Int.Bit1 w) = Int.Min*
 ⟨*proof*⟩

lemmas *sbintrunc-0-simps* =
 sbintrunc-Pls sbintrunc-Min sbintrunc-0-BIT-B0 sbintrunc-0-BIT-B1
 sbintrunc-0-Bit0 sbintrunc-0-Bit1

lemmas *bintrunc-simps* = *bintrunc.Z bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps sbintrunc-Sucs*

lemma *bintrunc-minus*:
 $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$
 ⟨*proof*⟩

lemma *sbintrunc-minus*:
 $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$
 ⟨*proof*⟩

lemmas *bintrunc-minus-simps* =
 bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemmas *sbintrunc-minus-simps* =
 sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemma *bintrunc-n-Pls* [*simp*]:
 $\text{bintrunc } n \ \text{Int.Pls} = \text{Int.Pls}$
 ⟨*proof*⟩

lemma *sbintrunc-n-PM* [*simp*]:
 $\text{sbintrunc } n \ \text{Int.Pls} = \text{Int.Pls}$

sbintrunc *n* *Int.Min* = *Int.Min*
 ⟨proof⟩

lemmas *thobini1* = *arg-cong* [where *f* = %*w*. *w* *BIT* *b*, *standard*]

lemmas *bintrunc-BIT-I* = *trans* [OF *bintrunc-BIT* *thobini1*]

lemmas *bintrunc-Min-I* = *trans* [OF *bintrunc-Min* *thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps*(1–3) [THEN *thobini1* [THEN [2] *trans*], *standard*]

lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)

lemmas *bintrunc-Min-minus-I* = *bmsts*(2)

lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-0-Min*: *bintrunc* 0 *Int.Min* = *Int.Pl*
 ⟨proof⟩

lemma *bintrunc-0-BIT*: *bintrunc* 0 (*w* *BIT* *b*) = *Int.Pl*
 ⟨proof⟩

lemma *bintrunc-Suc-lem*:

bintrunc (*Suc* *n*) *x* = *y* ==> *m* = *Suc* *n* ==> *bintrunc* *m* *x* = *y*
 ⟨proof⟩

lemmas *bintrunc-Suc-Ialts* =

bintrunc-Min-I [THEN *bintrunc-Suc-lem*, *standard*]

bintrunc-BIT-I [THEN *bintrunc-Suc-lem*, *standard*]

lemmas *sbintrunc-BIT-I* = *trans* [OF *sbintrunc-Suc-BIT* *thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs(1–3) [THEN *thobini1* [THEN [2] *trans*], *standard*]

lemmas *sbintrunc-Suc-minus-Is* =

sbintrunc-minus-simps(1–3) [THEN *thobini1* [THEN [2] *trans*], *standard*]

lemma *sbintrunc-Suc-lem*:

sbintrunc (*Suc* *n*) *x* = *y* ==> *m* = *Suc* *n* ==> *sbintrunc* *m* *x* = *y*
 ⟨proof⟩

lemmas *sbintrunc-Suc-Ialts* =

sbintrunc-Suc-Is [THEN *sbintrunc-Suc-lem*, *standard*]

lemma *sbintrunc-bintrunc-lt*:

m > *n* ==> *sbintrunc* *n* (*bintrunc* *m* *w*) = *sbintrunc* *n* *w*
 ⟨proof⟩

lemma *bintrunc-sbintrunc-le*:

m <= *Suc* *n* ==> *bintrunc* *m* (*sbintrunc* *n* *w*) = *bintrunc* *m* *w*
 ⟨proof⟩

lemmas *bintrunc-sbintrunc* [simp] = order-refl [THEN *bintrunc-sbintrunc-le*]
lemmas *sbintrunc-bintrunc* [simp] = lessI [THEN *sbintrunc-bintrunc-lt*]
lemmas *bintrunc-bintrunc* [simp] = order-refl [THEN *bintrunc-bintrunc-l*]
lemmas *sbintrunc-sbintrunc* [simp] = order-refl [THEN *sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc'* [simp]:
 $0 < n \implies \text{bintrunc } n (\text{sbintrunc } (n - 1) w) = \text{bintrunc } n w$
 <proof>

lemma *sbintrunc-bintrunc'* [simp]:
 $0 < n \implies \text{sbintrunc } (n - 1) (\text{bintrunc } n w) = \text{sbintrunc } (n - 1) w$
 <proof>

lemma *bin-sbin-eq-iff*:
 $\text{bintrunc } (\text{Suc } n) x = \text{bintrunc } (\text{Suc } n) y <->$
 $\text{sbintrunc } n x = \text{sbintrunc } n y$
 <proof>

lemma *bin-sbin-eq-iff'*:
 $0 < n \implies \text{bintrunc } n x = \text{bintrunc } n y <->$
 $\text{sbintrunc } (n - 1) x = \text{sbintrunc } (n - 1) y$
 <proof>

lemmas *bintrunc-sbintruncS0* [simp] = *bintrunc-sbintrunc'* [unfolded *One-nat-def*]
lemmas *sbintrunc-bintruncS0* [simp] = *sbintrunc-bintrunc'* [unfolded *One-nat-def*]

lemmas *bintrunc-bintrunc-l'* = le-add1 [THEN *bintrunc-bintrunc-l*]
lemmas *sbintrunc-sbintrunc-l'* = le-add1 [THEN *sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =
 trans [OF *iszero-def* [THEN *Not-eq-iff* [THEN *iffD2*]] refl, standard]

lemmas *bintrunc-pred-simps* [simp] =
bintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemmas *sbintrunc-pred-simps* [simp] =
sbintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemma *no-bintr-alt*:
 $\text{number-of } (\text{bintrunc } n w) = w \bmod 2^n$
 <proof>

lemma *no-bintr-alt1*: $\text{bintrunc } n = (\%w. w \bmod 2^n :: \text{int})$
 <proof>

lemma *range-bintrunc*: $\text{range } (\text{bintrunc } n) = \{i. 0 \leq i \ \& \ i < 2^n\}$

$\langle \text{proof} \rangle$

lemma *no-bintr*:

$\text{number-of } (\text{bintrunc } n \ w) = (\text{number-of } w \bmod 2^n :: \text{int})$

$\langle \text{proof} \rangle$

lemma *no-sbintr-alt2*:

$\text{sbintrunc } n = (\%w. (w + 2^n) \bmod 2^{\text{Suc } n} - 2^n :: \text{int})$

$\langle \text{proof} \rangle$

lemma *no-sbintr*:

$\text{number-of } (\text{sbintrunc } n \ w) =$

$((\text{number-of } w + 2^n) \bmod 2^{\text{Suc } n} - 2^n :: \text{int})$

$\langle \text{proof} \rangle$

lemma *range-sbintrunc*:

$\text{range } (\text{sbintrunc } n) = \{i. - (2^n) \leq i \ \& \ i < 2^n\}$

$\langle \text{proof} \rangle$

lemma *sb-inc-lem*:

$(a :: \text{int}) + 2^k < 0 \implies a + 2^k + 2^{\text{Suc } k} \leq (a + 2^k) \bmod 2^{\text{Suc } k}$

$\langle \text{proof} \rangle$

lemma *sb-inc-lem'*:

$(a :: \text{int}) < - (2^k) \implies a + 2^k + 2^{\text{Suc } k} \leq (a + 2^k) \bmod 2^{\text{Suc } k}$

$\langle \text{proof} \rangle$

lemma *sbintrunc-inc*:

$x < - (2^n) \implies x + 2^{\text{Suc } n} \leq \text{sbintrunc } n \ x$

$\langle \text{proof} \rangle$

lemma *sb-dec-lem*:

$(0 :: \text{int}) \leq - (2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$

$\langle \text{proof} \rangle$

lemma *sb-dec-lem'*:

$(2 :: \text{int})^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$

$\langle \text{proof} \rangle$

lemma *sbintrunc-dec*:

$x \geq (2^n) \implies x - 2^{\text{Suc } n} \geq \text{sbintrunc } n \ x$

$\langle \text{proof} \rangle$

lemmas *zmod-uminus'* = *zmod-uminus* [where $b=c$, standard]

lemmas *zpower-zmod'* = *zpower-zmod* [where $m=c$ and $y=k$, standard]

lemmas *brdmod1s'* [symmetric] =

mod-add-left-eq mod-add-right-eq

zmod-zsub-left-eq zmod-zsub-right-eq

zmod-zmult1-eq zmod-zmult1-eq-rev

lemmas *brdmods'* [*symmetric*] =
zpower-zmod' [*symmetric*]
trans [*OF mod-add-left-eq mod-add-right-eq*]
trans [*OF zmod-zsub-left-eq zmod-zsub-right-eq*]
trans [*OF zmod-zmult1-eq zmod-zmult1-eq-rev*]
zmod-uminus' [*symmetric*]
mod-add-left-eq [**where** *b* = 1::int]
zmod-zsub-left-eq [**where** *b* = 1]

lemmas *bintr-arith1s* =
brdmod1s' [**where** *c*=2^{*n*}::int, *folded pred-def succ-def bintrunc-mod2p, standard*]
lemmas *bintr-ariths* =
brdmods' [**where** *c*=2^{*n*}::int, *folded pred-def succ-def bintrunc-mod2p, standard*]

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [*OF zless2p, standard*]

lemma *bintr-ge0*: (*0* :: int) <= *number-of* (*bintrunc n w*)
 ⟨*proof*⟩

lemma *bintr-lt2p*: *number-of* (*bintrunc n w*) < (2^{*n*} :: int)
 ⟨*proof*⟩

lemma *bintr-Min*:
number-of (*bintrunc n Int.Min*) = (2^{*n*} :: int) − 1
 ⟨*proof*⟩

lemma *sbintr-ge*: (− (2^{*n*} :: int) <= *number-of* (*sbintrunc n w*)
 ⟨*proof*⟩

lemma *sbintr-lt*: *number-of* (*sbintrunc n w*) < (2^{*n*} :: int)
 ⟨*proof*⟩

lemma *bintrunc-Suc*:
bintrunc (*Suc n*) *bin* = *bintrunc n* (*bin-rest bin*) *BIT bin-last bin*
 ⟨*proof*⟩

lemma *sign-Pls-ge-0*:
(bin-sign bin = Int.Pl) = (*number-of bin* >= (*0* :: int))
 ⟨*proof*⟩

lemma *sign-Min-lt-0*:
(bin-sign bin = Int.Min) = (*number-of bin* < (*0* :: int))
 ⟨*proof*⟩

lemmas *sign-Min-neg* = *trans* [*OF sign-Min-lt-0 neg-def* [*symmetric*]]

lemma *bin-rest-trunc*:

!!bin. (bin-rest (bintrunc n bin)) = bintrunc (n - 1) (bin-rest bin)
 ⟨proof⟩

lemma bin-rest-power-trunc [rule-format] :
 (bin-rest ^ k) (bintrunc n bin) =
 bintrunc (n - k) ((bin-rest ^ k) bin)
 ⟨proof⟩

lemma bin-rest-trunc-i:
 bintrunc n (bin-rest bin) = bin-rest (bintrunc (Suc n) bin)
 ⟨proof⟩

lemma bin-rest-strunc:
 !!bin. bin-rest (sbintrunc (Suc n) bin) = sbintrunc n (bin-rest bin)
 ⟨proof⟩

lemma bintrunc-rest [simp]:
 !!bin. bintrunc n (bin-rest (bintrunc n bin)) = bin-rest (bintrunc n bin)
 ⟨proof⟩

lemma sbintrunc-rest [simp]:
 !!bin. sbintrunc n (bin-rest (sbintrunc n bin)) = bin-rest (sbintrunc n bin)
 ⟨proof⟩

lemma bintrunc-rest':
 bintrunc n o bin-rest o bintrunc n = bin-rest o bintrunc n
 ⟨proof⟩

lemma sbintrunc-rest':
 sbintrunc n o bin-rest o sbintrunc n = bin-rest o sbintrunc n
 ⟨proof⟩

lemma rco-lem:
 f o g o f = g o f ==> f o (g o f) ^ n = g ^ n o f
 ⟨proof⟩

lemma rco-alt: (f o g) ^ n o f = f o (g o f) ^ n
 ⟨proof⟩

lemmas rco-bintr = bintrunc-rest'
 [THEN rco-lem [THEN fun-cong], unfolded o-def]
lemmas rco-sbintr = sbintrunc-rest'
 [THEN rco-lem [THEN fun-cong], unfolded o-def]

4.6 Splitting and concatenation

primrec bin-split :: nat ⇒ int ⇒ int × int **where**
 Z: bin-split 0 w = (w, Int.Pls)
 | Suc: bin-split (Suc n) w = (let (w1, w2) = bin-split n (bin-rest w)

in (w1, w2 BIT bin-last w))

primrec *bin-cat* :: *int* \Rightarrow *nat* \Rightarrow *int* \Rightarrow *int* **where**

Z: *bin-cat* *w* 0 *v* = *w*

| *Suc*: *bin-cat* *w* (*Suc* *n*) *v* = *bin-cat* *w* *n* (*bin-rest* *v*) BIT *bin-last* *v*

4.7 Miscellaneous lemmas

lemma *funpow-minus-simp*:

$0 < n \implies f^{^^n} = f \circ f^{^^(n-1)}$

<proof>

lemmas *funpow-pred-simp* [*simp*] =

funpow-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *replicate-minus-simp* =

trans [*OF gen-minus* [**where** $f = \%n. \text{replicate } n \ x$] *replicate.replicate-Suc*,
standard]

lemmas *replicate-pred-simp* [*simp*] =

replicate-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *power-Suc-no* [*simp*] = *power-Suc* [*of number-of a, standard*]

lemmas *power-minus-simp* =

trans [*OF gen-minus* [**where** $f = \text{power } f$] *power-Suc*, *standard*]

lemmas *power-pred-simp* =

power-minus-simp [*of number-of bin, simplified nobm1, standard*]

lemmas *power-pred-simp-no* [*simp*] = *power-pred-simp* [**where** $f = \text{number-of } f$,
standard]

lemma *list-exhaust-size-gt0*:

assumes $y: \bigwedge a \text{ list. } y = a \# \text{list} \implies P$

shows $0 < \text{length } y \implies P$

<proof>

lemma *list-exhaust-size-eq0*:

assumes $y: y = [] \implies P$

shows $\text{length } y = 0 \implies P$

<proof>

lemma *size-Cons-lem-eq*:

$y = xa \# \text{list} \implies \text{size } y = \text{Suc } k \implies \text{size } \text{list} = k$

<proof>

lemma *size-Cons-lem-eq-bin*:

$y = xa \# \text{list} \implies \text{size } y = \text{number-of } (\text{Int.succ } k) \implies$

$\text{size } \text{list} = \text{number-of } k$

```

⟨proof⟩

lemmas ls-splits =
  prod.split split-split prod.split-asm split-split-asm split-if-asm

lemma not-B1-is-B0:  $y \neq \text{bit.B1} \implies y = \text{bit.B0}$ 
  ⟨proof⟩

lemma B1-ass-B0:
  assumes  $y: y = \text{bit.B0} \implies y = \text{bit.B1}$ 
  shows  $y = \text{bit.B1}$ 
  ⟨proof⟩
lemmas n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'

lemmas s2n-ths = n2s-ths [symmetric]

end

```

5 BitSyntax: Syntactic classes for bitwise operations

```

theory BitSyntax
imports BinGeneral
begin

```

```

class bit =
  fixes bitNOT :: 'a  $\Rightarrow$  'a      (NOT - [70] 71)
  and bitAND :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr AND 64)
  and bitOR  :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr OR 59)
  and bitXOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr XOR 59)

```

We want the bitwise operations to bind slightly weaker than $+$ and $-$, but $\sim\sim$ to bind slightly stronger than $*$.

Testing and shifting operations.

```

class bits = bit +
  fixes test-bit :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool (infixl !! 100)
  and lsb      :: 'a  $\Rightarrow$  bool
  and set-bit  :: 'a  $\Rightarrow$  nat  $\Rightarrow$  bool  $\Rightarrow$  'a
  and set-bits :: (nat  $\Rightarrow$  bool)  $\Rightarrow$  'a (binder BITS 10)
  and shiftrl :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl << 55)
  and shiftr  :: 'a  $\Rightarrow$  nat  $\Rightarrow$  'a (infixl >> 55)

class bitss = bits +
  fixes msb    :: 'a  $\Rightarrow$  bool

```

5.1 Bitwise operations on *bit*

instantiation *bit* :: *bit*

begin

primrec *bitNOT-bit* **where**

NOT bit.B0 = *bit.B1*
 | *NOT bit.B1* = *bit.B0*

primrec *bitAND-bit* **where**

bit.B0 AND y = *bit.B0*
 | *bit.B1 AND y* = *y*

primrec *bitOR-bit* **where**

bit.B0 OR y = *y*
 | *bit.B1 OR y* = *bit.B1*

primrec *bitXOR-bit* **where**

bit.B0 XOR y = *y*
 | *bit.B1 XOR y* = *NOT y*

instance $\langle proof \rangle$

end

lemmas *bit-simps* =

bitNOT-bit.simps bitAND-bit.simps bitOR-bit.simps bitXOR-bit.simps

lemma *bit-extra-simps* [*simp*]:

x AND bit.B0 = *bit.B0*
x AND bit.B1 = *x*
x OR bit.B1 = *bit.B1*
x OR bit.B0 = *x*
x XOR bit.B1 = *NOT x*
x XOR bit.B0 = *x*
 $\langle proof \rangle$

lemma *bit-ops-comm*:

$(x::bit) \text{ AND } y = y \text{ AND } x$
 $(x::bit) \text{ OR } y = y \text{ OR } x$
 $(x::bit) \text{ XOR } y = y \text{ XOR } x$
 $\langle proof \rangle$

lemma *bit-ops-same* [*simp*]:

$(x::bit) \text{ AND } x = x$
 $(x::bit) \text{ OR } x = x$
 $(x::bit) \text{ XOR } x = bit.B0$
 $\langle proof \rangle$

lemma *bit-not-not* [*simp*]: *NOT (NOT (x::bit))* = *x*

$\langle proof \rangle$

end

6 BinOperations: Bitwise Operations on Binary Integers

theory *BinOperations*
imports *BinGeneral BitSyntax*
begin

6.1 Logical operations

bit-wise logical operations on the int type

instantiation *int* :: *bit*
begin

definition

int-not-def [code del]: *bitNOT* = *bin-rec Int.Min Int.Pls*
 $(\lambda w \ b \ s. \ s \ \text{BIT} \ (\text{NOT } b))$

definition

int-and-def [code del]: *bitAND* = *bin-rec* $(\lambda x. \text{Int.Pls}) (\lambda y. \ y)$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{AND} \ \text{bin-last } y))$

definition

int-or-def [code del]: *bitOR* = *bin-rec* $(\lambda x. \ x) (\lambda y. \ \text{Int.Min})$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{OR} \ \text{bin-last } y))$

definition

int-xor-def [code del]: *bitXOR* = *bin-rec* $(\lambda x. \ x) \ \text{bitNOT}$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{XOR} \ \text{bin-last } y))$

instance $\langle proof \rangle$

end

lemma *int-not-simps* [simp]:

NOT Int.Pls = *Int.Min*
NOT Int.Min = *Int.Pls*
NOT (*Int.Bit0* *w*) = *Int.Bit1* (*NOT w*)
NOT (*Int.Bit1* *w*) = *Int.Bit0* (*NOT w*)
NOT (*w BIT b*) = (*NOT w*) *BIT* (*NOT b*)
 $\langle proof \rangle$

declare *int-not-simps*(1–4) [code]

lemma *int-xor-Pls* [*simp*, *code*]:

$$\text{Int.Pls XOR } x = x$$

<proof>

lemma *int-xor-Min* [*simp*, *code*]:

$$\text{Int.Min XOR } x = \text{NOT } x$$

<proof>

lemma *int-xor-Bits* [*simp*]:

$$(x \text{ BIT } b) \text{ XOR } (y \text{ BIT } c) = (x \text{ XOR } y) \text{ BIT } (b \text{ XOR } c)$$

<proof>

lemma *int-xor-Bits2* [*simp*, *code*]:

$$(\text{Int.Bit0 } x) \text{ XOR } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ XOR } y)$$

$$(\text{Int.Bit0 } x) \text{ XOR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ XOR } y)$$

$$(\text{Int.Bit1 } x) \text{ XOR } (\text{Int.Bit0 } y) = \text{Int.Bit1 } (x \text{ XOR } y)$$

$$(\text{Int.Bit1 } x) \text{ XOR } (\text{Int.Bit1 } y) = \text{Int.Bit0 } (x \text{ XOR } y)$$

<proof>

lemma *int-xor-x-simps'*:

$$w \text{ XOR } (\text{Int.Pls BIT bit.B0}) = w$$

$$w \text{ XOR } (\text{Int.Min BIT bit.B1}) = \text{NOT } w$$

<proof>

lemma *int-xor-extra-simps* [*simp*, *code*]:

$$w \text{ XOR } \text{Int.Pls} = w$$

$$w \text{ XOR } \text{Int.Min} = \text{NOT } w$$

<proof>

lemma *int-or-Pls* [*simp*, *code*]:

$$\text{Int.Pls OR } x = x$$

<proof>

lemma *int-or-Min* [*simp*, *code*]:

$$\text{Int.Min OR } x = \text{Int.Min}$$

<proof>

lemma *int-or-Bits* [*simp*]:

$$(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \text{ OR } c)$$

<proof>

lemma *int-or-Bits2* [*simp*, *code*]:

$$(\text{Int.Bit0 } x) \text{ OR } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ OR } y)$$

$$(\text{Int.Bit0 } x) \text{ OR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ OR } y)$$

$$(\text{Int.Bit1 } x) \text{ OR } (\text{Int.Bit0 } y) = \text{Int.Bit1 } (x \text{ OR } y)$$

$$(\text{Int.Bit1 } x) \text{ OR } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ OR } y)$$

<proof>

lemma *int-or-x-simps'*:

$w \text{ OR } (\text{Int.Pls BIT bit.B0}) = w$
 $w \text{ OR } (\text{Int.Min BIT bit.B1}) = \text{Int.Min}$
 ⟨proof⟩

lemma *int-or-extra-simps* [simp, code]:

$w \text{ OR Int.Pls} = w$
 $w \text{ OR Int.Min} = \text{Int.Min}$
 ⟨proof⟩

lemma *int-and-Pls* [simp, code]:

$\text{Int.Pls AND } x = \text{Int.Pls}$
 ⟨proof⟩

lemma *int-and-Min* [simp, code]:

$\text{Int.Min AND } x = x$
 ⟨proof⟩

lemma *int-and-Bits* [simp]:

$(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \text{ AND } c)$
 ⟨proof⟩

lemma *int-and-Bits2* [simp, code]:

$(\text{Int.Bit0 } x) \text{ AND } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit0 } x) \text{ AND } (\text{Int.Bit1 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit1 } x) \text{ AND } (\text{Int.Bit0 } y) = \text{Int.Bit0 } (x \text{ AND } y)$
 $(\text{Int.Bit1 } x) \text{ AND } (\text{Int.Bit1 } y) = \text{Int.Bit1 } (x \text{ AND } y)$
 ⟨proof⟩

lemma *int-and-x-simps'*:

$w \text{ AND } (\text{Int.Pls BIT bit.B0}) = \text{Int.Pls}$
 $w \text{ AND } (\text{Int.Min BIT bit.B1}) = w$
 ⟨proof⟩

lemma *int-and-extra-simps* [simp, code]:

$w \text{ AND Int.Pls} = \text{Int.Pls}$
 $w \text{ AND Int.Min} = w$
 ⟨proof⟩

lemma *bin-ops-comm*:

shows

int-and-comm: $!!y::\text{int}. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $!!y::\text{int}. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $!!y::\text{int}. x \text{ XOR } y = y \text{ XOR } x$
 ⟨proof⟩

lemma *bin-ops-same* [simp]:

$(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$

$(x::int) \text{ XOR } x = Int.Pls$
 $\langle proof \rangle$

lemma *int-not-not* [simp]: $NOT (NOT (x::int)) = x$
 $\langle proof \rangle$

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min

lemma *bbw-ao-absorb*:
 $!!y::int. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$
 $\langle proof \rangle$

lemma *bbw-ao-absorbs-other*:
 $x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::int)$
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::int)$
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::int)$
 $\langle proof \rangle$

lemmas *bbw-ao-absorbs* [simp] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:
 $!!y::int. (NOT \ x) \text{ XOR } y = NOT \ (x \text{ XOR } y) \ \&$
 $x \text{ XOR } (NOT \ y) = NOT \ (x \text{ XOR } y)$
 $\langle proof \rangle$

lemma *bbw-assocs'*:
 $!!y \ z::int. (x \text{ AND } y) \text{ AND } z = x \text{ AND } (y \text{ AND } z) \ \&$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } (y \text{ OR } z) \ \&$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } (y \text{ XOR } z)$
 $\langle proof \rangle$

lemma *int-and-assoc*:
 $(x \text{ AND } y) \text{ AND } (z::int) = x \text{ AND } (y \text{ AND } z)$
 $\langle proof \rangle$

lemma *int-or-assoc*:
 $(x \text{ OR } y) \text{ OR } (z::int) = x \text{ OR } (y \text{ OR } z)$
 $\langle proof \rangle$

lemma *int-xor-assoc*:
 $(x \text{ XOR } y) \text{ XOR } (z::int) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle proof \rangle$

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [simp]:

$(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle \text{proof} \rangle$

lemma *bbw-not-dist*:

$!!y::int. \text{ NOT } (x \text{ OR } y) = (\text{NOT } x) \text{ AND } (\text{NOT } y)$
 $!!y::int. \text{ NOT } (x \text{ AND } y) = (\text{NOT } x) \text{ OR } (\text{NOT } y)$
 $\langle \text{proof} \rangle$

lemma *bbw-oa-dist*:

$!!y \ z::int. (x \text{ AND } y) \text{ OR } z =$
 $(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 $\langle \text{proof} \rangle$

lemma *bbw-ao-dist*:

$!!y \ z::int. (x \text{ OR } y) \text{ AND } z =$
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$
 $\langle \text{proof} \rangle$

lemma *plus-and-or* [rule-format]:

$ALL \ y::int. (x \text{ AND } y) + (x \text{ OR } y) = x + y$
 $\langle \text{proof} \rangle$

lemma *le-int-or*:

$!!x. \text{ bin-sign } y = \text{Int.Pls} ==> x <= x \text{ OR } y$
 $\langle \text{proof} \rangle$

lemmas *int-and-le* =

xtr3 [OF *bbw-ao-absorbs* (2) [THEN *conjunct2*, *symmetric*] *le-int-or*]

lemma *bin-nth-ops*:

$!!x \ y. \text{ bin-nth } (x \text{ AND } y) \ n = (\text{bin-nth } x \ n \ \& \ \text{bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ OR } y) \ n = (\text{bin-nth } x \ n \ | \ \text{bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ XOR } y) \ n = (\text{bin-nth } x \ n \ \sim = \text{bin-nth } y \ n)$
 $!!x. \text{ bin-nth } (\text{NOT } x) \ n = (\sim \text{bin-nth } x \ n)$
 $\langle \text{proof} \rangle$

lemma *bin-add-not*: $x + \text{NOT } x = \text{Int.Min}$

$\langle \text{proof} \rangle$

lemma *bin-trunc-ao*:

$!!x \ y. (\text{bintrunc } n \ x) \text{ AND } (\text{bintrunc } n \ y) = \text{bintrunc } n \ (x \text{ AND } y)$

$!!x\ y. (\text{bintrunc } n\ x) \text{ OR } (\text{bintrunc } n\ y) = \text{bintrunc } n\ (x \text{ OR } y)$
 $\langle \text{proof} \rangle$

lemma *bin-trunc-xor*:

$!!x\ y. \text{bintrunc } n\ (\text{bintrunc } n\ x \text{ XOR } \text{bintrunc } n\ y) =$
 $\text{bintrunc } n\ (x \text{ XOR } y)$
 $\langle \text{proof} \rangle$

lemma *bin-trunc-not*:

$!!x. \text{bintrunc } n\ (\text{NOT } (\text{bintrunc } n\ x)) = \text{bintrunc } n\ (\text{NOT } x)$
 $\langle \text{proof} \rangle$

lemma *bintr-bintr-i*:

$x = \text{bintrunc } n\ y \implies \text{bintrunc } n\ x = \text{bintrunc } n\ y$
 $\langle \text{proof} \rangle$

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

6.2 Setting and clearing bits

primrec

$\text{bin-sc} :: \text{nat} \Rightarrow \text{bit} \Rightarrow \text{int} \Rightarrow \text{int}$

where

$Z: \text{bin-sc } 0\ b\ w = \text{bin-rest } w\ \text{BIT } b$

$| \text{Suc}: \text{bin-sc } (\text{Suc } n)\ b\ w = \text{bin-sc } n\ b\ (\text{bin-rest } w)\ \text{BIT } \text{bin-last } w$

lemma *bin-nth-sc* [simp]:

$!!w. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ n = (b = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *bin-sc-sc-same* [simp]:

