

# The Isabelle/HOL Algebra Library

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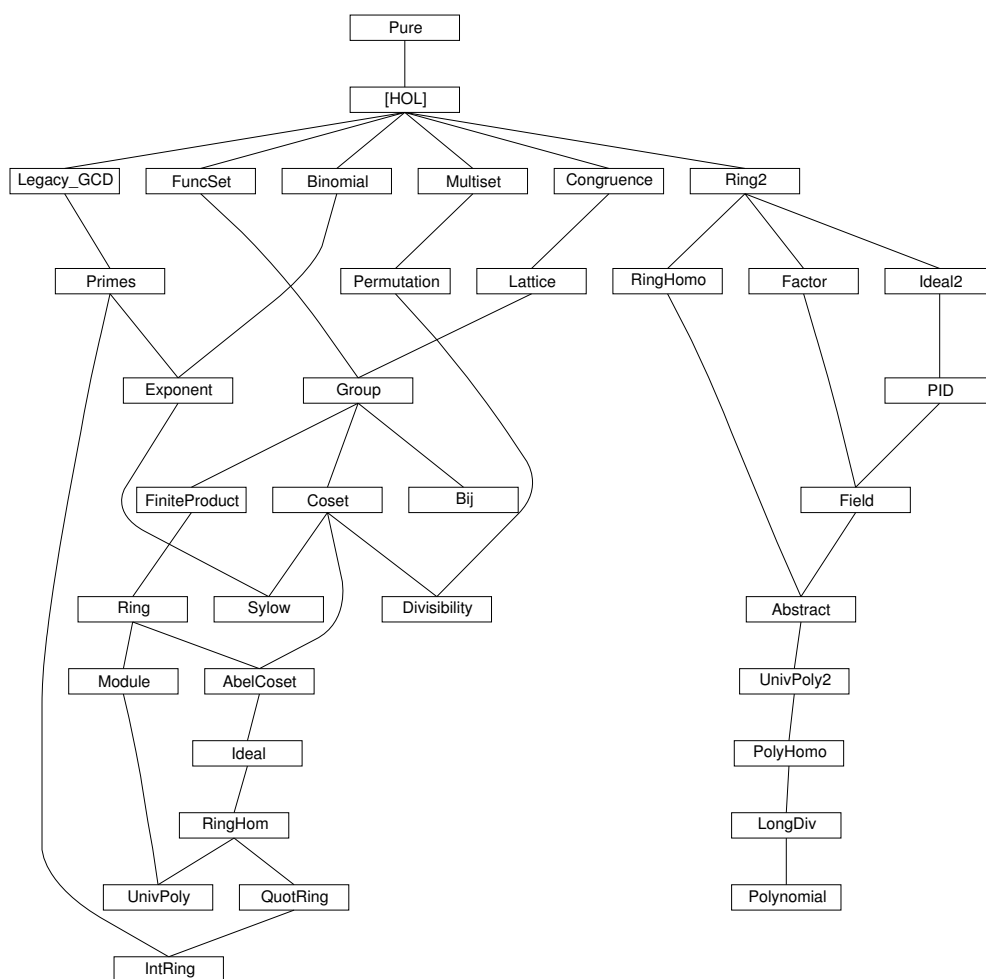
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```

theory Congruence
imports Main
begin

```

## 1 Objects

### 1.1 Structure with Carrier Set.

```

record 'a partial_object =
  carrier :: "'a set"

```

### 1.2 Structure with Carrier and Equivalence Relation eq

```

record 'a eq_object = "'a partial_object" +
  eq :: "'a  $\Rightarrow$  'a  $\Rightarrow$  bool" (infixl ".=ᵢ" 50)

```

#### definition

```

elem :: "_  $\Rightarrow$  'a  $\Rightarrow$  'a set  $\Rightarrow$  bool" (infixl ".∈ᵢ" 50)
where "x .∈S A  $\longleftrightarrow$  ( $\exists y \in A. x .=_S y$ )"

```

#### definition

```

set_eq :: "_  $\Rightarrow$  'a set  $\Rightarrow$  'a set  $\Rightarrow$  bool" (infixl "{.=}ᵢ" 50)
where "A {.=}S B  $\longleftrightarrow$  (( $\forall x \in A. x .∈_S B$ )  $\wedge$  ( $\forall x \in B. x .∈_S A$ ))"

```

#### definition

```

eq_class_of :: "_  $\Rightarrow$  'a  $\Rightarrow$  'a set" ("class'_ofᵢ _")
where "class_ofS x = {y  $\in$  carrier S. x .=_S y}"

```

#### definition

```

eq_closure_of :: "_  $\Rightarrow$  'a set  $\Rightarrow$  'a set" ("closure'_ofᵢ _")
where "closure_ofS A = {y  $\in$  carrier S. y .∈S A}"

```

#### definition

```

eq_is_closed :: "_  $\Rightarrow$  'a set  $\Rightarrow$  bool" ("is'_closedᵢ _")
where "is_closedS A  $\longleftrightarrow$  A  $\subseteq$  carrier S  $\wedge$  closure_ofS A = A"

```

#### abbreviation

```

not_eq :: "_  $\Rightarrow$  'a  $\Rightarrow$  'a  $\Rightarrow$  bool" (infixl ".≠ᵢ" 50)
where "x .≠S y ==  $\sim$ (x .=_S y)"

```

#### abbreviation

```

not_elem :: "_  $\Rightarrow$  'a  $\Rightarrow$  'a set  $\Rightarrow$  bool" (infixl ".∉ᵢ" 50)
where "x .∉S A ==  $\sim$ (x .∈S A)"

```

#### abbreviation

```

set_not_eq :: "_  $\Rightarrow$  'a set  $\Rightarrow$  'a set  $\Rightarrow$  bool" (infixl "{.≠}ᵢ" 50)
where "A {.=}S B ==  $\sim$ (A {.=}S B)"

```

```

locale equivalence =
  fixes S (structure)
    assumes refl [simp, intro]: "x ∈ carrier S  $\implies$  x .= x"
    and sym [sym]: "[ x .= y; x ∈ carrier S; y ∈ carrier S ]  $\implies$  y .= x"
    and trans [trans]: "[ x .= y; y .= z; x ∈ carrier S; y ∈ carrier S; z ∈ carrier S ]  $\implies$  x .= z"

```

```

lemma elemI:
  fixes R (structure)
  assumes "a' ∈ A" and "a .= a'"
  shows "a .∈ A"
  <proof>

```

```

lemma (in equivalence) elem_exact:
  assumes "a ∈ carrier S" and "a ∈ A"
  shows "a .∈ A"
  <proof>

```

```

lemma elemE:
  fixes S (structure)
  assumes "a .∈ A"
  and " $\bigwedge$ a'. [a' ∈ A; a .= a']  $\implies$  P"
  shows "P"
  <proof>

```

```

lemma (in equivalence) elem_cong_1 [trans]:
  assumes cong: "a' .= a"
  and a: "a .∈ A"
  and carr: "a ∈ carrier S" "a' ∈ carrier S"
  and Acarr: "A  $\subseteq$  carrier S"
  shows "a' .∈ A"
  <proof>

```

```

lemma (in equivalence) elem_subsetD:
  assumes "A  $\subseteq$  B"
  and aA: "a .∈ A"
  shows "a .∈ B"
  <proof>

```

```

lemma (in equivalence) mem_imp_elem [simp, intro]:
  "[| x ∈ A; x ∈ carrier S |]  $\implies$  x .∈ A"
  <proof>

```

```

lemma set_eqI:
  fixes R (structure)

```

```

    assumes ltr: " $\bigwedge a. a \in A \implies a \in B$ "
    and rtl: " $\bigwedge b. b \in B \implies b \in A$ "
    shows " $A \{.\} B$ "
  <proof>

```

```

lemma set_eqI2:
  fixes R (structure)
  assumes ltr: " $\bigwedge a b. a \in A \implies \exists b \in B. a \in b$ "
  and rtl: " $\bigwedge b. b \in B \implies \exists a \in A. b \in a$ "
  shows " $A \{.\} B$ "
  <proof>

```

```

lemma set_eqD1:
  fixes R (structure)
  assumes AA': " $A \{.\} A'$ "
  and "a  $\in A$ "
  shows " $\exists a' \in A'. a \in a'$ "
  <proof>

```

```

lemma set_eqD2:
  fixes R (structure)
  assumes AA': " $A \{.\} A'$ "
  and "a'  $\in A'$ "
  shows " $\exists a \in A. a' \in a$ "
  <proof>

```

```

lemma set_eqE:
  fixes R (structure)
  assumes AB: " $A \{.\} B$ "
  and r: " $\llbracket \forall a \in A. a \in B; \forall b \in B. b \in A \rrbracket \implies P$ "
  shows "P"
  <proof>

```

```

lemma set_eqE2:
  fixes R (structure)
  assumes AB: " $A \{.\} B$ "
  and r: " $\llbracket \forall a \in A. (\exists b \in B. a \in b); \forall b \in B. (\exists a \in A. b \in a) \rrbracket \implies P$ "
  shows "P"
  <proof>

```

```

lemma set_eqE':
  fixes R (structure)
  assumes AB: " $A \{.\} B$ "
  and aA: "a  $\in A$ " and bB: "b  $\in B$ "
  and r: " $\bigwedge a' b'. \llbracket a' \in A; b \in a'; b' \in B; a \in b' \rrbracket \implies P$ "
  shows "P"
  <proof>

```

```

lemma (in equivalence) eq_elem_cong_r [trans]:

```



```

    assumes a: "a .∈ A"
      and cong: "A {.=} A'"
      and carr: "a ∈ carrier S"
      and Carr: "A ⊆ carrier S" "A' ⊆ carrier S"
    shows "a .∈ A'"
  <proof>

```

```

lemma (in equivalence) set_eq_sym [sym]:
  assumes "A {.=} B"
    and "A ⊆ carrier S" "B ⊆ carrier S"
  shows "B {.=} A"
  <proof>

```

```

lemma (in equivalence) equal_set_eq_trans [trans]:
  assumes AB: "A = B" and BC: "B {.=} C"
  shows "A {.=} C"
  <proof>

```

```

lemma (in equivalence) set_eq_equal_trans [trans]:
  assumes AB: "A {.=} B" and BC: "B = C"
  shows "A {.=} C"
  <proof>

```

```

lemma (in equivalence) set_eq_trans [trans]:
  assumes AB: "A {.=} B" and BC: "B {.=} C"
    and carr: "A ⊆ carrier S" "B ⊆ carrier S" "C ⊆ carrier S"
  shows "A {.=} C"
  <proof>

```

```

lemma (in equivalence) set_eq_pairI:
  assumes xx': "x .= x'"
    and carr: "x ∈ carrier S" "x' ∈ carrier S" "y ∈ carrier S"
  shows "{x, y} {.=} {x', y}"
  <proof>

```

```

lemma (in equivalence) is_closedI:
  assumes closed: "!!x y. [| x .= y; x ∈ A; y ∈ carrier S |] ==> y ∈ A"
    and S: "A ⊆ carrier S"
  shows "is_closed A"
  <proof>

```

```

lemma (in equivalence) closure_of_eq:
  "[| x .= x'; A ⊆ carrier S; x ∈ closure_of A; x ∈ carrier S; x' ∈ carrier
S |] ==> x' ∈ closure_of A"
  <proof>

```

```

lemma (in equivalence) is_closed_eq [dest]:
  "[| x .= x'; x ∈ A; is_closed A; x ∈ carrier S; x' ∈ carrier S |] ==>
x' ∈ A"
  <proof>

```

```

lemma (in equivalence) is_closed_eq_rev [dest]:
  "[| x .= x'; x' ∈ A; is_closed A; x ∈ carrier S; x' ∈ carrier S |]
==> x ∈ A"
  <proof>

```

```

lemma closure_of_closed [simp, intro]:
  fixes S (structure)
  shows "closure_of A ⊆ carrier S"
  <proof>

```

```

lemma closure_of_memI:
  fixes S (structure)
  assumes "a .∈ A"
  and "a ∈ carrier S"
  shows "a ∈ closure_of A"
  <proof>

```

```

lemma closure_ofI2:
  fixes S (structure)
  assumes "a .= a'"
  and "a' ∈ A"
  and "a ∈ carrier S"
  shows "a ∈ closure_of A"
  <proof>

```

```

lemma closure_of_memE:
  fixes S (structure)
  assumes p: "a ∈ closure_of A"
  and r: "[| a ∈ carrier S; a .∈ A |] ==> P"
  shows "P"
  <proof>

```

```

lemma closure_ofE2:
  fixes S (structure)
  assumes p: "a ∈ closure_of A"
  and r: "⋀a'. [| a ∈ carrier S; a' ∈ A; a .= a' |] ==> P"
  shows "P"
  <proof>

```

end

```
theory Lattice
imports Congruence
begin
```

## 2 Orders and Lattices

### 2.1 Partial Orders

```
record 'a gorder = "'a eq_object" +
  le :: "[ 'a, 'a ] => bool" (infixl "⊆" 50)

locale weak_partial_order = equivalence L for L (structure) +
  assumes le_refl [intro, simp]:
    "x ∈ carrier L ==> x ⊆ x"
  and weak_le_antisym [intro]:
    "[| x ⊆ y; y ⊆ x; x ∈ carrier L; y ∈ carrier L |] ==> x .= y"
  and le_trans [trans]:
    "[| x ⊆ y; y ⊆ z; x ∈ carrier L; y ∈ carrier L; z ∈ carrier L
|] ==> x ⊆ z"
  and le_cong:
    "[| x .= y; z .= w; x ∈ carrier L; y ∈ carrier L; z ∈ carrier L;
w ∈ carrier L |] ==> x ⊆ z ⟷ y ⊆ w"
```

definition

```
lless :: "[_, 'a, 'a] => bool" (infixl "⊂" 50)
where "x ⊂L y ⟷ x ⊆L y & x ≠L y"
```

#### 2.1.1 The order relation

```
context weak_partial_order begin
```

```
lemma le_cong_l [intro, trans]:
  "[| x .= y; y ⊆ z; x ∈ carrier L; y ∈ carrier L; z ∈ carrier L |] ==>
x ⊆ z"
  <proof>
```

```
lemma le_cong_r [intro, trans]:
  "[| x ⊆ y; y .= z; x ∈ carrier L; y ∈ carrier L; z ∈ carrier L |] ==>
x ⊆ z"
  <proof>
```