$!!w. \text{bin-sc } n\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ c\ w$
 $\langle \text{proof} \rangle$

lemma *bin-sc-sc-diff*:

$!!w\ m. m \sim n \implies$
 $\text{bin-sc } m\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ b\ (\text{bin-sc } m\ c\ w)$
 $\langle \text{proof} \rangle$

lemma *bin-nth-sc-gen*:

$!!w\ m. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ m = (\text{if } m = n \text{ then } b = \text{bit.B1} \text{ else } \text{bin-nth } w\ m)$
 $\langle \text{proof} \rangle$

lemma *bin-sc-nth* [simp]:

$!!w. (\text{bin-sc } n\ (\text{If } (\text{bin-nth } w\ n)\ \text{bit.B1}\ \text{bit.B0})\ w) = w$

$\langle \text{proof} \rangle$

lemma *bin-sign-sc* [simp]:

$!!w. \text{bin-sign } (\text{bin-sc } n \ b \ w) = \text{bin-sign } w$

$\langle \text{proof} \rangle$

lemma *bin-sc-bintr* [simp]:

$!!w \ m. \text{bintrunc } m \ (\text{bin-sc } n \ x \ (\text{bintrunc } m \ (w))) = \text{bintrunc } m \ (\text{bin-sc } n \ x \ w)$

$\langle \text{proof} \rangle$

lemma *bin-clr-le*:

$!!w. \text{bin-sc } n \ \text{bit.B0 } w \leq w$

$\langle \text{proof} \rangle$

lemma *bin-set-ge*:

$!!w. \text{bin-sc } n \ \text{bit.B1 } w \geq w$

$\langle \text{proof} \rangle$

lemma *bintr-bin-clr-le*:

$!!w \ m. \text{bintrunc } n \ (\text{bin-sc } m \ \text{bit.B0 } w) \leq \text{bintrunc } n \ w$

$\langle \text{proof} \rangle$

lemma *bintr-bin-set-ge*:

$!!w \ m. \text{bintrunc } n \ (\text{bin-sc } m \ \text{bit.B1 } w) \geq \text{bintrunc } n \ w$

$\langle \text{proof} \rangle$

lemma *bin-sc-FP* [simp]: $\text{bin-sc } n \ \text{bit.B0 } \text{Int.Pls} = \text{Int.Pls}$

$\langle \text{proof} \rangle$

lemma *bin-sc-TM* [simp]: $\text{bin-sc } n \ \text{bit.B1 } \text{Int.Min} = \text{Int.Min}$

$\langle \text{proof} \rangle$

lemmas *bin-sc-simps* = *bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

lemma *bin-sc-minus*:

$0 < n \implies \text{bin-sc } (\text{Suc } (n - 1)) \ b \ w = \text{bin-sc } n \ b \ w$

$\langle \text{proof} \rangle$

lemmas *bin-sc-Suc-minus* =

trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]

lemmas *bin-sc-Suc-pred* [simp] =

bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]

6.3 Operations on lists of booleans

primrec *bl-to-bin-aux* :: *bool list* \Rightarrow *int* \Rightarrow *int* **where**

Nil: *bl-to-bin-aux* [] *w* = *w*

| *Cons*: *bl-to-bin-aux* (*b* # *bs*) *w* =

$bl\text{-}to\text{-}bin\text{-}aux\ bs\ (w\ BIT\ (if\ b\ then\ bit.B1\ else\ bit.B0))$

definition $bl\text{-}to\text{-}bin :: bool\ list \Rightarrow int$ **where**
 $bl\text{-}to\text{-}bin\text{-}def : bl\text{-}to\text{-}bin\ bs = bl\text{-}to\text{-}bin\text{-}aux\ bs\ Int.Pl$

primrec $bin\text{-}to\text{-}bl\text{-}aux :: nat \Rightarrow int \Rightarrow bool\ list \Rightarrow bool\ list$ **where**
 $Z: bin\text{-}to\text{-}bl\text{-}aux\ 0\ w\ bl = bl$
 $| Suc: bin\text{-}to\text{-}bl\text{-}aux\ (Suc\ n)\ w\ bl =$
 $bin\text{-}to\text{-}bl\text{-}aux\ n\ (bin\text{-}rest\ w)\ ((bin\text{-}last\ w = bit.B1) \# bl)$

definition $bin\text{-}to\text{-}bl :: nat \Rightarrow int \Rightarrow bool\ list$ **where**
 $bin\text{-}to\text{-}bl\text{-}def : bin\text{-}to\text{-}bl\ n\ w = bin\text{-}to\text{-}bl\text{-}aux\ n\ w\ []$

primrec $bl\text{-}of\text{-}nth :: nat \Rightarrow (nat \Rightarrow bool) \Rightarrow bool\ list$ **where**
 $Suc: bl\text{-}of\text{-}nth\ (Suc\ n)\ f = f\ n \# bl\text{-}of\text{-}nth\ n\ f$
 $| Z: bl\text{-}of\text{-}nth\ 0\ f = []$

primrec $takefill :: 'a \Rightarrow nat \Rightarrow 'a\ list \Rightarrow 'a\ list$ **where**
 $Z: takefill\ fill\ 0\ xs = []$
 $| Suc: takefill\ fill\ (Suc\ n)\ xs =$
 $case\ xs\ of\ [] \Rightarrow fill \# takefill\ fill\ n\ xs$
 $| y \# ys \Rightarrow y \# takefill\ fill\ n\ ys)$

definition $map2 :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'c\ list$ **where**
 $map2\ f\ as\ bs = map\ (split\ f)\ (zip\ as\ bs)$

6.4 Splitting and concatenation

definition $bin\text{-}rcat :: nat \Rightarrow int\ list \Rightarrow int$ **where**
 $bin\text{-}rcat\ n = foldl\ (\%u\ v.\ bin\text{-}cat\ u\ n\ v)\ Int.Pl$

fun $bin\text{-}rsplit\text{-}aux :: nat \Rightarrow nat \Rightarrow int \Rightarrow int\ list \Rightarrow int\ list$ **where**
 $bin\text{-}rsplit\text{-}aux\ n\ m\ c\ bs =$
 $(if\ m = 0 \mid n = 0\ then\ bs\ else$
 $let\ (a,\ b) = bin\text{-}split\ n\ c$
 $in\ bin\text{-}rsplit\text{-}aux\ n\ (m - n)\ a\ (b \# bs))$

definition $bin\text{-}rsplit :: nat \Rightarrow nat \times int \Rightarrow int\ list$ **where**
 $bin\text{-}rsplit\ n\ w = bin\text{-}rsplit\text{-}aux\ n\ (fst\ w)\ (snd\ w)\ []$

fun $bin\text{-}rsplitl\text{-}aux :: nat \Rightarrow nat \Rightarrow int \Rightarrow int\ list \Rightarrow int\ list$ **where**
 $bin\text{-}rsplitl\text{-}aux\ n\ m\ c\ bs =$
 $(if\ m = 0 \mid n = 0\ then\ bs\ else$
 $let\ (a,\ b) = bin\text{-}split\ (min\ m\ n)\ c$
 $in\ bin\text{-}rsplitl\text{-}aux\ n\ (m - n)\ a\ (b \# bs))$

definition $bin\text{-}rsplitl :: nat \Rightarrow nat \times int \Rightarrow int\ list$ **where**
 $bin\text{-}rsplitl\ n\ w = bin\text{-}rsplitl\text{-}aux\ n\ (fst\ w)\ (snd\ w)\ []$

declare *bin-rsplit-aux.simps* [*simp del*]
declare *bin-rsplitl-aux.simps* [*simp del*]

lemma *bin-sign-cat*:
 $!!y. \text{bin-sign } (\text{bin-cat } x \ n \ y) = \text{bin-sign } x$
 $\langle \text{proof} \rangle$

lemma *bin-cat-Suc-Bit*:
 $\text{bin-cat } w \ (\text{Suc } n) \ (v \ \text{BIT } b) = \text{bin-cat } w \ n \ v \ \text{BIT } b$
 $\langle \text{proof} \rangle$

lemma *bin-nth-cat*:
 $!!n \ y. \text{bin-nth } (\text{bin-cat } x \ k \ y) \ n =$
 $(\text{if } n < k \text{ then } \text{bin-nth } y \ n \text{ else } \text{bin-nth } x \ (n - k))$
 $\langle \text{proof} \rangle$

lemma *bin-nth-split*:
 $!!b \ c. \text{bin-split } n \ c = (a, b) ==>$
 $(\text{ALL } k. \text{bin-nth } a \ k = \text{bin-nth } c \ (n + k)) \ \&$
 $(\text{ALL } k. \text{bin-nth } b \ k = (k < n \ \& \ \text{bin-nth } c \ k))$
 $\langle \text{proof} \rangle$

lemma *bin-cat-assoc*:
 $!!z. \text{bin-cat } (\text{bin-cat } x \ m \ y) \ n \ z = \text{bin-cat } x \ (m + n) \ (\text{bin-cat } y \ n \ z)$
 $\langle \text{proof} \rangle$

lemma *bin-cat-assoc-sym*: $!!z \ m.$
 $\text{bin-cat } x \ m \ (\text{bin-cat } y \ n \ z) = \text{bin-cat } (\text{bin-cat } x \ (m - n) \ y) \ (\text{min } m \ n) \ z$
 $\langle \text{proof} \rangle$

lemma *bin-cat-Pls* [*simp*]:
 $!!w. \text{bin-cat } \text{Int.Pl} \ n \ w = \text{bintrunc } n \ w$
 $\langle \text{proof} \rangle$

lemma *bintr-cat1*:
 $!!b. \text{bintrunc } (k + n) \ (\text{bin-cat } a \ n \ b) = \text{bin-cat } (\text{bintrunc } k \ a) \ n \ b$
 $\langle \text{proof} \rangle$

lemma *bintr-cat*: $\text{bintrunc } m \ (\text{bin-cat } a \ n \ b) =$
 $\text{bin-cat } (\text{bintrunc } (m - n) \ a) \ n \ (\text{bintrunc } (\text{min } m \ n) \ b)$
 $\langle \text{proof} \rangle$

lemma *bintr-cat-same* [*simp*]:
 $\text{bintrunc } n \ (\text{bin-cat } a \ n \ b) = \text{bintrunc } n \ b$
 $\langle \text{proof} \rangle$

lemma *cat-bintr* [*simp*]:
 $!!b. \text{bin-cat } a \ n \ (\text{bintrunc } n \ b) = \text{bin-cat } a \ n \ b$
 $\langle \text{proof} \rangle$

lemma *split-bintrunc*:

!!*b c. bin-split n c = (a, b) ==> b = bintrunc n c*
 <proof>

lemma *bin-cat-split*:

!!*v w. bin-split n w = (u, v) ==> w = bin-cat u n v*
 <proof>

lemma *bin-split-cat*:

!!*w. bin-split n (bin-cat v n w) = (v, bintrunc n w)*
 <proof>

lemma *bin-split-Pls* [simp]:

bin-split n Int.Pls = (Int.Pls, Int.Pls)
 <proof>

lemma *bin-split-Min* [simp]:

bin-split n Int.Min = (Int.Min, bintrunc n Int.Min)
 <proof>

lemma *bin-split-trunc*:

!!*m b c. bin-split (min m n) c = (a, b) ==>*
bin-split n (bintrunc m c) = (bintrunc (m - n) a, b)
 <proof>

lemma *bin-split-trunc1*:

!!*m b c. bin-split n c = (a, b) ==>*
bin-split n (bintrunc m c) = (bintrunc (m - n) a, bintrunc m b)
 <proof>

lemma *bin-cat-num*:

!!*b. bin-cat a n b = a * 2 ^ n + bintrunc n b*
 <proof>

lemma *bin-split-num*:

!!*b. bin-split n b = (b div 2 ^ n, b mod 2 ^ n)*
 <proof>

6.5 Miscellaneous lemmas

lemma *nth-2p-bin*:

!!*m. bin-nth (2 ^ n) m = (m = n)*
 <proof>

lemma *ex-eq-or*:

$(EX m. n = Suc m \ \& \ (m = k \mid P m)) = (n = Suc k \mid (EX m. n = Suc m \ \& \ P$

$m))$
 $\langle \text{proof} \rangle$

end

7 BinBoolList: Bool lists and integers

theory *BinBoolList*
imports *BinOperations*
begin

7.1 Arithmetic in terms of bool lists

primrec *rbl-succ* :: *bool list* => *bool list* **where**
 $Nil: \text{rbl-succ } Nil = Nil$
 $| \text{ Cons: } \text{rbl-succ } (x \# xs) = (\text{if } x \text{ then } False \# \text{rbl-succ } xs \text{ else } True \# xs)$

primrec *rbl-pred* :: *bool list* => *bool list* **where**
 $Nil: \text{rbl-pred } Nil = Nil$
 $| \text{ Cons: } \text{rbl-pred } (x \# xs) = (\text{if } x \text{ then } False \# xs \text{ else } True \# \text{rbl-pred } xs)$

primrec *rbl-add* :: *bool list* => *bool list* => *bool list* **where**

$Nil: \text{rbl-add } Nil \ x = Nil$
 $| \text{ Cons: } \text{rbl-add } (y \# ys) \ x = (\text{let } ws = \text{rbl-add } ys \ (tl \ x) \ \text{in}$
 $\quad (y \sim = hd \ x) \# (\text{if } hd \ x \ \& \ y \text{ then } \text{rbl-succ } ws \text{ else } ws))$

primrec *rbl-mult* :: *bool list* => *bool list* => *bool list* **where**

$Nil: \text{rbl-mult } Nil \ x = Nil$
 $| \text{ Cons: } \text{rbl-mult } (y \# ys) \ x = (\text{let } ws = False \# \text{rbl-mult } ys \ x \ \text{in}$
 $\quad \text{if } y \text{ then } \text{rbl-add } ws \ x \text{ else } ws)$

lemma *butlast-power*:
 $(\text{butlast } ^{\wedge} n) \ bl = \text{take } (\text{length } bl - n) \ bl$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-Pls-minus-simp* [simp]:
 $0 < n ==> \text{bin-to-bl-aux } n \ \text{Int.Pls } bl =$
 $\quad \text{bin-to-bl-aux } (n - 1) \ \text{Int.Pls } (False \# bl)$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-Min-minus-simp* [simp]:
 $0 < n ==> \text{bin-to-bl-aux } n \ \text{Int.Min } bl =$
 $\quad \text{bin-to-bl-aux } (n - 1) \ \text{Int.Min } (True \# bl)$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-Bit-minus-simp* [simp]:
 $0 < n \implies \text{bin-to-bl-aux } n \ (w \text{ BIT } b) \ \text{bl} =$
 $\text{bin-to-bl-aux } (n - 1) \ w \ ((b = \text{bit.B1}) \# \text{bl})$
 ⟨proof⟩

lemma *bin-to-bl-aux-Bit0-minus-simp* [simp]:
 $0 < n \implies \text{bin-to-bl-aux } n \ (\text{Int.Bit0 } w) \ \text{bl} =$
 $\text{bin-to-bl-aux } (n - 1) \ w \ (\text{False} \# \text{bl})$
 ⟨proof⟩

lemma *bin-to-bl-aux-Bit1-minus-simp* [simp]:
 $0 < n \implies \text{bin-to-bl-aux } n \ (\text{Int.Bit1 } w) \ \text{bl} =$
 $\text{bin-to-bl-aux } (n - 1) \ w \ (\text{True} \# \text{bl})$
 ⟨proof⟩

lemma *bl-to-bin-aux-append*:
 $\text{bl-to-bin-aux } (bs @ cs) \ w = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin-aux } bs \ w)$
 ⟨proof⟩

lemma *bin-to-bl-aux-append*:
 $\text{bin-to-bl-aux } n \ w \ bs @ cs = \text{bin-to-bl-aux } n \ w \ (bs @ cs)$
 ⟨proof⟩

lemma *bl-to-bin-append*:
 $\text{bl-to-bin } (bs @ cs) = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin } bs)$
 ⟨proof⟩

lemma *bin-to-bl-aux-alt*:
 $\text{bin-to-bl-aux } n \ w \ bs = \text{bin-to-bl } n \ w @ bs$
 ⟨proof⟩

lemma *bin-to-bl-0*: $\text{bin-to-bl } 0 \ bs = []$
 ⟨proof⟩

lemma *size-bin-to-bl-aux*:
 $\text{size } (\text{bin-to-bl-aux } n \ w \ bs) = n + \text{length } bs$
 ⟨proof⟩

lemma *size-bin-to-bl*: $\text{size } (\text{bin-to-bl } n \ w) = n$
 ⟨proof⟩

lemma *bin-bl-bin'*:
 $\text{bl-to-bin } (\text{bin-to-bl-aux } n \ w \ bs) =$
 $\text{bl-to-bin-aux } bs \ (\text{bintrunc } n \ w)$
 ⟨proof⟩

lemma *bin-bl-bin*: $\text{bl-to-bin } (\text{bin-to-bl } n \ w) = \text{bintrunc } n \ w$

$\langle \text{proof} \rangle$

lemma *bl-bin-bl'*:

$\text{bin-to-bl } (n + \text{length } bs) \text{ (bl-to-bin-aux } bs \ w) =$
 $\text{bin-to-bl-aux } n \ w \ bs$
 $\langle \text{proof} \rangle$

lemma *bl-bin-bl*: $\text{bin-to-bl } (\text{length } bs) \text{ (bl-to-bin } bs) = bs$
 $\langle \text{proof} \rangle$

declare

bin-to-bl-0 [simp]
size-bin-to-bl [simp]
bin-bl-bin [simp]
bl-bin-bl [simp]

lemma *bl-to-bin-inj*:

$\text{bl-to-bin } bs = \text{bl-to-bin } cs \implies \text{length } bs = \text{length } cs \implies bs = cs$
 $\langle \text{proof} \rangle$

lemma *bl-to-bin-False*: $\text{bl-to-bin } (\text{False} \# \text{bl}) = \text{bl-to-bin } bl$
 $\langle \text{proof} \rangle$

lemma *bl-to-bin-Nil*: $\text{bl-to-bin } [] = \text{Int.Pls}$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-Pls-aux*:

$\text{bin-to-bl-aux } n \ \text{Int.Pls } bl = \text{replicate } n \ \text{False} \ @ \ bl$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-Pls*: $\text{bin-to-bl } n \ \text{Int.Pls} = \text{replicate } n \ \text{False}$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-Min-aux* [rule-format] :

$\text{ALL } bl. \text{bin-to-bl-aux } n \ \text{Int.Min } bl = \text{replicate } n \ \text{True} \ @ \ bl$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-Min*: $\text{bin-to-bl } n \ \text{Int.Min} = \text{replicate } n \ \text{True}$
 $\langle \text{proof} \rangle$

lemma *bl-to-bin-rep-F*:

$\text{bl-to-bin } (\text{replicate } n \ \text{False} \ @ \ bl) = \text{bl-to-bin } bl$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-trunc*:

$n \leq m \implies \text{bin-to-bl } n \ (\text{bintrunc } m \ w) = \text{bin-to-bl } n \ w$
 $\langle \text{proof} \rangle$

declare

bin-to-bl-trunc [simp]
bl-to-bin-False [simp]
bl-to-bin-Nil [simp]

lemma *bin-to-bl-aux-bintr* [rule-format] :
 ALL *m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =*
replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl
 ⟨proof⟩

lemmas *bin-to-bl-bintr* =
bin-to-bl-aux-bintr [where *bl* = [], folded *bin-to-bl-def*]

lemma *bl-to-bin-rep-False*: *bl-to-bin (replicate n False) = Int.Pls*
 ⟨proof⟩

lemma *len-bin-to-bl-aux*:
length (bin-to-bl-aux n w bs) = n + length bs
 ⟨proof⟩

lemma *len-bin-to-bl* [simp]: *length (bin-to-bl n w) = n*
 ⟨proof⟩

lemma *sign-bl-bin'*:
bin-sign (bl-to-bin-aux bs w) = bin-sign w
 ⟨proof⟩

lemma *sign-bl-bin*: *bin-sign (bl-to-bin bs) = Int.Pls*
 ⟨proof⟩

lemma *bl-sbin-sign-aux*:
hd (bin-to-bl-aux (Suc n) w bs) =
(bin-sign (sbintrunc n w) = Int.Min)
 ⟨proof⟩

lemma *bl-sbin-sign*:
hd (bin-to-bl (Suc n) w) = (bin-sign (sbintrunc n w) = Int.Min)
 ⟨proof⟩

lemma *bin-nth-of-bl-aux* [rule-format]:
 $\forall w. \text{bin-nth } (\text{bl-to-bin-aux } \text{bl } w) \text{ } n =$
 $(n < \text{size bl} \ \& \ \text{rev bl ! } n \mid n \geq \text{length bl} \ \& \ \text{bin-nth } w \ (n - \text{size bl}))$
 ⟨proof⟩

lemma *bin-nth-of-bl*: *bin-nth (bl-to-bin bl) n = (n < length bl & rev bl ! n)*
 ⟨proof⟩

lemma *bin-nth-bl* [rule-format] : ALL *m w. n < m -->*
bin-nth w n = nth (rev (bin-to-bl m w)) n
 ⟨proof⟩

lemma *nth-rev* [rule-format] :

$n < \text{length } xs \dashrightarrow \text{rev } xs ! n = xs ! (\text{length } xs - 1 - n)$
 ⟨proof⟩

lemmas *nth-rev-alt* = *nth-rev* [where $xs = \text{rev } ys$, simplified, standard]

lemma *nth-bin-to-bl-aux* [rule-format] :

ALL $w \ n \ bl. n < m + \text{length } bl \dashrightarrow (\text{bin-to-bl-aux } m \ w \ bl) ! n =$
 (if $n < m$ then $\text{bin-nth } w \ (m - 1 - n)$ else $bl ! (n - m)$)
 ⟨proof⟩

lemma *nth-bin-to-bl*: $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p-aux* [rule-format]:

$\forall w. \text{bl-to-bin-aux } bs \ w < (w + 1) * (2 ^ \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p*: $\text{bl-to-bin } bs < (2 ^ \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge2p-aux* [rule-format] :

$\forall w. \text{bl-to-bin-aux } bs \ w \geq w * (2 ^ \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge0*: $\text{bl-to-bin } bs \geq 0$
 ⟨proof⟩

lemma *butlast-rest-bin*:

$\text{butlast } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } (n - 1) \ (\text{bin-rest } w)$
 ⟨proof⟩

lemmas *butlast-bin-rest* = *butlast-rest-bin*

[where $w = \text{bl-to-bin } bl$ and $n = \text{length } bl$, simplified, standard]

lemma *butlast-rest-bl2bin-aux*:

$bl \sim [] \implies$
 $\text{bl-to-bin-aux } (\text{butlast } bl) \ w = \text{bin-rest } (\text{bl-to-bin-aux } bl \ w)$
 ⟨proof⟩

lemma *butlast-rest-bl2bin*:

$\text{bl-to-bin } (\text{butlast } bl) = \text{bin-rest } (\text{bl-to-bin } bl)$
 ⟨proof⟩

lemma *trunc-bl2bin-aux* [rule-format]:

ALL $w. \text{bintrunc } m \ (\text{bl-to-bin-aux } bl \ w) =$
 $\text{bl-to-bin-aux } (\text{drop } (\text{length } bl - m) \ bl) \ (\text{bintrunc } (m - \text{length } bl) \ w)$
 ⟨proof⟩

lemma *trunc-bl2bin*:

$\text{bintrunc } m \text{ (bl-to-bin bl)} = \text{bl-to-bin (drop (length bl - m) bl)}$
 $\langle \text{proof} \rangle$

lemmas *trunc-bl2bin-len [simp]* =

trunc-bl2bin [of length bl bl, simplified, standard]

lemma *bl2bin-drop*:

$\text{bl-to-bin (drop k bl)} = \text{bintrunc (length bl - k) (bl-to-bin bl)}$
 $\langle \text{proof} \rangle$

lemma *nth-rest-power-bin [rule-format]* :

ALL n. bin-nth ((bin-rest ^ k) w) n = bin-nth w (n + k)
 $\langle \text{proof} \rangle$

lemma *take-rest-power-bin*:

$m \leq n \implies \text{take } m \text{ (bin-to-bl } n \text{ w)} = \text{bin-to-bl } m \text{ ((bin-rest ^ (n - m)) w)}$
 $\langle \text{proof} \rangle$

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd (butlast xs)} = \text{hd } xs$

$\langle \text{proof} \rangle$

lemma *last-bin-last'*:

$\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last (bl-to-bin-aux xs w)} = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *last-bin-last*:

$\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last (bl-to-bin xs)} = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *bin-last-last*:

$\text{bin-last } w = (\text{if last (bin-to-bl (Suc n) w) then bit.B1 else bit.B0})$
 $\langle \text{proof} \rangle$

lemma *map2-Nil [simp]*: $\text{map2 } f \ [] \ ys = []$

$\langle \text{proof} \rangle$

lemma *map2-Cons [simp]*:

$\text{map2 } f \ (x \# xs) \ (y \# ys) = f \ x \ y \ \# \ \text{map2 } f \ xs \ ys$
 $\langle \text{proof} \rangle$

lemma *bl-xor-aux-bin [rule-format]* : *ALL v w bs cs.*

$\text{map2 } (\%x \ y. x \sim y) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ XOR \ w) \ (\text{map2 } (\%x \ y. x \sim y) \ bs \ cs)$
 $\langle \text{proof} \rangle$

lemma *bl-or-aux-bin* [rule-format] : *ALL v w bs cs.*

$$\text{map2 } (op \mid) (bin\text{-}to\text{-}bl\text{-}aux \ n \ v \ bs) (bin\text{-}to\text{-}bl\text{-}aux \ n \ w \ cs) =$$

$$bin\text{-}to\text{-}bl\text{-}aux \ n \ (v \ OR \ w) (\text{map2 } (op \mid) \ bs \ cs)$$
<proof>

lemma *bl-and-aux-bin* [rule-format] : *ALL v w bs cs.*

$$\text{map2 } (op \ \&) (bin\text{-}to\text{-}bl\text{-}aux \ n \ v \ bs) (bin\text{-}to\text{-}bl\text{-}aux \ n \ w \ cs) =$$

$$bin\text{-}to\text{-}bl\text{-}aux \ n \ (v \ AND \ w) (\text{map2 } (op \ \&) \ bs \ cs)$$
<proof>

lemma *bl-not-aux-bin* [rule-format] :
ALL w cs. map Not (bin-to-bl-aux n w cs) =

$$bin\text{-}to\text{-}bl\text{-}aux \ n \ (NOT \ w) (\text{map } Not \ cs)$$
<proof>

lemmas *bl-not-bin = bl-not-aux-bin*
 [where *cs = []*, *unfolded bin-to-bl-def* [symmetric] *map.simps*]

lemmas *bl-and-bin = bl-and-aux-bin* [where *bs=[]* and *cs=[]*,
unfolded map2-Nil, folded bin-to-bl-def]

lemmas *bl-or-bin = bl-or-aux-bin* [where *bs=[]* and *cs=[]*,
unfolded map2-Nil, folded bin-to-bl-def]

lemmas *bl-xor-bin = bl-xor-aux-bin* [where *bs=[]* and *cs=[]*,
unfolded map2-Nil, folded bin-to-bl-def]

lemma *drop-bin2bl-aux* [rule-format] :
ALL m bin bs. drop m (bin-to-bl-aux n bin bs) =

$$bin\text{-}to\text{-}bl\text{-}aux \ (n - m) \ bin \ (drop \ (m - n) \ bs)$$
<proof>

lemma *drop-bin2bl*: *drop m (bin-to-bl n bin) = bin-to-bl (n - m) bin*
<proof>

lemma *take-bin2bl-lem1* [rule-format] :
ALL w bs. take m (bin-to-bl-aux m w bs) = bin-to-bl m w
<proof>

lemma *take-bin2bl-lem* [rule-format] :
ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =

$$take \ m \ (bin\text{-}to\text{-}bl \ (m + n) \ w)$$
<proof>

lemma *bin-split-take* [rule-format] :
ALL b c. bin-split n c = (a, b) -->

$$bin\text{-}to\text{-}bl \ m \ a = take \ m \ (bin\text{-}to\text{-}bl \ (m + n) \ c)$$
<proof>

lemma *bin-split-take1*:

$k = m + n \implies \text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-to-bl } m \ a = \text{take } m \ (\text{bin-to-bl } k \ c)$
 $\langle \text{proof} \rangle$

lemma *nth-takefill* [rule-format] : $ALL \ m \ l. \ m < n \implies$

$\text{takefill fill } n \ l \ ! \ m = (\text{if } m < \text{length } l \text{ then } l \ ! \ m \text{ else fill})$
 $\langle \text{proof} \rangle$

lemma *takefill-alt* [rule-format] :

$ALL \ l. \ \text{takefill fill } n \ l = \text{take } n \ l \ @ \ \text{replicate } (n - \text{length } l) \ \text{fill}$
 $\langle \text{proof} \rangle$

lemma *takefill-replicate* [simp]:

$\text{takefill fill } n \ (\text{replicate } m \ \text{fill}) = \text{replicate } n \ \text{fill}$
 $\langle \text{proof} \rangle$

lemma *takefill-le'* [rule-format] :

$ALL \ l \ n. \ n = m + k \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
 $\langle \text{proof} \rangle$

lemma *length-takefill* [simp]: $\text{length } (\text{takefill fill } n \ l) = n$

$\langle \text{proof} \rangle$

lemma *take-takefill'*:

$!!w \ n. \ n = k + m \implies \text{take } k \ (\text{takefill fill } n \ w) = \text{takefill fill } k \ w$
 $\langle \text{proof} \rangle$

lemma *drop-takefill*:

$!!w. \ \text{drop } k \ (\text{takefill fill } (m + k) \ w) = \text{takefill fill } m \ (\text{drop } k \ w)$
 $\langle \text{proof} \rangle$

lemma *takefill-le* [simp]:

$m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
 $\langle \text{proof} \rangle$

lemma *take-takefill* [simp]:

$m \leq n \implies \text{take } m \ (\text{takefill fill } n \ w) = \text{takefill fill } m \ w$
 $\langle \text{proof} \rangle$

lemma *takefill-append*:

$\text{takefill fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill fill } m \ w)$
 $\langle \text{proof} \rangle$

lemma *takefill-same'*:

$l = \text{length } xs \implies \text{takefill fill } l \ xs = xs$
 $\langle \text{proof} \rangle$

lemmas *takefill-same* [simp] = *takefill-same'* [OF refl]

lemma *takefill-bintrunc*:

takefill False n bl = rev (bin-to-bl n (bl-to-bin (rev bl)))
<proof>

lemma *bl-bin-bl-rtf*:

bin-to-bl n (bl-to-bin bl) = rev (takefill False n (rev bl))
<proof>

lemmas *bl-bin-bl-rep-drop =*

bl-bin-bl-rtf [simplified takefill-alt,
simplified, simplified rev-take, simplified]

lemma *tf-rev*:

n + k = m + length bl ==> takefill x m (rev (takefill y n bl)) =
rev (takefill y m (rev (takefill x k (rev bl))))
<proof>

lemma *takefill-minus*:

0 < n ==> takefill fill (Suc (n - 1)) w = takefill fill n w
<proof>

lemmas *takefill-Suc-cases =*

list.cases [THEN takefill.Suc [THEN trans], standard]

lemmas *takefill-Suc-Nil = takefill-Suc-cases (1)*

lemmas *takefill-Suc-Cons = takefill-Suc-cases (2)*

lemmas *takefill-minus-simps = takefill-Suc-cases [THEN [2]*

takefill-minus [symmetric, THEN trans], standard]

lemmas *takefill-pred-simps [simp] =*

takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]

lemma *bl-to-bin-aux-cat*:

!!nv v. bl-to-bin-aux bs (bin-cat w nv v) =
bin-cat w (nv + length bs) (bl-to-bin-aux bs v)
<proof>

lemma *bin-to-bl-aux-cat*:

!!w bs. bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs =
bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)
<proof>

lemmas *bl-to-bin-aux-alt =*

bl-to-bin-aux-cat [where nv = 0 and v = Int.Pl,
simplified bl-to-bin-def [symmetric], simplified]

lemmas *bin-to-bl-cat* =
bin-to-bl-aux-cat [where *bs* = [], folded *bin-to-bl-def*]
lemmas *bl-to-bin-aux-app-cat* =
trans [OF *bl-to-bin-aux-append bl-to-bin-aux-alt*]
lemmas *bin-to-bl-aux-cat-app* =
trans [OF *bin-to-bl-aux-cat bin-to-bl-aux-alt*]
lemmas *bl-to-bin-app-cat* = *bl-to-bin-aux-app-cat*
[where *w* = *Int.Pls*, folded *bl-to-bin-def*]
lemmas *bin-to-bl-cat-app* = *bin-to-bl-aux-cat-app*
[where *bs* = [], folded *bin-to-bl-def*]

lemma *bl-to-bin-app-cat-alt*:
bin-cat (*bl-to-bin cs*) *n w* = *bl-to-bin* (*cs @ bin-to-bl n w*)
⟨*proof*⟩

lemma *mask-lem*: (*bl-to-bin* (*True # replicate n False*)) =
Int.succ (*bl-to-bin* (*replicate n True*))
⟨*proof*⟩

lemma *length-bl-of-nth* [*simp*]: *length* (*bl-of-nth n f*) = *n*
⟨*proof*⟩

lemma *nth-bl-of-nth* [*simp*]:
m < *n* \implies *rev* (*bl-of-nth n f*) ! *m* = *f m*
⟨*proof*⟩

lemma *bl-of-nth-inj*:
(!*k*. *k* < *n* \implies *f k* = *g k*) \implies *bl-of-nth n f* = *bl-of-nth n g*
⟨*proof*⟩

lemma *bl-of-nth-nth-le* [*rule-format*] : ALL *xs*.
length xs $\geq n \implies$ *bl-of-nth n* (*nth* (*rev xs*)) = *drop* (*length xs* - *n*) *xs*
⟨*proof*⟩

lemmas *bl-of-nth-nth* [*simp*] = *order-refl* [THEN *bl-of-nth-nth-le*, *simplified*]

lemma *size-rbl-pred*: *length* (*rbl-pred bl*) = *length bl*
⟨*proof*⟩

lemma *size-rbl-succ*: *length* (*rbl-succ bl*) = *length bl*
⟨*proof*⟩

lemma *size-rbl-add*:

!!*cl*. *length* (*rbl-add* *bl cl*) = *length* *bl*
 ⟨*proof*⟩

lemma *size-rbl-mult*:

!!*cl*. *length* (*rbl-mult* *bl cl*) = *length* *bl*
 ⟨*proof*⟩

lemmas *rbl-sizes* [*simp*] =

size-rbl-pred *size-rbl-succ* *size-rbl-add* *size-rbl-mult*

lemmas *rbl-Nils* =

rbl-pred.Nil *rbl-succ.Nil* *rbl-add.Nil* *rbl-mult.Nil*

lemma *rbl-pred*:

!!*bin*. *rbl-pred* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Int.pred* *bin*))
 ⟨*proof*⟩

lemma *rbl-succ*:

!!*bin*. *rbl-succ* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Int.succ* *bin*))
 ⟨*proof*⟩

lemma *rbl-add*:

!!*bina binb*. *rbl-add* (*rev* (*bin-to-bl* *n bina*)) (*rev* (*bin-to-bl* *n binb*)) =
rev (*bin-to-bl* *n* (*bina* + *binb*))
 ⟨*proof*⟩

lemma *rbl-add-app2*:

!!*blb*. *length* *blb* >= *length* *bla* ==>
rbl-add *bla* (*blb* @ *blc*) = *rbl-add* *bla* *blb*
 ⟨*proof*⟩

lemma *rbl-add-take2*:

!!*blb*. *length* *blb* >= *length* *bla* ==>
rbl-add *bla* (*take* (*length* *bla*) *blb*) = *rbl-add* *bla* *blb*
 ⟨*proof*⟩

lemma *rbl-add-long*:

m >= *n* ==> *rbl-add* (*rev* (*bin-to-bl* *n bina*)) (*rev* (*bin-to-bl* *m binb*)) =
rev (*bin-to-bl* *n* (*bina* + *binb*))
 ⟨*proof*⟩

lemma *rbl-mult-app2*:

!!*blb*. *length* *blb* >= *length* *bla* ==>
rbl-mult *bla* (*blb* @ *blc*) = *rbl-mult* *bla* *blb*
 ⟨*proof*⟩

lemma *rbl-mult-take2*:

length *blb* >= *length* *bla* ==>

$rbl_mult\ bla\ (take\ (length\ bla)\ bbl) = rbl_mult\ bla\ bbl$
 $\langle proof \rangle$

lemma *rbl-mult-gt1*:

$m \geq length\ bl \implies rbl_mult\ bl\ (rev\ (bin_to_bl\ m\ binb)) =$
 $rbl_mult\ bl\ (rev\ (bin_to_bl\ (length\ bl)\ binb))$
 $\langle proof \rangle$

lemma *rbl-mult-gt*:

$m > n \implies rbl_mult\ (rev\ (bin_to_bl\ n\ bina))\ (rev\ (bin_to_bl\ m\ binb)) =$
 $rbl_mult\ (rev\ (bin_to_bl\ n\ bina))\ (rev\ (bin_to_bl\ n\ binb))$
 $\langle proof \rangle$

lemmas *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

lemma *rbbL-Cons*:

$b \# rev\ (bin_to_bl\ n\ x) = rev\ (bin_to_bl\ (Suc\ n)\ (x\ BIT\ If\ b\ bit.B1\ bit.B0))$
 $\langle proof \rangle$

lemma *rbl-mult: !!bina binb.*

$rbl_mult\ (rev\ (bin_to_bl\ n\ bina))\ (rev\ (bin_to_bl\ n\ binb)) =$
 $rev\ (bin_to_bl\ n\ (bina * binb))$
 $\langle proof \rangle$

lemma *rbl-add-split*:

$P\ (rbl_add\ (y\ \# \ ys)\ (x\ \# \ xs)) =$
 $(ALL\ ws.\ length\ ws = length\ ys \longrightarrow ws = rbl_add\ ys\ xs \longrightarrow$
 $(y \longrightarrow ((x \longrightarrow P\ (False\ \# \ rbl_succ\ ws)) \ \&\ (\sim x \longrightarrow P\ (True\ \# \ ws)))) \ \&$
 $(\sim y \longrightarrow P\ (x\ \# \ ws)))$
 $\langle proof \rangle$

lemma *rbl-mult-split*:

$P\ (rbl_mult\ (y\ \# \ ys)\ xs) =$
 $(ALL\ ws.\ length\ ws = Suc\ (length\ ys) \longrightarrow ws = False\ \# \ rbl_mult\ ys\ xs \longrightarrow$
 $(y \longrightarrow P\ (rbl_add\ ws\ xs)) \ \&\ (\sim y \longrightarrow P\ ws))$
 $\langle proof \rangle$

lemma *and-len*: $xs = ys \implies xs = ys \ \&\ length\ xs = length\ ys$

$\langle proof \rangle$

lemma *size-if*: $size\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ size\ xs\ else\ size\ ys)$

$\langle proof \rangle$

lemma *tl-if*: $tl\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ tl\ xs\ else\ tl\ ys)$

$\langle proof \rangle$

lemma *hd-if*: $hd\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ hd\ xs\ else\ hd\ ys)$

$\langle proof \rangle$

lemma *if-Not-x*: $(\text{if } p \text{ then } \sim x \text{ else } x) = (p = (\sim x))$
 $\langle \text{proof} \rangle$

lemma *if-x-Not*: $(\text{if } p \text{ then } x \text{ else } \sim x) = (p = x)$
 $\langle \text{proof} \rangle$

lemma *if-same-and*: $(\text{If } p \ x \ y \ \& \ \text{If } p \ u \ v) = (\text{if } p \text{ then } x \ \& \ u \text{ else } y \ \& \ v)$
 $\langle \text{proof} \rangle$

lemma *if-same-eq*: $(\text{If } p \ x \ y = (\text{If } p \ u \ v)) = (\text{if } p \text{ then } x = (u) \text{ else } y = (v))$
 $\langle \text{proof} \rangle$

lemma *if-same-eq-not*:
 $(\text{If } p \ x \ y = (\sim \text{If } p \ u \ v)) = (\text{if } p \text{ then } x = (\sim u) \text{ else } y = (\sim v))$
 $\langle \text{proof} \rangle$

lemma *if-Cons*: $(\text{if } p \text{ then } x \ \# \ xs \text{ else } y \ \# \ ys) = \text{If } p \ x \ y \ \# \ \text{If } p \ xs \ ys$
 $\langle \text{proof} \rangle$

lemma *if-single*:
 $(\text{if } xc \text{ then } [xab] \text{ else } [an]) = [\text{if } xc \text{ then } xab \text{ else } an]$
 $\langle \text{proof} \rangle$

lemma *if-bool-simps*:
 $\text{If } p \ \text{True} \ y = (p \mid y) \ \& \ \text{If } p \ \text{False} \ y = (\sim p \ \& \ y) \ \& \$
 $\text{If } p \ y \ \text{True} = (p \dashrightarrow y) \ \& \ \text{If } p \ y \ \text{False} = (p \ \& \ y)$
 $\langle \text{proof} \rangle$

lemmas *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *seqr* = *eq-reflection* [where $x = \text{size } w, \text{ standard}$]

lemmas *tl-Nil* = *tl.simps* (1)
lemmas *tl-Cons* = *tl.simps* (2)

7.2 Repeated splitting or concatenation

lemma *sclem*:
 $\text{size } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ xs)) = \text{length } xs * n$
 $\langle \text{proof} \rangle$

lemma *bin-cat-foldl-lem* [rule-format] :
 $\text{ALL } x. \text{foldl } (\%u. \text{bin-cat } u \ n) \ x \ xs =$
 $\text{bin-cat } x \ (\text{size } xs * n) \ (\text{foldl } (\%u. \text{bin-cat } u \ n) \ y \ xs)$
 $\langle \text{proof} \rangle$

lemma *bin-rcat-bl*:

$(\text{bin-rcat } n \text{ } wl) = \text{bl-to-bin } (\text{concat } (\text{map } (\text{bin-to-bl } n) \text{ } wl))$
 $\langle \text{proof} \rangle$

lemmas $\text{bin-rsplit-aux-simps} = \text{bin-rsplit-aux.simps } \text{bin-rsplittl-aux.simps}$

lemmas $\text{rsplit-aux-simps} = \text{bin-rsplit-aux-simps}$

lemmas $\text{th-if-simp1} = \text{split-if } [\text{where } P = \text{op} = l,$
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct1}, \text{ THEN } \text{mp}, \text{ standard}]$

lemmas $\text{th-if-simp2} = \text{split-if } [\text{where } P = \text{op} = l,$
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct2}, \text{ THEN } \text{mp}, \text{ standard}]$

lemmas $\text{rsplit-aux-simp1s} = \text{rsplit-aux-simps } [\text{THEN } \text{th-if-simp1}]$

lemmas $\text{rsplit-aux-simp2ls} = \text{rsplit-aux-simps } [\text{THEN } \text{th-if-simp2}]$

lemmas $\text{bin-rsplit-aux-simp2s } [\text{simp}] = \text{rsplit-aux-simp2ls } [\text{unfolded Let-def}]$

lemmas $\text{rbscl} = \text{bin-rsplit-aux-simp2s } (2)$

lemmas $\text{rsplit-aux-0-simps } [\text{simp}] =$
 $\text{rsplit-aux-simp1s } [\text{OF } \text{disjI1}] \text{ rsplit-aux-simp1s } [\text{OF } \text{disjI2}]$

lemma $\text{bin-rsplit-aux-append}:$
 $\text{bin-rsplit-aux } n \text{ } m \text{ } c \text{ } (bs \text{ @ } cs) = \text{bin-rsplit-aux } n \text{ } m \text{ } c \text{ } bs \text{ @ } cs$
 $\langle \text{proof} \rangle$

lemma $\text{bin-rsplittl-aux-append}:$
 $\text{bin-rsplittl-aux } n \text{ } m \text{ } c \text{ } (bs \text{ @ } cs) = \text{bin-rsplittl-aux } n \text{ } m \text{ } c \text{ } bs \text{ @ } cs$
 $\langle \text{proof} \rangle$

lemmas $\text{rsplit-aux-apps } [\text{where } bs = []] =$
 $\text{bin-rsplit-aux-append } \text{bin-rsplittl-aux-append}$

lemmas $\text{rsplit-def-auxs} = \text{bin-rsplit-def } \text{bin-rsplittl-def}$

lemmas $\text{rsplit-aux-alts} = \text{rsplit-aux-apps}$
 $[\text{unfolded append-Nil } \text{rsplit-def-auxs } [\text{symmetric}]]$

lemma $\text{bin-split-minus}: 0 < n ==> \text{bin-split } (\text{Suc } (n - 1)) \text{ } w = \text{bin-split } n \text{ } w$
 $\langle \text{proof} \rangle$

lemmas $\text{bin-split-minus-simp} =$
 $\text{bin-split.Suc } [\text{THEN } [2] \text{ bin-split-minus } [\text{symmetric}, \text{ THEN } \text{trans}], \text{ standard}]$

lemma $\text{bin-split-pred-simp } [\text{simp}]:$
 $(0::\text{nat}) < \text{number-of bin} \implies$
 $\text{bin-split } (\text{number-of bin}) \text{ } w =$
 $(\text{let } (w1, w2) = \text{bin-split } (\text{number-of } (\text{Int.pred bin})) \text{ } (\text{bin-rest } w)$
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$
 $\langle \text{proof} \rangle$

declare *bin-split-pred-simp* [*simp*]

lemma *bin-rsplit-aux-simp-alt*:

bin-rsplit-aux *n m c bs* =
 (if *m* = 0 \vee *n* = 0
 then *bs*
 else let (*a*, *b*) = *bin-split* *n c* in *bin-rsplit* *n* (*m* - *n*, *a*) @ *b* # *bs*)
 ⟨proof⟩

lemmas *bin-rsplit-simp-alt* =

trans [*OF bin-rsplit-def*
bin-rsplit-aux-simp-alt, *standard*]

lemmas *bthrs* = *bin-rsplit-simp-alt* [*THEN* [2] *trans*]

lemma *bin-rsplit-size-sign'* [*rule-format*] :

n > 0 ==> (ALL *nw w*. *rev sw* = *bin-rsplit* *n* (*nw*, *w*) -->
 (ALL *v*: set *sw*. *bintrunc* *n v* = *v*))
 ⟨proof⟩

lemmas *bin-rsplit-size-sign* = *bin-rsplit-size-sign'* [*OF asm-rl*

rev-rev-ident [*THEN trans*] *set-rev* [*THEN equalityD2* [*THEN subsetD*]],
standard]

lemma *bin-nth-rsplit* [*rule-format*] :

n > 0 ==> *m* < *n* ==> (ALL *w k nw*. *rev sw* = *bin-rsplit* *n* (*nw*, *w*) -->
k < size *sw* --> *bin-nth* (*sw* ! *k*) *m* = *bin-nth* *w* (*k* * *n* + *m*))
 ⟨proof⟩

lemma *bin-rsplit-all*:

0 < *nw* ==> *nw* <= *n* ==> *bin-rsplit* *n* (*nw*, *w*) = [*bintrunc* *n w*]
 ⟨proof⟩

lemma *bin-rsplit-l* [*rule-format*] :

ALL *bin*. *bin-rsplitl* *n* (*m*, *bin*) = *bin-rsplit* *n* (*m*, *bintrunc* *m bin*)
 ⟨proof⟩

lemma *bin-rsplit-rcat* [*rule-format*] :

n > 0 --> *bin-rsplit* *n* (*n* * size *ws*, *bin-rcat* *n ws*) = map (*bintrunc* *n*) *ws*
 ⟨proof⟩

lemma *bin-rsplit-aux-len-le* [*rule-format*] :

\forall *ws m*. *n* \neq 0 \longrightarrow *ws* = *bin-rsplit-aux* *n nw w bs* \longrightarrow
length ws \leq *m* \longleftrightarrow *nw* + *length bs* * *n* \leq *m* * *n*
 ⟨proof⟩

lemma *bin-rsplit-len-le*:

n \neq 0 --> *ws* = *bin-rsplit* *n* (*nw*, *w*) --> (*length ws* <= *m*) = (*nw* <= *m* *

n)
 $\langle proof \rangle$

lemma *bin-rsplit-aux-len* [rule-format] :
 $n \neq 0 \rightarrow \text{length } (\text{bin-rsplit-aux } n \text{ } nw \text{ } w \text{ } cs) =$
 $(nw + n - 1) \text{ div } n + \text{length } cs$
 $\langle proof \rangle$

lemma *bin-rsplit-len*:
 $n \neq 0 \Rightarrow \text{length } (\text{bin-rsplit } n \text{ } (nw, w)) = (nw + n - 1) \text{ div } n$
 $\langle proof \rangle$

lemma *bin-rsplit-aux-len-indep*:
 $n \neq 0 \Rightarrow \text{length } bs = \text{length } cs \Rightarrow$
 $\text{length } (\text{bin-rsplit-aux } n \text{ } nw \text{ } v \text{ } bs) =$
 $\text{length } (\text{bin-rsplit-aux } n \text{ } nw \text{ } w \text{ } cs)$
 $\langle proof \rangle$

lemma *bin-rsplit-len-indep*:
 $n \neq 0 \Rightarrow \text{length } (\text{bin-rsplit } n \text{ } (nw, v)) = \text{length } (\text{bin-rsplit } n \text{ } (nw, w))$
 $\langle proof \rangle$

end

8 TdThs: Type Definition Theorems

theory *TdThs*
imports *Main*
begin

9 More lemmas about normal type definitions

lemma
 $tdD1: \text{type-definition } Rep \text{ } Abs \text{ } A \Rightarrow \forall x. Rep \text{ } x \in A \text{ and}$
 $tdD2: \text{type-definition } Rep \text{ } Abs \text{ } A \Rightarrow \forall x. Abs \text{ } (Rep \text{ } x) = x \text{ and}$
 $tdD3: \text{type-definition } Rep \text{ } Abs \text{ } A \Rightarrow \forall y. y \in A \rightarrow Rep \text{ } (Abs \text{ } y) = y$
 $\langle proof \rangle$

lemma *td-nat-int*:
 $\text{type-definition } int \text{ } nat \text{ } (Collect \text{ } (op \leq 0))$
 $\langle proof \rangle$

context *type-definition*
begin

lemmas $Rep' \text{ } [iff] = Rep \text{ } [simplified]$

declare *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

lemma *Abs-eqD*: $Abs\ x = Abs\ y \implies x \in A \implies y \in A \implies x = y$
 $\langle proof \rangle$

lemma *Abs-inverse'*:
 $r : A \implies Abs\ r = a \implies Rep\ a = r$
 $\langle proof \rangle$

lemma *Rep-comp-inverse*:
 $Rep\ o\ f = g \implies Abs\ o\ g = f$
 $\langle proof \rangle$

lemma *Rep-eqD* [*elim!*]: $Rep\ x = Rep\ y \implies x = y$
 $\langle proof \rangle$

lemma *Rep-inverse'*: $Rep\ a = r \implies Abs\ r = a$
 $\langle proof \rangle$

lemma *comp-Abs-inverse*:
 $f\ o\ Abs = g \implies g\ o\ Rep = f$
 $\langle proof \rangle$

lemma *set-Rep*:
 $A = range\ Rep$
 $\langle proof \rangle$

lemma *set-Rep-Abs*: $A = range\ (Rep\ o\ Abs)$
 $\langle proof \rangle$

lemma *Abs-inj-on*: *inj-on* *Abs* *A*
 $\langle proof \rangle$

lemma *image*: $Abs\ ` A = UNIV$
 $\langle proof \rangle$

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:
 $Rep\ o\ fa = fr\ o\ Rep \mid fa\ o\ Abs = Abs\ o\ fr \implies Abs\ o\ fr\ o\ Rep = fa$
 $\langle proof \rangle$

lemmas *fns1a* = *disjI1* [*THEN* *fns1*]

lemmas *fns1b* = *disjI2* [*THEN* *fns1*]

lemma *fns4*:
 $Rep\ o\ fa\ o\ Abs = fr \implies$
 $Rep\ o\ fa = fr\ o\ Rep \ \&\ fa\ o\ Abs = Abs\ o\ fr$
 $\langle proof \rangle$

end

interpretation *nat-int*: *type-definition int nat Collect (op <= 0)*
<proof>

declare

nat-int.Rep-cases [*cases del*]
nat-int.Abs-cases [*cases del*]
nat-int.Rep-induct [*induct del*]
nat-int.Abs-induct [*induct del*]

9.1 Extended form of type definition predicate

lemma *td-conds*:

norm o norm = norm ==> (fr o norm = norm o fr) =
(norm o fr o norm = fr o norm & norm o fr o norm = norm o fr)
<proof>

lemma *fn-comm-power*:

fa o tr = tr o fr ==> fa ^^ n o tr = tr o fr ^^ n
<proof>

lemmas *fn-comm-power'* =

ext [THEN fn-comm-power, THEN fun-cong, unfolded o-def, standard]

locale *td-ext* = *type-definition* +

fixes *norm*

assumes *eq-norm*: $\bigwedge x. \text{Rep } (\text{Abs } x) = \text{norm } x$

begin

lemma *Abs-norm* [*simp*]:

Abs (norm x) = Abs x
<proof>

lemma *td-th*:

g o Abs = f ==> f (Rep x) = g x
<proof>

lemma *eq-norm'*: *Rep o Abs = norm*

<proof>

lemma *norm-Rep* [*simp*]: *norm (Rep x) = Rep x*

<proof>

lemmas *td* = *td-thm*

lemma *set-iff-norm*: *w : A <-> w = norm w*

$\langle \text{proof} \rangle$

lemma *inverse-norm*:

$(\text{Abs } n = w) = (\text{Rep } w = \text{norm } n)$
 $\langle \text{proof} \rangle$

lemma *norm-eq-iff*:

$(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$
 $\langle \text{proof} \rangle$

lemma *norm-comps*:

$\text{Abs } o \text{ norm} = \text{Abs}$
 $\text{norm } o \text{ Rep} = \text{Rep}$
 $\text{norm } o \text{ norm} = \text{norm}$
 $\langle \text{proof} \rangle$

lemmas *norm-norm* [*simp*] = *norm-comps*

lemma *fns5*:

$\text{Rep } o \text{ fa } o \text{ Abs} = \text{fr} ==>$
 $\text{fr } o \text{ norm} = \text{fr} \ \& \ \text{norm } o \text{ fr} = \text{fr}$
 $\langle \text{proof} \rangle$

lemma *fns2*:

$\text{Abs } o \text{ fr } o \text{ Rep} = \text{fa} ==>$
 $(\text{norm } o \text{ fr } o \text{ norm} = \text{fr } o \text{ norm}) = (\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep})$
 $\langle \text{proof} \rangle$

lemma *fns3*:

$\text{Abs } o \text{ fr } o \text{ Rep} = \text{fa} ==>$
 $(\text{norm } o \text{ fr } o \text{ norm} = \text{norm } o \text{ fr}) = (\text{fa } o \text{ Abs} = \text{Abs } o \text{ fr})$
 $\langle \text{proof} \rangle$

lemma *fns*:

$\text{fr } o \text{ norm} = \text{norm } o \text{ fr} ==>$
 $(\text{fa } o \text{ Abs} = \text{Abs } o \text{ fr}) = (\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep})$
 $\langle \text{proof} \rangle$

lemma *range-norm*:

$\text{range } (\text{Rep } o \text{ Abs}) = A$
 $\langle \text{proof} \rangle$

end

lemmas *td-ext-def'* =

td-ext-def [*unfolded type-definition-def* *td-ext-axioms-def*]

end

10 WordDefinition: Definition of Word Type

```
theory WordDefinition
imports Size BinBoolList TdThs
begin
```

10.1 Type definition

```
typedef (open word) 'a word = {(0::int) ..< 2^len-of TYPE('a::len0)}
morphisms uint Abs-word <proof>
```

definition *word-of-int* :: *int* \Rightarrow *'a::len0 word* **where**
— representation of words using unsigned or signed bins, only difference in these
is the type class
word-of-int *w* = *Abs-word* (*bintrunc* (*len-of TYPE* ('*a*)) *w*)

lemma *uint-word-of-int* [*code*]: *uint* (*word-of-int* *w* :: '*a::len0 word*) = *w mod 2 ^*
len-of TYPE ('*a*)
<proof>

```
code-datatype word-of-int
```

```
notation fcomp (infixl o> 60)
notation scomp (infixl o→ 60)
```

```
instantiation word :: ({len0, typerep}) random
begin
```

definition
random-word *i* = *Random.range* (*max* *i* (*2 ^ len-of TYPE* ('*a*))) *o*→ (λk . *Pair* (
let *j* = *word-of-int* (*Code-Numeral.int-of* *k*) :: '*a word*
in (*j*, λ -.unit. *Code-Evaluation.term-of* *j*)))

```
instance <proof>
```

```
end
```

```
no-notation fcomp (infixl o> 60)
no-notation scomp (infixl o→ 60)
```