```

lemma weak_refl [intro, simp]: "[ x .= y; x ∈ carrier L; y ∈ carrier
L ] ⇒ x ⊆ y"
  ⟨proof⟩

```

```

end

```

```

lemma weak_llessI:
  fixes R (structure)
  assumes "x ⊆ y" and "~(x .= y)"
  shows "x ⊂ y"
  ⟨proof⟩

```

```

lemma lless_imp_le:
  fixes R (structure)
  assumes "x ⊂ y"
  shows "x ⊆ y"
  ⟨proof⟩

```

```

lemma weak_lless_imp_not_eq:
  fixes R (structure)
  assumes "x ⊂ y"
  shows "¬ (x .= y)"
  ⟨proof⟩

```

```

lemma weak_llessE:
  fixes R (structure)
  assumes p: "x ⊂ y" and e: "[x ⊆ y; ¬ (x .= y)] ⇒ P"
  shows "P"
  ⟨proof⟩

```

```

lemma (in weak_partial_order) lless_cong_l [trans]:
  assumes xx': "x .= x'"
  and xy: "x' ⊂ y"
  and carr: "x ∈ carrier L" "x' ∈ carrier L" "y ∈ carrier L"
  shows "x ⊂ y"
  ⟨proof⟩

```

```

lemma (in weak_partial_order) lless_cong_r [trans]:
  assumes xy: "x ⊂ y"
  and yy': "y .= y'"
  and carr: "x ∈ carrier L" "y ∈ carrier L" "y' ∈ carrier L"
  shows "x ⊂ y'"
  ⟨proof⟩

```

```

lemma (in weak_partial_order) lless_antisym:
  assumes "a ∈ carrier L" "b ∈ carrier L"
  and "a ⊂ b" "b ⊂ a"
  shows "P"

```

*<proof>*

```
lemma (in weak_partial_order) lless_trans [trans]:
  assumes "a  $\sqsubset$  b" "b  $\sqsubset$  c"
    and carr[simp]: "a  $\in$  carrier L" "b  $\in$  carrier L" "c  $\in$  carrier L"
  shows "a  $\sqsubset$  c"
<proof>
```

### 2.1.2 Upper and lower bounds of a set

**definition**

```
Upper :: "[_, 'a set] => 'a set"
where "Upper L A = {u. (ALL x. x  $\in$  A  $\cap$  carrier L --> x  $\sqsubseteq_L$  u)}  $\cap$  carrier L"
```

**definition**

```
Lower :: "[_, 'a set] => 'a set"
where "Lower L A = {l. (ALL x. x  $\in$  A  $\cap$  carrier L --> l  $\sqsubseteq_L$  x)}  $\cap$  carrier L"
```

```
lemma Upper_closed [intro!, simp]:
  "Upper L A  $\subseteq$  carrier L"
<proof>
```

```
lemma Upper_memD [dest]:
  fixes L (structure)
  shows "[| u  $\in$  Upper L A; x  $\in$  A; A  $\subseteq$  carrier L |] ==> x  $\sqsubseteq$  u  $\wedge$  u  $\in$  carrier L"
<proof>
```

```
lemma (in weak_partial_order) Upper_elemD [dest]:
  "[| u  $\in$  Upper L A; u  $\in$  carrier L; x  $\in$  A; A  $\subseteq$  carrier L |] ==> x  $\sqsubseteq$  u"
<proof>
```

```
lemma Upper_memI:
  fixes L (structure)
  shows "[| !! y. y  $\in$  A ==> y  $\sqsubseteq$  x; x  $\in$  carrier L |] ==> x  $\in$  Upper L A"
<proof>
```

```
lemma (in weak_partial_order) Upper_elemI:
  "[| !! y. y  $\in$  A ==> y  $\sqsubseteq$  x; x  $\in$  carrier L |] ==> x  $\in$  Upper L A"
<proof>
```

```
lemma Upper_antimono:
  "A  $\subseteq$  B ==> Upper L B  $\subseteq$  Upper L A"
<proof>
```

```
lemma (in weak_partial_order) Upper_is_closed [simp]:
  "A  $\subseteq$  carrier L  $\implies$  is_closed (Upper L A)"
  <proof>
```

```
lemma (in weak_partial_order) Upper_mem_cong:
  assumes a'carr: "a'  $\in$  carrier L" and Acarr: "A  $\subseteq$  carrier L"
    and aa': "a .= a'"
    and aelem: "a  $\in$  Upper L A"
  shows "a'  $\in$  Upper L A"
  <proof>
```

```
lemma (in weak_partial_order) Upper_cong:
  assumes Acarr: "A  $\subseteq$  carrier L" and A'carr: "A'  $\subseteq$  carrier L"
    and AA': "A  $\{.\} A'"
  shows "Upper L A = Upper L A'"
  <proof>$ 
```

```
lemma Lower_closed [intro!, simp]:
  "Lower L A  $\subseteq$  carrier L"
  <proof>
```

```
lemma Lower_memD [dest]:
  fixes L (structure)
  shows "[| l  $\in$  Lower L A; x  $\in$  A; A  $\subseteq$  carrier L |]  $\implies$  l  $\sqsubseteq$  x  $\wedge$  l  $\in$ 
  carrier L"
  <proof>
```

```
lemma Lower_memI:
  fixes L (structure)
  shows "[| !! y. y  $\in$  A  $\implies$  x  $\sqsubseteq$  y; x  $\in$  carrier L |]  $\implies$  x  $\in$  Lower L
  A"
  <proof>
```

```
lemma Lower_antimono:
  "A  $\subseteq$  B  $\implies$  Lower L B  $\subseteq$  Lower L A"
  <proof>
```

```
lemma (in weak_partial_order) Lower_is_closed [simp]:
  "A  $\subseteq$  carrier L  $\implies$  is_closed (Lower L A)"
  <proof>
```

```
lemma (in weak_partial_order) Lower_mem_cong:
  assumes a'carr: "a'  $\in$  carrier L" and Acarr: "A  $\subseteq$  carrier L"
    and aa': "a .= a'"
    and aelem: "a  $\in$  Lower L A"
  shows "a'  $\in$  Lower L A"
  <proof>
```

```
lemma (in weak_partial_order) Lower_cong:
```

```

    assumes Acarr: "A  $\subseteq$  carrier L" and A'carr: "A'  $\subseteq$  carrier L"
    and AA': "A  $\{.\} A'$ "
    shows "Lower L A = Lower L A'"
  <proof>

```

### 2.1.3 Least and greatest, as predicate

definition

```

least :: "[_, 'a, 'a set] => bool"
where "least L l A  $\longleftrightarrow$  A  $\subseteq$  carrier L & l  $\in$  A & (ALL x : A. l  $\sqsubseteq_L$  x)"

```

definition

```

greatest :: "[_, 'a, 'a set] => bool"
where "greatest L g A  $\longleftrightarrow$  A  $\subseteq$  carrier L & g  $\in$  A & (ALL x : A. x  $\sqsubseteq_L$  g)"

```

Could weaken these to  $l \in \text{carrier L} \wedge l \in A$  and  $g \in \text{carrier L} \wedge g \in A$ .

```

lemma least_closed [intro, simp]:
  "least L l A  $\implies$  l  $\in$  carrier L"
  <proof>

```

```

lemma least_mem:
  "least L l A  $\implies$  l  $\in$  A"
  <proof>

```

```

lemma (in weak_partial_order) weak_least_unique:
  "[| least L x A; least L y A |]  $\implies$  x  $\{.\}$  y"
  <proof>

```

```

lemma least_le:
  fixes L (structure)
  shows "[| least L x A; a  $\in$  A |]  $\implies$  x  $\sqsubseteq$  a"
  <proof>

```

```

lemma (in weak_partial_order) least_cong:
  "[| x  $\{.\}$  x'; x  $\in$  carrier L; x'  $\in$  carrier L; is_closed A |]  $\implies$  least
  L x A = least L x' A"
  <proof>

```

least is not congruent in the second parameter for A  $\{.\}$  A'

```

lemma (in weak_partial_order) least_Upper_cong_l:
  assumes "x  $\{.\}$  x'"
  and "x  $\in$  carrier L" "x'  $\in$  carrier L"
  and "A  $\subseteq$  carrier L"
  shows "least L x (Upper L A) = least L x' (Upper L A)"
  <proof>

```

```

lemma (in weak_partial_order) least_Upper_cong_r:

```

```

    assumes Acarrs: "A  $\subseteq$  carrier L" "A'  $\subseteq$  carrier L"
    and AA': "A  $\{.\} A'$ "
    shows "least L x (Upper L A) = least L x (Upper L A')"
  <proof>

```

```

lemma least_UpperI:
  fixes L (structure)
  assumes above: "!! x. x  $\in$  A  $\implies$  x  $\sqsubseteq$  s"
    and below: "!! y. y  $\in$  Upper L A  $\implies$  s  $\sqsubseteq$  y"
    and L: "A  $\subseteq$  carrier L" "s  $\in$  carrier L"
  shows "least L s (Upper L A)"
  <proof>

```

```

lemma least_Upper_above:
  fixes L (structure)
  shows "[| least L s (Upper L A); x  $\in$  A; A  $\subseteq$  carrier L |]  $\implies$  x  $\sqsubseteq$  s"
  <proof>

```

```

lemma greatest_closed [intro, simp]:
  "greatest L l A  $\implies$  l  $\in$  carrier L"
  <proof>

```

```

lemma greatest_mem:
  "greatest L l A  $\implies$  l  $\in$  A"
  <proof>

```

```

lemma (in weak_partial_order) weak_greatest_unique:
  "[| greatest L x A; greatest L y A |]  $\implies$  x  $\{.\} y$ "
  <proof>

```

```

lemma greatest_le:
  fixes L (structure)
  shows "[| greatest L x A; a  $\in$  A |]  $\implies$  a  $\sqsubseteq$  x"
  <proof>

```

```

lemma (in weak_partial_order) greatest_cong:
  "[| x  $\{.\} x'$ ; x  $\in$  carrier L; x'  $\in$  carrier L; is_closed A |]  $\implies$ 
  greatest L x A = greatest L x' A"
  <proof>

```

greatest is not congruent in the second parameter for A  $\{.\} A'$

```

lemma (in weak_partial_order) greatest_Lower_cong_l:
  assumes "x  $\{.\} x'"
    and "x  $\in$  carrier L" "x'  $\in$  carrier L"
    and "A  $\subseteq$  carrier L"
  shows "greatest L x (Lower L A) = greatest L x' (Lower L A)"
  <proof>$ 
```

```

lemma (in weak_partial_order) greatest_Lower_cong_r:

```



```

    assumes Acarrs: "A  $\subseteq$  carrier L" "A'  $\subseteq$  carrier L"
    and AA': "A  $\{.\} A'$ "
    shows "greatest L x (Lower L A) = greatest L x (Lower L A')"
```

*<proof>*

```

lemma greatest_LowerI:
  fixes L (structure)
  assumes below: "!! x. x  $\in$  A  $\Rightarrow$  i  $\sqsubseteq$  x"
    and above: "!! y. y  $\in$  Lower L A  $\Rightarrow$  y  $\sqsubseteq$  i"
    and L: "A  $\subseteq$  carrier L" "i  $\in$  carrier L"
  shows "greatest L i (Lower L A)"
```

*<proof>*

```

lemma greatest_Lower_below:
  fixes L (structure)
  shows "[| greatest L i (Lower L A); x  $\in$  A; A  $\subseteq$  carrier L |]  $\Rightarrow$  i  $\sqsubseteq$ 
x"
```

*<proof>*

Supremum and infimum

**definition**

```

  sup :: "[_, 'a set] => 'a" ("⋒" [90] 90)
  where "⋒L A = (SOME x. least L x (Upper L A))"
```

**definition**

```

  inf :: "[_, 'a set] => 'a" ("⋓" [90] 90)
  where "⋓L A = (SOME x. greatest L x (Lower L A))"
```

**definition**

```

  join :: "[_, 'a, 'a] => 'a" (infixl "⋒" 65)
  where "x ⋒L y = ⋒L {x, y}"
```

**definition**

```

  meet :: "[_, 'a, 'a] => 'a" (infixl "⋓" 70)
  where "x ⋓L y = ⋓L {x, y}"
```

## 2.2 Lattices

```

locale weak_upper_semilattice = weak_partial_order +
  assumes sup_of_two_exists:
    "[| x  $\in$  carrier L; y  $\in$  carrier L |]  $\Rightarrow$  EX s. least L s (Upper L {x,
y})"
```

```

locale weak_lower_semilattice = weak_partial_order +
  assumes inf_of_two_exists:
    "[| x  $\in$  carrier L; y  $\in$  carrier L |]  $\Rightarrow$  EX s. greatest L s (Lower
L {x, y})"
```

```

locale weak_lattice = weak_upper_semilattice + weak_lower_semilattice
```

### 2.2.1 Supremum

```
lemma (in weak_upper_semilattice) joinI:
  "[| !!l. least L l (Upper L {x, y}) ==> P l; x ∈ carrier L; y ∈ carrier
L |]
  ==> P (x ⊔ y)"
⟨proof⟩
```

```
lemma (in weak_upper_semilattice) join_closed [simp]:
  "[| x ∈ carrier L; y ∈ carrier L |] ==> x ⊔ y ∈ carrier L"
⟨proof⟩
```

```
lemma (in weak_upper_semilattice) join_cong_l:
  assumes carr: "x ∈ carrier L" "x' ∈ carrier L" "y ∈ carrier L"
  and xx': "x .= x'"
  shows "x ⊔ y .= x' ⊔ y"
⟨proof⟩
```

```
lemma (in weak_upper_semilattice) join_cong_r:
  assumes carr: "x ∈ carrier L" "y ∈ carrier L" "y' ∈ carrier L"
  and yy': "y .= y'"
  shows "x ⊔ y .= x ⊔ y'"
⟨proof⟩
```

```
lemma (in weak_partial_order) sup_of_singletonI:
  "x ∈ carrier L ==> least L x (Upper L {x})"
⟨proof⟩
```

```
lemma (in weak_partial_order) weak_sup_of_singleton [simp]:
  "x ∈ carrier L ==> ⊔{x} .= x"
⟨proof⟩
```

```
lemma (in weak_partial_order) sup_of_singleton_closed [simp]:
  "x ∈ carrier L ==> ⊔{x} ∈ carrier L"
⟨proof⟩
```

Condition on A: supremum exists.

```
lemma (in weak_upper_semilattice) sup_insertI:
  "[| !!s. least L s (Upper L (insert x A)) ==> P s;
  least L a (Upper L A); x ∈ carrier L; A ⊆ carrier L |]
  ==> P (⊔(insert x A))"
⟨proof⟩
```

```
lemma (in weak_upper_semilattice) finite_sup_least:
  "[| finite A; A ⊆ carrier L; A ~= {} |] ==> least L (⊔A) (Upper L A)"
⟨proof⟩
```

```
lemma (in weak_upper_semilattice) finite_sup_insertI:
  assumes P: "!!l. least L l (Upper L (insert x A)) ==> P l"
  and xA: "finite A" "x ∈ carrier L" "A ⊆ carrier L"
```

```

    shows "P ( $\sqcup$  (insert x A))"
  <proof>

```

```

lemma (in weak_upper_semilattice) finite_sup_closed [simp]:
  "[| finite A; A  $\subseteq$  carrier L; A  $\sim$  {} |] ==>  $\sqcup$  A  $\in$  carrier L"
  <proof>

```

```

lemma (in weak_upper_semilattice) join_left:
  "[| x  $\in$  carrier L; y  $\in$  carrier L |] ==> x  $\sqsubseteq$  x  $\sqcup$  y"
  <proof>

```

```

lemma (in weak_upper_semilattice) join_right:
  "[| x  $\in$  carrier L; y  $\in$  carrier L |] ==> y  $\sqsubseteq$  x  $\sqcup$  y"
  <proof>

```

```

lemma (in weak_upper_semilattice) sup_of_two_least:
  "[| x  $\in$  carrier L; y  $\in$  carrier L |] ==> least L ( $\sqcup$  {x, y}) (Upper L
{x, y})"
  <proof>

```

```

lemma (in weak_upper_semilattice) join_le:
  assumes sub: "x  $\sqsubseteq$  z" "y  $\sqsubseteq$  z"
  and x: "x  $\in$  carrier L" and y: "y  $\in$  carrier L" and z: "z  $\in$  carrier
L"
  shows "x  $\sqcup$  y  $\sqsubseteq$  z"
  <proof>

```

```

lemma (in weak_upper_semilattice) weak_join_assoc_lemma:
  assumes L: "x  $\in$  carrier L" "y  $\in$  carrier L" "z  $\in$  carrier L"
  shows "x  $\sqcup$  (y  $\sqcup$  z)  $\sqsubseteq$   $\sqcup$  {x, y, z}"
  <proof>

```

Commutativity holds for  $\sqcup$ .

```

lemma join_comm:
  fixes L (structure)
  shows "x  $\sqcup$  y = y  $\sqcup$  x"
  <proof>

```

```

lemma (in weak_upper_semilattice) weak_join_assoc:
  assumes L: "x  $\in$  carrier L" "y  $\in$  carrier L" "z  $\in$  carrier L"
  shows "(x  $\sqcup$  y)  $\sqcup$  z  $\sqsubseteq$  x  $\sqcup$  (y  $\sqcup$  z)"
  <proof>

```

### 2.2.2 Infimum

```

lemma (in weak_lower_semilattice) meetI:
  "[| !!i. greatest L i (Lower L {x, y}) ==> P i;
x  $\in$  carrier L; y  $\in$  carrier L |]
==> P (x  $\sqcap$  y)"

```

*<proof>*

```
lemma (in weak_lower_semilattice) meet_closed [simp]:
  "[| x ∈ carrier L; y ∈ carrier L |] ==> x ⊓ y ∈ carrier L"
<proof>
```

```
lemma (in weak_lower_semilattice) meet_cong_l:
  assumes carr: "x ∈ carrier L" "x' ∈ carrier L" "y ∈ carrier L"
  and xx': "x .= x'"
  shows "x ⊓ y .= x' ⊓ y"
<proof>
```

```
lemma (in weak_lower_semilattice) meet_cong_r:
  assumes carr: "x ∈ carrier L" "y ∈ carrier L" "y' ∈ carrier L"
  and yy': "y .= y'"
  shows "x ⊓ y .= x ⊓ y'"
<proof>
```

```
lemma (in weak_partial_order) inf_of_singletonI:
  "x ∈ carrier L ==> greatest L x (Lower L {x})"
<proof>
```

```
lemma (in weak_partial_order) weak_inf_of_singleton [simp]:
  "x ∈ carrier L ==> ⊓ {x} .= x"
<proof>
```

```
lemma (in weak_partial_order) inf_of_singleton_closed:
  "x ∈ carrier L ==> ⊓ {x} ∈ carrier L"
<proof>
```

Condition on A: infimum exists.

```
lemma (in weak_lower_semilattice) inf_insertI:
  "[| !!i. greatest L i (Lower L (insert x A)) ==> P i;
    greatest L a (Lower L A); x ∈ carrier L; A ⊆ carrier L |]
  ==> P (⊓ (insert x A))"
<proof>
```

```
lemma (in weak_lower_semilattice) finite_inf_greatest:
  "[| finite A; A ⊆ carrier L; A ~= {} |] ==> greatest L (⊓ A) (Lower
L A)"
<proof>
```

```
lemma (in weak_lower_semilattice) finite_inf_insertI:
  assumes P: "!!i. greatest L i (Lower L (insert x A)) ==> P i"
  and xA: "finite A" "x ∈ carrier L" "A ⊆ carrier L"
  shows "P (⊓ (insert x A))"
<proof>
```

```
lemma (in weak_lower_semilattice) finite_inf_closed [simp]:
```

```
"[| finite A; A ⊆ carrier L; A ~= {} |] ==> ⋂ A ∈ carrier L"
⟨proof⟩
```

```
lemma (in weak_lower_semilattice) meet_left:
  "[| x ∈ carrier L; y ∈ carrier L |] ==> x ⊓ y ⊆ x"
  ⟨proof⟩
```

```
lemma (in weak_lower_semilattice) meet_right:
  "[| x ∈ carrier L; y ∈ carrier L |] ==> x ⊓ y ⊆ y"
  ⟨proof⟩
```

```
lemma (in weak_lower_semilattice) inf_of_two_greatest:
  "[| x ∈ carrier L; y ∈ carrier L |] ==>
   greatest L (⋂ {x, y}) (Lower L {x, y})"
  ⟨proof⟩
```

```
lemma (in weak_lower_semilattice) meet_le:
  assumes sub: "z ⊆ x" "z ⊆ y"
  and x: "x ∈ carrier L" and y: "y ∈ carrier L" and z: "z ∈ carrier
  L"
  shows "z ⊆ x ⊓ y"
  ⟨proof⟩
```

```
lemma (in weak_lower_semilattice) weak_meet_assoc_lemma:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "x ⊓ (y ⊓ z) .= ⋂ {x, y, z}"
  ⟨proof⟩
```

```
lemma meet_comm:
  fixes L (structure)
  shows "x ⊓ y = y ⊓ x"
  ⟨proof⟩
```

```
lemma (in weak_lower_semilattice) weak_meet_assoc:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "(x ⊓ y) ⊓ z .= x ⊓ (y ⊓ z)"
  ⟨proof⟩
```

## 2.3 Total Orders

```
local weak_total_order = weak_partial_order +
  assumes total: "[| x ∈ carrier L; y ∈ carrier L |] ==> x ⊆ y | y ⊆
  x"
```

Introduction rule: the usual definition of total order

```
lemma (in weak_partial_order) weak_total_orderI:
  assumes total: "!!x y. [| x ∈ carrier L; y ∈ carrier L |] ==> x ⊆
  y | y ⊆ x"
  shows "weak_total_order L"
```

*<proof>*

Total orders are lattices.

```
sublocale weak_total_order < weak: weak_lattice
<proof>
```

## 2.4 Complete Lattices

```
locale weak_complete_lattice = weak_lattice +
  assumes sup_exists:
    "[| A ⊆ carrier L |] ==> EX s. least L s (Upper L A)"
  and inf_exists:
    "[| A ⊆ carrier L |] ==> EX i. greatest L i (Lower L A)"
```

Introduction rule: the usual definition of complete lattice

```
lemma (in weak_partial_order) weak_complete_latticeI:
  assumes sup_exists:
    "!!A. [| A ⊆ carrier L |] ==> EX s. least L s (Upper L A)"
  and inf_exists:
    "!!A. [| A ⊆ carrier L |] ==> EX i. greatest L i (Lower L A)"
  shows "weak_complete_lattice L"
<proof>
```

**definition**

```
top :: "_ => 'a" ("⊤")
where "⊤L = sup L (carrier L)"
```

**definition**

```
bottom :: "_ => 'a" ("⊥")
where "⊥L = inf L (carrier L)"
```

```
lemma (in weak_complete_lattice) supI:
  "[| !!l. least L l (Upper L A) ==> P l; A ⊆ carrier L |]
  ==> P (⊔ A)"
<proof>
```

```
lemma (in weak_complete_lattice) sup_closed [simp]:
  "A ⊆ carrier L ==> ⊔ A ∈ carrier L"
<proof>
```

```
lemma (in weak_complete_lattice) top_closed [simp, intro]:
  "⊤ ∈ carrier L"
<proof>
```

```
lemma (in weak_complete_lattice) infI:
  "[| !!i. greatest L i (Lower L A) ==> P i; A ⊆ carrier L |]
  ==> P (⊓ A)"
<proof>
```

```

lemma (in weak_complete_lattice) inf_closed [simp]:
  "A  $\subseteq$  carrier L ==>  $\bigcap$  A  $\in$  carrier L"
  <proof>

lemma (in weak_complete_lattice) bottom_closed [simp, intro]:
  " $\perp \in$  carrier L"
  <proof>

Jacobson: Theorem 8.1

lemma Lower_empty [simp]:
  "Lower L {} = carrier L"
  <proof>

lemma Upper_empty [simp]:
  "Upper L {} = carrier L"
  <proof>

theorem (in weak_partial_order) weak_complete_lattice_criterion1:
  assumes top_exists: "EX g. greatest L g (carrier L)"
  and inf_exists:
    "!!A. [| A  $\subseteq$  carrier L; A  $\sim$  {} |] ==> EX i. greatest L i (Lower
L A)"
  shows "weak_complete_lattice L"
  <proof>

```

## 2.5 Orders and Lattices where eq is the Equality

```

locale partial_order = weak_partial_order +
  assumes eq_is_equal: "op .= = op ="
begin

declare weak_le_antisym [rule del]

lemma le_antisym [intro]:
  "[| x  $\sqsubseteq$  y; y  $\sqsubseteq$  x; x  $\in$  carrier L; y  $\in$  carrier L |] ==> x = y"
  <proof>

lemma lless_eq:
  "x  $\sqsubset$  y  $\longleftrightarrow$  x  $\sqsubseteq$  y & x  $\neq$  y"
  <proof>

lemma lless_asym:
  assumes "a  $\in$  carrier L" "b  $\in$  carrier L"
  and "a  $\sqsubset$  b" "b  $\sqsubset$  a"
  shows "P"
  <proof>

end

```

Least and greatest, as predicate

```
lemma (in partial_order) least_unique:
  "[| least L x A; least L y A |] ==> x = y"
  <proof>
```

```
lemma (in partial_order) greatest_unique:
  "[| greatest L x A; greatest L y A |] ==> x = y"
  <proof>
```

Lattices

```
locale upper_semilattice = partial_order +
  assumes sup_of_two_exists:
    "[| x ∈ carrier L; y ∈ carrier L |] ==> EX s. least L s (Upper L {x,
y})"
```

```
sublocale upper_semilattice < weak: weak_upper_semilattice
  <proof>
```

```
locale lower_semilattice = partial_order +
  assumes inf_of_two_exists:
    "[| x ∈ carrier L; y ∈ carrier L |] ==> EX s. greatest L s (Lower
L {x, y})"
```

```
sublocale lower_semilattice < weak: weak_lower_semilattice
  <proof>
```

```
locale lattice = upper_semilattice + lower_semilattice
```

Supremum

```
declare (in partial_order) weak_sup_of_singleton [simp del]
```

```
lemma (in partial_order) sup_of_singleton [simp]:
  "x ∈ carrier L ==>  $\sqcup$ {x} = x"
  <proof>
```

```
lemma (in upper_semilattice) join_assoc_lemma:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "x  $\sqcup$  (y  $\sqcup$  z) =  $\sqcup$ {x, y, z}"
  <proof>
```

```
lemma (in upper_semilattice) join_assoc:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "(x  $\sqcup$  y)  $\sqcup$  z = x  $\sqcup$  (y  $\sqcup$  z)"
  <proof>
```

Infimum

```
declare (in partial_order) weak_inf_of_singleton [simp del]
```



```

lemma (in partial_order) inf_of_singleton [simp]:
  "x ∈ carrier L ==>  $\bigcap \{x\} = x$ "
  <proof>

```

Condition on A: infimum exists.

```

lemma (in lower_semilattice) meet_assoc_lemma:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "x  $\sqcap$  (y  $\sqcap$  z) =  $\bigcap \{x, y, z\}$ "
  <proof>

```

```

lemma (in lower_semilattice) meet_assoc:
  assumes L: "x ∈ carrier L" "y ∈ carrier L" "z ∈ carrier L"
  shows "(x  $\sqcap$  y)  $\sqcap$  z = x  $\sqcap$  (y  $\sqcap$  z)"
  <proof>

```

Total Orders

```

locale total_order = partial_order +
  assumes total_order_total: "[| x ∈ carrier L; y ∈ carrier L |] ==>
x  $\sqsubseteq$  y | y  $\sqsubseteq$  x"

```

```

sublocale total_order < weak: weak_total_order
  <proof>

```

Introduction rule: the usual definition of total order

```

lemma (in partial_order) total_orderI:
  assumes total: "!!x y. [| x ∈ carrier L; y ∈ carrier L |] ==> x  $\sqsubseteq$ 
y | y  $\sqsubseteq$  x"
  shows "total_order L"
  <proof>

```

Total orders are lattices.