10.2 Type conversions and casting

definition *sint* :: '*a* :: *len word* => *int* **where**
— treats the most-significant-bit as a sign bit
sint-uint: *sint* *w* = *sbintrunc* (*len-of TYPE* ('*a*) - 1) (*uint* *w*)

definition *unat* :: 'a :: len0 word => nat **where**
unat w = nat (uint w)

definition *uints* :: nat => int set **where**
 — the sets of integers representing the words
uints n = range (bintrunc n)

definition *sints* :: nat => int set **where**
sints n = range (sbintrunc (n - 1))

definition *unats* :: nat => nat set **where**
unats n = {i. i < 2 ^ n}

definition *norm-sint* :: nat => int => int **where**
norm-sint n w = (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)

definition *scast* :: 'a :: len word => 'b :: len word **where**
 — cast a word to a different length
scast w = word-of-int (sint w)

definition *ucast* :: 'a :: len0 word => 'b :: len0 word **where**
ucast w = word-of-int (uint w)

instantiation *word* :: (len0) size
begin

definition
word-size: size (w :: 'a word) = len-of TYPE('a)

instance <proof>

end

definition *source-size* :: ('a :: len0 word => 'b) => nat **where**
 — whether a cast (or other) function is to a longer or shorter length
source-size c = (let arb = undefined ; x = c arb in size arb)

definition *target-size* :: ('a => 'b :: len0 word) => nat **where**
target-size c = size (c undefined)

definition *is-up* :: ('a :: len0 word => 'b :: len0 word) => bool **where**
is-up c <=> source-size c <= target-size c

definition *is-down* :: ('a :: len0 word => 'b :: len0 word) => bool **where**
is-down c <=> target-size c <= source-size c

definition *of-bl* :: bool list => 'a :: len0 word **where**
of-bl bl = word-of-int (bl-to-bin bl)

definition *to-bl* :: 'a :: len0 word => bool list **where**
to-bl w = *bin-to-bl* (*len-of TYPE* ('a)) (*uint w*)

definition *word-reverse* :: 'a :: len0 word => 'a word **where**
word-reverse w = *of-bl* (*rev* (*to-bl w*))

definition *word-int-case* :: (int => 'b) => ('a :: len0 word) => 'b **where**
word-int-case f w = *f* (*uint w*)

syntax

of-int :: int => 'a

translations

case x of CONST of-int y => b == CONST word-int-case (%y. b) x

10.3 Arithmetic operations

instantiation *word* :: (len0) {*number*, *uminus*, *minus*, *plus*, *one*, *zero*, *times*,
Divides.div, *ord*, *bit*}
begin

definition

word-0-wi: 0 = *word-of-int* 0

definition

word-1-wi: 1 = *word-of-int* 1

definition

word-add-def: $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$

definition

word-sub-wi: $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$

definition

word-minus-def: $- a = \text{word-of-int } (- \text{uint } a)$

definition

word-mult-def: $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$

definition

word-div-def: $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div } \text{uint } b)$

definition

word-mod-def: $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod } \text{uint } b)$

definition

word-number-of-def: *number-of w* = *word-of-int w*

definition

word-le-def: $a \leq b \longleftrightarrow \text{uint } a \leq \text{uint } b$

definition

word-less-def: $x < y \iff x \leq y \wedge x \neq (y :: 'a \text{ word})$

definition

word-and-def:
 $(a :: 'a \text{ word}) \text{ AND } b = \text{word-of-int } (\text{uint } a \text{ AND } \text{uint } b)$

definition

word-or-def:
 $(a :: 'a \text{ word}) \text{ OR } b = \text{word-of-int } (\text{uint } a \text{ OR } \text{uint } b)$

definition

word-xor-def:
 $(a :: 'a \text{ word}) \text{ XOR } b = \text{word-of-int } (\text{uint } a \text{ XOR } \text{uint } b)$

definition

word-not-def:
 $\text{NOT } (a :: 'a \text{ word}) = \text{word-of-int } (\text{NOT } (\text{uint } a))$

instance $\langle \text{proof} \rangle$

end

definition

word-succ :: $'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

where

word-succ $a = \text{word-of-int } (\text{Int.succ } (\text{uint } a))$

definition

word-pred :: $'a :: \text{len0 word} \Rightarrow 'a \text{ word}$

where

word-pred $a = \text{word-of-int } (\text{Int.pred } (\text{uint } a))$

definition *udvd* :: $'a :: \text{len word} \Rightarrow 'a :: \text{len word} \Rightarrow \text{bool}$ (**infixl** *udvd* 50) **where**

$a \text{ udvd } b == \text{EX } n \geq 0. \text{uint } b = n * \text{uint } a$

definition *word-sle* :: $'a :: \text{len word} \Rightarrow 'a \text{ word} \Rightarrow \text{bool}$ ((-/ <=s -) [50, 51] 50)

where

$a <=s b == \text{sint } a \leq \text{sint } b$

definition *word-sless* :: $'a :: \text{len word} \Rightarrow 'a \text{ word} \Rightarrow \text{bool}$ ((-/ <s -) [50, 51] 50)

where

$(x <s y) == (x <=s y \ \& \ x \sim= y)$

10.4 Bit-wise operations

instantiation *word* :: $(\text{len0}) \text{ bits}$

begin

definition

word-test-bit-def: $\text{test-bit } a = \text{bin-nth } (\text{uint } a)$

definition

word-set-bit-def: $\text{set-bit } a \ n \ x =$
word-of-int ($\text{bin-sc } n \ (\text{If } x \ \text{bit.B1} \ \text{bit.B0}) \ (\text{uint } a)$)

definition

word-set-bits-def: $(\text{BITS } n. \ f \ n) = \text{of-bl } (\text{bl-of-nth } (\text{len-of } \text{TYPE } ('a)) \ f)$

definition

word-lsb-def: $\text{lsb } a \longleftrightarrow \text{bin-last } (\text{uint } a) = \text{bit.B1}$

definition *shiftrl1* :: 'a word \Rightarrow 'a word **where**

shiftrl1 $w = \text{word-of-int } (\text{uint } w \ \text{BIT} \ \text{bit.B0})$

definition *shiftr1* :: 'a word \Rightarrow 'a word **where**

— shift right as unsigned or as signed, ie logical or arithmetic
shiftr1 $w = \text{word-of-int } (\text{bin-rest } (\text{uint } w))$

definition

shiftrl-def: $w << n = (\text{shiftrl1 } \wedge^{\wedge} n) \ w$

definition

shiftr-def: $w >> n = (\text{shiftr1 } \wedge^{\wedge} n) \ w$

instance $\langle \text{proof} \rangle$

end

instantiation *word* :: (len) bitss

begin

definition

word-msb-def:
 $\text{msb } a \longleftrightarrow \text{bin-sign } (\text{sint } a) = \text{Int.Min}$

instance $\langle \text{proof} \rangle$

end

definition *setBit* :: 'a :: len0 word \Rightarrow nat \Rightarrow 'a word **where**

setBit $w \ n == \text{set-bit } w \ n \ \text{True}$

definition *clearBit* :: 'a :: len0 word \Rightarrow nat \Rightarrow 'a word **where**

clearBit $w \ n == \text{set-bit } w \ n \ \text{False}$

10.5 Shift operations

definition *sshiftr1* :: 'a :: len word => 'a word **where**
sshiftr1 w == word-of-int (bin-rest (sint w))

definition *bshiftr1* :: bool => 'a :: len word => 'a word **where**
bshiftr1 b w == of-bl (b # butlast (to-bl w))

definition *sshiftr* :: 'a :: len word => nat => 'a word (infixl >>> 55) **where**
w >>> *n* == (*sshiftr1* ^^ *n*) *w*

definition *mask* :: nat => 'a::len word **where**
mask n == (1 << n) - 1

definition *revcast* :: 'a :: len0 word => 'b :: len0 word **where**
revcast w == of-bl (takefill False (len-of TYPE('b)) (to-bl w))

definition *slice1* :: nat => 'a :: len0 word => 'b :: len0 word **where**
slice1 n w == of-bl (takefill False n (to-bl w))

definition *slice* :: nat => 'a :: len0 word => 'b :: len0 word **where**
slice n w == *slice1* (size w - n) w

10.6 Rotation

definition *rotater1* :: 'a list => 'a list **where**
rotater1 ys ==
 case ys of [] => [] | x # xs => last ys # butlast ys

definition *rotater* :: nat => 'a list => 'a list **where**
rotater n == *rotater1* ^^ n

definition *word-rotr* :: nat => 'a :: len0 word => 'a :: len0 word **where**
word-rotr n w == of-bl (*rotater* n (to-bl w))

definition *word-rotl* :: nat => 'a :: len0 word => 'a :: len0 word **where**
word-rotl n w == of-bl (rotate n (to-bl w))

definition *word-roti* :: int => 'a :: len0 word => 'a :: len0 word **where**
word-roti i w == if i >= 0 then *word-rotr* (nat i) w
 else *word-rotl* (nat (- i)) w

10.7 Split and cat operations

definition *word-cat* :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word **where**
word-cat a b == word-of-int (bin-cat (uint a) (len-of TYPE ('b)) (uint b))

definition *word-split* :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)
where
word-split a ==

case bin-split (len-of TYPE ('c)) (uint a) of
(u, v) => (word-of-int u, word-of-int v)

definition *word-rcat* :: 'a :: len0 word list => 'b :: len0 word **where**
word-rcat ws ==
word-of-int (bin-rcat (len-of TYPE ('a)) (map uint ws))

definition *word-rsplit* :: 'a :: len0 word => 'b :: len word list **where**
word-rsplit w ==
map word-of-int (bin-rsplit (len-of TYPE ('b)) (len-of TYPE ('a), uint w))

definition *max-word* :: 'a::len word — Largest representable machine integer.
where
max-word ≡ word-of-int (2 ^ len-of TYPE('a) - 1)

primrec *of-bool* :: bool => 'a::len word **where**
of-bool False = 0
| of-bool True = 1

lemmas *of-nth-def* = *word-set-bits-def*

lemmas *word-size-gt-0 [iff]* =
xtr1 [OF word-size len-gt-0, standard]
lemmas *lens-gt-0* = *word-size-gt-0 len-gt-0*
lemmas *lens-not-0 [iff]* = *lens-gt-0 [THEN gr-implies-not0, standard]*

lemma *uints-num*: *uints n = {i. 0 ≤ i ∧ i < 2 ^ n}*
<proof>

lemma *sints-num*: *sints n = {i. - (2 ^ (n - 1)) ≤ i ∧ i < 2 ^ (n - 1)}*
<proof>

lemmas *atLeastLessThan-alt* = *atLeastLessThan-def [unfolded*
atLeast-def lessThan-def Collect-conj-eq [symmetric]]

lemma *mod-in-reps*: *m > 0 ==> y mod m : {0::int ..< m}*
<proof>

lemma
uint-0:0 <= uint x and
uint-lt: uint (x::'a::len0 word) < 2 ^ len-of TYPE('a)
<proof>

lemma *uint-mod-same*:
uint x mod 2 ^ len-of TYPE('a) = uint (x::'a::len0 word)
<proof>

lemma *td-ext-uint*:

td-ext (*uint* :: 'a word => int) *word-of-int* (*uints* (*len-of TYPE*('a::len0)))
 (%w::int. w mod 2 ^ *len-of TYPE*('a))
 <proof>

lemmas *int-word-uint* = *td-ext-uint* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-uint*:
td-ext uint::'a::len0 word ⇒ *int*
word-of-int
uints (*len-of TYPE*('a::len0))
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE}('a::\text{len0})$
 <proof>

lemmas *td-uint* = *word-uint.td-thm*

lemmas *td-ext-ubin* = *td-ext-uint*
 [simplified *len-gt-0* *no-bintr-alt1* [symmetric]]

interpretation *word-ubin*:
td-ext uint::'a::len0 word ⇒ *int*
word-of-int
uints (*len-of TYPE*('a::len0))
bintrunc (*len-of TYPE*('a::len0))
 <proof>

lemma *sint-sbintrunc'*:
sint (*word-of-int bin* :: 'a word) =
 (*sbintrunc* (*len-of TYPE* ('a :: len) - 1) *bin*)
 <proof>

lemma *uint-sint*:
uint w = *bintrunc* (*len-of TYPE*('a)) (*sint* (*w* :: 'a :: len word))
 <proof>

lemma *bintr-uint'*:
 $n \geq \text{size } w \implies \text{bintrunc } n (\text{uint } w) = \text{uint } w$
 <proof>

lemma *wi-bintr'*:
 $wb = \text{word-of-int bin} \implies n \geq \text{size } wb \implies$
 $\text{word-of-int } (\text{bintrunc } n \text{ bin}) = wb$
 <proof>

lemmas *bintr-uint* = *bintr-uint'* [unfolded *word-size*]

lemmas *wi-bintr* = *wi-bintr'* [unfolded *word-size*]

lemma *td-ext-sbin*:
td-ext (*sint* :: 'a word => int) *word-of-int* (*sints* (*len-of TYPE*('a::len)))
 (*sbintrunc* (*len-of TYPE*('a) - 1))

$\langle \text{proof} \rangle$

lemmas $td\text{-}ext\text{-}sint = td\text{-}ext\text{-}sbin$
 $[simplified\ len\text{-}gt\text{-}0\ no\text{-}sbintr\text{-}alt2\ Suc\text{-}pred'\ [symmetric]]$

interpretation $word\text{-}sint$:
 $td\text{-}ext\ sint :: 'a::len\ word \Rightarrow int$
 $word\text{-}of\text{-}int$
 $sints\ (len\text{-}of\ TYPE('a::len))$
 $\%w. (w + 2^{(len\text{-}of\ TYPE('a::len) - 1)}) \bmod 2^{len\text{-}of\ TYPE('a::len) - 1}$
 $2^{(len\text{-}of\ TYPE('a::len) - 1)}$
 $\langle \text{proof} \rangle$

interpretation $word\text{-}sbin$:
 $td\text{-}ext\ sint :: 'a::len\ word \Rightarrow int$
 $word\text{-}of\text{-}int$
 $sints\ (len\text{-}of\ TYPE('a::len))$
 $sbintrunc\ (len\text{-}of\ TYPE('a::len) - 1)$
 $\langle \text{proof} \rangle$

lemmas $int\text{-}word\text{-}sint = td\text{-}ext\text{-}sint\ [THEN\ td\text{-}ext.\text{eq}\text{-}norm,\ standard]$

lemmas $td\text{-}sint = word\text{-}sint.td$

lemma $word\text{-}number\text{-}of\text{-}alt$: $number\text{-}of\ b == word\text{-}of\text{-}int\ (number\text{-}of\ b)$
 $\langle \text{proof} \rangle$

lemma $word\text{-}no\text{-}wi$: $number\text{-}of = word\text{-}of\text{-}int$
 $\langle \text{proof} \rangle$

lemma $to\text{-}bl\text{-}def'$:
 $(to\text{-}bl :: 'a :: len0\ word \Rightarrow bool\ list) =$
 $bin\text{-}to\text{-}bl\ (len\text{-}of\ TYPE('a))\ o\ uint$
 $\langle \text{proof} \rangle$

lemmas $word\text{-}reverse\text{-}no\text{-}def\ [simp] = word\text{-}reverse\text{-}def\ [of\ number\text{-}of\ w,\ standard]$

lemmas $uints\text{-}mod = uints\text{-}def\ [unfolded\ no\text{-}bintr\text{-}alt1]$

lemma $uint\text{-}bintrunc$: $uint\ (number\text{-}of\ bin :: 'a\ word) =$
 $number\text{-}of\ (bintrunc\ (len\text{-}of\ TYPE('a :: len0))\ bin)$
 $\langle \text{proof} \rangle$

lemma $sint\text{-}sbintrunc$: $sint\ (number\text{-}of\ bin :: 'a\ word) =$
 $number\text{-}of\ (sbintrunc\ (len\text{-}of\ TYPE('a :: len) - 1)\ bin)$
 $\langle \text{proof} \rangle$

lemma $unat\text{-}bintrunc$:

unat (*number-of* *bin* :: 'a :: len0 word) =
number-of (*bintrunc* (*len-of* *TYPE*('a)) *bin*)
 ⟨*proof*⟩

declare

uint-bintrunc [*simp*]
sint-sbintrunc [*simp*]
unat-bintrunc [*simp*]

lemma *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*
 ⟨*proof*⟩

lemmas *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]
lemmas *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]
lemmas *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]
lemmas *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]
lemmas *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]
lemmas *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

lemma *sign-uint-Pls* [*simp*]:
bin-sign (*uint* *x*) = *Int.Pl*
 ⟨*proof*⟩

lemmas *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]
lemmas *uint-m2p-not-non-neg* =
iffD2 [*OF linorder-not-le uint-m2p-neg, standard*]

lemma *lt2p-lem*:
len-of TYPE('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*
 ⟨*proof*⟩

lemmas *uint-le-0-iff* [*simp*] =
uint-ge-0 [*THEN leD, THEN linorder-antisym-conv1, standard*]

lemma *uint-nat*: *uint* *w* == *int* (*unat* *w*)
 ⟨*proof*⟩

lemma *uint-number-of*:
uint (*number-of* *b* :: 'a :: len0 word) = *number-of* *b* mod 2 ^ *len-of TYPE*('a)
 ⟨*proof*⟩

lemma *unat-number-of*:
bin-sign *b* = *Int.Pl* ==>
unat (*number-of* *b*::'a::len0 word) = *number-of* *b* mod 2 ^ *len-of TYPE* ('a)
 ⟨*proof*⟩

lemma *sint-number-of*: *sint* (*number-of* *b* :: 'a :: len word) = (*number-of* *b* +
 2 ^ (*len-of TYPE*('a) - 1)) mod 2 ^ *len-of TYPE*('a) -

$2 \wedge (\text{len-of TYPE}('a) - 1)$
 $\langle \text{proof} \rangle$

lemma *word-of-int-bin [simp]* :
 $(\text{word-of-int} (\text{number-of bin}) :: 'a :: \text{len0 word}) = (\text{number-of bin})$
 $\langle \text{proof} \rangle$

lemma *word-int-case-wi*:
 $\text{word-int-case } f (\text{word-of-int } i :: 'b \text{ word}) =$
 $f (i \bmod 2 \wedge \text{len-of TYPE}('b::\text{len0}))$
 $\langle \text{proof} \rangle$

lemma *word-int-split*:
 $P (\text{word-int-case } f x) =$
 $(\text{ALL } i. x = (\text{word-of-int } i :: 'b :: \text{len0 word}) \ \&$
 $0 \leq i \ \& \ i < 2 \wedge \text{len-of TYPE}('b) \longrightarrow P (f i))$
 $\langle \text{proof} \rangle$

lemma *word-int-split-asm*:
 $P (\text{word-int-case } f x) =$
 $(\sim (EX \ n. x = (\text{word-of-int } n :: 'b::\text{len0 word}) \ \&$
 $0 \leq n \ \& \ n < 2 \wedge \text{len-of TYPE}('b::\text{len0}) \ \& \ \sim P (f n)))$
 $\langle \text{proof} \rangle$

lemmas *uint-range'* =
 $\text{word-uint.Rep} [\text{unfolded uints-num mem-Collect-eq, standard}]$
lemmas *sint-range'* = $\text{word-sint.Rep} [\text{unfolded One-nat-def}$
 $\text{sints-num mem-Collect-eq, standard}]$

lemma *uint-range-size*: $0 \leq \text{uint } w \ \& \ \text{uint } w < 2 \wedge \text{size } w$
 $\langle \text{proof} \rangle$

lemma *sint-range-size*:
 $-(2 \wedge (\text{size } w - \text{Suc } 0)) \leq \text{sint } w \ \& \ \text{sint } w < 2 \wedge (\text{size } w - \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemmas *sint-above-size* = *sint-range-size*
 $[\text{THEN conjunct2, THEN } [2] \text{ xtr8, folded One-nat-def, standard}]$

lemmas *sint-below-size* = *sint-range-size*
 $[\text{THEN conjunct1, THEN } [2] \text{ order-trans, folded One-nat-def, standard}]$

lemma *test-bit-eq-iff*: $(\text{test-bit } (u::'a::\text{len0 word}) = \text{test-bit } v) = (u = v)$
 $\langle \text{proof} \rangle$

lemma *test-bit-size [rule-format]* : $(w::'a::\text{len0 word}) !! n \longrightarrow n < \text{size } w$
 $\langle \text{proof} \rangle$

lemma *word-eqI [rule-format]* :

fixes $u :: 'a::len0 \text{ word}$
shows $(\text{ALL } n. n < \text{size } u \dashv\vdash u !! n = v !! n) \implies u = v$
 $\langle \text{proof} \rangle$

lemmas $\text{word-eqD} = \text{test-bit-eq-iff} \text{ [THEN iffD2, THEN fun-cong, standard]}$

lemma test-bit-bin' : $w !! n = (n < \text{size } w \ \& \ \text{bin-nth} (\text{uint } w) \ n)$
 $\langle \text{proof} \rangle$

lemmas $\text{test-bit-bin} = \text{test-bit-bin'}$ $[\text{unfolded word-size}]$

lemma bin-nth-uint-imp' : $\text{bin-nth} (\text{uint } w) \ n \dashv\vdash n < \text{size } w$
 $\langle \text{proof} \rangle$

lemma bin-nth-sint' :
 $n \geq \text{size } w \dashv\vdash \text{bin-nth} (\text{sint } w) \ n = \text{bin-nth} (\text{sint } w) (\text{size } w - 1)$
 $\langle \text{proof} \rangle$

lemmas $\text{bin-nth-uint-imp} = \text{bin-nth-uint-imp'}$ $[\text{rule-format, unfolded word-size}]$
lemmas $\text{bin-nth-sint} = \text{bin-nth-sint'}$ $[\text{rule-format, unfolded word-size}]$

lemma td-bl :
 $\text{type-definition } (\text{to-bl} :: 'a::len0 \text{ word} \Rightarrow \text{bool list})$
 of-bl
 $\{\text{bl. length bl} = \text{len-of TYPE('a)}\}$
 $\langle \text{proof} \rangle$

interpretation word-bl :
 $\text{type-definition } \text{to-bl} :: 'a::len0 \text{ word} \Rightarrow \text{bool list}$
 of-bl
 $\{\text{bl. length bl} = \text{len-of TYPE('a::len0)}\}$
 $\langle \text{proof} \rangle$

lemma word-size-bl : $\text{size } w == \text{size} (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma to-bl-use-of-bl :
 $(\text{to-bl } w = \text{bl}) = (w = \text{of-bl bl} \wedge \text{length bl} = \text{length} (\text{to-bl } w))$
 $\langle \text{proof} \rangle$

lemma to-bl-word-rev : $\text{to-bl} (\text{word-reverse } w) = \text{rev} (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma word-rev-rev $[\text{simp}]$: $\text{word-reverse} (\text{word-reverse } w) = w$
 $\langle \text{proof} \rangle$

lemma word-rev-gal : $\text{word-reverse } w = u \implies \text{word-reverse } u = w$
 $\langle \text{proof} \rangle$

lemmas *word-rev-gal'* = *sym* [*THEN word-rev-gal, symmetric, standard*]

lemmas *length-bl-gt-0* [*iff*] = *xtr1* [*OF word-bl.Rep' len-gt-0, standard*]

lemmas *bl-not-Nil* [*iff*] =

length-bl-gt-0 [*THEN length-greater-0-conv* [*THEN iffD1*], *standard*]

lemmas *length-bl-neq-0* [*iff*] = *length-bl-gt-0* [*THEN gr-implies-not0*]

lemma *hd-bl-sign-sint*: *hd (to-bl w) = (bin-sign (sint w) = Int.Min)*
<proof>

lemma *of-bl-drop'*:

lend = length bl - len-of TYPE ('a :: len0) ==>

of-bl (drop lend bl) = (of-bl bl :: 'a word)

<proof>

lemmas *of-bl-no* = *of-bl-def* [*folded word-number-of-def*]

lemma *test-bit-of-bl*:

(of-bl bl :: 'a :: len0 word) !! n = (rev bl ! n ∧ n < len-of TYPE('a) ∧ n < length bl)

<proof>

lemma *no-of-bl*:

(number-of bin :: 'a :: len0 word) = of-bl (bin-to-bl (len-of TYPE ('a)) bin)

<proof>

lemma *uint-bl*: *to-bl w == bin-to-bl (size w) (uint w)*

<proof>

lemma *to-bl-bin*: *bl-to-bin (to-bl w) = uint w*

<proof>

lemma *to-bl-of-bin*:

to-bl (word-of-int bin :: 'a :: len0 word) = bin-to-bl (len-of TYPE('a)) bin

<proof>

lemmas *to-bl-no-bin* [*simp*] = *to-bl-of-bin* [*folded word-number-of-def*]

lemma *to-bl-to-bin* [*simp*] : *bl-to-bin (to-bl w) = uint w*

<proof>

lemmas *uint-bl-bin* [*simp*] = *trans* [*OF bin-bl-bin word-ubin.norm-Rep, standard*]

lemmas *num-AB-u* [*simp*] = *word-uint.Rep-inverse*

[*unfolded o-def word-number-of-def* [*symmetric*], *standard*]

lemmas *num-AB-s* [*simp*] = *word-sint.Rep-inverse*

[*unfolded o-def word-number-of-def* [*symmetric*], *standard*]

lemma *uints-unats*: *uints n = int ‘ unats n*
 ⟨*proof*⟩

lemma *unats-uints*: *unats n = nat ‘ uints n*
 ⟨*proof*⟩

lemmas *bintr-num = word-ubin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *sbintr-num = word-sbin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *num-of-bintr = word-ubin.Abs-norm* [*folded word-number-of-def, standard*]

lemmas *num-of-sbintr = word-sbin.Abs-norm* [*folded word-number-of-def, standard*]

lemma *num-of-bintr'*:
bintrunc (len-of TYPE('a :: len0)) a = b ==>
number-of a = (number-of b :: 'a word)
 ⟨*proof*⟩

lemma *num-of-sbintr'*:
sbintrunc (len-of TYPE('a :: len) - 1) a = b ==>
number-of a = (number-of b :: 'a word)
 ⟨*proof*⟩

lemmas *num-abs-bintr = sym [THEN trans,*
OF num-of-bintr word-number-of-def, standard]

lemmas *num-abs-sbintr = sym [THEN trans,*
OF num-of-sbintr word-number-of-def, standard]

lemma *ucast-id*: *ucast w = w*
 ⟨*proof*⟩

lemma *scast-id*: *scast w = w*
 ⟨*proof*⟩

lemma *ucast-bl*: *ucast w == of-bl (to-bl w)*
 ⟨*proof*⟩

lemma *nth-ucast*:
(ucast w :: 'a :: len0 word) !! n = (w !! n & n < len-of TYPE('a))
 ⟨*proof*⟩

lemma *ucast-bintr* [simp]:

ucast (*number-of* *w* :: 'a::len0 *word*) =
number-of (*bintrunc* (*len-of TYPE*('a)) *w*)
 ⟨*proof*⟩

lemma *scast-sbintr* [simp]:

scast (*number-of* *w* :: 'a::len *word*) =
number-of (*sbintrunc* (*len-of TYPE*('a) - Suc 0) *w*)
 ⟨*proof*⟩

lemmas *source-size* = *source-size-def* [unfolded *Let-def word-size*]

lemmas *target-size* = *target-size-def* [unfolded *Let-def word-size*]

lemmas *is-down* = *is-down-def* [unfolded *source-size target-size*]

lemmas *is-up* = *is-up-def* [unfolded *source-size target-size*]

lemmas *is-up-down* = *trans* [*OF is-up is-down* [symmetric], *standard*]

lemma *down-cast-same'*: *uc* = *ucast* ==> *is-down uc* ==> *uc* = *scast*
 ⟨*proof*⟩

lemma *word-rev-tf'*:

r = *to-bl* (*of-bl bl*) ==> *r* = *rev* (*takefill False* (*length r*) (*rev bl*))
 ⟨*proof*⟩

lemmas *word-rev-tf* = *refl* [*THEN word-rev-tf'*, *unfolded word-bl.Rep'*, *standard*]

lemmas *word-rep-drop* = *word-rev-tf* [simplified *takefill-alt*,
simplified, *simplified rev-take*, *simplified*]

lemma *to-bl-ucast*:

to-bl (*ucast* (*w*::'b::len0 *word*) :: 'a::len0 *word*) =
replicate (*len-of TYPE*('a) - *len-of TYPE*('b)) *False* @
drop (*len-of TYPE*('b) - *len-of TYPE*('a)) (*to-bl w*)
 ⟨*proof*⟩

lemma *ucast-up-app'*:

uc = *ucast* ==> *source-size uc* + *n* = *target-size uc* ==>
to-bl (*uc w*) = *replicate n False* @ (*to-bl w*)
 ⟨*proof*⟩

lemma *ucast-down-drop'*:

uc = *ucast* ==> *source-size uc* = *target-size uc* + *n* ==>
to-bl (*uc w*) = *drop n* (*to-bl w*)
 ⟨*proof*⟩

lemma *scast-down-drop'*:

sc = *scast* ==> *source-size sc* = *target-size sc* + *n* ==>

$to-bl (sc\ w) = drop\ n\ (to-bl\ w)$
 ⟨proof⟩

lemma *sint-up-scast'*:

$sc = scast ==> is-up\ sc ==> sint\ (sc\ w) = sint\ w$
 ⟨proof⟩

lemma *uint-up-ucast'*:

$uc = ucast ==> is-up\ uc ==> uint\ (uc\ w) = uint\ w$
 ⟨proof⟩

lemmas *down-cast-same* = refl [THEN down-cast-same']

lemmas *ucast-up-app* = refl [THEN ucast-up-app']

lemmas *ucast-down-drop* = refl [THEN ucast-down-drop']

lemmas *scast-down-drop* = refl [THEN scast-down-drop']

lemmas *uint-up-ucast* = refl [THEN uint-up-ucast']

lemmas *sint-up-scast* = refl [THEN sint-up-scast']

lemma *ucast-up-ucast'*: $uc = ucast ==> is-up\ uc ==> ucast\ (uc\ w) = ucast\ w$
 ⟨proof⟩

lemma *scast-up-scast'*: $sc = scast ==> is-up\ sc ==> scast\ (sc\ w) = scast\ w$
 ⟨proof⟩

lemma *ucast-of-bl-up'*:

$w = of-bl\ bl ==> size\ bl <= size\ w ==> ucast\ w = of-bl\ bl$
 ⟨proof⟩

lemmas *ucast-up-ucast* = refl [THEN ucast-up-ucast']

lemmas *scast-up-scast* = refl [THEN scast-up-scast']

lemmas *ucast-of-bl-up* = refl [THEN ucast-of-bl-up']

lemmas *ucast-up-ucast-id* = trans [OF ucast-up-ucast ucast-id]

lemmas *scast-up-scast-id* = trans [OF scast-up-scast scast-id]

lemmas *isduu* = is-up-down [where $c = ucast$, THEN iffD2]

lemmas *isdus* = is-up-down [where $c = scast$, THEN iffD2]

lemmas *ucast-down-ucast-id* = isduu [THEN ucast-up-ucast-id]

lemmas *scast-down-scast-id* = isdus [THEN ucast-up-ucast-id]

lemma *up-ucast-surj*:

$is-up\ (ucast :: 'b::len0\ word ==> 'a::len0\ word) ==>$
 $surj\ (ucast :: 'a\ word ==> 'b\ word)$
 ⟨proof⟩

lemma *up-scast-surj*:

$is-up\ (scast :: 'b::len\ word ==> 'a::len\ word) ==>$
 $surj\ (scast :: 'a\ word ==> 'b\ word)$
 ⟨proof⟩

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) ==>
inj-on (*ucast* :: 'a word => 'b word) *A*
 ⟨*proof*⟩

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) ==>
inj-on (*ucast* :: 'a word => 'b word) *A*
 ⟨*proof*⟩

lemma *of-bl-append-same*: *of-bl* (*X* @ *to-bl w*) = *w*

⟨*proof*⟩

lemma *ucast-down-no'*:

uc = *ucast* ==> *is-down uc* ==> *uc (number-of bin) = number-of bin*
 ⟨*proof*⟩

lemmas *ucast-down-no* = *ucast-down-no'* [*OF refl*]

lemma *ucast-down-bl'*: *uc* = *ucast* ==> *is-down uc* ==> *uc (of-bl bl) = of-bl bl*

⟨*proof*⟩

lemmas *ucast-down-bl* = *ucast-down-bl'* [*OF refl*]

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def word-or-def word-xor-def word-not-def*

lemmas *word-log-bin-defs* = *word-log-defs*

Executable equality

instantiation *word* :: ({*len0*}) *eq*

begin

definition *eq-word* :: 'a word ⇒ 'a word ⇒ bool **where**

eq-word k l ⇔ *HOL.eq (uint k) (uint l)*

instance ⟨*proof*⟩

end

end

11 WordArith: Word Arithmetic

theory *WordArith*

imports *WordDefinition*

begin

lemma *word-less-alt*: $(a < b) = (uint\ a < uint\ b)$
 $\langle proof \rangle$

lemma *signed-linorder*: *class.linorder word-sle word-sless*
 $\langle proof \rangle$

interpretation *signed*: *linorder word-sle word-sless*
 $\langle proof \rangle$

lemmas *word-arith-wis* =
word-add-def word-mult-def word-minus-def
word-succ-def word-pred-def word-0-wi word-1-wi

lemma *udvdI*:
 $0 \leq n \implies uint\ b = n * uint\ a \implies a\ udvd\ b$
 $\langle proof \rangle$

lemmas *word-div-no* [*simp*] =
word-div-def [*of number-of a number-of b, standard*]

lemmas *word-mod-no* [*simp*] =
word-mod-def [*of number-of a number-of b, standard*]

lemmas *word-less-no* [*simp*] =
word-less-def [*of number-of a number-of b, standard*]

lemmas *word-le-no* [*simp*] =
word-le-def [*of number-of a number-of b, standard*]

lemmas *word-sless-no* [*simp*] =
word-sless-def [*of number-of a number-of b, standard*]

lemmas *word-sle-no* [*simp*] =
word-sle-def [*of number-of a number-of b, standard*]

lemmas *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]

lemmas *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

lemma *int-one-bin*: $(1 :: int) == (Int.Pls\ BIT\ bit.B1)$
 $\langle proof \rangle$

lemma *word-1-no*:
 $(1 :: 'a :: len0\ word) == number-of\ (Int.Pls\ BIT\ bit.B1)$
 $\langle proof \rangle$

lemma *word-m1-wi*: $-1 == word-of-int\ -1$

$\langle proof \rangle$

lemma *word-m1-wi-Min*: $-1 = \text{word-of-int Int.Min}$
 $\langle proof \rangle$

lemma *word-0-bl*: $\text{of-bl } [] = 0$
 $\langle proof \rangle$

lemma *word-1-bl*: $\text{of-bl } [True] = 1$
 $\langle proof \rangle$

lemma *uint-0 [simp]*: $(\text{uint } 0 = 0)$
 $\langle proof \rangle$

lemma *of-bl-0 [simp]*: $\text{of-bl } (\text{replicate } n \text{ False}) = 0$
 $\langle proof \rangle$

lemma *to-bl-0*:
 $\text{to-bl } (0 :: 'a :: \text{len0 word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ False}$
 $\langle proof \rangle$

lemma *uint-0-iff*: $(\text{uint } x = 0) = (x = 0)$
 $\langle proof \rangle$

lemma *unat-0-iff*: $(\text{unat } x = 0) = (x = 0)$
 $\langle proof \rangle$

lemma *unat-0 [simp]*: $\text{unat } 0 = 0$
 $\langle proof \rangle$

lemma *size-0-same'*: $\text{size } w = 0 ==> w = (v :: 'a :: \text{len0 word})$
 $\langle proof \rangle$

lemmas *size-0-same* = *size-0-same'* [folded word-size]

lemmas *unat-eq-0* = *unat-0-iff*
lemmas *unat-eq-zero* = *unat-0-iff*

lemma *unat-gt-0*: $(0 < \text{unat } x) = (x \sim= 0)$
 $\langle proof \rangle$

lemma *ucast-0 [simp]*: $\text{ucast } 0 = 0$
 $\langle proof \rangle$

lemma *sint-0 [simp]*: $\text{sint } 0 = 0$
 $\langle proof \rangle$

lemma *scast-0 [simp]*: $\text{scast } 0 = 0$
 $\langle proof \rangle$

lemma *sint-n1* [*simp*] : *sint* $-1 = -1$
 ⟨*proof*⟩

lemma *scast-n1* [*simp*] : *scast* $-1 = -1$
 ⟨*proof*⟩

lemma *uint-1* [*simp*] : *uint* ($1 :: 'a :: \text{len word}$) = 1
 ⟨*proof*⟩

lemma *unat-1* [*simp*] : *unat* ($1 :: 'a :: \text{len word}$) = 1
 ⟨*proof*⟩

lemma *ucast-1* [*simp*] : *ucast* ($1 :: 'a :: \text{len word}$) = 1
 ⟨*proof*⟩

lemmas *ariths* =
bintr-ariths [*THEN* *word-ubin.norm-eq-iff* [*THEN iffD1*],
folded word-ubin.eq-norm, standard]

lemma *wi-homs*:
shows
wi-hom-add: *word-of-int* $a + \text{word-of-int } b = \text{word-of-int } (a + b)$ **and**
wi-hom-mult: *word-of-int* $a * \text{word-of-int } b = \text{word-of-int } (a * b)$ **and**
wi-hom-neg: $-\text{word-of-int } a = \text{word-of-int } (-a)$ **and**
wi-hom-succ: *word-succ* (*word-of-int* a) = *word-of-int* (*Int.succ* a) **and**
wi-hom-pred: *word-pred* (*word-of-int* a) = *word-of-int* (*Int.pred* a)
 ⟨*proof*⟩

lemmas *wi-hom-syms* = *wi-homs* [*symmetric*]

lemma *word-sub-def*: $a - b == a + - (b :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemmas *word-diff-minus* = *word-sub-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-of-int-sub-hom*:
 (*word-of-int* a) $- \text{word-of-int } b = \text{word-of-int } (a - b)$
 ⟨*proof*⟩

lemmas *new-word-of-int-homs* =
word-of-int-sub-hom wi-homs word-0-wi word-1-wi

lemmas *new-word-of-int-hom-syms* = *new-word-of-int-homs* [*symmetric, standard*]

lemmas *word-of-int-hom-syms* =
new-word-of-int-hom-syms [*unfolded succ-def pred-def*]

lemmas *word-of-int-homs* =
new-word-of-int-homs [*unfolded succ-def pred-def*]

lemmas *word-of-int-add-hom* = *word-of-int-homs* (2)
lemmas *word-of-int-mult-hom* = *word-of-int-homs* (3)
lemmas *word-of-int-minus-hom* = *word-of-int-homs* (4)
lemmas *word-of-int-succ-hom* = *word-of-int-homs* (5)
lemmas *word-of-int-pred-hom* = *word-of-int-homs* (6)
lemmas *word-of-int-0-hom* = *word-of-int-homs* (7)
lemmas *word-of-int-1-hom* = *word-of-int-homs* (8)

lemmas *word-arith-alt*s =
word-sub-wi [*unfolded succ-def pred-def, standard*]
word-arith-wis [*unfolded succ-def pred-def, standard*]

lemmas *word-sub-alt* = *word-arith-alt*s (1)
lemmas *word-add-alt* = *word-arith-alt*s (2)
lemmas *word-mult-alt* = *word-arith-alt*s (3)
lemmas *word-minus-alt* = *word-arith-alt*s (4)
lemmas *word-succ-alt* = *word-arith-alt*s (5)
lemmas *word-pred-alt* = *word-arith-alt*s (6)
lemmas *word-0-alt* = *word-arith-alt*s (7)
lemmas *word-1-alt* = *word-arith-alt*s (8)

11.1 Transferring goals from words to ints

lemma *word-ths*:
shows
word-succ-p1: *word-succ a* = *a* + 1 **and**
word-pred-m1: *word-pred a* = *a* - 1 **and**
word-pred-succ: *word-pred (word-succ a)* = *a* **and**
word-succ-pred: *word-succ (word-pred a)* = *a* **and**
word-mult-succ: *word-succ a* * *b* = *b* + *a* * *b*
 ⟨*proof*⟩

lemmas *uint-cong* = *arg-cong* [**where** *f* = *uint*]

lemmas *uint-word-ariths* =
*word-arith-alt*s [*THEN trans* [*OF uint-cong int-word-uint*], *standard*]

lemmas *uint-word-arith-bintr*s = *uint-word-ariths* [*folded bintrunc-mod2p*]

lemmas *sint-word-ariths* = *uint-word-arith-bintr*s
 [*THEN uint-sint* [*symmetric, THEN trans*],
unfolded uint-sint bintr-arith1s bintr-ariths]

len-gt-0 [THEN *bin-sbin-eq-iff*] *word-sbin.norm-Rep*, *standard*]

lemmas *uint-div-alt* = *word-div-def*
 [THEN *trans* [OF *uint-cong int-word-uint*], *standard*]
lemmas *uint-mod-alt* = *word-mod-def*
 [THEN *trans* [OF *uint-cong int-word-uint*], *standard*]

lemma *word-pred-0-n1*: *word-pred 0 = word-of-int -1*
 ⟨*proof*⟩

lemma *word-pred-0-Min*: *word-pred 0 = word-of-int Int.Min*
 ⟨*proof*⟩

lemma *word-m1-Min*: *- 1 = word-of-int Int.Min*
 ⟨*proof*⟩

lemma *succ-pred-no* [simp]:
word-succ (number-of bin) = number-of (Int.succ bin) &
word-pred (number-of bin) = number-of (Int.pred bin)
 ⟨*proof*⟩

lemma *word-sp-01* [simp] :
word-succ -1 = 0 & word-succ 0 = 1 & word-pred 0 = -1 & word-pred 1 = 0
 ⟨*proof*⟩

lemma *word-of-int-Ex*:
 $\exists y. x = \text{word-of-int } y$
 ⟨*proof*⟩

lemma *word-arith-eqs*:
fixes *a* :: 'a::len0 word
fixes *b* :: 'a::len0 word
shows
word-add-0: $0 + a = a$ **and**
word-add-0-right: $a + 0 = a$ **and**
word-mult-1: $1 * a = a$ **and**
word-mult-1-right: $a * 1 = a$ **and**
word-add-commute: $a + b = b + a$ **and**
word-add-assoc: $a + b + c = a + (b + c)$ **and**
word-add-left-commute: $a + (b + c) = b + (a + c)$ **and**
word-mult-commute: $a * b = b * a$ **and**
word-mult-assoc: $a * b * c = a * (b * c)$ **and**
word-mult-left-commute: $a * (b * c) = b * (a * c)$ **and**
word-left-distrib: $(a + b) * c = a * c + b * c$ **and**
word-right-distrib: $a * (b + c) = a * b + a * c$ **and**
word-left-minus: $- a + a = 0$ **and**
word-diff-0-right: $a - 0 = a$ **and**
word-diff-self: $a - a = 0$

<proof>

lemmas *word-add-ac* = *word-add-commute word-add-assoc word-add-left-commute*

lemmas *word-mult-ac* = *word-mult-commute word-mult-assoc word-mult-left-commute*

lemmas *word-plus-ac0* = *word-add-0 word-add-0-right word-add-ac*

lemmas *word-times-ac1* = *word-mult-1 word-mult-1-right word-mult-ac*

11.2 Order on fixed-length words

lemma *word-order-trans*: $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$
<proof>

lemma *word-order-refl*: $z \leq (z :: 'a :: \text{len0 word})$
<proof>

lemma *word-order-antisym*: $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$
<proof>

lemma *word-order-linear*:
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$
<proof>

lemma *word-zero-le* [*simp*] :
 $0 \leq (y :: 'a :: \text{len0 word})$
<proof>

instance *word* :: (*len0*) *semigroup-add*
<proof>

instance *word* :: (*len0*) *linorder*
<proof>

instance *word* :: (*len0*) *ring*
<proof>

lemma *word-m1-ge* [*simp*] : *word-pred* 0 $\geq y$
<proof>

lemmas *word-n1-ge* [*simp*] = *word-m1-ge* [*simplified word-sp-01*]

lemmas *word-not-simps* [*simp*] =
word-zero-le [*THEN leD*] *word-m1-ge* [*THEN leD*] *word-n1-ge* [*THEN leD*]

lemma *word-gt-0*: $0 < y = (0 \sim (y :: 'a :: \text{len0 word}))$
<proof>

lemmas *word-gt-0-no* [*simp*] = *word-gt-0* [*of number-of y, standard*]

lemma *word-sless-alt*: $(a <_s b) == (sint\ a < sint\ b)$
 $\langle proof \rangle$

lemma *word-le-nat-alt*: $(a <= b) = (unat\ a <= unat\ b)$
 $\langle proof \rangle$

lemma *word-less-nat-alt*: $(a < b) = (unat\ a < unat\ b)$
 $\langle proof \rangle$

lemma *wi-less*:
 $(word-of-int\ n < (word-of-int\ m :: 'a :: len0\ word)) =$
 $(n\ mod\ 2 \wedge len-of\ TYPE('a) < m\ mod\ 2 \wedge len-of\ TYPE('a))$
 $\langle proof \rangle$

lemma *wi-le*:
 $(word-of-int\ n <= (word-of-int\ m :: 'a :: len0\ word)) =$
 $(n\ mod\ 2 \wedge len-of\ TYPE('a) <= m\ mod\ 2 \wedge len-of\ TYPE('a))$
 $\langle proof \rangle$

lemma *udvd-nat-alt*: $a\ udvd\ b = (EX\ n \geq 0. unat\ b = n * unat\ a)$
 $\langle proof \rangle$

lemma *udvd-iff-dvd*: $x\ udvd\ y \iff unat\ x\ dvd\ unat\ y$
 $\langle proof \rangle$

lemmas *unat-mono* = *word-less-nat-alt* [THEN *iffD1*, *standard*]

lemma *word-zero-neq-one*: $0 < len-of\ TYPE\ ('a :: len0) ==> (0 :: 'a\ word) \sim=$
 1
 $\langle proof \rangle$

lemmas *lenw1-zero-neq-one* = *len-gt-0* [THEN *word-zero-neq-one*]

lemma *no-no* [simp]: $number-of\ (number-of\ b) = number-of\ b$
 $\langle proof \rangle$

lemma *unat-minus-one*: $x \sim= 0 ==> unat\ (x - 1) = unat\ x - 1$
 $\langle proof \rangle$

lemma *measure-unat*: $p \sim= 0 ==> unat\ (p - 1) < unat\ p$
 $\langle proof \rangle$

lemmas *uint-add-ge0* [simp] =
add-nonneg-nonneg [OF *uint-ge-0 uint-ge-0*, *standard*]

lemmas *uint-mult-ge0* [simp] =
mult-nonneg-nonneg [OF *uint-ge-0 uint-ge-0*, *standard*]

lemma *uint-sub-lt2p* [simp]:
 $uint\ (x :: 'a :: len0\ word) - uint\ (y :: 'b :: len0\ word) <$

$2 \wedge \text{len-of TYPE('a)}$
 $\langle \text{proof} \rangle$

11.3 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:

$(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y)$
 $\langle \text{proof} \rangle$

lemma *uint-mult-lem*:

$(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$
 $\langle \text{proof} \rangle$

lemma *uint-sub-lem*:

$(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$
 $\langle \text{proof} \rangle$

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
 $\langle \text{proof} \rangle$

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
 $\langle \text{proof} \rangle$

lemmas *uint-sub-if'* =

trans [*OF uint-word-ariths*(1) *mod-sub-if-z, simplified, standard*]

lemmas *uint-plus-if'* =

trans [*OF uint-word-ariths*(2) *mod-add-if-z, simplified, standard*]

11.4 Definition of uint_arith

lemma *word-of-int-inverse*:

$\text{word-of-int } r = a \implies 0 \leq r \implies r < 2 \wedge \text{len-of TYPE('a)} \implies$
 $\text{uint } (a :: 'a :: \text{len0 word}) = r$
 $\langle \text{proof} \rangle$

lemma *uint-split*:

fixes $x :: 'a :: \text{len0 word}$

shows $P (\text{uint } x) =$

$(\text{ALL } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2 \wedge \text{len-of TYPE('a)} \implies P \ i)$

$\langle \text{proof} \rangle$

lemma *uint-split-asm*:

fixes $x :: 'a :: \text{len0 word}$

shows $P (\text{uint } x) =$

$(\sim (\text{EX } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2 \wedge \text{len-of TYPE('a)} \ \& \ \sim P \ i))$

$\langle \text{proof} \rangle$

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =
word-le-def word-less-alt
word-uint.Rep-inject [symmetric]
uint-sub-if' uint-plus-if'

lemma *power-False-cong*: $\text{False} \implies a \wedge b = c \wedge d$
 $\langle \text{proof} \rangle$

$\langle ML \rangle$

11.5 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*:
 $((x :: 'a :: \text{len0 word}) \leq x + y) = (\text{uint } x + \text{uint } y < 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add* = *no-plus-overflow-uint-size [unfolded word-size]*

lemma *no-ulen-sub*: $((x :: 'a :: \text{len0 word}) \geq x - y) = (\text{uint } y \leq \text{uint } x)$
 $\langle \text{proof} \rangle$

lemma *no-olen-add'*:
fixes $x :: 'a :: \text{len0 word}$
shows $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)})}$
 $\langle \text{proof} \rangle$

lemmas *olen-add-eqv* = *trans [OF no-olen-add no-olen-add' [symmetric], standard]*

lemmas *uint-plus-simple-iff* = *trans [OF no-olen-add uint-add-lem, standard]*
lemmas *uint-plus-simple* = *uint-plus-simple-iff [THEN iffD1, standard]*
lemmas *uint-minus-simple-iff* = *trans [OF no-ulen-sub uint-sub-lem, standard]*
lemmas *uint-minus-simple-alt* = *uint-sub-lem [folded word-le-def]*
lemmas *word-sub-le-iff* = *no-ulen-sub [folded word-le-def]*
lemmas *word-sub-le* = *word-sub-le-iff [THEN iffD2, standard]*

lemma *word-less-sub1*:
 $(x :: 'a :: \text{len word}) \sim 0 \implies (1 < x) = (0 < x - 1)$
 $\langle \text{proof} \rangle$

lemma *word-le-sub1*:
 $(x :: 'a :: \text{len word}) \sim 0 \implies (1 \leq x) = (0 \leq x - 1)$
 $\langle \text{proof} \rangle$

lemma *sub-wrap-lt*:

$$((x :: 'a :: \text{len0 word}) < x - z) = (x < z)$$

<proof>

lemma *sub-wrap*:

$$((x :: 'a :: \text{len0 word}) \leq x - z) = (z = 0 \mid x < z)$$

<proof>

lemma *plus-minus-not-NULL-ab*:

$$(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies c \sim 0 \implies x + c \sim 0$$

<proof>

lemma *plus-minus-no-overflow-ab*:

$$(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies x \leq x + c$$

<proof>

lemma *le-minus'*:

$$(a :: 'a :: \text{len0 word}) + c \leq b \implies a \leq a + c \implies c \leq b - a$$

<proof>

lemma *le-plus'*:

$$(a :: 'a :: \text{len0 word}) \leq b \implies c \leq b - a \implies a + c \leq b$$

<proof>

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*, standard]

lemma *word-plus-mono-right*:

$$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq x + z \implies x + y \leq x + z$$

<proof>

lemma *word-less-minus-cancel*:

$$y - x < z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) < z$$

<proof>

lemma *word-less-minus-mono-left*:

$$(y :: 'a :: \text{len0 word}) < z \implies x \leq y \implies y - x < z - x$$

<proof>

lemma *word-less-minus-mono*:

$$a < c \implies d < b \implies a - b < a \implies c - d < c$$

$$\implies a - b < c - (d :: 'a :: \text{len word})$$

<proof>

lemma *word-le-minus-cancel*:

$$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) \leq z$$

<proof>

lemma *word-le-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) <= z ==> x <= y ==> y - x <= z - x$
 $\langle \text{proof} \rangle$

lemma *word-le-minus-mono*:

$a <= c ==> d <= b ==> a - b <= a ==> c - d <= c$
 $==> a - b <= c - (d :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma *plus-le-left-cancel-wrap*:

$(x :: 'a :: \text{len0 word}) + y' < x ==> x + y < x ==> (x + y' < x + y) = (y' < y)$
 $\langle \text{proof} \rangle$

lemma *plus-le-left-cancel-nowrap*:

$(x :: 'a :: \text{len0 word}) <= x + y' ==> x <= x + y ==>$
 $(x + y' < x + y) = (y' < y)$
 $\langle \text{proof} \rangle$

lemma *word-plus-mono-right2*:

$(a :: 'a :: \text{len0 word}) <= a + b ==> c <= b ==> a <= a + c$
 $\langle \text{proof} \rangle$

lemma *word-less-add-right*:

$(x :: 'a :: \text{len0 word}) < y - z ==> z <= y ==> x + z < y$
 $\langle \text{proof} \rangle$

lemma *word-less-sub-right*:

$(x :: 'a :: \text{len0 word}) < y + z ==> y <= x ==> x - y < z$
 $\langle \text{proof} \rangle$

lemma *word-le-plus-either*:

$(x :: 'a :: \text{len0 word}) <= y \mid x <= z ==> y <= y + z ==> x <= y + z$
 $\langle \text{proof} \rangle$

lemma *word-less-nowrapI*:

$(x :: 'a :: \text{len0 word}) < z - k ==> k <= z ==> 0 < k ==> x < x + k$
 $\langle \text{proof} \rangle$

lemma *inc-le*: $(i :: 'a :: \text{len word}) < m ==> i + 1 <= m$

$\langle \text{proof} \rangle$

lemma *inc-i*:

$(1 :: 'a :: \text{len word}) <= i ==> i < m ==> 1 <= (i + 1) \ \& \ i + 1 <= m$
 $\langle \text{proof} \rangle$

lemma *udvd-incr-lem*:

$up < uq ==> up = ua + n * \text{uint } K ==>$
 $uq = ua + n' * \text{uint } K ==> up + \text{uint } K <= uq$

$\langle proof \rangle$

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K \leq q$
 $\langle proof \rangle$

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p \leq q - K$
 $\langle proof \rangle$

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [where *ua=0, simplified*]

lemmas *udvd-incr0* = *udvd-incr'* [where *ua=0, simplified*]

lemmas *udvd-decr0* = *udvd-decr'* [where *ua=0, simplified*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy \leq k - z$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq$
 $p \implies$
 $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *word-succ-rbl*:

$to\text{-}bl \ w = bl \implies to\text{-}bl \ (\text{word-succ } w) = (\text{rev } (\text{rbl-succ } (\text{rev } bl)))$
 $\langle proof \rangle$

lemma *word-pred-rbl*:

$to\text{-}bl \ w = bl \implies to\text{-}bl \ (\text{word-pred } w) = (\text{rev } (\text{rbl-pred } (\text{rev } bl)))$
 $\langle proof \rangle$

lemma *word-add-rbl*:

$to\text{-}bl \ v = vbl \implies to\text{-}bl \ w = wbl \implies$
 $to\text{-}bl \ (v + w) = (\text{rev } (\text{rbl-add } (\text{rev } vbl) (\text{rev } wbl)))$
 $\langle proof \rangle$

lemma *word-mult-rbl*:

$to\text{-}bl \ v = vbl \implies to\text{-}bl \ w = wbl \implies$
 $to\text{-}bl \ (v * w) = (\text{rev } (\text{rbl-mult } (\text{rev } vbl) (\text{rev } wbl)))$
 $\langle proof \rangle$

lemma *rtb-rbl-ariths*:

$$\text{rev } (\text{to-bl } w) = ys \implies \text{rev } (\text{to-bl } (\text{word-succ } w)) = \text{rbl-succ } ys$$

$$\text{rev } (\text{to-bl } w) = ys \implies \text{rev } (\text{to-bl } (\text{word-pred } w)) = \text{rbl-pred } ys$$

$$\begin{aligned} &[\text{rev } (\text{to-bl } v) = ys; \text{rev } (\text{to-bl } w) = xs] \\ \implies &\text{rev } (\text{to-bl } (v * w)) = \text{rbl-mult } ys \ xs \end{aligned}$$

$$\begin{aligned} &[\text{rev } (\text{to-bl } v) = ys; \text{rev } (\text{to-bl } w) = xs] \\ \implies &\text{rev } (\text{to-bl } (v + w)) = \text{rbl-add } ys \ xs \\ &\langle \text{proof} \rangle \end{aligned}$$

11.6 Arithmetic type class instantiations

instance *word* :: (len0) comm-monoid-add $\langle \text{proof} \rangle$

instance *word* :: (len0) comm-monoid-mult $\langle \text{proof} \rangle$

instance *word* :: (len0) comm-semiring $\langle \text{proof} \rangle$

instance *word* :: (len0) ab-group-add $\langle \text{proof} \rangle$

instance *word* :: (len0) comm-ring $\langle \text{proof} \rangle$

instance *word* :: (len) comm-semiring-1 $\langle \text{proof} \rangle$

instance *word* :: (len) comm-ring-1 $\langle \text{proof} \rangle$

instance *word* :: (len0) comm-semiring-0 $\langle \text{proof} \rangle$

instance *word* :: (len0) order $\langle \text{proof} \rangle$

lemma *zero-bintrunc*:

$$\begin{aligned} &\text{iszero } (\text{number-of } x :: 'a :: \text{len } \text{word}) = \\ &(\text{bintrunc } (\text{len-of TYPE('a)}) \ x = \text{Int.Pl}) \\ &\langle \text{proof} \rangle \end{aligned}$$

lemmas *word-le-0-iff* [simp] =
word-zero-le [THEN leD, THEN linorder-antisym-conv1]

lemma *word-of-nat*: of-nat *n* = word-of-int (int *n*) $\langle \text{proof} \rangle$

lemma *word-of-int*: of-int = word-of-int $\langle \text{proof} \rangle$

lemma *word-of-int-nat*:
 $0 \leq x \implies \text{word-of-int } x = \text{of-nat } (\text{nat } x)$
 ⟨proof⟩

lemma *word-number-of-eq*:
 $\text{number-of } w = (\text{of-int } w :: 'a :: \text{len word})$
 ⟨proof⟩

instance *word* :: (len) number-ring
 ⟨proof⟩

lemma *iszero-word-no* [simp] :
 $\text{iszero } (\text{number-of bin} :: 'a :: \text{len word}) =$
 $\text{iszero } (\text{number-of } (\text{bintrunc } (\text{len-of TYPE } ('a)) \text{ bin}) :: \text{int})$
 ⟨proof⟩

11.7 Word and nat

lemma *td-ext-unat'*:
 $n = \text{len-of TYPE } ('a :: \text{len}) \implies$
 $\text{td-ext } (\text{unat} :: 'a \text{ word} \Rightarrow \text{nat}) \text{ of-nat}$
 $(\text{unats } n) (\%i. i \bmod 2^n)$
 ⟨proof⟩

lemmas *td-ext-unat* = refl [THEN *td-ext-unat'*]
lemmas *unat-of-nat* = *td-ext-unat* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-unat*:
 $\text{td-ext unat} :: 'a :: \text{len word} \Rightarrow \text{nat}$
 of-nat
 $\text{unats } (\text{len-of TYPE } ('a :: \text{len}))$
 $\%i. i \bmod 2^{\text{len-of TYPE } ('a :: \text{len})}$
 ⟨proof⟩

lemmas *td-unat* = *word-unat.td-thm*

lemmas *unat-lt2p* [iff] = *word-unat.Rep* [unfolded *unats-def mem-Collect-eq*]

lemma *unat-le*: $y \leq \text{unat } (z :: 'a :: \text{len word}) \implies y : \text{unats } (\text{len-of TYPE } ('a))$
 ⟨proof⟩

lemma *word-nchotomy*:
 $\text{ALL } w. \text{EX } n. (w :: 'a :: \text{len word}) = \text{of-nat } n \ \& \ n < 2^{\text{len-of TYPE } ('a)}$
 ⟨proof⟩

lemma *of-nat-eq*:
fixes $w :: 'a :: \text{len word}$

shows $(\text{of-nat } n = w) = (\exists q. n = \text{unat } w + q * 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemma *of-nat-eq-size*:
 $(\text{of-nat } n = w) = (EX q. n = \text{unat } w + q * 2^{\text{size } w})$
 $\langle \text{proof} \rangle$

lemma *of-nat-0*:
 $(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemmas *of-nat-2p = mult-1 [symmetric, THEN iffD2 [OF of-nat-0 exI]]*

lemma *of-nat-gt-0*: $\text{of-nat } k \sim = 0 ==> 0 < k$
 $\langle \text{proof} \rangle$

lemma *of-nat-neq-0*:
 $0 < k ==> k < 2^{\text{len-of TYPE ('a :: len)}} ==> \text{of-nat } k \sim = (0 :: 'a \text{ word})$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-add*:
 $\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-mult*:
 $\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-Suc*:
 $\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-0*: $(0 :: 'a :: \text{len word}) = \text{of-nat } 0$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-1*: $(1 :: 'a :: \text{len word}) = \text{of-nat } (\text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemmas *Abs-fnat-homs =*
 $\text{Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc}$
 $\text{Abs-fnat-hom-0 Abs-fnat-hom-1}$

lemma *word-arith-nat-add*:
 $a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-mult*:
 $a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-arith-nat-Suc*:

word-succ a = of-nat (Suc (unat a))
<proof>

lemma *word-arith-nat-div*:

a div b = of-nat (unat a div unat b)
<proof>

lemma *word-arith-nat-mod*:

a mod b = of-nat (unat a mod unat b)
<proof>

lemmas *word-arith-nat-defs =*

word-arith-nat-add word-arith-nat-mult
word-arith-nat-Suc Abs-fnat-hom-0
Abs-fnat-hom-1 word-arith-nat-div
word-arith-nat-mod

lemmas *unat-cong = arg-cong [where f = unat]*

lemmas *unat-word-ariths = word-arith-nat-defs*

[THEN trans [OF unat-cong unat-of-nat], standard]

lemmas *word-sub-less-iff = word-sub-le-iff*

[simplified linorder-not-less [symmetric], simplified]

lemma *unat-add-lem*:

(unat x + unat y < 2 ^ len-of TYPE('a)) =
(unat (x + y :: 'a :: len word) = unat x + unat y)
<proof>

lemma *unat-mult-lem*:

*(unat x * unat y < 2 ^ len-of TYPE('a)) =*
*(unat (x * y :: 'a :: len word) = unat x * unat y)*
<proof>

lemmas *unat-plus-if' =*

trans [OF unat-word-ariths(1) mod-nat-add, simplified, standard]

lemma *le-no-overflow*:

x <= b ==> a <= a + b ==> x <= a + (b :: 'a :: len0 word)
<proof>

lemmas *un-ui-le = trans*

[OF word-le-nat-alt [symmetric]
word-le-def,
standard]

lemma *unat-sub-if-size*:

unat ($x - y$) = (if *unat* $y \leq$ *unat* x
 then *unat* $x -$ *unat* y
 else *unat* $x + 2^{\text{size } x} -$ *unat* y)
 ⟨proof⟩

lemmas *unat-sub-if' = unat-sub-if-size* [unfolded word-size]

lemma *unat-div*: *unat* (($x :: 'a :: \text{len word}$) *div* y) = *unat* x *div* *unat* y
 ⟨proof⟩

lemma *unat-mod*: *unat* (($x :: 'a :: \text{len word}$) *mod* y) = *unat* x *mod* *unat* y
 ⟨proof⟩

lemma *uint-div*: *uint* (($x :: 'a :: \text{len word}$) *div* y) = *uint* x *div* *uint* y
 ⟨proof⟩

lemma *uint-mod*: *uint* (($x :: 'a :: \text{len word}$) *mod* y) = *uint* x *mod* *uint* y
 ⟨proof⟩

11.8 Definition of unat_arith tactic

lemma *unat-split*:

fixes $x :: 'a :: \text{len word}$
shows P (*unat* x) =
 (ALL n . of-nat $n = x \ \& \ n < 2^{\text{len-of TYPE('a)}}$ $\longrightarrow P \ n$)
 ⟨proof⟩

lemma *unat-split-asm*:

fixes $x :: 'a :: \text{len word}$
shows P (*unat* x) =
 (\sim (EX n . of-nat $n = x \ \& \ n < 2^{\text{len-of TYPE('a)}}$ $\& \ \sim P \ n$))
 ⟨proof⟩

lemmas *of-nat-inverse* =

word-unat.Abs-inverse' [rotated, unfolded unats-def, simplified]

lemmas *unat-splits = unat-split unat-split-asm*

lemmas *unat-arith-simps* =

word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [symmetric]
unat-sub-if' unat-plus-if' unat-div unat-mod

⟨ML⟩

lemma *no-plus-overflow-unat-size*:

(($x :: 'a :: \text{len word}$) \leq $x + y$) = (*unat* $x +$ *unat* $y < 2^{\text{size } x}$)

<proof>

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$
<proof>

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

lemmas *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem, standard*]

lemma *word-div-mult*:
 $(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$
 $x * y \text{ div } y = x$
<proof>

lemma *div-lt'*: $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$
 $\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$
<proof>

lemmas *div-lt''* = *order-less-imp-le* [*THEN div-lt'*]

lemma *div-lt-mult*: $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$
<proof>

lemma *div-le-mult*:
 $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$
<proof>

lemma *div-lt-uint'*:
 $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$
<proof>

lemmas *div-lt-uint''* = *order-less-imp-le* [*THEN div-lt-uint'*]

lemma *word-le-exists'*:
 $(x :: 'a :: \text{len0 word}) \leq y \implies$
 $(\exists z. y = x + z \ \& \ \text{uint } x + \text{uint } z < 2^{\text{len-of TYPE('a)}})$
<proof>

lemmas *plus-minus-not-NULL* = *order-less-imp-le* [*THEN plus-minus-not-NULL-ab*]

lemmas *plus-minus-no-overflow* =
order-less-imp-le [*THEN plus-minus-no-overflow-ab*]

lemmas *mcs* = *word-less-minus-cancel word-less-minus-mono-left*
word-le-minus-cancel word-le-minus-mono-left

lemmas *word-l-diffs* = *mcs* [**where** $y = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-diff-ls* = *mcs* [**where** $z = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-plus-mcs* = *word-diff-ls*

[**where** $y = v + x$, *unfolded add-diff-cancel, standard*]

lemmas $le-unat-voi = unat-le$ [*THEN word-unat.Abs-inverse*]

lemmas $thd = refl$ [*THEN* [2] *split-div-lemma* [*THEN iffD2*], *THEN conjunct1*]

lemma $thd1$:
 $a \text{ div } b * b \leq (a :: nat)$
<proof>

lemmas $uno-simps$ [*THEN le-unat-voi, standard*] =
 $mod-le-divisor \text{ div-le-dividend } thd1$

lemma $word-mod-div-equality$:
 $(n \text{ div } b) * b + (n \text{ mod } b) = (n :: 'a :: len \text{ word})$
<proof>

lemma $word-div-mult-le$: $a \text{ div } b * b \leq (a :: 'a :: len \text{ word})$
<proof>

lemma $word-mod-less-divisor$: $0 < n ==> m \text{ mod } n < (n :: 'a :: len \text{ word})$
<proof>

lemma $word-of-int-power-hom$:
 $word-of-int \ a \ ^n = (word-of-int \ (a \ ^n) :: 'a :: len \text{ word})$
<proof>

lemma $word-arith-power-alt$:
 $a \ ^n = (word-of-int \ (uint \ a \ ^n) :: 'a :: len \text{ word})$
<proof>

lemma $of-bl-length-less$:
 $length \ x = k ==> k < len-of \ TYPE('a) ==> (of-bl \ x :: 'a :: len \text{ word}) < 2 \ ^k$
<proof>

11.9 Cardinality, finiteness of set of words

lemmas $card-lessThan' = card-lessThan$ [*unfolded lessThan-def*]

lemmas $card-eq = word-unat.Abs-inj-on$ [*THEN card-image, unfolded word-unat.image, unfolded unats-def, standard*]

lemmas $card-word = trans$ [*OF card-eq card-lessThan', standard*]

lemma $finite-word-UNIV$: $finite \ (UNIV :: 'a :: len \text{ word} \text{ set})$
<proof>

lemma $card-word-size$:
 $card \ (UNIV :: 'a :: len \text{ word} \text{ set}) = (2 \ ^{size \ (x :: 'a \text{ word})})$

<proof>

end

12 WordBitwise: Bitwise Operations on Words

theory *WordBitwise*

imports *WordArith*

begin

lemmas *bin-log-bintrs = bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

lemmas *wils1 = bin-log-bintrs [THEN word-ubin.norm-eq-iff [THEN iffD1],
folded word-ubin.eq-norm, THEN eq-reflection, standard]*

lemmas *word-log-binary-defs =
word-and-def word-or-def word-xor-def*

lemmas *word-no-log-defs [simp] =
word-not-def [where a=number-of a,
unfolded word-no-wi wils1, folded word-no-wi, standard]
word-log-binary-defs [where a=number-of a and b=number-of b,
unfolded word-no-wi wils1, folded word-no-wi, standard]*

lemmas *word-wi-log-defs = word-no-log-defs [unfolded word-no-wi]*

lemma *uint-or: uint (x OR y) = (uint x) OR (uint y)*
<proof>

lemma *uint-and: uint (x AND y) = (uint x) AND (uint y)*
<proof>

lemma *word-ops-nth-size:*
n < size (x::'a::len0 word) ==>
(x OR y) !! n = (x !! n | y !! n) &
(x AND y) !! n = (x !! n & y !! n) &
(x XOR y) !! n = (x !! n ~ y !! n) &
(NOT x) !! n = (~ x !! n)
<proof>

lemma *word-ao-nth:*
fixes *x :: 'a::len0 word*
shows *(x OR y) !! n = (x !! n | y !! n) &*
(x AND y) !! n = (x !! n & y !! n)

$\langle proof \rangle$

lemmas *bwsimps* =
word-of-int-homs(2)
word-0-wi-Pls
word-m1-wi-Min
word-wi-log-defs

lemma *word-bw-assocs*:
fixes *x* :: 'a::len0 word
shows
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle proof \rangle$

lemma *word-bw-comms*:
fixes *x* :: 'a::len0 word
shows
 $x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
 $\langle proof \rangle$

lemma *word-bw-lcs*:
fixes *x* :: 'a::len0 word
shows
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle proof \rangle$

lemma *word-log-esimps* [*simp*]:
fixes *x* :: 'a::len0 word
shows
 $x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$
 $0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$

$\langle proof \rangle$

lemma *word-not-dist*:

fixes $x :: 'a::len0$ word

shows

$NOT (x OR y) = NOT x AND NOT y$

$NOT (x AND y) = NOT x OR NOT y$

$\langle proof \rangle$

lemma *word-bw-same*:

fixes $x :: 'a::len0$ word

shows

$x AND x = x$

$x OR x = x$

$x XOR x = 0$

$\langle proof \rangle$

lemma *word-ao-absorbs* [simp]:

fixes $x :: 'a::len0$ word

shows

$x AND (y OR x) = x$

$x OR y AND x = x$

$x AND (x OR y) = x$

$y AND x OR x = x$

$(y OR x) AND x = x$

$x OR x AND y = x$

$(x OR y) AND x = x$

$x AND y OR x = x$

$\langle proof \rangle$

lemma *word-not-not* [simp]:

$NOT NOT (x::'a::len0 \text{ word}) = x$

$\langle proof \rangle$

lemma *word-ao-dist*:

fixes $x :: 'a::len0$ word

shows $(x OR y) AND z = x AND z OR y AND z$

$\langle proof \rangle$

lemma *word-oa-dist*:

fixes $x :: 'a::len0$ word

shows $x AND y OR z = (x OR z) AND (y OR z)$

$\langle proof \rangle$

lemma *word-add-not* [simp]:

fixes $x :: 'a::len0$ word

shows $x + NOT x = -1$

$\langle proof \rangle$

lemma *word-plus-and-or* [simp]:

fixes $x :: 'a::len0 \text{ word}$

shows $(x \text{ AND } y) + (x \text{ OR } y) = x + y$

$\langle \text{proof} \rangle$

lemma *leoa*:

fixes $x :: 'a::len0 \text{ word}$

shows $(w = (x \text{ OR } y)) \implies (y = (w \text{ AND } x)) \langle \text{proof} \rangle$

lemma *leao*:

fixes $x' :: 'a::len0 \text{ word}$

shows $(w' = (x' \text{ AND } y')) \implies (x' = (x' \text{ OR } w')) \langle \text{proof} \rangle$

lemmas *word-ao-equiv* = *leao* [COMP *leoa* [COMP *iffI*]]

lemma *le-word-or2*: $x \leq x \text{ OR } (y::'a::len0 \text{ word})$

$\langle \text{proof} \rangle$

lemmas *le-word-or1* = *xtr3* [OF *word-bw-comms* (2) *le-word-or2*, *standard*]

lemmas *word-and-le1* =

xtr3 [OF *word-ao-absorbs* (4) [symmetric] *le-word-or2*, *standard*]

lemmas *word-and-le2* =

xtr3 [OF *word-ao-absorbs* (8) [symmetric] *le-word-or2*, *standard*]

lemma *bl-word-not*: $\text{to-bl } (\text{NOT } w) = \text{map Not } (\text{to-bl } w)$

$\langle \text{proof} \rangle$

lemma *bl-word-xor*: $\text{to-bl } (v \text{ XOR } w) = \text{map2 op } \sim = (\text{to-bl } v) (\text{to-bl } w)$

$\langle \text{proof} \rangle$

lemma *bl-word-or*: $\text{to-bl } (v \text{ OR } w) = \text{map2 op } | (\text{to-bl } v) (\text{to-bl } w)$

$\langle \text{proof} \rangle$

lemma *bl-word-and*: $\text{to-bl } (v \text{ AND } w) = \text{map2 op } \& (\text{to-bl } v) (\text{to-bl } w)$

$\langle \text{proof} \rangle$

lemma *word-lsb-alt*: $\text{lsb } (w::'a::len0 \text{ word}) = \text{test-bit } w \ 0$

$\langle \text{proof} \rangle$

lemma *word-lsb-1-0*: $\text{lsb } (1::'a::len \text{ word}) \& \sim \text{lsb } (0::'b::len0 \text{ word})$

$\langle \text{proof} \rangle$

lemma *word-lsb-last*: $\text{lsb } (w::'a::len \text{ word}) = \text{last } (\text{to-bl } w)$

$\langle \text{proof} \rangle$

lemma *word-lsb-int*: $\text{lsb } w = (\text{uint } w \bmod 2 = 1)$

$\langle \text{proof} \rangle$

lemma *word-msb-sint*: $\text{msb } w = (\text{sint } w < 0)$

$\langle \text{proof} \rangle$

lemma *word-msb-no'*:

$w = \text{number-of bin} \implies \text{msb } (w :: 'a :: \text{len word}) = \text{bin-nth bin } (\text{size } w - 1)$
 ⟨proof⟩

lemmas *word-msb-no* = *refl* [THEN *word-msb-no'*, *unfolded word-size*]

lemma *word-msb-nth'*: $\text{msb } (w :: 'a :: \text{len word}) = \text{bin-nth } (\text{uint } w) (\text{size } w - 1)$
 ⟨proof⟩

lemmas *word-msb-nth* = *word-msb-nth'* [*unfolded word-size*]

lemma *word-msb-alt*: $\text{msb } (w :: 'a :: \text{len word}) = \text{hd } (\text{to-bl } w)$
 ⟨proof⟩

lemma *word-set-nth*:

$\text{set-bit } w \ n \ (\text{test-bit } w \ n) = (w :: 'a :: \text{len0 word})$
 ⟨proof⟩

lemma *bin-nth-uint'*:

$\text{bin-nth } (\text{uint } w) \ n = (\text{rev } (\text{bin-to-bl } (\text{size } w) (\text{uint } w))) ! \ n \ \& \ n < \text{size } w$
 ⟨proof⟩

lemmas *bin-nth-uint* = *bin-nth-uint'* [*unfolded word-size*]

lemma *test-bit-bl*: $w !! \ n = (\text{rev } (\text{to-bl } w)) ! \ n \ \& \ n < \text{size } w$
 ⟨proof⟩

lemma *to-bl-nth*: $n < \text{size } w \implies \text{to-bl } w ! \ n = w !! (\text{size } w - \text{Suc } n)$
 ⟨proof⟩

lemma *test-bit-set*:

fixes $w :: 'a :: \text{len0 word}$
shows $(\text{set-bit } w \ n \ x) !! \ n = (n < \text{size } w \ \& \ x)$
 ⟨proof⟩

lemma *test-bit-set-gen*:

fixes $w :: 'a :: \text{len0 word}$
shows $\text{test-bit } (\text{set-bit } w \ n \ x) \ m =$
 $(\text{if } m = n \text{ then } n < \text{size } w \ \& \ x \text{ else } \text{test-bit } w \ m)$
 ⟨proof⟩

lemma *of-bl-rep-False*: $\text{of-bl } (\text{replicate } n \ \text{False} \ @ \ bs) = \text{of-bl } bs$
 ⟨proof⟩

lemma *msb-nth'*:

fixes $w :: 'a :: \text{len word}$
shows $\text{msb } w = w !! (\text{size } w - 1)$
 ⟨proof⟩

lemmas $msb\text{-}nth = msb\text{-}nth'$ [unfolded word-size]

lemmas $msb0 = len\text{-}gt\text{-}0$ [THEN diff-Suc-less, THEN
word-ops-nth-size [unfolded word-size], standard]

lemmas $msb1 = msb0$ [where $i = 0$]

lemmas $word\text{-}ops\text{-}msb = msb1$ [unfolded msb-nth [symmetric, unfolded One-nat-def]]

lemmas $lsb0 = len\text{-}gt\text{-}0$ [THEN word-ops-nth-size [unfolded word-size], standard]

lemmas $word\text{-}ops\text{-}lsb = lsb0$ [unfolded word-lsb-alt]

lemma $td\text{-}ext\text{-}nth'$:

$n = size\ (w :: 'a :: len0\ word) ==> ofn = set\text{-}bits ==> [w, ofn\ g] = l ==>$
 $td\text{-}ext\ test\text{-}bit\ ofn\ \{f.\ ALL\ i.\ f\ i \longrightarrow i < n\}\ (\%h\ i.\ h\ i \ \&\ i < n)$
 <proof>

lemmas $td\text{-}ext\text{-}nth = td\text{-}ext\text{-}nth'$ [OF refl refl refl, unfolded word-size]

interpretation $test\text{-}bit$:

$td\text{-}ext\ op\ !! :: 'a :: len0\ word ==> nat ==> bool$
 $set\text{-}bits$
 $\{f.\ \forall i.\ f\ i \longrightarrow i < len\text{-}of\ TYPE('a :: len0)\}$
 $(\lambda h\ i.\ h\ i \wedge i < len\text{-}of\ TYPE('a :: len0))$
 <proof>

declare $test\text{-}bit.Rep'$ [simp del]

declare $test\text{-}bit.Rep'$ [rule del]

lemmas $td\text{-}nth = test\text{-}bit.td\text{-}thm$

lemma $word\text{-}set\text{-}set\text{-}same$:

fixes $w :: 'a :: len0\ word$
shows $set\text{-}bit\ (set\text{-}bit\ w\ n\ x)\ n\ y = set\text{-}bit\ w\ n\ y$
 <proof>

lemma $word\text{-}set\text{-}set\text{-}diff$:

fixes $w :: 'a :: len0\ word$
assumes $m \sim n$
shows $set\text{-}bit\ (set\text{-}bit\ w\ m\ x)\ n\ y = set\text{-}bit\ (set\text{-}bit\ w\ n\ y)\ m\ x$
 <proof>

lemma $test\text{-}bit\text{-}no'$:

fixes $w :: 'a :: len0\ word$
shows $w = number\text{-}of\ bin ==> test\text{-}bit\ w\ n = (n < size\ w \ \&\ bin\text{-}nth\ bin\ n)$
 <proof>

lemmas $test\text{-}bit\text{-}no =$

$refl$ [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

lemma *nth-0*: $\sim (0::'a::len0 \text{ word}) !! n$
 ⟨proof⟩

lemma *nth-sint*:
 fixes $w :: 'a::len \text{ word}$
 defines $l \equiv \text{len-of } TYPE ('a)$
 shows $\text{bin-nth } (sint \ w) \ n = (\text{if } n < l - 1 \text{ then } w !! n \text{ else } w !! (l - 1))$
 ⟨proof⟩

lemma *word-lsb-no*:
 $\text{lsb } (\text{number-of bin} :: 'a :: len \text{ word}) = (\text{bin-last bin} = \text{bit.B1})$
 ⟨proof⟩

lemma *word-set-no*:
 $\text{set-bit } (\text{number-of bin} :: 'a::len0 \text{ word}) \ n \ b =$
 $\text{number-of } (\text{bin-sc } n \ (\text{if } b \text{ then bit.B1 else bit.B0}) \ \text{bin})$
 ⟨proof⟩

lemmas *setBit-no* = *setBit-def* [THEN trans [OF meta-eq-to-obj-eq word-set-no],
 simplified if-simps, THEN eq-reflection, standard]
lemmas *clearBit-no* = *clearBit-def* [THEN trans [OF meta-eq-to-obj-eq word-set-no],
 simplified if-simps, THEN eq-reflection, standard]

lemma *to-bl-n1*:
 $\text{to-bl } (-1::'a::len0 \text{ word}) = \text{replicate } (\text{len-of } TYPE ('a)) \ \text{True}$
 ⟨proof⟩

lemma *word-msb-n1*: $\text{msb } (-1::'a::len \text{ word})$
 ⟨proof⟩

declare *word-set-set-same* [simp] *word-set-nth* [simp]
test-bit-no [simp] *word-set-no* [simp] *nth-0* [simp]
setBit-no [simp] *clearBit-no* [simp]
word-lsb-no [simp] *word-msb-no* [simp] *word-msb-n1* [simp] *word-lsb-1-0* [simp]

lemma *word-set-nth-iff*:
 $(\text{set-bit } w \ n \ b = w) = (w !! n = b \mid n >= \text{size } (w::'a::len0 \text{ word}))$
 ⟨proof⟩

lemma *test-bit-2p'*:
 $w = \text{word-of-int } (2 \wedge n) ==>$
 $w !! m = (m = n \ \& \ m < \text{size } (w :: 'a :: len \text{ word}))$
 ⟨proof⟩

lemmas *test-bit-2p* = *refl* [THEN *test-bit-2p'*, unfolded word-size]

lemma *nth-w2p*:
 $((2::'a::len \text{ word}) \wedge n) !! m \longleftrightarrow m = n \wedge m < \text{len-of } TYPE('a::len)$
 ⟨proof⟩

```

lemma uint-2p:
   $(0 :: 'a :: \text{len word}) < 2^{\wedge} n ==> \text{uint } (2^{\wedge} n :: 'a :: \text{len word}) = 2^{\wedge} n$ 
  <proof>

lemma word-of-int-2p:  $(\text{word-of-int } (2^{\wedge} n) :: 'a :: \text{len word}) = 2^{\wedge} n$ 
  <proof>

lemma bang-is-le:  $x \text{ !! } m ==> 2^{\wedge} m \leq (x :: 'a :: \text{len word})$ 
  <proof>

lemma word-clr-le:
  fixes  $w :: 'a :: \text{len0 word}$ 
  shows  $w \geq \text{set-bit } w \ n \ \text{False}$ 
  <proof>

lemma word-set-ge:
  fixes  $w :: 'a :: \text{len word}$ 
  shows  $w \leq \text{set-bit } w \ n \ \text{True}$ 
  <proof>

end

```

13 WordShift: Shifting, Rotating, and Splitting Words

```

theory WordShift
imports WordBitwise
begin

```

13.1 Bit shifting

```

lemma shiffl1-number [simp] :
   $\text{shiffl1 } (\text{number-of } w) = \text{number-of } (w \ \text{BIT } \text{bit.B0})$ 
  <proof>

lemma shiffl1-0 [simp] :  $\text{shiffl1 } 0 = 0$ 
  <proof>

lemmas shiffl1-def-u = shiffl1-def [folded word-number-of-def]

lemma shiffl1-def-s:  $\text{shiffl1 } w = \text{number-of } (\text{sint } w \ \text{BIT } \text{bit.B0})$ 
  <proof>

lemma shiftr1-0 [simp] :  $\text{shiftr1 } 0 = 0$ 
  <proof>

```

lemma *sshiftr1-0* [*simp*] : *sshiftr1 0 = 0*
 ⟨*proof*⟩

lemma *sshiftr1-n1* [*simp*] : *sshiftr1 -1 = -1*
 ⟨*proof*⟩

lemma *shiftr1-0* [*simp*] : *(0::'a::len0 word) << n = 0*
 ⟨*proof*⟩

lemma *shiftr1-0* [*simp*] : *(0::'a::len0 word) >> n = 0*
 ⟨*proof*⟩

lemma *sshiftr1-0* [*simp*] : *0 >>> n = 0*
 ⟨*proof*⟩

lemma *sshiftr1-n1* [*simp*] : *-1 >>> n = -1*
 ⟨*proof*⟩

lemma *nth-shiftr1*: *shiftr1 w !! n = (n < size w & n > 0 & w !! (n - 1))*
 ⟨*proof*⟩

lemma *nth-shiftr1'* [*rule-format*]:
ALL n. ((w::'a::len0 word) << m) !! n = (n < size w & n >= m & w !! (n - m))
 ⟨*proof*⟩

lemmas *nth-shiftr1 = nth-shiftr1'* [*unfolded word-size*]

lemma *nth-shiftr1*: *shiftr1 w !! n = w !! Suc n*
 ⟨*proof*⟩

lemma *nth-shiftr*:
∧ n. ((w::'a::len0 word) >> m) !! n = w !! (n + m)
 ⟨*proof*⟩

lemma *uint-shiftr1*: *uint (shiftr1 w) = bin-rest (uint w)*
 ⟨*proof*⟩

lemma *nth-sshiftr1*:
sshiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
 ⟨*proof*⟩

lemma *nth-sshiftr* [*rule-format*] :
ALL n. sshiftr w m !! n = (n < size w & (if n + m >= size w then w !! (size w - 1) else w !! (n + m)))
 ⟨*proof*⟩

lemma *shiftr1-div-2*: $\text{uint } (\text{shiftr1 } w) = \text{uint } w \text{ div } 2$
 $\langle \text{proof} \rangle$

lemma *sshiftr1-div-2*: $\text{sint } (\text{sshiftr1 } w) = \text{sint } w \text{ div } 2$
 $\langle \text{proof} \rangle$

lemma *shiftr-div-2n*: $\text{uint } (\text{shiftr } w \ n) = \text{uint } w \text{ div } 2 \wedge n$
 $\langle \text{proof} \rangle$

lemma *sshiftr-div-2n*: $\text{sint } (\text{sshiftr } w \ n) = \text{sint } w \text{ div } 2 \wedge n$
 $\langle \text{proof} \rangle$