```

sublocale total_order < weak: lattice
  <proof>

```

Complete lattices

```

locale complete_lattice = lattice +
  assumes sup_exists:
    "[| A  $\sqsubseteq$  carrier L |] ==> EX s. least L s (Upper L A)"
  and inf_exists:
    "[| A  $\sqsubseteq$  carrier L |] ==> EX i. greatest L i (Lower L A)"

```

```

sublocale complete_lattice < weak: weak_complete_lattice
  <proof>

```

Introduction rule: the usual definition of complete lattice

```

lemma (in partial_order) complete_latticeI:
  assumes sup_exists:
    "!!A. [| A  $\sqsubseteq$  carrier L |] ==> EX s. least L s (Upper L A)"

```

```

    and inf_exists:
      "!!A. [| A ⊆ carrier L |] ==> EX i. greatest L i (Lower L A)"
    shows "complete_lattice L"
    <proof>

theorem (in partial_order) complete_lattice_criterion1:
  assumes top_exists: "EX g. greatest L g (carrier L)"
  and inf_exists:
    "!!A. [| A ⊆ carrier L; A ~= {} |] ==> EX i. greatest L i (Lower
L A)"
  shows "complete_lattice L"
  <proof>

```

## 2.6 Examples

### 2.6.1 The Powerset of a Set is a Complete Lattice

```

theorem powerset_is_complete_lattice:
  "complete_lattice (| carrier = Pow A, eq = op =, le = op ⊆ |)"
  (is "complete_lattice ?L")
  <proof>

```

Another example, that of the lattice of subgroups of a group, can be found in Group theory (Section 3.8).

end

```

theory Group
imports Lattice FuncSet
begin

```

## 3 Monoids and Groups

### 3.1 Definitions

Definitions follow [2].

```

record 'a monoid = "'a partial_object" +
  mult      :: "'a, 'a] ⇒ 'a" (infixl "⊗" 70)
  one       :: 'a ("1")

```

**definition**

```

m_inv :: "('a, 'b) monoid_scheme => 'a => 'a" ("inv" [81] 80)
where "invG x = (THE y. y ∈ carrier G & x ⊗G y = 1G & y ⊗G x = 1G)"

```

**definition**

```

Units :: "_ => 'a set"
— The set of invertible elements

```

```

    where "Units G = {y. y ∈ carrier G & (∃x ∈ carrier G. x ⊗G y = 1G
    & y ⊗G x = 1G)}"

consts
  pow :: "('a, 'm) monoid_scheme, 'a, 'b::number] => 'a" (infixr "'^(^')_z"
75)

overloading nat_pow == "pow :: [_, 'a, nat] => 'a"
begin
  definition "nat_pow G a n = nat_rec 1G (%u b. b ⊗G a) n"
end

overloading int_pow == "pow :: [_, 'a, int] => 'a"
begin
  definition "int_pow G a z =
    (let p = nat_rec 1G (%u b. b ⊗G a)
    in if neg z then invG (p (nat (-z))) else p (nat z))"
end

locale monoid =
  fixes G (structure)
  assumes m_closed [intro, simp]:
    "[x ∈ carrier G; y ∈ carrier G] ⇒ x ⊗ y ∈ carrier G"
  and m_assoc:
    "[x ∈ carrier G; y ∈ carrier G; z ∈ carrier G]
    ⇒ (x ⊗ y) ⊗ z = x ⊗ (y ⊗ z)"
  and one_closed [intro, simp]: "1 ∈ carrier G"
  and l_one [simp]: "x ∈ carrier G ⇒ 1 ⊗ x = x"
  and r_one [simp]: "x ∈ carrier G ⇒ x ⊗ 1 = x"

lemma monoidI:
  fixes G (structure)
  assumes m_closed:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y ∈ carrier
G"
  and one_closed: "1 ∈ carrier G"
  and m_assoc:
    "!!x y z. [| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>
(x ⊗ y) ⊗ z = x ⊗ (y ⊗ z)"
  and l_one: "!!x. x ∈ carrier G ==> 1 ⊗ x = x"
  and r_one: "!!x. x ∈ carrier G ==> x ⊗ 1 = x"
  shows "monoid G"
  <proof>

lemma (in monoid) Units_closed [dest]:
  "x ∈ Units G ==> x ∈ carrier G"
  <proof>

lemma (in monoid) inv_unique:

```

```

    assumes eq: "y  $\otimes$  x = 1"  "x  $\otimes$  y' = 1"
    and G: "x  $\in$  carrier G"  "y  $\in$  carrier G"  "y'  $\in$  carrier G"
    shows "y = y'"
  <proof>

```

```

lemma (in monoid) Units_m_closed [intro, simp]:
  assumes x: "x  $\in$  Units G" and y: "y  $\in$  Units G"
  shows "x  $\otimes$  y  $\in$  Units G"
  <proof>

```

```

lemma (in monoid) Units_one_closed [intro, simp]:
  "1  $\in$  Units G"
  <proof>

```

```

lemma (in monoid) Units_inv_closed [intro, simp]:
  "x  $\in$  Units G  $\implies$  inv x  $\in$  carrier G"
  <proof>

```

```

lemma (in monoid) Units_l_inv_ex:
  "x  $\in$  Units G  $\implies \exists y \in$  carrier G. y  $\otimes$  x = 1"
  <proof>

```

```

lemma (in monoid) Units_r_inv_ex:
  "x  $\in$  Units G  $\implies \exists y \in$  carrier G. x  $\otimes$  y = 1"
  <proof>

```

```

lemma (in monoid) Units_l_inv [simp]:
  "x  $\in$  Units G  $\implies$  inv x  $\otimes$  x = 1"
  <proof>

```

```

lemma (in monoid) Units_r_inv [simp]:
  "x  $\in$  Units G  $\implies$  x  $\otimes$  inv x = 1"
  <proof>

```

```

lemma (in monoid) Units_inv_Units [intro, simp]:
  "x  $\in$  Units G  $\implies$  inv x  $\in$  Units G"
  <proof>

```

```

lemma (in monoid) Units_l_cancel [simp]:
  "[| x  $\in$  Units G; y  $\in$  carrier G; z  $\in$  carrier G |]  $\implies$ 
  (x  $\otimes$  y = x  $\otimes$  z) = (y = z)"
  <proof>

```

```

lemma (in monoid) Units_inv_inv [simp]:
  "x  $\in$  Units G  $\implies$  inv (inv x) = x"
  <proof>

```

```

lemma (in monoid) inv_inj_on_Units:
  "inj_on (m_inv G) (Units G)"

```

*<proof>*

```
lemma (in monoid) Units_inv_comm:
  assumes inv: "x  $\otimes$  y = 1"
  and G: "x  $\in$  Units G" "y  $\in$  Units G"
  shows "y  $\otimes$  x = 1"
<proof>
```

Power

```
lemma (in monoid) nat_pow_closed [intro, simp]:
  "x  $\in$  carrier G ==> x ( $\wedge$ ) (n::nat)  $\in$  carrier G"
<proof>
```

```
lemma (in monoid) nat_pow_0 [simp]:
  "x ( $\wedge$ ) (0::nat) = 1"
<proof>
```

```
lemma (in monoid) nat_pow_Suc [simp]:
  "x ( $\wedge$ ) (Suc n) = x ( $\wedge$ ) n  $\otimes$  x"
<proof>
```

```
lemma (in monoid) nat_pow_one [simp]:
  "1 ( $\wedge$ ) (n::nat) = 1"
<proof>
```

```
lemma (in monoid) nat_pow_mult:
  "x  $\in$  carrier G ==> x ( $\wedge$ ) (n::nat)  $\otimes$  x ( $\wedge$ ) m = x ( $\wedge$ ) (n + m)"
<proof>
```

```
lemma (in monoid) nat_pow_pow:
  "x  $\in$  carrier G ==> (x ( $\wedge$ ) n) ( $\wedge$ ) m = x ( $\wedge$ ) (n * m::nat)"
<proof>
```

### 3.2 Groups

A group is a monoid all of whose elements are invertible.

```
locale group = monoid +
  assumes Units: "carrier G <= Units G"
```

```
lemma (in group) is_group: "group G" <proof>
```

```
theorem groupI:
  fixes G (structure)
  assumes m_closed [simp]:
    "!!x y. [| x  $\in$  carrier G; y  $\in$  carrier G |] ==> x  $\otimes$  y  $\in$  carrier G"
  and one_closed [simp]: "1  $\in$  carrier G"
  and m_assoc:
    "!!x y z. [| x  $\in$  carrier G; y  $\in$  carrier G; z  $\in$  carrier G |] ==>
```

```

      (x ⊗ y) ⊗ z = x ⊗ (y ⊗ z)"
    and l_one [simp]: "!!x. x ∈ carrier G ==> 1 ⊗ x = x"
    and l_inv_ex: "!!x. x ∈ carrier G ==> ∃y ∈ carrier G. y ⊗ x = 1"
  shows "group G"
<proof>

```

```

lemma (in monoid) group_l_invI:
  assumes l_inv_ex:
    "!!x. x ∈ carrier G ==> ∃y ∈ carrier G. y ⊗ x = 1"
  shows "group G"
<proof>

```

```

lemma (in group) Units_eq [simp]:
  "Units G = carrier G"
<proof>

```

```

lemma (in group) inv_closed [intro, simp]:
  "x ∈ carrier G ==> inv x ∈ carrier G"
<proof>

```

```

lemma (in group) l_inv_ex [simp]:
  "x ∈ carrier G ==> ∃y ∈ carrier G. y ⊗ x = 1"
<proof>

```

```

lemma (in group) r_inv_ex [simp]:
  "x ∈ carrier G ==> ∃y ∈ carrier G. x ⊗ y = 1"
<proof>

```

```

lemma (in group) l_inv [simp]:
  "x ∈ carrier G ==> inv x ⊗ x = 1"
<proof>

```

### 3.3 Cancellation Laws and Basic Properties

```

lemma (in group) l_cancel [simp]:
  "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>
   (x ⊗ y = x ⊗ z) = (y = z)"
<proof>

```

```

lemma (in group) r_inv [simp]:
  "x ∈ carrier G ==> x ⊗ inv x = 1"
<proof>

```

```

lemma (in group) r_cancel [simp]:
  "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>
   (y ⊗ x = z ⊗ x) = (y = z)"
<proof>

```

```

lemma (in group) inv_one [simp]:

```

```
"inv 1 = 1"
⟨proof⟩
```

```
lemma (in group) inv_inv [simp]:
  "x ∈ carrier G ==> inv (inv x) = x"
⟨proof⟩
```

```
lemma (in group) inv_inj:
  "inj_on (m_inv G) (carrier G)"
⟨proof⟩
```

```
lemma (in group) inv_mult_group:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> inv (x ⊗ y) = inv y ⊗ inv x"
⟨proof⟩
```

```
lemma (in group) inv_comm:
  "[| x ⊗ y = 1; x ∈ carrier G; y ∈ carrier G |] ==> y ⊗ x = 1"
⟨proof⟩
```

```
lemma (in group) inv_equality:
  "[| y ⊗ x = 1; x ∈ carrier G; y ∈ carrier G |] ==> inv x = y"
⟨proof⟩
```

Power

```
lemma (in group) int_pow_def2:
  "a (^) (z::int) = (if neg z then inv (a (^) (nat (-z))) else a (^) (nat z))"
⟨proof⟩
```

```
lemma (in group) int_pow_0 [simp]:
  "x (^) (0::int) = 1"
⟨proof⟩
```

```
lemma (in group) int_pow_one [simp]:
  "1 (^) (z::int) = 1"
⟨proof⟩
```

### 3.4 Subgroups

```
locale subgroup =
  fixes H and G (structure)
  assumes subset: "H ⊆ carrier G"
    and m_closed [intro, simp]: "[| x ∈ H; y ∈ H |] ==> x ⊗ y ∈ H"
    and one_closed [simp]: "1 ∈ H"
    and m_inv_closed [intro, simp]: "x ∈ H ==> inv x ∈ H"
```

```
lemma (in subgroup) is_subgroup:
  "subgroup H G" ⟨proof⟩
```

```
declare (in subgroup) group.intro [intro]
```

```
lemma (in subgroup) mem_carrier [simp]:
  "x ∈ H ⇒ x ∈ carrier G"
  ⟨proof⟩
```

```
lemma subgroup_imp_subset:
  "subgroup H G ⇒ H ⊆ carrier G"
  ⟨proof⟩
```

```
lemma (in subgroup) subgroup_is_group [intro]:
  assumes "group G"
  shows "group (G⟦carrier := H⟧)"
  ⟨proof⟩
```

Since  $H$  is nonempty, it contains some element  $x$ . Since it is closed under inverse, it contains  $\text{inv } x$ . Since it is closed under product, it contains  $x \otimes \text{inv } x = 1$ .

```
lemma (in group) one_in_subset:
  "[| H ⊆ carrier G; H ≠ {}; ∀ a ∈ H. inv a ∈ H; ∀ a ∈ H. ∀ b ∈ H. a ⊗ b
   ∈ H |]
   ==> 1 ∈ H"
  ⟨proof⟩
```

A characterization of subgroups: closed, non-empty subset.

```
lemma (in group) subgroupI:
  assumes subset: "H ⊆ carrier G" and non_empty: "H ≠ {}"
  and inv: "!!a. a ∈ H ⇒ inv a ∈ H"
  and mult: "!!a b. [a ∈ H; b ∈ H] ⇒ a ⊗ b ∈ H"
  shows "subgroup H G"
  ⟨proof⟩
```

```
declare monoid.one_closed [iff] group.inv_closed [simp]
  monoid.l_one [simp] monoid.r_one [simp] group.inv_inv [simp]
```

```
lemma subgroup_nonempty:
  "~ subgroup {} G"
  ⟨proof⟩
```

```
lemma (in subgroup) finite_imp_card_positive:
  "finite (carrier G) ==> 0 < card H"
  ⟨proof⟩
```

### 3.5 Direct Products

**definition**

```
DirProd :: "'a ⇒ _ ⇒ ('a × 'b) monoid" (infixr "××" 80) where
  "G ×× H =
    (carrier = carrier G × carrier H,
```



```

mult = (λ(g, h) (g', h'). (g ⊗G g', h ⊗H h')),
one = (1G, 1H)"

```

```

lemma DirProd_monoid:
  assumes "monoid G" and "monoid H"
  shows "monoid (G ×× H)"
<proof>

```

Does not use the previous result because it's easier just to use auto.