13.1.1 shift functions in terms of lists of bools

lemmas *bshiftr1-no-bin* [simp] =
bshiftr1-def [where $w = \text{number-of } w$, *unfolded to-bl-no-bin*, *standard*]

lemma *bshiftr1-bl*: $\text{to-bl } (\text{bshiftr1 } b \ w) = b \ \# \ \text{butlast } (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *shiftrl1-of-bl*: $\text{shiftrl1 } (\text{of-bl } bl) = \text{of-bl } (bl \ @ \ [\text{False}])$
 $\langle \text{proof} \rangle$

lemma *shiftrl1-bl*: $\text{shiftrl1 } (w :: 'a :: \text{len0 word}) = \text{of-bl } (\text{to-bl } w \ @ \ [\text{False}])$
 $\langle \text{proof} \rangle$

lemma *bl-shiftrl1*:
 $\text{to-bl } (\text{shiftrl1 } (w :: 'a :: \text{len word})) = \text{tl } (\text{to-bl } w) \ @ \ [\text{False}]$
 $\langle \text{proof} \rangle$

lemma *shiftr1-bl*: $\text{shiftr1 } w = \text{of-bl } (\text{butlast } (\text{to-bl } w))$
 $\langle \text{proof} \rangle$

lemma *bl-shiftr1*:
 $\text{to-bl } (\text{shiftr1 } (w :: 'a :: \text{len word})) = \text{False} \ \# \ \text{butlast } (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *shiftrl1-rev*:
 $\text{shiftrl1 } (w :: 'a :: \text{len word}) = \text{word-reverse } (\text{shiftr1 } (\text{word-reverse } w))$
 $\langle \text{proof} \rangle$

lemma *shiftrl-rev*:
 $\text{shiftrl } (w :: 'a :: \text{len word}) \ n = \text{word-reverse } (\text{shiftr } (\text{word-reverse } w) \ n)$
 $\langle \text{proof} \rangle$

lemmas *rev-shiftrl* =

shiffl-rev [where $w = \text{word-reverse } w$, *simplified*, *standard*]

lemmas *shiftr-rev* = *rev-shiffl* [THEN *word-rev-gal'*, *standard*]

lemmas *rev-shiftr* = *shiffl-rev* [THEN *word-rev-gal'*, *standard*]

lemma *bl-sshiftr1*:

$\text{to-bl } (\text{sshiftr1 } (w :: 'a :: \text{len word})) = \text{hd } (\text{to-bl } w) \# \text{butlast } (\text{to-bl } w)$
 ⟨proof⟩

lemma *drop-shiftr*:

$\text{drop } n \ (\text{to-bl } ((w :: 'a :: \text{len word}) >> n)) = \text{take } (\text{size } w - n) \ (\text{to-bl } w)$
 ⟨proof⟩

lemma *drop-sshiftr*:

$\text{drop } n \ (\text{to-bl } ((w :: 'a :: \text{len word}) >>> n)) = \text{take } (\text{size } w - n) \ (\text{to-bl } w)$
 ⟨proof⟩

lemma *take-shiftr* [rule-format] :

$n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{take } n \ (\text{to-bl } (w >> n)) =$
 $\text{replicate } n \ \text{False}$
 ⟨proof⟩

lemma *take-sshiftr'* [rule-format] :

$n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{hd } (\text{to-bl } (w >>> n)) = \text{hd } (\text{to-bl } w) \ \&$
 $\text{take } n \ (\text{to-bl } (w >>> n)) = \text{replicate } n \ (\text{hd } (\text{to-bl } w))$
 ⟨proof⟩

lemmas *hd-sshiftr* = *take-sshiftr'* [THEN *conjunct1*, *standard*]

lemmas *take-sshiftr* = *take-sshiftr'* [THEN *conjunct2*, *standard*]

lemma *atd-lem*: $\text{take } n \ xs = t \implies \text{drop } n \ xs = d \implies xs = t @ d$
 ⟨proof⟩

lemmas *bl-shiftr* = *atd-lem* [OF *take-shiftr drop-shiftr*]

lemmas *bl-sshiftr* = *atd-lem* [OF *take-sshiftr drop-sshiftr*]

lemma *shiffl-of-bl*: $\text{of-bl } bl \ll n = \text{of-bl } (bl @ \text{replicate } n \ \text{False})$
 ⟨proof⟩

lemma *shiffl-bl*:

$(w :: 'a :: \text{len0 word}) \ll (n :: \text{nat}) = \text{of-bl } (\text{to-bl } w @ \text{replicate } n \ \text{False})$
 ⟨proof⟩

lemmas *shiffl-number* [simp] = *shiffl-def* [where $w = \text{number-of } w$, *standard*]

lemma *bl-shiffl*:

$\text{to-bl } (w \ll n) = \text{drop } n \ (\text{to-bl } w) @ \text{replicate } (\min (\text{size } w) \ n) \ \text{False}$
 ⟨proof⟩

lemma *shiffl-zero-size*:

fixes $x :: 'a::len0 \text{ word}$

shows $\text{size } x \leq n \implies x \ll n = 0$

<proof>

lemma *shiffl1-2t*: $\text{shiffl1 } (w :: 'a :: len \text{ word}) = 2 * w$

<proof>

lemma *shiffl1-p*: $\text{shiffl1 } (w :: 'a :: len \text{ word}) = w + w$

<proof>

lemma *shiffl-t2n*: $\text{shiffl } (w :: 'a :: len \text{ word}) \ n = 2 ^ n * w$

<proof>

lemma *shiftr1-bintr* [simp]:

$(\text{shiftr1 } (\text{number-of } w) :: 'a :: len0 \text{ word}) =$

$\text{number-of } (\text{bin-rest } (\text{bintrunc } (\text{len-of TYPE } ('a)) \ w))$

<proof>

lemma *sshiftr1-sbintr* [simp] :

$(\text{sshiftr1 } (\text{number-of } w) :: 'a :: len \text{ word}) =$

$\text{number-of } (\text{bin-rest } (\text{sbintrunc } (\text{len-of TYPE } ('a) - 1) \ w))$

<proof>

lemma *shiftr-no'*:

$w = \text{number-of bin} \implies$

$(w :: 'a :: len0 \text{ word}) \gg n = \text{number-of } ((\text{bin-rest } ^n) (\text{bintrunc } (\text{size } w) \ \text{bin}))$

<proof>

lemma *sshiftr-no'*:

$w = \text{number-of bin} \implies w \ggg n = \text{number-of } ((\text{bin-rest } ^n)$

$(\text{sbintrunc } (\text{size } w - 1) \ \text{bin}))$

<proof>

lemmas *sshiftr-no* [simp] =

sshiftr-no' [where $w = \text{number-of } w$, OF *refl*, *unfolded word-size*, *standard*]

lemmas *shiftr-no* [simp] =

shiftr-no' [where $w = \text{number-of } w$, OF *refl*, *unfolded word-size*, *standard*]

lemma *shiftr1-bl-of'*:

$us = \text{shiftr1 } (\text{of-bl } bl) \implies \text{length } bl \leq \text{size } us \implies$

$us = \text{of-bl } (\text{butlast } bl)$

<proof>

lemmas *shiftr1-bl-of* = *refl* [THEN *shiftr1-bl-of'*, *unfolded word-size*]

lemma *shiftr-bl-of'* [rule-format]:

us = of-bl bl >> n ==> length bl <= size us -->
us = of-bl (take (length bl - n) bl)
 ⟨proof⟩

lemmas *shiftr-bl-of* = *refl* [THEN *shiftr-bl-of'*, *unfolded word-size*]

lemmas *shiftr-bl* = *word-bl.Rep'* [THEN *eq-imp-le*, THEN *shiftr-bl-of*,
simplified word-size, *simplified*, THEN *eq-reflection*, *standard*]

lemma *msb-shift'*: *msb (w::'a::len word) <-> (w >> (size w - 1)) ~ = 0*
 ⟨proof⟩

lemmas *msb-shift* = *msb-shift'* [*unfolded word-size*]

lemma *align-lem-or* [rule-format] :

ALL x m. length x = n + m --> length y = n + m -->
drop m x = replicate n False --> take m y = replicate m False -->
map2 op | x y = take m x @ drop m y
 ⟨proof⟩

lemma *align-lem-and* [rule-format] :

ALL x m. length x = n + m --> length y = n + m -->
drop m x = replicate n False --> take m y = replicate m False -->
map2 op & x y = replicate (n + m) False
 ⟨proof⟩

lemma *aligned-bl-add-size'*:

size x - n = m ==> n <= size x ==> drop m (to-bl x) = replicate n False
==>
take m (to-bl y) = replicate m False ==>
to-bl (x + y) = take m (to-bl x) @ drop m (to-bl y)
 ⟨proof⟩

lemmas *aligned-bl-add-size* = *refl* [THEN *aligned-bl-add-size'*]

13.1.2 Mask

lemma *nth-mask'*: *m = mask n ==> test-bit m i = (i < n & i < size m)*
 ⟨proof⟩

lemmas *nth-mask* [*simp*] = *refl* [THEN *nth-mask'*]

lemma *mask-bl*: *mask n = of-bl (replicate n True)*
 ⟨proof⟩

lemma *mask-bin*: *mask n = number-of (bintrunc n Int.Min)*
 ⟨proof⟩

lemma *and-mask-bintr*: $w \text{ AND } \text{mask } n = \text{number-of } (\text{bintrunc } n \text{ (uint } w))$
 ⟨proof⟩

lemma *and-mask-no*: $\text{number-of } i \text{ AND } \text{mask } n = \text{number-of } (\text{bintrunc } n \text{ } i)$
 ⟨proof⟩

lemmas *and-mask-wi* = *and-mask-no* [unfolded word-number-of-def]

lemma *bl-and-mask*:
 $\text{to-bl } (w \text{ AND } \text{mask } n :: 'a :: \text{len word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{drop } (\text{len-of TYPE('a)} - n) (\text{to-bl } w)$
 ⟨proof⟩

lemmas *and-mask-mod-2p* =
and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma *and-mask-lt-2p*: $\text{uint } (w \text{ AND } \text{mask } n) < 2 \wedge n$
 ⟨proof⟩

lemmas *eq-mod-iff* = *trans* [symmetric, OF int-mod-lem eq-sym-conv]

lemma *mask-eq-iff*: $(w \text{ AND } \text{mask } n) = w \iff \text{uint } w < 2 \wedge n$
 ⟨proof⟩

lemma *and-mask-dvd*: $2 \wedge n \text{ dvd uint } w = (w \text{ AND } \text{mask } n = 0)$
 ⟨proof⟩

lemma *and-mask-dvd-nat*: $2 \wedge n \text{ dvd unat } w = (w \text{ AND } \text{mask } n = 0)$
 ⟨proof⟩

lemma *word-2p-lem*:
 $n < \text{size } w \implies w < 2 \wedge n = (\text{uint } (w :: 'a :: \text{len word}) < 2 \wedge n)$
 ⟨proof⟩

lemma *less-mask-eq*: $x < 2 \wedge n \implies x \text{ AND } \text{mask } n = (x :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *mask-eq-iff-w2p* =
trans [OF mask-eq-iff word-2p-lem [symmetric], standard]

lemmas *and-mask-less'* =
iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND } \text{mask } n < 2 \wedge n$
 ⟨proof⟩

lemma *word-mod-2p-is-mask'*:
 $c = 2 \wedge n \implies c > 0 \implies x \text{ mod } c = (x :: 'a :: \text{len word}) \text{ AND } \text{mask } n$

<proof>

lemmas *word-mod-2p-is-mask* = refl [THEN *word-mod-2p-is-mask'*]

lemma *mask-egs*:

$(a \text{ AND } \text{mask } n) + b \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$
 $a + (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) - b \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$
 $a - (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$
 $a * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$
 $(b \text{ AND } \text{mask } n) * a \text{ AND } \text{mask } n = b * a \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) + (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) - (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$
 $(a \text{ AND } \text{mask } n) * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$
 $-(a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = -a \text{ AND } \text{mask } n$
 $\text{word-succ } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-succ } a \text{ AND } \text{mask } n$
 $\text{word-pred } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-pred } a \text{ AND } \text{mask } n$
<proof>

lemma *mask-power-eq*:

$(x \text{ AND } \text{mask } n) \wedge^k \text{ AND } \text{mask } n = x \wedge^k \text{ AND } \text{mask } n$
<proof>

13.1.3 Revcast

lemmas *revcast-def'* = *revcast-def* [*simplified*]

lemmas *revcast-def''* = *revcast-def'* [*simplified word-size*]

lemmas *revcast-no-def* [*simp*] =
revcast-def' [**where** *w=number-of w, unfolded word-size, standard*]

lemma *to-bl-revcast*:

$\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
<proof>

lemma *revcast-rev-ucast'*:

$cs = [rc, uc] ==> rc = \text{revcast } (\text{word-reverse } w) ==> uc = \text{ucast } w ==>$
 $rc = \text{word-reverse } uc$
<proof>

lemmas *revcast-rev-ucast* = *revcast-rev-ucast'* [*OF refl refl refl*]

lemmas *revcast-ucast* = *revcast-rev-ucast*

[**where** *w = word-reverse w, simplified word-rev-rev, standard*]

lemmas *ucast-revcast* = *revcast-rev-ucast* [*THEN word-rev-gal', standard*]

lemmas *ucast-rev-revcast* = *revcast-ucast* [*THEN word-rev-gal', standard*]

— linking revcast and cast via shift

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *revcast-down-uu'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc \ (w :: 'a :: \text{len word}) = \text{ucast } (w >> n)$
 ⟨proof⟩

lemma *revcast-down-us'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc \ (w :: 'a :: \text{len word}) = \text{ucast } (w >>> n)$
 ⟨proof⟩

lemma *revcast-down-su'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc \ (w :: 'a :: \text{len word}) = \text{scast } (w >> n)$
 ⟨proof⟩

lemma *revcast-down-ss'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc \ (w :: 'a :: \text{len word}) = \text{scast } (w >>> n)$
 ⟨proof⟩

lemmas *revcast-down-uu* = *refl [THEN revcast-down-uu']*

lemmas *revcast-down-us* = *refl [THEN revcast-down-us']*

lemmas *revcast-down-su* = *refl [THEN revcast-down-su']*

lemmas *revcast-down-ss* = *refl [THEN revcast-down-ss']*

lemma *cast-down-rev*:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies$
 $uc \ w = \text{revcast } ((w :: 'a :: \text{len word}) << n)$
 ⟨proof⟩

lemma *revcast-up'*:

$rc = \text{revcast} \implies \text{source-size } rc + n = \text{target-size } rc \implies$
 $rc \ w = (\text{ucast } w :: 'a :: \text{len word}) << n$
 ⟨proof⟩

lemmas *revcast-up* = *refl [THEN revcast-up']*

lemmas *rc1* = *revcast-up [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *rc2* = *revcast-down-uu [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *ucast-up* =

rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

lemmas *ucast-down* =

rc2 [*simplified rev-shiftr revcast-ucast* [*symmetric*]]

13.1.4 Slices

lemmas *slice1-no-bin* [*simp*] =
slice1-def [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

lemmas *slice-no-bin* [*simp*] =
trans [*OF slice-def* [*THEN meta-eq-to-obj-eq*]
slice1-no-bin [*THEN meta-eq-to-obj-eq*],
unfolded word-size, standard]

lemma *slice1-0* [*simp*] : *slice1 n 0 = 0*
 ⟨*proof*⟩

lemma *slice-0* [*simp*] : *slice n 0 = 0*
 ⟨*proof*⟩

lemma *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*
 ⟨*proof*⟩

lemmas *slice-take = slice-take'* [*unfolded word-size*]

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

lemmas *shiftr-slice = trans*
 [*OF shiftr-bl* [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*], *standard*]

lemma *slice-shiftr*: *slice n w = ucast (w >> n)*
 ⟨*proof*⟩

lemma *nth-slice*:
 (*slice n w :: 'a :: len0 word*) !! *m* =
 (*w* !! (*m* + *n*) & *m* < *len-of TYPE ('a)*)
 ⟨*proof*⟩

lemma *slice1-down-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>
to-bl sl = takefill False fs (drop k (to-bl w))
 ⟨*proof*⟩

lemma *slice1-up-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>
to-bl sl = takefill False fs (replicate k False @ (to-bl w))
 ⟨*proof*⟩

lemmas *sd1 = slice1-down-alt'* [*OF refl refl, unfolded word-size*]

lemmas *su1 = slice1-up-alt'* [*OF refl refl, unfolded word-size*]

lemmas *slice1-down-alt = le-add-diff-inverse* [*THEN sd1*]

lemmas *slice1-up-alt =*

le-add-diff-inverse [symmetric, THEN *su1*]
le-add-diff-inverse2 [symmetric, THEN *su1*]

lemma *ucast-slice1*: *ucast w = slice1 (size w) w*
 ⟨proof⟩

lemma *ucast-slice*: *ucast w = slice 0 w*
 ⟨proof⟩

lemmas *slice-id = trans [OF ucast-slice [symmetric] ucast-id]*

lemma *revcast-slice1'*:
rc = revcast w ==> slice1 (size rc) w = rc
 ⟨proof⟩

lemmas *revcast-slice1 = refl [THEN revcast-slice1']*

lemma *slice1-tf-tf'*:
to-bl (slice1 n w :: 'a :: len0 word) =
rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))
 ⟨proof⟩

lemmas *slice1-tf-tf' = slice1-tf-tf'*
 [THEN *word-bl.Rep-inverse'*, symmetric, standard]

lemma *rev-slice1*:
n + k = len-of TYPE('a) + len-of TYPE('b) ==>
slice1 n (word-reverse w :: 'b :: len0 word) =
word-reverse (slice1 k w :: 'a :: len0 word)
 ⟨proof⟩

lemma *rev-slice'*:
res = slice n (word-reverse w) ==> n + k + size res = size w ==>
res = word-reverse (slice k w)
 ⟨proof⟩

lemmas *rev-slice = refl [THEN rev-slice', unfolded word-size]*

lemmas *sym-notr =*
not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:
(sint (x :: 'a :: len word) + sint y = sint (x + y)) =
((((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)
 ⟨proof⟩

13.2 Split and cat

lemmas *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

lemmas *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-rsplit-no*:

(*word-rsplit* (*number-of* *bin* :: 'b :: len0 word) :: 'a word list) =
 map *number-of* (*bin-rsplit* (*len-of TYPE*('a :: len))
 (*len-of TYPE*('b), *bintrunc* (*len-of TYPE*('b)) *bin*))
 ⟨*proof*⟩

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*
 [*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

lemma *test-bit-cat*:

wc = *word-cat* *a b* ==> *wc* !! *n* = (*n* < *size wc* &
 (*if n* < *size b* then *b* !! *n* else *a* !! (*n* - *size b*)))
 ⟨*proof*⟩

lemma *word-cat-bl*: *word-cat* *a b* = *of-bl* (*to-bl* *a* @ *to-bl* *b*)
 ⟨*proof*⟩

lemma *of-bl-append*:

(*of-bl* (*xs* @ *ys*) :: 'a :: len word) = *of-bl* *xs* * 2^(*length ys*) + *of-bl* *ys*
 ⟨*proof*⟩

lemma *of-bl-False* [*simp*]:

of-bl (*False*#*xs*) = *of-bl* *xs*
 ⟨*proof*⟩

lemma *of-bl-True*:

(*of-bl* (*True*#*xs*)::'a::len word) = 2^{*length xs*} + *of-bl* *xs*
 ⟨*proof*⟩

lemma *of-bl-Cons*:

of-bl (*x*#*xs*) = *of-bool* *x* * 2^{*length xs*} + *of-bl* *xs*
 ⟨*proof*⟩

lemma *split-uint-lem*: *bin-split* *n* (*uint* (*w* :: 'a :: len0 word)) = (*a*, *b*) ==>
a = *bintrunc* (*len-of TYPE*('a) - *n*) *a* & *b* = *bintrunc* (*len-of TYPE*('a)) *b*
 ⟨*proof*⟩

lemma *word-split-bl'*:

std = *size c* - *size b* ==> (*word-split* *c* = (*a*, *b*)) ==>
 (*a* = *of-bl* (*take std* (*to-bl* *c*)) & *b* = *of-bl* (*drop std* (*to-bl* *c*)))
 ⟨*proof*⟩

lemma *word-split-bl*: *std* = *size c* - *size b* ==>

(*a* = *of-bl* (*take std* (*to-bl* *c*)) & *b* = *of-bl* (*drop std* (*to-bl* *c*))) <->
word-split *c* = (*a*, *b*)

$\langle \text{proof} \rangle$

lemma *word-split-bl-eq*:

$(\text{word-split } (c :: 'a :: \text{len word}) :: ('c :: \text{len0 word} * 'd :: \text{len0 word})) =$
 $(\text{of-bl } (\text{take } (\text{len-of TYPE('a :: len)} - \text{len-of TYPE('d :: len0)}) (\text{to-bl } c)),$
 $\text{of-bl } (\text{drop } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)}) (\text{to-bl } c)))$
 $\langle \text{proof} \rangle$

lemma *test-bit-split'*:

$\text{word-split } c = (a, b) \dashv\dashv (ALL\ n\ m.\ b\ !!\ n = (n < \text{size } b \ \&\ c\ !!\ n) \ \&$
 $a\ !!\ m = (m < \text{size } a \ \&\ c\ !!\ (m + \text{size } b)))$
 $\langle \text{proof} \rangle$

lemmas *test-bit-split* =

test-bit-split' [THEN mp, simplified all-simps, standard]

lemma *test-bit-split-eq*: $\text{word-split } c = (a, b) \dashv\dashv$

$((ALL\ n :: \text{nat}.\ b\ !!\ n = (n < \text{size } b \ \&\ c\ !!\ n)) \ \&$
 $(ALL\ m :: \text{nat}.\ a\ !!\ m = (m < \text{size } a \ \&\ c\ !!\ (m + \text{size } b))))$
 $\langle \text{proof} \rangle$

lemma *word-cat-id*: $\text{word-cat } a\ b = b$

$\langle \text{proof} \rangle$

lemma *word-cat-hom*:

$\text{len-of TYPE('a :: len0)} \leq \text{len-of TYPE('b :: len0)} + \text{len-of TYPE('c :: len0)}$
 \implies
 $(\text{word-cat } (\text{word-of-int } w :: 'b\ \text{word})\ (b :: 'c\ \text{word}) :: 'a\ \text{word}) =$
 $\text{word-of-int } (\text{bin-cat } w\ (\text{size } b)\ (\text{uint } b))$
 $\langle \text{proof} \rangle$

lemma *word-cat-split-alt*:

$\text{size } w \leq \text{size } u + \text{size } v \implies \text{word-split } w = (u, v) \implies \text{word-cat } u\ v = w$
 $\langle \text{proof} \rangle$

lemmas *word-cat-split-size* =

sym [THEN [2] word-cat-split-alt [symmetric], standard]

13.2.1 Split and slice

lemma *split-slices*:

$\text{word-split } w = (u, v) \implies u = \text{slice } (\text{size } v)\ w \ \&\ v = \text{slice } 0\ w$
 $\langle \text{proof} \rangle$

lemma *slice-cat1'*:

$wc = \text{word-cat } a\ b \implies \text{size } wc \geq \text{size } a + \text{size } b \implies \text{slice } (\text{size } b)\ wc = a$
 $\langle \text{proof} \rangle$

lemmas *slice-cat1* = *refl [THEN slice-cat1']*

lemmas *slice-cat2* = *trans [OF slice-id word-cat-id]*

lemma *cat-slices*:

$a = \text{slice } n \ c ==> b = \text{slice } 0 \ c ==> n = \text{size } b ==>$
 $\text{size } a + \text{size } b >= \text{size } c ==> \text{word-cat } a \ b = c$
 ⟨proof⟩

lemma *word-split-cat-alt*:

$w = \text{word-cat } u \ v ==> \text{size } u + \text{size } v <= \text{size } w ==> \text{word-split } w = (u, v)$
 ⟨proof⟩

lemmas *word-cat-bl-no-bin* [simp] =

word-cat-bl [where $a = \text{number-of } a$
 and $b = \text{number-of } b$,
 unfolded to-bl-no-bin, standard]

lemmas *word-split-bl-no-bin* [simp] =

word-split-bl-eq [where $c = \text{number-of } c$, unfolded to-bl-no-bin, standard]

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

lemma *word-rsplit-same*: $\text{word-rsplit } w = [w]$

⟨proof⟩

lemma *word-rsplit-empty-iff-size*:

$(\text{word-rsplit } w = []) = (\text{size } w = 0)$
 ⟨proof⟩

lemma *test-bit-rsplit*:

$sw = \text{word-rsplit } w ==> m < \text{size } (\text{hd } sw :: 'a :: \text{len word}) ==>$
 $k < \text{length } sw ==> (\text{rev } sw ! k) !! m = (w !! (k * \text{size } (\text{hd } sw) + m))$
 ⟨proof⟩

lemma *word-rcat-bl*: $\text{word-rcat } wl == \text{of-bl } (\text{concat } (\text{map to-bl } wl))$

⟨proof⟩

lemma *size-rcat-lem'*:

$\text{size } (\text{concat } (\text{map to-bl } wl)) = \text{length } wl * \text{size } (\text{hd } wl)$
 ⟨proof⟩

lemmas *size-rcat-lem* = *size-rcat-lem'* [unfolded word-size]

lemmas *td-gal-lt-len* = *len-gt-0* [THEN *td-gal-lt*, standard]

lemma *nth-rcat-lem'* [rule-format] :

$sw = \text{size } (\text{hd } wl :: 'a :: \text{len word}) ==> (\text{ALL } n. n < \text{size } wl * sw \longrightarrow$
 $\text{rev } (\text{concat } (\text{map to-bl } wl)) ! n =$
 $\text{rev } (\text{to-bl } (\text{rev } wl ! (n \text{ div } sw))) ! (n \text{ mod } sw))$
 ⟨proof⟩

lemmas *nth-rcat-lem* = refl [THEN *nth-rcat-lem'*, unfolded *word-size*]

lemma *test-bit-rcat*:

sw = size (hd *wl* :: 'a :: len *word*) ==> *rc* = word-rcat *wl* ==> *rc* !! *n* =
 (*n* < size *rc* & *n* div *sw* < size *wl* & (rev *wl*) ! (*n* div *sw*) !! (*n* mod *sw*))
 <proof>

lemma *foldl-eq-foldr* [rule-format] :

ALL *x*. foldl *op* + *x* *xs* = foldr *op* + (*x* # *xs*) (0 :: 'a :: comm-monoid-add)
 <proof>

lemmas *test-bit-cong* = arg-cong [where *f* = *test-bit*, THEN *fun-cong*]

lemmas *test-bit-rsplit-alt* =

trans [OF *nth-rev-alt* [THEN *test-bit-cong*]
test-bit-rsplit [OF refl *asm-rl diff-Suc-less*]]

— lazy way of expressing that *u* and *v*, and *su* and *sv*, have same types

lemma *word-rsplit-len-indep'*:

[*u,v*] = *p* ==> [*su,sv*] = *q* ==> word-rsplit *u* = *su* ==>
 word-rsplit *v* = *sv* ==> length *su* = length *sv*
 <proof>

lemmas *word-rsplit-len-indep* = *word-rsplit-len-indep'* [OF refl refl refl refl]

lemma *length-word-rsplit-size*:

n = len-of TYPE ('a :: len) ==>
 (length (word-rsplit *w* :: 'a word list) <= *m*) = (size *w* <= *m* * *n*)
 <proof>

lemmas *length-word-rsplit-lt-size* =

length-word-rsplit-size [unfolded *Not-eq-iff linorder-not-less* [symmetric]]

lemma *length-word-rsplit-exp-size*:

n = len-of TYPE ('a :: len) ==>
 length (word-rsplit *w* :: 'a word list) = (size *w* + *n* - 1) div *n*
 <proof>

lemma *length-word-rsplit-even-size*:

n = len-of TYPE ('a :: len) ==> size *w* = *m* * *n* ==>
 length (word-rsplit *w* :: 'a word list) = *m*
 <proof>

lemmas *length-word-rsplit-exp-size'* = refl [THEN *length-word-rsplit-exp-size*]

lemmas *tdle* = iffD2 [OF *split-div-lemma* refl, THEN *conjunct1*]

lemmas *dtle* = xtr4 [OF *tdle* mult-commute]

lemma *word-rcat-rsplit*: *word-rcat (word-rsplit w) = w*
 ⟨*proof*⟩

lemma *size-word-rsplit-rcat-size'*:
word-rcat (ws :: 'a :: len word list) = frcw ==>
*size frcw = length ws * len-of TYPE ('a) ==>*
size (hd [word-rsplit frcw, ws]) = size ws
 ⟨*proof*⟩

lemmas *size-word-rsplit-rcat-size =*
size-word-rsplit-rcat-size' [simplified]

lemma *msrevs*:
fixes *n::nat*
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
 ⟨*proof*⟩

lemma *word-rsplit-rcat-size'*:
word-rcat (ws :: 'a :: len word list) = frcw ==>
*size frcw = length ws * len-of TYPE ('a) ==> word-rsplit frcw = ws*
 ⟨*proof*⟩

lemmas *word-rsplit-rcat-size = refl [THEN word-rsplit-rcat-size']*

13.3 Rotation

lemmas *rotater-0'* [*simp*] = *rotater-def [where n = 0, simplified]*

lemmas *word-rot-defs = word-roti-def word-rotr-def word-rotl-def*

lemma *rotate-eq-mod*:
m mod length xs = n mod length xs ==> rotate m xs = rotate n xs
 ⟨*proof*⟩

lemmas *rotate-egs [standard] =*
trans [OF rotate0 [THEN fun-cong] id-apply]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