```

lemma DirProd_group:
  assumes "group G" and "group H"
  shows "group (G ×× H)"
<proof>

```

```

lemma carrier_DirProd [simp]:
  "carrier (G ×× H) = carrier G × carrier H"
<proof>

```

```

lemma one_DirProd [simp]:
  "1G ×× H = (1G, 1H)"
<proof>

```

```

lemma mult_DirProd [simp]:
  "(g, h) ⊗(G ×× H) (g', h') = (g ⊗G g', h ⊗H h')"
<proof>

```

```

lemma inv_DirProd [simp]:
  assumes "group G" and "group H"
  assumes g: "g ∈ carrier G"
  and h: "h ∈ carrier H"
  shows "m_inv (G ×× H) (g, h) = (invG g, invH h)"
<proof>

```

### 3.6 Homomorphisms and Isomorphisms

```

definition
  hom :: "_ => _ => ('a => 'b) set" where
    "hom G H =
      {h. h ∈ carrier G -> carrier H &
        (∀ x ∈ carrier G. ∀ y ∈ carrier G. h (x ⊗G y) = h x ⊗H h y)}"

```

```

lemma (in group) hom_compose:
  "[| h ∈ hom G H; i ∈ hom H I |] ==> compose (carrier G) i h ∈ hom G I"
<proof>

```

```

definition
  iso :: "_ => _ => ('a => 'b) set" (infixr "≅" 60)
  where "G ≅ H = {h. h ∈ hom G H & bij_betw h (carrier G) (carrier H)}"

```

```
lemma iso_refl: "(%x. x) ∈ G ≅ G"
⟨proof⟩
```

```
lemma (in group) iso_sym:
  "h ∈ G ≅ H ==> inv_into (carrier G) h ∈ H ≅ G"
⟨proof⟩
```

```
lemma (in group) iso_trans:
  "[| h ∈ G ≅ H; i ∈ H ≅ I |] ==> (compose (carrier G) i h) ∈ G ≅ I"
⟨proof⟩
```

```
lemma DirProd_commute_iso:
  shows "(λ(x,y). (y,x)) ∈ (G ×× H) ≅ (H ×× G)"
⟨proof⟩
```

```
lemma DirProd_assoc_iso:
  shows "(λ(x,y,z). (x,(y,z))) ∈ (G ×× H ×× I) ≅ (G ×× (H ×× I))"
⟨proof⟩
```

Basis for homomorphism proofs: we assume two groups  $G$  and  $H$ , with a homomorphism  $h$  between them

```
locale group_hom = G: group G + H: group H for G (structure) and H (structure) +
  fixes h
  assumes homh: "h ∈ hom G H"
```

```
lemma (in group_hom) hom_mult [simp]:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> h (x ⊗G y) = h x ⊗H h y"
⟨proof⟩
```

```
lemma (in group_hom) hom_closed [simp]:
  "x ∈ carrier G ==> h x ∈ carrier H"
⟨proof⟩
```

```
lemma (in group_hom) one_closed [simp]:
  "h 1 ∈ carrier H"
⟨proof⟩
```

```
lemma (in group_hom) hom_one [simp]:
  "h 1 = 1H"
⟨proof⟩
```

```
lemma (in group_hom) inv_closed [simp]:
  "x ∈ carrier G ==> h (inv x) ∈ carrier H"
⟨proof⟩
```

```
lemma (in group_hom) hom_inv [simp]:
  "x ∈ carrier G ==> h (inv x) = invH (h x)"
⟨proof⟩
```

### 3.7 Commutative Structures

Naming convention: multiplicative structures that are commutative are called *commutative*, additive structures are called *Abelian*.

```

locale comm_monoid = monoid +
  assumes m_comm: "[x ∈ carrier G; y ∈ carrier G] ==> x ⊗ y = y ⊗ x"

```

```

lemma (in comm_monoid) m_lcomm:
  "[x ∈ carrier G; y ∈ carrier G; z ∈ carrier G] ==>
   x ⊗ (y ⊗ z) = y ⊗ (x ⊗ z)"
  <proof>

```

```

lemmas (in comm_monoid) m_ac = m_assoc m_comm m_lcomm

```

```

lemma comm_monoidI:
  fixes G (structure)
  assumes m_closed:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y ∈ carrier
G"
  and one_closed: "1 ∈ carrier G"
  and m_assoc:
    "!!x y z. [| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>
    (x ⊗ y) ⊗ z = x ⊗ (y ⊗ z)"
  and l_one: "!!x. x ∈ carrier G ==> 1 ⊗ x = x"
  and m_comm:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y = y ⊗ x"
  shows "comm_monoid G"
  <proof>

```

```

lemma (in monoid) monoid_comm_monoidI:
  assumes m_comm:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y = y ⊗ x"
  shows "comm_monoid G"
  <proof>

```

```

lemma (in comm_monoid) nat_pow_distr:
  "[| x ∈ carrier G; y ∈ carrier G |] ==>
  (x ⊗ y) (^) (n::nat) = x (^) n ⊗ y (^) n"
  <proof>

```

```

locale comm_group = comm_monoid + group

```

```

lemma (in group) group_comm_groupI:
  assumes m_comm: "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==>
    x ⊗ y = y ⊗ x"
  shows "comm_group G"
  <proof>

```

```

lemma comm_groupI:
  fixes G (structure)
  assumes m_closed:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y ∈ carrier
G"
  and one_closed: "1 ∈ carrier G"
  and m_assoc:
    "!!x y z. [| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>
(x ⊗ y) ⊗ z = x ⊗ (y ⊗ z)"
  and m_comm:
    "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊗ y = y ⊗ x"
  and l_one: "!!x. x ∈ carrier G ==> 1 ⊗ x = x"
  and l_inv_ex: "!!x. x ∈ carrier G ==> ∃y ∈ carrier G. y ⊗ x = 1"
  shows "comm_group G"
  <proof>

lemma (in comm_group) inv_mult:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> inv (x ⊗ y) = inv x ⊗ inv y"
  <proof>

```

### 3.8 The Lattice of Subgroups of a Group

```

theorem (in group) subgroups_partial_order:
  "partial_order (| carrier = {H. subgroup H G}, eq = op =, le = op ⊆
|)"
  <proof>

lemma (in group) subgroup_self:
  "subgroup (carrier G) G"
  <proof>

lemma (in group) subgroup_imp_group:
  "subgroup H G ==> group (G(| carrier := H |))"
  <proof>

lemma (in group) is_monoid [intro, simp]:
  "monoid G"
  <proof>

lemma (in group) subgroup_inv_equality:
  "[| subgroup H G; x ∈ H |] ==> m_inv (G (| carrier := H |)) x = inv
x"
  <proof>

theorem (in group) subgroups_Inter:
  assumes subgr: "(!!H. H ∈ A ==> subgroup H G)"
  and not_empty: "A ~= {}"
  shows "subgroup (⋂A) G"

```

*<proof>*

```

theorem (in group) subgroups_complete_lattice:
  "complete_lattice (| carrier = {H. subgroup H G}, eq = op =, le = op
  ⊆ |)"
  (is "complete_lattice ?L")
<proof>

end

```

```

theory FiniteProduct
imports Group
begin

```

### 3.9 Product Operator for Commutative Monoids

#### 3.9.1 Inductive Definition of a Relation for Products over Sets

Instantiation of locale LC of theory `Finite_Set` is not possible, because here we have explicit typing rules like  $x \in \text{carrier } G$ . We introduce an explicit argument for the domain  $D$ .

```

inductive_set
  foldSetD :: "[ 'a set, 'b => 'a => 'a, 'a ] => ( 'b set * 'a ) set"
  for D :: "'a set" and f :: "'b => 'a => 'a" and e :: 'a
  where
    emptyI [intro]: "e ∈ D ==> ({}, e) ∈ foldSetD D f e"
    | insertI [intro]: "[| x ~: A; f x y ∈ D; (A, y) ∈ foldSetD D f e |]
    ==>
      (insert x A, f x y) ∈ foldSetD D f e"

inductive_cases empty_foldSetDE [elim!]: "({}, x) ∈ foldSetD D f e"

```

**definition**

```

foldD :: "[ 'a set, 'b => 'a => 'a, 'a, 'b set ] => 'a"
where "foldD D f e A = (THE x. (A, x) ∈ foldSetD D f e)"

```

**lemma** `foldSetD_closed`:

```

"[| (A, z) ∈ foldSetD D f e ; e ∈ D; !!x y. [| x ∈ A; y ∈ D |] ==>
f x y ∈ D
|] ==> z ∈ D"
<proof>

```

**lemma** `Diff1_foldSetD`:

```

"[| (A - {x}, y) ∈ foldSetD D f e; x ∈ A; f x y ∈ D |] ==>
(A, f x y) ∈ foldSetD D f e"
<proof>

```

```
lemma foldSetD_imp_finite [simp]: "(A, x) ∈ foldSetD D f e ==> finite A"
  <proof>
```

```
lemma finite_imp_foldSetD:
  "[| finite A; e ∈ D; !!x y. [| x ∈ A; y ∈ D |] ==> f x y ∈ D |] ==>
   EX x. (A, x) ∈ foldSetD D f e"
  <proof>
```

Left-Commutative Operations

```
locale LCD =
  fixes B :: "'b set"
  and D :: "'a set"
  and f :: "'b => 'a => 'a"      (infixl "." 70)
  assumes left_commute:
    "[| x ∈ B; y ∈ B; z ∈ D |] ==> x · (y · z) = y · (x · z)"
  and f_closed [simp, intro!]: "!!x y. [| x ∈ B; y ∈ D |] ==> f x y ∈ D"
  <proof>
```

```
lemma (in LCD) foldSetD_closed [dest]:
  "(A, z) ∈ foldSetD D f e ==> z ∈ D"
  <proof>
```

```
lemma (in LCD) Diff1_foldSetD:
  "[| (A - {x}, y) ∈ foldSetD D f e; x ∈ A; A ⊆ B |] ==>
   (A, f x y) ∈ foldSetD D f e"
  <proof>
```

```
lemma (in LCD) foldSetD_imp_finite [simp]:
  "(A, x) ∈ foldSetD D f e ==> finite A"
  <proof>
```

```
lemma (in LCD) finite_imp_foldSetD:
  "[| finite A; A ⊆ B; e ∈ D |] ==> EX x. (A, x) ∈ foldSetD D f e"
  <proof>
```

```
lemma (in LCD) foldSetD_determ_aux:
  "e ∈ D ==> ∀A x. A ⊆ B & card A < n --> (A, x) ∈ foldSetD D f e -->
   (∀y. (A, y) ∈ foldSetD D f e --> y = x)"
  <proof>
```

```
lemma (in LCD) foldSetD_determ:
  "[| (A, x) ∈ foldSetD D f e; (A, y) ∈ foldSetD D f e; e ∈ D; A ⊆ B |]
   ==> y = x"
  <proof>
```

```
lemma (in LCD) foldD_equality:
  "[| (A, y) ∈ foldSetD D f e; e ∈ D; A ⊆ B |] ==> foldD D f e A = y"
  <proof>
```

*<proof>*

```
lemma foldD_empty [simp]:
  "e ∈ D ==> foldD D f e {} = e"
<proof>
```

```
lemma (in LCD) foldD_insert_aux:
  "[| x ~: A; x ∈ B; e ∈ D; A ⊆ B |] ==>
    ((insert x A, v) ∈ foldSetD D f e) =
    (EX y. (A, y) ∈ foldSetD D f e & v = f x y)"
<proof>
```

```
lemma (in LCD) foldD_insert:
  "[| finite A; x ~: A; x ∈ B; e ∈ D; A ⊆ B |] ==>
    foldD D f e (insert x A) = f x (foldD D f e A)"
<proof>
```

```
lemma (in LCD) foldD_closed [simp]:
  "[| finite A; e ∈ D; A ⊆ B |] ==> foldD D f e A ∈ D"
<proof>
```

```
lemma (in LCD) foldD_commute:
  "[| finite A; x ∈ B; e ∈ D; A ⊆ B |] ==>
    f x (foldD D f e A) = foldD D f (f x e) A"
<proof>
```

```
lemma Int_mono2:
  "[| A ⊆ C; B ⊆ C |] ==> A Int B ⊆ C"
<proof>
```

```
lemma (in LCD) foldD_nest_Un_Int:
  "[| finite A; finite C; e ∈ D; A ⊆ B; C ⊆ B |] ==>
    foldD D f (foldD D f e C) A = foldD D f (foldD D f e (A Int C)) (A
    Un C)"
<proof>
```

```
lemma (in LCD) foldD_nest_Un_disjoint:
  "[| finite A; finite B; A Int B = {}; e ∈ D; A ⊆ B; C ⊆ B |]
    ==> foldD D f e (A Un B) = foldD D f (foldD D f e B) A"
<proof>
```

```
declare foldSetD_imp_finite [simp del]
empty_foldSetDE [rule del]
foldSetD.intros [rule del]
declare (in LCD)
foldSetD_closed [rule del]
```

Commutative Monoids

We enter a more restrictive context, with  $f :: 'a \Rightarrow 'a \Rightarrow 'a$  instead of  $'b$

```

=> 'a => 'a.

locale ACeD =
  fixes D :: "'a set"
  and f :: "'a => 'a => 'a"      (infixl "." 70)
  and e :: 'a
  assumes ident [simp]: "x ∈ D ==> x · e = x"
  and commute: "[| x ∈ D; y ∈ D |] ==> x · y = y · x"
  and assoc: "[| x ∈ D; y ∈ D; z ∈ D |] ==> (x · y) · z = x · (y · z)"
  and e_closed [simp]: "e ∈ D"
  and f_closed [simp]: "[| x ∈ D; y ∈ D |] ==> x · y ∈ D"

lemma (in ACeD) left_commute:
  "[| x ∈ D; y ∈ D; z ∈ D |] ==> x · (y · z) = y · (x · z)"
  <proof>

lemmas (in ACeD) AC = assoc commute left_commute

lemma (in ACeD) left_ident [simp]: "x ∈ D ==> e · x = x"
  <proof>

lemma (in ACeD) foldD_Un_Int:
  "[| finite A; finite B; A ⊆ D; B ⊆ D |] ==>
    foldD D f e A · foldD D f e B =
    foldD D f e (A Un B) · foldD D f e (A Int B)"
  <proof>

lemma (in ACeD) foldD_Un_disjoint:
  "[| finite A; finite B; A Int B = {}; A ⊆ D; B ⊆ D |] ==>
    foldD D f e (A Un B) = foldD D f e A · foldD D f e B"
  <proof>

```

### 3.9.2 Products over Finite Sets

#### definition

```

finprod :: "[( 'b, 'm) monoid_scheme, 'a => 'b, 'a set] => 'b"
where "finprod G f A =
  (if finite A
   then foldD (carrier G) (mult G o f) 1G A
   else undefined)"

```

#### syntax

```

"_finprod" :: "index => idt => 'a set => 'b => 'b"
  ("(3⊗ __:_. _)" [1000, 0, 51, 10] 10)

```

#### syntax (xsymbols)

```

"_finprod" :: "index => idt => 'a set => 'b => 'b"
  ("(3⊗ __∈_. _)" [1000, 0, 51, 10] 10)

```

#### syntax (HTML output)

```

"_finprod" :: "index => idt => 'a set => 'b => 'b"
  ("(3⊗ __∈_. _)" [1000, 0, 51, 10] 10)

```



```

translations
  " $\bigotimes_{i:A} b$ " == "CONST finprod  $\diamond z$  (%i. b) A"
  — Beware of argument permutation!

lemma (in comm_monoid) finprod_empty [simp]:
  "finprod G f {} = 1"
  <proof>

declare funcsetI [intro]
  funcset_mem [dest]

context comm_monoid begin

lemma finprod_insert [simp]:
  "[| finite F; a  $\notin$  F; f  $\in$  F  $\rightarrow$  carrier G; f a  $\in$  carrier G |] ==>
    finprod G f (insert a F) = f a  $\otimes$  finprod G f F"
  <proof>

lemma finprod_one [simp]:
  "finite A ==> ( $\bigotimes_{i:A} 1$ ) = 1"
  <proof>

lemma finprod_closed [simp]:
  fixes A
  assumes fin: "finite A" and f: "f  $\in$  A  $\rightarrow$  carrier G"
  shows "finprod G f A  $\in$  carrier G"
  <proof>

lemma funcset_Int_left [simp, intro]:
  "[| f  $\in$  A  $\rightarrow$  C; f  $\in$  B  $\rightarrow$  C |] ==> f  $\in$  A Int B  $\rightarrow$  C"
  <proof>

lemma funcset_Un_left [iff]:
  "(f  $\in$  A Un B  $\rightarrow$  C) = (f  $\in$  A  $\rightarrow$  C & f  $\in$  B  $\rightarrow$  C)"
  <proof>

lemma finprod_Un_Int:
  "[| finite A; finite B; g  $\in$  A  $\rightarrow$  carrier G; g  $\in$  B  $\rightarrow$  carrier G |] ==>
    finprod G g (A Un B)  $\otimes$  finprod G g (A Int B) =
    finprod G g A  $\otimes$  finprod G g B"
  — The reversed orientation looks more natural, but LOOPS as a simprule!
  <proof>

lemma finprod_Un_disjoint:
  "[| finite A; finite B; A Int B = {};
    g  $\in$  A  $\rightarrow$  carrier G; g  $\in$  B  $\rightarrow$  carrier G |]
  ==> finprod G g (A Un B) = finprod G g A  $\otimes$  finprod G g B"
  <proof>

```

```

lemma finprod_multf:
  "[| finite A; f ∈ A -> carrier G; g ∈ A -> carrier G |] ==>
    finprod G (%x. f x ⊗ g x) A = (finprod G f A ⊗ finprod G g A)"
  <proof>

lemma finprod_cong':
  "[| A = B; g ∈ B -> carrier G;
    !!i. i ∈ B ==> f i = g i |] ==> finprod G f A = finprod G g B"
  <proof>

lemma finprod_cong:
  "[| A = B; f ∈ B -> carrier G = True;
    !!i. i ∈ B ==> f i = g i |] ==> finprod G f A = finprod G g B"
  <proof>

Usually, if this rule causes a failed congruence proof error, the reason is that
the premise  $g \in B \rightarrow \text{carrier } G$  cannot be shown. Adding Pi_def to the
simpset is often useful. For this reason, comm_monoid.finprod_cong is not
added to the simpset by default.

end

declare funcsetI [rule del]
  funcset_mem [rule del]

context comm_monoid begin

lemma finprod_0 [simp]:
  "f ∈ {0::nat} -> carrier G ==> finprod G f {..0} = f 0"
  <proof>

lemma finprod_Suc [simp]:
  "f ∈ {..Suc n} -> carrier G ==>
    finprod G f {..Suc n} = (f (Suc n) ⊗ finprod G f {..n})"
  <proof>

lemma finprod_Suc2:
  "f ∈ {..Suc n} -> carrier G ==>
    finprod G f {..Suc n} = (finprod G (%i. f (Suc i)) {..n} ⊗ f 0)"
  <proof>

lemma finprod_mult [simp]:
  "[| f ∈ {..n} -> carrier G; g ∈ {..n} -> carrier G |] ==>
    finprod G (%i. f i ⊗ g i) {..n::nat} =
    finprod G f {..n} ⊗ finprod G g {..n}"
  <proof>

```

```

lemma finprod_reindex:
  assumes fin: "finite A"
  shows "f : (h ' A) → carrier G ==>
    inj_on h A ==> finprod G f (h ' A) = finprod G (%x. f (h x)) A"
  <proof>

lemma finprod_const:
  assumes fin [simp]: "finite A"
  and a [simp]: "a : carrier G"
  shows "finprod G (%x. a) A = a (^) card A"
  <proof>

lemma finprod_singleton:
  assumes i_in_A: "i ∈ A" and fin_A: "finite A" and f_Pi: "f ∈ A →
  carrier G"
  shows "(⊗j∈A. if i = j then f j else 1) = f i"
  <proof>

end

end

```

```

theory Coset
imports Group
begin

```

## 4 Cosets and Quotient Groups

```

definition
  r_coset    :: "[_, 'a set, 'a] ⇒ 'a set"      (infixl ">ₐ" 60)
  where "H #>ₐ a = (⋃ h∈H. {h ⊗ₐ a})"

definition
  l_coset    :: "[_, 'a, 'a set] ⇒ 'a set"      (infixl "<ₐ" 60)
  where "a <#ₐ H = (⋃ h∈H. {a ⊗ₐ h})"

definition
  RCOSSETS   :: "[_, 'a set] ⇒ ('a set)set"      ("rcosetsₐ _" [81] 80)
  where "rcosetsₐ H = (⋃ a∈carrier G. {H #>ₐ a})"

definition
  set_mult    :: "[_, 'a set, 'a set] ⇒ 'a set" (infixl "<#ₐ" 60)
  where "H <#ₐ K = (⋃ h∈H. ⋃ k∈K. {h ⊗ₐ k})"

definition

```

```
SET_INV :: "[_, 'a set] ⇒ 'a set"  ("set'_invz _" [81] 80)
where "set_invG H = (⋃ h∈H. {invG h})"
```

```
locale normal = subgroup + group +
  assumes coset_eq: "(∀ x ∈ carrier G. H #> x = x <# H)"
```

abbreviation

```
normal_rel :: "[ 'a set, ( 'a, 'b) monoid_scheme] ⇒ bool"  (infixl "<"
60) where
  "H < G ≡ normal H G"
```

## 4.1 Basic Properties of Cosets

```
lemma (in group) coset_mult_assoc:
  "[| M ⊆ carrier G; g ∈ carrier G; h ∈ carrier G |]
  ==> (M #> g) #> h = M #> (g ⊗ h)"
⟨proof⟩
```

```
lemma (in group) coset_mult_one [simp]: "M ⊆ carrier G ==> M #> 1 =
M"
⟨proof⟩
```

```
lemma (in group) coset_mult_inv1:
  "[| M #> (x ⊗ (inv y)) = M; x ∈ carrier G ; y ∈ carrier G;
  M ⊆ carrier G |] ==> M #> x = M #> y"
⟨proof⟩
```

```
lemma (in group) coset_mult_inv2:
  "[| M #> x = M #> y; x ∈ carrier G; y ∈ carrier G; M ⊆ carrier
G |]
  ==> M #> (x ⊗ (inv y)) = M "
⟨proof⟩
```

```
lemma (in group) coset_join1:
  "[| H #> x = H; x ∈ carrier G; subgroup H G |] ==> x ∈ H"
⟨proof⟩
```

```
lemma (in group) solve_equation:
  "[subgroup H G; x ∈ H; y ∈ H] ⇒ ∃ h∈H. y = h ⊗ x"
⟨proof⟩
```

```
lemma (in group) repr_independence:
  "[y ∈ H #> x; x ∈ carrier G; subgroup H G] ⇒ H #> x = H #> y"
⟨proof⟩
```

```
lemma (in group) coset_join2:
  "[x ∈ carrier G; subgroup H G; x∈H] ⇒ H #> x = H"
— Alternative proof is to put x = 1 in repr_independence.
```

*<proof>*

**lemma** (in monoid) r\_coset\_subset\_G:  
 "[| H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> H #> x  $\subseteq$  carrier G"  
*<proof>*

**lemma** (in group) rcosI:  
 "[| h  $\in$  H; H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> h  $\otimes$  x  $\in$  H #> x"  
*<proof>*

**lemma** (in group) rcosetsI:  
 "[| H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> H #> x  $\in$  rcosets H"  
*<proof>*

Really needed?

**lemma** (in group) transpose\_inv:  
 "[| x  $\otimes$  y = z; x  $\in$  carrier G; y  $\in$  carrier G; z  $\in$  carrier G |]  
 ==> (inv x)  $\otimes$  z = y"  
*<proof>*

**lemma** (in group) rcos\_self: "[| x  $\in$  carrier G; subgroup H G |] ==> x  
 $\in$  H #> x"  
*<proof>*

Opposite of "repr\_independence"

**lemma** (in group) repr\_independenceD:  
 assumes "subgroup H G"  
 assumes ycarr: "y  $\in$  carrier G"  
 and repr: "H #> x = H #> y"  
 shows "y  $\in$  H #> x"  
*<proof>*

Elements of a right coset are in the carrier

**lemma** (in subgroup) elemrcos\_carrier:  
 assumes "group G"  
 assumes acarr: "a  $\in$  carrier G"  
 and a': "a'  $\in$  H #> a"  
 shows "a'  $\in$  carrier G"  
*<proof>*

**lemma** (in subgroup) rcos\_const:  
 assumes "group G"  
 assumes hH: "h  $\in$  H"  
 shows "H #> h = H"  
*<proof>*

Step one for lemma rcos\_module

**lemma** (in subgroup) rcos\_module\_imp:

```

    assumes "group G"
    assumes xcarr: "x ∈ carrier G"
      and x'cos: "x' ∈ H #> x"
    shows "(x' ⊗ inv x) ∈ H"
  <proof>

```

Step two for lemma rcos\_module

```

lemma (in subgroup) rcos_module_rev:
  assumes "group G"
  assumes carr: "x ∈ carrier G" "x' ∈ carrier G"
    and xixH: "(x' ⊗ inv x) ∈ H"
  shows "x' ∈ H #> x"
  <proof>

```

Module property of right cosets

```

lemma (in subgroup) rcos_module:
  assumes "group G"
  assumes carr: "x ∈ carrier G" "x' ∈ carrier G"
  shows "(x' ∈ H #> x) = (x' ⊗ inv x ∈ H)"
  <proof>

```

Right cosets are subsets of the carrier.

```

lemma (in subgroup) rcosets_carrier:
  assumes "group G"
  assumes XH: "X ∈ rcosets H"
  shows "X ⊆ carrier G"
  <proof>

```

Multiplication of general subsets

```

lemma (in monoid) set_mult_closed:
  assumes Acarr: "A ⊆ carrier G"
    and Bcarr: "B ⊆ carrier G"
  shows "A <#> B ⊆ carrier G"
  <proof>

```

```

lemma (in comm_group) mult_subgroups:
  assumes subH: "subgroup H G"
    and subK: "subgroup K G"
  shows "subgroup (H <#> K) G"
  <proof>

```

```

lemma (in subgroup) lcos_module_rev:
  assumes "group G"
  assumes carr: "x ∈ carrier G" "x' ∈ carrier G"
    and xixH: "(inv x ⊗ x') ∈ H"
  shows "x' ∈ x <# H"
  <proof>

```

## 4.2 Normal subgroups

**lemma** normal\_imp\_subgroup: " $H \triangleleft G \implies \text{subgroup } H \ G$ "  
*<proof>*

**lemma** (in group) normalI:  
 " $\text{subgroup } H \ G \implies (\forall x \in \text{carrier } G. H \#> x = x <\# H) \implies H \triangleleft G$ "  
*<proof>*

**lemma** (in normal) inv\_op\_closed1:  
 " $\llbracket x \in \text{carrier } G; h \in H \rrbracket \implies (\text{inv } x) \otimes h \otimes x \in H$ "  
*<proof>*

**lemma** (in normal) inv\_op\_closed2:  
 " $\llbracket x \in \text{carrier } G; h \in H \rrbracket \implies x \otimes h \otimes (\text{inv } x) \in H$ "  
*<proof>*

Alternative characterization of normal subgroups

**lemma** (in group) normal\_inv\_iff:  
 " $(N \triangleleft G) =$   
 $(\text{subgroup } N \ G \ \& \ (\forall x \in \text{carrier } G. \forall h \in N. x \otimes h \otimes (\text{inv } x) \in N))$ "  
 (is " $\_ = ?\text{rhs}$ ")  
*<proof>*

## 4.3 More Properties of Cosets

**lemma** (in group) lcos\_m\_assoc:  
 " $\llbracket M \subseteq \text{carrier } G; g \in \text{carrier } G; h \in \text{carrier } G \rrbracket$   
 $\implies g <\# (h <\# M) = (g \otimes h) <\# M$ "  
*<proof>*

**lemma** (in group) lcos\_mult\_one: " $M \subseteq \text{carrier } G \implies 1 <\# M = M$ "  
*<proof>*

**lemma** (in group) l\_coset\_subset\_G:  
 " $\llbracket H \subseteq \text{carrier } G; x \in \text{carrier } G \rrbracket \implies x <\# H \subseteq \text{carrier } G$ "  
*<proof>*

**lemma** (in group) l\_coset\_swap:  
 " $\llbracket y \in x <\# H; x \in \text{carrier } G; \text{subgroup } H \ G \rrbracket \implies x \in y <\# H$ "  
*<proof>*

**lemma** (in group) l\_coset\_carrier:  
 " $\llbracket y \in x <\# H; x \in \text{carrier } G; \text{subgroup } H \ G \rrbracket \implies y \in \text{carrier } G$ "  
*<proof>*

**lemma** (in group) l\_repr\_imp\_subset:  
 assumes y: " $y \in x <\# H$ " and x: " $x \in \text{carrier } G$ " and sb: " $\text{subgroup } H \ G$ "  
 G"

shows "y <# H  $\subseteq$  x <# H"  
 <proof>

lemma (in group) l\_repr\_independence:  
 assumes y: "y  $\in$  x <# H" and x: "x  $\in$  carrier G" and sb: "subgroup H G"  
 shows "x <# H = y <# H"  
 <proof>

lemma (in group) setmult\_subset\_G:  
 "[H  $\subseteq$  carrier G; K  $\subseteq$  carrier G]  $\implies$  H <#> K  $\subseteq$  carrier G"  
 <proof>

lemma (in group) subgroup\_mult\_id: "subgroup H G  $\implies$  H <#> H = H"  
 <proof>

#### 4.3.1 Set of Inverses of an r\_coset.

lemma (in normal) rcos\_inv:  
 assumes x: "x  $\in$  carrier G"  
 shows "set\_inv (H #> x) = H #> (inv x)"  
 <proof>

#### 4.3.2 Theorems for <#> with #> or <#.

lemma (in group) setmult\_rcos\_assoc:  
 "[H  $\subseteq$  carrier G; K  $\subseteq$  carrier G; x  $\in$  carrier G]  
 $\implies$  H <#> (K #> x) = (H <#> K) #> x"  
 <proof>

lemma (in group) rcos\_assoc\_lcos:  
 "[H  $\subseteq$  carrier G; K  $\subseteq$  carrier G; x  $\in$  carrier G]  
 $\implies$  (H #> x) <#> K = H <#> (x <# K)"  
 <proof>

lemma (in normal) rcos\_mult\_step1:  
 "[x  $\in$  carrier G; y  $\in$  carrier G]  
 $\implies$  (H #> x) <#> (H #> y) = (H <#> (x <# H)) #> y"  
 <proof>

lemma (in normal) rcos\_mult\_step2:  
 "[x  $\in$  carrier G; y  $\in$  carrier G]  
 $\implies$  (H <#> (x <# H)) #> y = (H <#> (H #> x)) #> y"  
 <proof>

lemma (in normal) rcos\_mult\_step3:  
 "[x  $\in$  carrier G; y  $\in$  carrier G]  
 $\implies$  (H <#> (H #> x)) #> y = H #> (x  $\otimes$  y)"  
 <proof>