13.3.1 Rotation of list to right

lemma *rotate1-rl'*: *rotater1 (l @ [a]) = a # l*
 ⟨*proof*⟩

lemma *rotate1-rl [simp]* : *rotater1 (rotate1 l) = l*
 ⟨*proof*⟩

lemma *rotate1-lr* [*simp*] : *rotate1* (*rotater1* *l*) = *l*
 ⟨*proof*⟩

lemma *rotater1-rev'*: *rotater1* (*rev xs*) = *rev* (*rotate1 xs*)
 ⟨*proof*⟩

lemma *rotater-rev'*: *rotater* *n* (*rev xs*) = *rev* (*rotate n xs*)
 ⟨*proof*⟩

lemmas *rotater-rev* = *rotater-rev'* [where *xs* = *rev ys*, *simplified*, *standard*]

lemma *rotater-drop-take*:
rotater *n xs* =
 drop (*length xs* − *n mod length xs*) *xs* @
 take (*length xs* − *n mod length xs*) *xs*
 ⟨*proof*⟩

lemma *rotater-Suc* [*simp*] :
rotater (*Suc n*) *xs* = *rotater1* (*rotater n xs*)
 ⟨*proof*⟩

lemma *rotate-inv-plus* [*rule-format*] :
 ALL *k*. *k* = *m* + *n* --> *rotater k* (*rotate n xs*) = *rotater m xs* &
 rotate k (*rotater n xs*) = *rotate m xs* &
 rotater n (*rotate k xs*) = *rotate m xs* &
 rotate n (*rotater k xs*) = *rotater m xs*
 ⟨*proof*⟩

lemmas *rotate-inv-rel* = *le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

lemmas *rotate-inv-eq* = *order-refl* [*THEN rotate-inv-rel*, *simplified*]

lemmas *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*, *standard*]

lemmas *rotate-rl* [*simp*] =
rotate-inv-eq [*THEN conjunct2*, *THEN conjunct1*, *standard*]

lemma *rotate-gal*: (*rotater n xs* = *ys*) = (*rotate n ys* = *xs*)
 ⟨*proof*⟩

lemma *rotate-gal'*: (*ys* = *rotater n xs*) = (*xs* = *rotate n ys*)
 ⟨*proof*⟩

lemma *length-rotater* [*simp*]:
length (*rotater n xs*) = *length xs*
 ⟨*proof*⟩

lemmas *rrs0* = *rotate-eqs* [*THEN restrict-to-left*,
simplified rotate-gal [*symmetric*] *rotate-gal'* [*symmetric*], *standard*]

lemmas $rrs1 = rrs0$ [THEN refl [THEN rev-iffD1]]
lemmas $rotater-eqs = rrs1$ [simplified length-rotater, standard]
lemmas $rotater-0 = rotater-eqs$ (1)
lemmas $rotater-add = rotater-eqs$ (2)

13.3.2 map, map2, commuting with rotate(r)

lemma $last-map$: $xs \sim [] ==> last (map f xs) = f (last xs)$
 <proof>

lemma $butlast-map$:
 $xs \sim [] ==> butlast (map f xs) = map f (butlast xs)$
 <proof>

lemma $rotater1-map$: $rotater1 (map f xs) = map f (rotater1 xs)$
 <proof>

lemma $rotater-map$:
 $rotater n (map f xs) = map f (rotater n xs)$
 <proof>

lemma $but-last-zip$ [rule-format] :
 ALL ys . $length xs = length ys --> xs \sim [] -->$
 $last (zip xs ys) = (last xs, last ys) \&$
 $butlast (zip xs ys) = zip (butlast xs) (butlast ys)$
 <proof>

lemma $but-last-map2$ [rule-format] :
 ALL ys . $length xs = length ys --> xs \sim [] -->$
 $last (map2 f xs ys) = f (last xs) (last ys) \&$
 $butlast (map2 f xs ys) = map2 f (butlast xs) (butlast ys)$
 <proof>

lemma $rotater1-zip$:
 $length xs = length ys ==>$
 $rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)$
 <proof>

lemma $rotater1-map2$:
 $length xs = length ys ==>$
 $rotater1 (map2 f xs ys) = map2 f (rotater1 xs) (rotater1 ys)$
 <proof>

lemmas $lrth =$
 $box-equals$ [OF $asm-rl$ length-rotater [symmetric]
 $length-rotater$ [symmetric],
 THEN $rotater1-map2$]

lemma $rotater-map2$:

$length\ xs = length\ ys ==>$
 $rotater\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotater\ n\ xs)\ (rotater\ n\ ys)$
 $\langle proof \rangle$

lemma *rotate1-map2*:
 $length\ xs = length\ ys ==>$
 $rotate1\ (map2\ f\ xs\ ys) = map2\ f\ (rotate1\ xs)\ (rotate1\ ys)$
 $\langle proof \rangle$

lemmas *lth = box-equals* [*OF asm-rl length-rotate* [*symmetric*]
 $length-rotate$ [*symmetric*], *THEN rotate1-map2*]

lemma *rotate-map2*:
 $length\ xs = length\ ys ==>$
 $rotate\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotate\ n\ xs)\ (rotate\ n\ ys)$
 $\langle proof \rangle$

lemma *to-bl-rotl*:
 $to-bl\ (word-rotl\ n\ w) = rotate\ n\ (to-bl\ w)$
 $\langle proof \rangle$

lemmas *blrs0 = rotate-egs* [*THEN to-bl-rotl* [*THEN trans*]]

lemmas *word-rotl-egs =*
 $blrs0$ [*simplified word-bl.Rep'* *word-bl.Rep-inject to-bl-rotl* [*symmetric*]]

lemma *to-bl-rotr*:
 $to-bl\ (word-rotr\ n\ w) = rotater\ n\ (to-bl\ w)$
 $\langle proof \rangle$

lemmas *brrs0 = rotater-egs* [*THEN to-bl-rotr* [*THEN trans*]]

lemmas *word-rotr-egs =*
 $brrs0$ [*simplified word-bl.Rep'* *word-bl.Rep-inject to-bl-rotr* [*symmetric*]]

declare *word-rotr-egs* (1) [*simp*]

declare *word-rotl-egs* (1) [*simp*]

lemma
 $word-rot-rl$ [*simp*]:
 $word-rotl\ k\ (word-rotr\ k\ v) = v$ **and**
 $word-rot-lr$ [*simp*]:
 $word-rotr\ k\ (word-rotl\ k\ v) = v$
 $\langle proof \rangle$

lemma
 $word-rot-gal$:
 $(word-rotr\ n\ v = w) = (word-rotl\ n\ w = v)$ **and**
 $word-rot-gal'$:

$(w = \text{word-rottr } n \ v) = (v = \text{word-rotl } n \ w)$
 $\langle \text{proof} \rangle$

lemma *word-rottr-rev*:

$\text{word-rottr } n \ w = \text{word-reverse } (\text{word-rotl } n \ (\text{word-reverse } w))$
 $\langle \text{proof} \rangle$

lemma *word-roti-0* [simp]: $\text{word-roti } 0 \ w = w$
 $\langle \text{proof} \rangle$

lemmas *abl-cong* = *arg-cong* [where $f = \text{of-bl}$]

lemma *word-roti-add*:

$\text{word-roti } (m + n) \ w = \text{word-roti } m \ (\text{word-roti } n \ w)$
 $\langle \text{proof} \rangle$

lemma *word-roti-conv-mod'*: $\text{word-roti } n \ w = \text{word-roti } (n \bmod \text{int } (\text{size } w)) \ w$
 $\langle \text{proof} \rangle$

lemmas *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded *word-size*]

13.3.3 Word rotation commutes with bit-wise operations

locale *word-rotate*

context *word-rotate*

begin

lemmas *word-rot-defs'* = *to-bl-rotl to-bl-rottr*

lemmas *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

lemmas *ths-map2* [OF *lbl-lbl*] = *rotate-map2 rotater-map2*

lemmas *ths-map* [where $xs = \text{to-bl } v, \text{ standard}$] = *rotate-map rotater-map*

lemmas *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-map2 ths-map*

lemma *word-rot-logs*:

$\text{word-rotl } n \ (\text{NOT } v) = \text{NOT } \text{word-rotl } n \ v$
 $\text{word-rottr } n \ (\text{NOT } v) = \text{NOT } \text{word-rottr } n \ v$
 $\text{word-rotl } n \ (x \ \text{AND } y) = \text{word-rotl } n \ x \ \text{AND } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{AND } y) = \text{word-rottr } n \ x \ \text{AND } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{OR } y) = \text{word-rotl } n \ x \ \text{OR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{OR } y) = \text{word-rottr } n \ x \ \text{OR } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{XOR } y) = \text{word-rotl } n \ x \ \text{XOR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{XOR } y) = \text{word-rottr } n \ x \ \text{XOR } \text{word-rottr } n \ y$


```

    <proof>
  end

  lemmas word-rot-logs = word-rotate.word-rot-logs

  lemmas bl-word-rotl-dt = trans [OF to-bl-rotl rotate-drop-take,
    simplified word-bl.Rep', standard]

  lemmas bl-word-rotr-dt = trans [OF to-bl-rotr rotater-drop-take,
    simplified word-bl.Rep', standard]

  lemma bl-word-roti-dt':
    n = nat ((- i) mod int (size (w :: 'a :: len word))) ==>
      to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)
    <proof>

  lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

  lemmas word-rotl-dt = bl-word-rotl-dt
    [THEN word-bl.Rep-inverse' [symmetric], standard]
  lemmas word-rotr-dt = bl-word-rotr-dt
    [THEN word-bl.Rep-inverse' [symmetric], standard]
  lemmas word-roti-dt = bl-word-roti-dt
    [THEN word-bl.Rep-inverse' [symmetric], standard]

  lemma word-rotx-0 [simp] : word-rotr i 0 = 0 & word-rotl i 0 = 0
    <proof>

  lemma word-roti-0' [simp] : word-roti n 0 = 0
    <proof>

  lemmas word-rotr-dt-no-bin' [simp] =
    word-rotr-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

  lemmas word-rotl-dt-no-bin' [simp] =
    word-rotl-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

  declare word-roti-def [simp]

  end

```

14 Boolean-Algebra: Boolean Algebras

```

theory Boolean-Algebra
imports Main
begin

```

```

locale boolean =
  fixes conj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcap$  70)
  fixes disj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcup$  65)
  fixes compl :: 'a  $\Rightarrow$  'a ( $\sim$  - [81] 80)
  fixes zero :: 'a (0)
  fixes one  :: 'a (1)
  assumes conj-assoc:  $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$ 
  assumes disj-assoc:  $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$ 
  assumes conj-commute:  $x \sqcap y = y \sqcap x$ 
  assumes disj-commute:  $x \sqcup y = y \sqcup x$ 
  assumes conj-disj-distrib:  $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$ 
  assumes disj-conj-distrib:  $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$ 
  assumes conj-one-right [simp]:  $x \sqcap \mathbf{1} = x$ 
  assumes disj-zero-right [simp]:  $x \sqcup \mathbf{0} = x$ 
  assumes conj-cancel-right [simp]:  $x \sqcap \sim x = \mathbf{0}$ 
  assumes disj-cancel-right [simp]:  $x \sqcup \sim x = \mathbf{1}$ 

```

```

sublocale boolean < conj!: abel-semigroup conj <proof>

```

```

sublocale boolean < disj!: abel-semigroup disj <proof>

```

```

context boolean
begin

```

```

lemmas conj-left-commute = conj.left-commute

```

```

lemmas disj-left-commute = disj.left-commute

```

```

lemmas conj-ac = conj.assoc conj.commute conj.left-commute

```

```

lemmas disj-ac = disj.assoc disj.commute disj.left-commute

```

```

lemma dual: boolean disj conj compl one zero
  <proof>

```

14.1 Complement

```

lemma complement-unique:

```

```

  assumes 1:  $a \sqcap x = \mathbf{0}$ 

```

```

  assumes 2:  $a \sqcup x = \mathbf{1}$ 

```

```

  assumes 3:  $a \sqcap y = \mathbf{0}$ 

```

```

  assumes 4:  $a \sqcup y = \mathbf{1}$ 

```

```

  shows  $x = y$ 

```

```

  <proof>

```

```

lemma compl-unique:  $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \Longrightarrow \sim x = y$ 

```

```

  <proof>

```

```

lemma double-compl [simp]:  $\sim (\sim x) = x$ 

```

```

  <proof>

```

lemma *compl-eq-compl-iff* [simp]: $(\sim x = \sim y) = (x = y)$
 $\langle proof \rangle$

14.2 Conjunction

lemma *conj-absorb* [simp]: $x \sqcap x = x$
 $\langle proof \rangle$

lemma *conj-zero-right* [simp]: $x \sqcap \mathbf{0} = \mathbf{0}$
 $\langle proof \rangle$

lemma *compl-one* [simp]: $\sim \mathbf{1} = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-zero-left* [simp]: $\mathbf{0} \sqcap x = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-one-left* [simp]: $\mathbf{1} \sqcap x = x$
 $\langle proof \rangle$

lemma *conj-cancel-left* [simp]: $\sim x \sqcap x = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-left-absorb* [simp]: $x \sqcap (x \sqcap y) = x \sqcap y$
 $\langle proof \rangle$

lemma *conj-disj-distrib2*:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
 $\langle proof \rangle$

lemmas *conj-disj-distrib* =
conj-disj-distrib conj-disj-distrib2

14.3 Disjunction

lemma *disj-absorb* [simp]: $x \sqcup x = x$
 $\langle proof \rangle$

lemma *disj-one-right* [simp]: $x \sqcup \mathbf{1} = \mathbf{1}$
 $\langle proof \rangle$

lemma *compl-zero* [simp]: $\sim \mathbf{0} = \mathbf{1}$
 $\langle proof \rangle$

lemma *disj-zero-left* [simp]: $\mathbf{0} \sqcup x = x$
 $\langle proof \rangle$

lemma *disj-one-left* [simp]: $\mathbf{1} \sqcup x = \mathbf{1}$
 $\langle proof \rangle$

lemma *disj-cancel-left* [simp]: $\sim x \sqcup x = \mathbf{1}$
 $\langle \text{proof} \rangle$

lemma *disj-left-absorb* [simp]: $x \sqcup (x \sqcup y) = x \sqcup y$
 $\langle \text{proof} \rangle$

lemma *disj-conj-distrib2*:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
 $\langle \text{proof} \rangle$

lemmas *disj-conj-distrib* =
disj-conj-distrib disj-conj-distrib2

14.4 De Morgan’s Laws

lemma *de-Morgan-conj* [simp]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
 $\langle \text{proof} \rangle$

lemma *de-Morgan-disj* [simp]: $\sim (x \sqcup y) = \sim x \sqcap \sim y$
 $\langle \text{proof} \rangle$

end

14.5 Symmetric Difference

locale *boolean-xor* = *boolean* +
fixes *xor* :: 'a => 'a => 'a (**infixr** \oplus 65)
assumes *xor-def*: $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$

sublocale *boolean-xor* < *xor*!: *abel-semigroup xor* $\langle \text{proof} \rangle$

context *boolean-xor*
begin

lemmas *xor-assoc* = *xor.assoc*
lemmas *xor-commute* = *xor.commute*
lemmas *xor-left-commute* = *xor.left-commute*

lemmas *xor-ac* = *xor.assoc xor.commute xor.left-commute*

lemma *xor-def2*:
 $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$
 $\langle \text{proof} \rangle$

lemma *xor-zero-right* [simp]: $x \oplus \mathbf{0} = x$
 $\langle \text{proof} \rangle$

lemma *xor-zero-left* [simp]: $\mathbf{0} \oplus x = x$
 $\langle \text{proof} \rangle$

lemma *xor-one-right* [*simp*]: $x \oplus \mathbf{1} = \sim x$
 $\langle proof \rangle$

lemma *xor-one-left* [*simp*]: $\mathbf{1} \oplus x = \sim x$
 $\langle proof \rangle$

lemma *xor-self* [*simp*]: $x \oplus x = \mathbf{0}$
 $\langle proof \rangle$

lemma *xor-left-self* [*simp*]: $x \oplus (x \oplus y) = y$
 $\langle proof \rangle$

lemma *xor-compl-left* [*simp*]: $\sim x \oplus y = \sim (x \oplus y)$
 $\langle proof \rangle$

lemma *xor-compl-right* [*simp*]: $x \oplus \sim y = \sim (x \oplus y)$
 $\langle proof \rangle$

lemma *xor-cancel-right*: $x \oplus \sim x = \mathbf{1}$
 $\langle proof \rangle$

lemma *xor-cancel-left*: $\sim x \oplus x = \mathbf{1}$
 $\langle proof \rangle$

lemma *conj-xor-distrib*: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$
 $\langle proof \rangle$

lemma *conj-xor-distrib2*:
 $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
 $\langle proof \rangle$

lemmas *conj-xor-distrib* =
conj-xor-distrib conj-xor-distrib2

end

end

15 WordGenLib: Miscellaneous Library for Words

theory *WordGenLib*
imports *WordShift Boolean-Algebra*
begin

declare *of-nat-2p* [*simp*]

lemma *word-int-cases*:

$$\llbracket \bigwedge n. \llbracket (x :: 'a::len0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$$

$$\implies P$$

$$\langle \text{proof} \rangle$$

lemma *word-nat-cases* [*cases type: word*]:

$$\llbracket \bigwedge n. \llbracket (x :: 'a::len \text{ word}) = \text{of-nat } n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$$

$$\implies P$$

$$\langle \text{proof} \rangle$$

lemma *max-word-eq*:

$$(\text{max-word} :: 'a::len \text{ word}) = 2^{\text{len-of TYPE('a)}} - 1$$

$$\langle \text{proof} \rangle$$

lemma *max-word-max* [*simp,intro!*]:

$$n \leq \text{max-word}$$

$$\langle \text{proof} \rangle$$

lemma *word-of-int-2p-len*:

$$\text{word-of-int } (2^{\text{len-of TYPE('a)}}) = (0 :: 'a::len0 \text{ word})$$

$$\langle \text{proof} \rangle$$

lemma *word-pow-0*:

$$(2 :: 'a::len \text{ word})^{\text{len-of TYPE('a)}} = 0$$

$$\langle \text{proof} \rangle$$

lemma *max-word-wrap*: $x + 1 = 0 \implies x = \text{max-word}$

$$\langle \text{proof} \rangle$$

lemma *max-word-minus*:

$$\text{max-word} = (-1 :: 'a::len \text{ word})$$

$$\langle \text{proof} \rangle$$

lemma *max-word-bl* [*simp*]:

$$\text{to-bl } (\text{max-word} :: 'a::len \text{ word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ True}$$

$$\langle \text{proof} \rangle$$

lemma *max-test-bit* [*simp*]:

$$(\text{max-word} :: 'a::len \text{ word}) !! n = (n < \text{len-of TYPE('a)})$$

$$\langle \text{proof} \rangle$$

lemma *word-and-max* [*simp*]:

$$x \text{ AND } \text{max-word} = x$$

$$\langle \text{proof} \rangle$$

lemma *word-or-max* [*simp*]:

$$x \text{ OR } \text{max-word} = \text{max-word}$$

$$\langle \text{proof} \rangle$$

lemma *word-ao-dist2*:

$x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

lemma *word-oa-dist2*:

$x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::len0 \text{ word}))$
 $\langle \text{proof} \rangle$

lemma *word-and-not [simp]*:

$x \text{ AND NOT } x = (0::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

lemma *word-or-not [simp]*:

$x \text{ OR NOT } x = \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *word-boolean*:

$\text{boolean } (op \text{ AND}) (op \text{ OR}) \text{ bitNOT } 0 \text{ max-word}$
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*:

$\text{boolean } op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word}$
 $\langle \text{proof} \rangle$

lemma *word-xor-and-or*:

$x \text{ XOR } y = x \text{ AND NOT } y \text{ OR NOT } x \text{ AND } (y::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*:

$\text{boolean-xor } op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word } op \text{ XOR}$
 $\langle \text{proof} \rangle$

lemma *shiftr-0 [iff]*:

$(x::'a::len0 \text{ word}) >> 0 = x$
 $\langle \text{proof} \rangle$

lemma *shiftr-0 [simp]*:

$(x :: 'a :: len \text{ word}) << 0 = x$
 $\langle \text{proof} \rangle$

lemma *shiftr-1 [simp]*:

$(1::'a::len \text{ word}) << n = 2^n$
 $\langle \text{proof} \rangle$

lemma *uint-lt-0 [simp]*:

$\text{uint } x < 0 = \text{False}$
 $\langle \text{proof} \rangle$

lemma *shiftr1-1 [simp]*:

shiftr1 (*1::'a::len word*) = 0
 ⟨proof⟩

lemma *shiftr-1* [*simp*]:
 (*1::'a::len word*) >> *n* = (if *n* = 0 then 1 else 0)
 ⟨proof⟩

lemma *word-less-1* [*simp*]:
 ((*x::'a::len word*) < 1) = (*x* = 0)
 ⟨proof⟩

lemma *to-bl-mask*:
to-bl (*mask n :: 'a::len word*) =
replicate (*len-of TYPE('a) - n*) *False* @
replicate (*min (len-of TYPE('a)) n*) *True*
 ⟨proof⟩

lemma *map-replicate-True*:
n = *length xs* ==>
map ($\lambda(x,y). x \& y$) (*zip xs (replicate n True)*) = *xs*
 ⟨proof⟩

lemma *map-replicate-False*:
n = *length xs* ==> *map* ($\lambda(x,y). x \& y$)
 (*zip xs (replicate n False)*) = *replicate n False*
 ⟨proof⟩

lemma *bl-and-mask*:
fixes *w :: 'a::len word*
fixes *n*
defines *n' ≡ len-of TYPE('a) - n*
shows *to-bl (w AND mask n) = replicate n' False @ drop n' (to-bl w)*
 ⟨proof⟩

lemma *drop-rev-takefill*:
length xs ≤ *n* ==>
drop (*n - length xs*) (*rev (takefill False n (rev xs))*) = *xs*
 ⟨proof⟩

lemma *map-nth-0* [*simp*]:
map (*op* !! (*0::'a::len0 word*)) *xs* = *replicate (length xs) False*
 ⟨proof⟩

lemma *uint-plus-if-size*:
uint (*x + y*) =
 (if *uint x + uint y* < 2^{size} *x* then

uint x + uint y

else

uint x + uint y - 2^{size} x)

$\langle \text{proof} \rangle$

lemma *unat-plus-if-size*:
 $\text{unat } (x + (y :: 'a :: \text{len word})) =$
 $(\text{if } \text{unat } x + \text{unat } y < 2^{\text{size } x} \text{ then}$
 $\quad \text{unat } x + \text{unat } y$
 else
 $\quad \text{unat } x + \text{unat } y - 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemma *word-neq-0-conv* [simp]:
fixes $w :: 'a :: \text{len word}$
shows $(w \neq 0) = (0 < w)$
 $\langle \text{proof} \rangle$

lemma *max-lt*:
 $\text{unat } (\max a b \text{ div } c) = \text{unat } (\max a b) \text{ div } \text{unat } (c :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma *uint-sub-if-size*:
 $\text{uint } (x - y) =$
 $(\text{if } \text{uint } y \leq \text{uint } x \text{ then}$
 $\quad \text{uint } x - \text{uint } y$
 else
 $\quad \text{uint } x - \text{uint } y + 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemma *unat-sub-simple*:
 $x \leq y \implies \text{unat } (y - x) = \text{unat } y - \text{unat } x$
 $\langle \text{proof} \rangle$

lemmas *unat-sub = unat-sub-simple*

lemma *word-less-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 < x = (0 < x - 1)$
 $\langle \text{proof} \rangle$

lemma *word-le-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 \leq x = (0 \leq x - 1)$
 $\langle \text{proof} \rangle$

lemmas *word-less-sub1-numberof* [simp] =
 $\text{word-less-sub1 } [\text{of number-of } w, \text{ standard}]$
lemmas *word-le-sub1-numberof* [simp] =
 $\text{word-le-sub1 } [\text{of number-of } w, \text{ standard}]$

lemma *word-of-int-minus*:

$\text{word-of-int } (2^{\text{len-of TYPE}('a)} - i) = (\text{word-of-int } (-i)::'a::\text{len word})$
 $\langle \text{proof} \rangle$

lemmas $\text{word-of-int-inj} =$
 $\text{word-uint.Abs-inject } [\text{unfolded uints-num, simplified}]$

lemma word-le-less-eq :
 $(x :: 'z::\text{len word}) \leq y = (x = y \vee x < y)$
 $\langle \text{proof} \rangle$

lemma mod-plus-cong :
assumes $1: (b::\text{int}) = b'$
and $2: x \bmod b' = x' \bmod b'$
and $3: y \bmod b' = y' \bmod b'$
and $4: x' + y' = z'$
shows $(x + y) \bmod b = z' \bmod b'$
 $\langle \text{proof} \rangle$

lemma mod-minus-cong :
assumes $1: (b::\text{int}) = b'$
and $2: x \bmod b' = x' \bmod b'$
and $3: y \bmod b' = y' \bmod b'$
and $4: x' - y' = z'$
shows $(x - y) \bmod b = z' \bmod b'$
 $\langle \text{proof} \rangle$

lemma word-induct-less :
 $\llbracket P (0::'a::\text{len word}); \bigwedge n. \llbracket n < m; P n \rrbracket \implies P (1 + n) \rrbracket \implies P m$
 $\langle \text{proof} \rangle$

lemma word-induct :
 $\llbracket P (0::'a::\text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$
 $\langle \text{proof} \rangle$

lemma word-induct2 $[\text{induct type}]$:
 $\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n::'b::\text{len word})$
 $\langle \text{proof} \rangle$

definition $\text{word-rec} :: 'a \Rightarrow ('b::\text{len word} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b \text{ word} \Rightarrow 'a$ **where**
 $\text{word-rec forZero forSuc } n \equiv \text{nat-rec forZero (forSuc } \circ \text{ of-nat) (unat } n)$

lemma word-rec-0 : $\text{word-rec } z \text{ } s \text{ } 0 = z$
 $\langle \text{proof} \rangle$

lemma word-rec-Suc :
 $1 + n \neq (0::'a::\text{len word}) \implies \text{word-rec } z \text{ } s \text{ } (1 + n) = s \text{ } n \text{ } (\text{word-rec } z \text{ } s \text{ } n)$
 $\langle \text{proof} \rangle$

lemma word-rec-Pred :

$n \neq 0 \implies \text{word-rec } z \ s \ n = s \ (n - 1) \ (\text{word-rec } z \ s \ (n - 1))$
 $\langle \text{proof} \rangle$

lemma *word-rec-in*:

$f \ (\text{word-rec } z \ (\lambda \cdot. f) \ n) = \text{word-rec } (f \ z) \ (\lambda \cdot. f) \ n$
 $\langle \text{proof} \rangle$

lemma *word-rec-in2*:

$f \ n \ (\text{word-rec } z \ f \ n) = \text{word-rec } (f \ 0 \ z) \ (f \circ \text{op} + 1) \ n$
 $\langle \text{proof} \rangle$

lemma *word-rec-twice*:

$m \leq n \implies \text{word-rec } z \ f \ n = \text{word-rec } (\text{word-rec } z \ f \ (n - m)) \ (f \circ \text{op} + (n - m)) \ m$
 $\langle \text{proof} \rangle$

lemma *word-rec-id*: $\text{word-rec } z \ (\lambda \cdot. \text{id}) \ n = z$

$\langle \text{proof} \rangle$

lemma *word-rec-id-eq*: $\forall m < n. f \ m = \text{id} \implies \text{word-rec } z \ f \ n = z$

$\langle \text{proof} \rangle$

lemma *word-rec-max*:

$\forall m \geq n. m \neq -1 \longrightarrow f \ m = \text{id} \implies \text{word-rec } z \ f \ -1 = \text{word-rec } z \ f \ n$
 $\langle \text{proof} \rangle$

lemma *unatSuc*:

$1 + n \neq (0::'a::\text{len } \text{word}) \implies \text{unat } (1 + n) = \text{Suc } (\text{unat } n)$
 $\langle \text{proof} \rangle$

lemmas *word-no-1* [simp] = *word-1-no* [symmetric, unfolded BIT-simps]

lemmas *word-no-0* [simp] = *word-0-no* [symmetric]

declare *word-0-bl* [simp]

declare *bin-to-bl-def* [simp]

declare *to-bl-0* [simp]

declare *of-bl-True* [simp]

end

16 Word: Word Library interface

theory *Word*

imports *WordGenLib*

uses $\sim\sim / \text{src} / \text{HOL} / \text{Tools} / \text{SMT} / \text{smt-word.ML}$

begin

$\langle ML \rangle$

see *Examples/WordExamples.thy* for examples

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.