```

lemma (in normal) rcos_sum:
  "[x ∈ carrier G; y ∈ carrier G]
  ⇒ (H #> x) <#> (H #> y) = H #> (x ⊗ y)"
⟨proof⟩

```

```

lemma (in normal) rcosets_mult_eq: "M ∈ rcosets H ⇒ H <#> M = M"
  — generalizes subgroup_mult_id
⟨proof⟩

```

### 4.3.3 An Equivalence Relation

**definition**

```

r_congruent :: "[('a,'b)monoid_scheme, 'a set] ⇒ ('a*'a)set" ("rcong2
_")
where "rcong G H = {(x,y). x ∈ carrier G & y ∈ carrier G & invG x ⊗G
y ∈ H}"

```

```

lemma (in subgroup) equiv_rcong:
  assumes "group G"
  shows "equiv (carrier G) (rcong H)"
⟨proof⟩

```

Equivalence classes of `rcong` correspond to left cosets. Was there a mistake in the definitions? I'd have expected them to correspond to right cosets.

```

lemma (in subgroup) l_coset_eq_rcong:
  assumes "group G"
  assumes a: "a ∈ carrier G"
  shows "a <# H = rcong H `` {a}"
⟨proof⟩

```

### 4.3.4 Two Distinct Right Cosets are Disjoint

```

lemma (in group) rcos_equation:
  assumes "subgroup H G"
  assumes p: "ha ⊗ a = h ⊗ b" "a ∈ carrier G" "b ∈ carrier G" "h ∈ H"
  "ha ∈ H" "hb ∈ H"
  shows "hb ⊗ a ∈ (⋃ h∈H. {h ⊗ b})"
⟨proof⟩

```

```

lemma (in group) rcos_disjoint:
  assumes "subgroup H G"
  assumes p: "a ∈ rcosets H" "b ∈ rcosets H" "a ≠ b"
  shows "a ∩ b = {}"
⟨proof⟩

```

## 4.4 Further lemmas for `r_congruent`

The relation is a congruence

```

lemma (in normal) congruent_rcong:
  shows "congruent2 (rcong H) (rcong H) ( $\lambda a b. a \otimes b <\# H$ )"
<proof>

```

## 4.5 Order of a Group and Lagrange's Theorem

```

definition
  order :: "('a, 'b) monoid_scheme  $\Rightarrow$  nat"
  where "order S = card (carrier S)"

```

```

lemma (in group) rcosets_part_G:
  assumes "subgroup H G"
  shows " $\bigcup$  (rcosets H) = carrier G"
<proof>

```

```

lemma (in group) cosets_finite:
  "[ $c \in$  rcosets H;  $H \subseteq$  carrier G; finite (carrier G)]  $\Longrightarrow$  finite c"
<proof>

```

The next two lemmas support the proof of card\_cosets\_equal.

```

lemma (in group) inj_on_f:
  "[ $H \subseteq$  carrier G;  $a \in$  carrier G]  $\Longrightarrow$  inj_on ( $\lambda y. y \otimes \text{inv } a$ ) (H #> a)"
<proof>

```

```

lemma (in group) inj_on_g:
  "[ $H \subseteq$  carrier G;  $a \in$  carrier G]  $\Longrightarrow$  inj_on ( $\lambda y. y \otimes a$ ) H"
<proof>

```

```

lemma (in group) card_cosets_equal:
  "[ $c \in$  rcosets H;  $H \subseteq$  carrier G; finite(carrier G)]
 $\Longrightarrow$  card c = card H"
<proof>

```

```

lemma (in group) rcosets_subset_PowG:
  "subgroup H G  $\Longrightarrow$  rcosets H  $\subseteq$  Pow(carrier G)"
<proof>

```

```

theorem (in group) lagrange:
  "[finite(carrier G); subgroup H G]
 $\Longrightarrow$  card(rcosets H) * card(H) = order(G)"
<proof>

```

## 4.6 Quotient Groups: Factorization of a Group

```

definition
  FactGroup :: "[('a, 'b) monoid_scheme, 'a set]  $\Rightarrow$  ('a set) monoid" (infixl "Mod" 65)

```

— Actually defined for groups rather than monoids  
 where "FactGroup G H = (⟦carrier = rcosets<sub>G</sub> H, mult = set\_mult G, one = H⟧)"

**lemma** (in normal) setmult\_closed:  
 "⟦K1 ∈ rcosets H; K2 ∈ rcosets H⟧ ⇒ K1 <#> K2 ∈ rcosets H"  
 ⟨proof⟩

**lemma** (in normal) setinv\_closed:  
 "K ∈ rcosets H ⇒ set\_inv K ∈ rcosets H"  
 ⟨proof⟩

**lemma** (in normal) rcosets\_assoc:  
 "⟦M1 ∈ rcosets H; M2 ∈ rcosets H; M3 ∈ rcosets H⟧  
 ⇒ M1 <#> M2 <#> M3 = M1 <#> (M2 <#> M3)"  
 ⟨proof⟩

**lemma** (in subgroup) subgroup\_in\_rcosets:  
 assumes "group G"  
 shows "H ∈ rcosets H"  
 ⟨proof⟩

**lemma** (in normal) rcosets\_inv\_mult\_group\_eq:  
 "M ∈ rcosets H ⇒ set\_inv M <#> M = H"  
 ⟨proof⟩

**theorem** (in normal) factorgroup\_is\_group:  
 "group (G Mod H)"  
 ⟨proof⟩

**lemma** mult\_FactGroup [simp]: "X ⊗<sub>(G Mod H)</sub> X' = X <#><sub>G</sub> X'"  
 ⟨proof⟩

**lemma** (in normal) inv\_FactGroup:  
 "X ∈ carrier (G Mod H) ⇒ inv<sub>G Mod H</sub> X = set\_inv X"  
 ⟨proof⟩

The coset map is a homomorphism from G to the quotient group G Mod H

**lemma** (in normal) r\_coset\_hom\_Mod:  
 "(λa. H #> a) ∈ hom G (G Mod H)"  
 ⟨proof⟩

## 4.7 The First Isomorphism Theorem

The quotient by the kernel of a homomorphism is isomorphic to the range of that homomorphism.

**definition**

kernel :: "('a, 'm) monoid\_scheme ⇒ ('b, 'n) monoid\_scheme ⇒ ('a ⇒ 'b) ⇒ 'a set"

— the kernel of a homomorphism

**where** "kernel G H h = {x. x ∈ carrier G & h x = 1<sub>H</sub>}"

**lemma** (in group\_hom) subgroup\_kernel: "subgroup (kernel G H h) G"  
*<proof>*

The kernel of a homomorphism is a normal subgroup

**lemma** (in group\_hom) normal\_kernel: "(kernel G H h) ◁ G"  
*<proof>*

**lemma** (in group\_hom) FactGroup\_nonempty:  
 assumes X: "X ∈ carrier (G Mod kernel G H h)"  
 shows "X ≠ {}"  
*<proof>*

**lemma** (in group\_hom) FactGroup\_contents\_mem:  
 assumes X: "X ∈ carrier (G Mod (kernel G H h))"  
 shows "contents (h'X) ∈ carrier H"  
*<proof>*

**lemma** (in group\_hom) FactGroup\_hom:  
 "(λX. contents (h'X)) ∈ hom (G Mod (kernel G H h)) H"  
*<proof>*

Lemma for the following injectivity result

**lemma** (in group\_hom) FactGroup\_subset:  
 "[g ∈ carrier G; g' ∈ carrier G; h g = h g']  
 ⇒ kernel G H h #> g ⊆ kernel G H h #> g'"  
*<proof>*

**lemma** (in group\_hom) FactGroup\_inj\_on:  
 "inj\_on (λX. contents (h'X)) (carrier (G Mod kernel G H h))"  
*<proof>*

If the homomorphism h is onto H, then so is the homomorphism from the quotient group

**lemma** (in group\_hom) FactGroup\_onto:  
 assumes h: "h ' carrier G = carrier H"  
 shows "(λX. contents (h'X)) ' carrier (G Mod kernel G H h) = carrier H"  
*<proof>*

If h is a homomorphism from G onto H, then the quotient group G Mod kernel G H h is isomorphic to H.

**theorem** (in group\_hom) FactGroup\_iso:  
 "h ' carrier G = carrier H  
 ⇒ (λX. contents (h'X)) ∈ (G Mod (kernel G H h)) ≅ H"

*<proof>*

**end**

```
theory Exponent
imports Main "~~/src/HOL/Old_Number_Theory/Primes" Binomial
begin
```

## 5 Sylow's Theorem

### 5.1 The Combinatorial Argument Underlying the First Sylow Theorem

**definition**

```
exponent :: "nat => nat => nat"
where "exponent p s = (if prime p then (GREATEST r. p^r dvd s) else
0)"
```

Prime Theorems

**lemma** prime\_imp\_one\_less: "prime p ==> Suc 0 < p"

*<proof>*

**lemma** prime\_iff:

```
"(prime p) = (Suc 0 < p & (∀ a b. p dvd a*b --> (p dvd a) | (p dvd b)))"
```

*<proof>*

**lemma** zero\_less\_prime\_power: "prime p ==> 0 < p^a"

*<proof>*

**lemma** zero\_less\_card\_empty: "[| finite S; S ≠ {} |] ==> 0 < card(S)"

*<proof>*

**lemma** prime\_dvd\_cases:

```
"[| p*k dvd m*n; prime p |]
==> (∃ x. k dvd x*n & m = p*x) | (∃ y. k dvd m*y & n = p*y)"
```

*<proof>*

**lemma** prime\_power\_dvd\_cases [rule\_format (no\_asm)]: "prime p

```
==> ∀ m n. p^c dvd m*n -->
(∀ a b. a+b = Suc c --> p^a dvd m | p^b dvd n)"
```

*<proof>*

```

lemma div_combine:
  "[| prime p; ~ (p ^ (Suc r) dvd n); p^(a+r) dvd n*k |]
   ==> p ^ a dvd k"
<proof>

```

```

lemma Suc_le_power: "Suc 0 < p ==> Suc n <= p^n"
<proof>

```

```

lemma power_dvd_bound: "[| p^n dvd a; Suc 0 < p; a > 0 |] ==> n < a"
<proof>

```

### Exponent Theorems

```

lemma exponent_ge [rule_format]:
  "[| p^k dvd n; prime p; 0 < n |] ==> k <= exponent p n"
<proof>

```

```

lemma power_exponent_dvd: "s > 0 ==> (p ^ exponent p s) dvd s"
<proof>

```

```

lemma power_Suc_exponent_Not_dvd:
  "[| (p * p ^ exponent p s) dvd s; prime p |] ==> s = 0"
<proof>

```

```

lemma exponent_power_eq [simp]: "prime p ==> exponent p (p^a) = a"
<proof>

```

```

lemma exponent_equalityI:
  "!r::nat. (p^r dvd a) = (p^r dvd b) ==> exponent p a = exponent p b"
<proof>

```

```

lemma exponent_eq_0 [simp]: "~ prime p ==> exponent p s = 0"
<proof>

```

```

lemma exponent_mult_add1: "[| a > 0; b > 0 |]
  ==> (exponent p a) + (exponent p b) <= exponent p (a * b)"
<proof>

```

```

lemma exponent_mult_add2: "[| a > 0; b > 0 |]
  ==> exponent p (a * b) <= (exponent p a) + (exponent p b)"
<proof>

```

```

lemma exponent_mult_add: "[| a > 0; b > 0 |]
  ==> exponent p (a * b) = (exponent p a) + (exponent p b)"
<proof>

```

**lemma** not\_divides\_exponent\_0: "~ (p dvd n) ==> exponent p n = 0"  
*<proof>*

**lemma** exponent\_1\_eq\_0 [simp]: "exponent p (Suc 0) = 0"  
*<proof>*

Main Combinatorial Argument

**lemma** le\_extend\_mult: "[| c > 0; a <= b |] ==> a <= b \* (c::nat)"  
*<proof>*

**lemma** p\_fac\_forw\_lemma:  
 "[| (m::nat) > 0; k > 0; k < p^a; (p^r) dvd (p^a)\* m - k |] ==> r <= a"  
*<proof>*

**lemma** p\_fac\_forw: "[| (m::nat) > 0; k>0; k < p^a; (p^r) dvd (p^a)\* m - k |]  
 ==> (p^r) dvd (p^a) - k"  
*<proof>*

**lemma** r\_le\_a\_forw:  
 "[| (k::nat) > 0; k < p^a; p>0; (p^r) dvd (p^a) - k |] ==> r <= a"  
*<proof>*

**lemma** p\_fac\_backw: "[| m>0; k>0; (p::nat)≠0; k < p^a; (p^r) dvd p^a - k |]  
 ==> (p^r) dvd (p^a)\*m - k"  
*<proof>*

**lemma** exponent\_p\_a\_m\_k\_equation: "[| m>0; k>0; (p::nat)≠0; k < p^a |]  
 ==> exponent p (p^a \* m - k) = exponent p (p^a - k)"  
*<proof>*

Suc rules that we have to delete from the simpset

**lemmas** bad\_Sucs = binomial\_Suc\_Suc mult\_Suc mult\_Suc\_right

**lemma** p\_not\_div\_choose\_lemma [rule\_format]:  
 "[| ∀i. Suc i < K --> exponent p (Suc i) = exponent p (Suc(j+i)) |]  
 ==> k<K --> exponent p ((j+k) choose k) = 0"  
*<proof>*

**lemma** p\_not\_div\_choose:  
 "[| k<K; k<=n;

```

       $\forall j. 0 < j \ \& \ j < K \ \rightarrow \text{exponent } p \ (n - k + (K - j)) = \text{exponent } p \ (K - j) \mid]$ 
    ==> exponent p (n choose k) = 0"
  <proof>

```

```

lemma const_p_fac_right:
  "m>0 ==> exponent p ((p^a * m - Suc 0) choose (p^a - Suc 0)) = 0"
  <proof>

```

```

lemma const_p_fac:
  "m>0 ==> exponent p (((p^a) * m) choose p^a) = exponent p m"
  <proof>

```

end

```

theory Sylow
imports Coset Exponent
begin

```

See also [3].

The combinatorial argument is in theory Exponent

```

locale sylow = group +
  fixes p and a and m and calM and RelM
  assumes prime_p: "prime p"
    and order_G: "order(G) = (p^a) * m"
    and finite_G [iff]: "finite (carrier G)"
  defines "calM == {s. s  $\subseteq$  carrier(G) & card(s) = p^a}"
    and "RelM == {(N1,N2). N1  $\in$  calM & N2  $\in$  calM &
      ( $\exists g \in$  carrier(G). N1 = (N2 #> g) )}"

```

```

lemma (in sylow) RelM_refl_on: "refl_on calM RelM"
  <proof>

```

```

lemma (in sylow) RelM_sym: "sym RelM"
  <proof>

```

```

lemma (in sylow) RelM_trans: "trans RelM"
  <proof>

```

```

lemma (in sylow) RelM_equiv: "equiv calM RelM"
  <proof>

```

```

lemma (in sylow) M_subset_calM_prep: "M'  $\in$  calM // RelM ==> M'  $\subseteq$  calM"
  <proof>

```



## 5.2 Main Part of the Proof

```

locale sylow_central = sylow +
  fixes H and M1 and M
  assumes M_in_quot: "M ∈ calM // RelM"
    and not_dvd_M: "~(p ^ Suc(exponent p m) dvd card(M))"
    and M1_in_M: "M1 ∈ M"
  defines "H == {g. g ∈ carrier G & M1 #> g = M1}"

lemma (in sylow_central) M_subset_calM: "M ⊆ calM"
  <proof>

lemma (in sylow_central) card_M1: "card(M1) = p^a"
  <proof>

lemma card_nonempty: "0 < card(S) ==> S ≠ {}"
  <proof>

lemma (in sylow_central) exists_x_in_M1: "∃x. x ∈ M1"
  <proof>

lemma (in sylow_central) M1_subset_G [simp]: "M1 ⊆ carrier G"
  <proof>

lemma (in sylow_central) M1_inj_H: "∃f ∈ H → M1. inj_on f H"
  <proof>

```

## 5.3 Discharging the Assumptions of sylow\_central

```

lemma (in sylow) EmptyNotInEquivSet: "{} ∉ calM // RelM"
  <proof>

lemma (in sylow) existsM1inM: "M ∈ calM // RelM ==> ∃M1. M1 ∈ M"
  <proof>

lemma (in sylow) zero_less_o_G: "0 < order(G)"
  <proof>

lemma (in sylow) zero_less_m: "m > 0"
  <proof>

lemma (in sylow) card_calM: "card(calM) = (p^a) * m choose p^a"
  <proof>

lemma (in sylow) zero_less_card_calM: "card calM > 0"
  <proof>

lemma (in sylow) max_p_div_calM:
  "~ (p ^ Suc(exponent p m) dvd card(calM))"
  <proof>

```

**lemma** (in sylow) finite\_calM: "finite calM"

*<proof>*

**lemma** (in sylow) lemma\_A1:

" $\exists M \in \text{calM} // \text{RelM}. \sim (p \wedge \text{Suc}(\text{exponent } p \ m) \text{ dvd } \text{card}(M))$ "

*<proof>*

### 5.3.1 Introduction and Destruct Rules for H

**lemma** (in sylow\_central) H\_I: "[|g ∈ carrier G; M1 #> g = M1|] ==> g ∈ H"

*<proof>*

**lemma** (in sylow\_central) H\_into\_carrier\_G: "x ∈ H ==> x ∈ carrier G"

*<proof>*

**lemma** (in sylow\_central) in\_H\_imp\_eq: "g : H ==> M1 #> g = M1"

*<proof>*

**lemma** (in sylow\_central) H\_m\_closed: "[| x∈H; y∈H|] ==> x ⊗ y ∈ H"

*<proof>*

**lemma** (in sylow\_central) H\_not\_empty: "H ≠ {}"

*<proof>*

**lemma** (in sylow\_central) H\_is\_subgroup: "subgroup H G"

*<proof>*

**lemma** (in sylow\_central) rcosetGM1g\_subset\_G:

"[| g ∈ carrier G; x ∈ M1 #> g |] ==> x ∈ carrier G"

*<proof>*

**lemma** (in sylow\_central) finite\_M1: "finite M1"

*<proof>*

**lemma** (in sylow\_central) finite\_rcosetGM1g: "g∈carrier G ==> finite (M1 #> g)"

*<proof>*

**lemma** (in sylow\_central) M1\_cardeq\_rcosetGM1g:

"g ∈ carrier G ==> card(M1 #> g) = card(M1)"

*<proof>*

**lemma** (in sylow\_central) M1\_RelM\_rcosetGM1g:

"g ∈ carrier G ==> (M1, M1 #> g) ∈ RelM"

*<proof>*

## 5.4 Equal Cardinalities of $M$ and the Set of Cosets

Injectons between  $M$  and  $\text{rcosets}_G H$  show that their cardinalities are equal.

**lemma** ElemClassEquiv:

"[| equiv A r; C ∈ A // r |] ==> ∀x ∈ C. ∀y ∈ C. (x,y)∈r"  
*<proof>*

**lemma** (in sylow\_central) M\_elem\_map:

"M2 ∈ M ==> ∃g. g ∈ carrier G & M1 #> g = M2"  
*<proof>*

**lemmas** (in sylow\_central) M\_elem\_map\_carrier =

M\_elem\_map [THEN someI\_ex, THEN conjunct1]

**lemmas** (in sylow\_central) M\_elem\_map\_eq =

M\_elem\_map [THEN someI\_ex, THEN conjunct2]

**lemma** (in sylow\_central) M\_funcset\_rcosets\_H:

"(%x:M. H #> (SOME g. g ∈ carrier G & M1 #> g = x)) ∈ M → rcosets H"  
*<proof>*

**lemma** (in sylow\_central) inj\_M\_GmodH: "∃f ∈ M→rcosets H. inj\_on f M"

*<proof>*

### 5.4.1 The Opposite Injection

**lemma** (in sylow\_central) H\_elem\_map:

"H1 ∈ rcosets H ==> ∃g. g ∈ carrier G & H #> g = H1"  
*<proof>*

**lemmas** (in sylow\_central) H\_elem\_map\_carrier =

H\_elem\_map [THEN someI\_ex, THEN conjunct1]

**lemmas** (in sylow\_central) H\_elem\_map\_eq =

H\_elem\_map [THEN someI\_ex, THEN conjunct2]

**lemma** EquivElemClass:

"[|equiv A r; M ∈ A//r; M1∈M; (M1,M2) ∈ r |] ==> M2 ∈ M"  
*<proof>*

**lemma** (in sylow\_central) rcosets\_H\_funcset\_M:

"(λC ∈ rcosets H. M1 #> (@g. g ∈ carrier G ∧ H #> g = C)) ∈ rcosets H → M"  
*<proof>*

close to a duplicate of inj\_M\_GmodH

```
lemma (in sylow_central) inj_GmodH_M:
  "∃ g ∈ rcosets H → M. inj_on g (rcosets H)"
<proof>
```

```
lemma (in sylow_central) calM_subset_PowG: "calM ⊆ Pow(carrier G)"
<proof>
```

```
lemma (in sylow_central) finite_M: "finite M"
<proof>
```

```
lemma (in sylow_central) cardMeqIndexH: "card(M) = card(rcosets H)"
<proof>
```

```
lemma (in sylow_central) index_lem: "card(M) * card(H) = order(G)"
<proof>
```

```
lemma (in sylow_central) lemma_leq1: "p^a ≤ card(H)"
<proof>
```

```
lemma (in sylow_central) lemma_leq2: "card(H) ≤ p^a"
<proof>
```

```
lemma (in sylow_central) card_H_eq: "card(H) = p^a"
<proof>
```

```
lemma (in sylow) sylow_thm: "∃ H. subgroup H G & card(H) = p^a"
<proof>
```

Needed because the locale's automatic definition refers to `semigroup G` and `group_axioms G` rather than simply to `group G`.

```
lemma sylow_eq: "syLOW G p a m = (group G & sylow_axioms G p a m)"
<proof>
```

## 5.5 Sylow's Theorem

```
theorem sylow_thm:
  "[| prime p; group(G); order(G) = (p^a) * m; finite (carrier G) |]
  ==> ∃ H. subgroup H G & card(H) = p^a"
<proof>
```

```
end
```

```
theory Bij
imports Group
begin
```

## 6 Bijections of a Set, Permutation and Automorphism Groups

### definition

```
Bij :: "'a set ⇒ ('a ⇒ 'a) set"
  — Only extensional functions, since otherwise we get too many.
  where "Bij S = extensional S ∩ {f. bij_betw f S S}"
```

### definition

```
BijGroup :: "'a set ⇒ ('a ⇒ 'a) monoid"
  where "BijGroup S =
    (|carrier = Bij S,
      mult = λg ∈ Bij S. λf ∈ Bij S. compose S g f,
      one = λx ∈ S. x|)"
```

```
declare Id_compose [simp] compose_Id [simp]
```

```
lemma Bij_imp_extensional: "f ∈ Bij S ⇒ f ∈ extensional S"
  <proof>
```

```
lemma Bij_imp_funcset: "f ∈ Bij S ⇒ f ∈ S → S"
  <proof>
```

### 6.1 Bijections Form a Group

```
lemma restrict_inv_into_Bij: "f ∈ Bij S ⇒ (λx ∈ S. (inv_into S f)
x) ∈ Bij S"
  <proof>
```

```
lemma id_Bij: "(λx ∈ S. x) ∈ Bij S"
  <proof>
```

```
lemma compose_Bij: "[x ∈ Bij S; y ∈ Bij S] ⇒ compose S x y ∈ Bij S"
  <proof>
```

```
lemma Bij_compose_restrict_eq:
  "f ∈ Bij S ⇒ compose S (restrict (inv_into S f) S) f = (λx ∈ S.
x)"
  <proof>
```

```
theorem group_BijGroup: "group (BijGroup S)"
  <proof>
```

### 6.2 Automorphisms Form a Group

```
lemma Bij_inv_into_mem: "[f ∈ Bij S; x ∈ S] ⇒ inv_into S f x ∈ S"
  <proof>
```

```

lemma Bij_inv_into_lemma:
  assumes eq: " $\bigwedge x y. [x \in S; y \in S] \implies h(g\ x\ y) = g\ (h\ x)\ (h\ y)$ "
  shows "[ $h \in \text{Bij } S; g \in S \rightarrow S \rightarrow S; x \in S; y \in S$ ]"
     $\implies \text{inv\_into } S\ h\ (g\ x\ y) = g\ (\text{inv\_into } S\ h\ x)\ (\text{inv\_into } S\ h\ y)$ "
  <proof>

```

**definition**

```

auto :: "('a, 'b) monoid_scheme  $\Rightarrow$  ('a  $\Rightarrow$  'a) set"
where "auto G = hom G G  $\cap$  Bij (carrier G)"

```

**definition**

```

AutoGroup :: "('a, 'c) monoid_scheme  $\Rightarrow$  ('a  $\Rightarrow$  'a) monoid"
where "AutoGroup G = BijGroup (carrier G) ( $\llcorner$ carrier := auto G)"

```

```

lemma (in group) id_in_auto: " $(\lambda x \in \text{carrier } G. x) \in \text{auto } G$ "
  <proof>

```

```

lemma (in group) mult_funcset: "mult G  $\in$  carrier G  $\rightarrow$  carrier G  $\rightarrow$  carrier G"
  <proof>

```

```

lemma (in group) restrict_inv_into_hom:
  "[ $h \in \text{hom } G\ G; h \in \text{Bij } (\text{carrier } G)$ ]"
     $\implies \text{restrict } (\text{inv\_into } (\text{carrier } G)\ h)\ (\text{carrier } G) \in \text{hom } G\ G$ "
  <proof>

```

```

lemma inv_BijGroup:
  " $f \in \text{Bij } S \implies \text{m\_inv } (\text{BijGroup } S)\ f = (\lambda x \in S. (\text{inv\_into } S\ f)\ x)$ "
  <proof>

```

```

lemma (in group) subgroup_auto:
  "subgroup (auto G) (BijGroup (carrier G))"
  <proof>

```

```

theorem (in group) AutoGroup: "group (AutoGroup G)"
  <proof>

```

**end**

```

theory Divisibility
imports Permutation Coset Group
begin

```

## 7 Factorial Monoids

### 7.1 Monoids with Cancellation Law

```

locale monoid_cancel = monoid +
  assumes l_cancel:
    " $\llbracket c \otimes a = c \otimes b; a \in \text{carrier } G; b \in \text{carrier } G; c \in \text{carrier } G \rrbracket \implies a = b$ "
  and r_cancel:
    " $\llbracket a \otimes c = b \otimes c; a \in \text{carrier } G; b \in \text{carrier } G; c \in \text{carrier } G \rrbracket \implies a = b$ "

```

```

lemma (in monoid) monoid_cancelI:
  assumes l_cancel:
    " $\bigwedge a b c. \llbracket c \otimes a = c \otimes b; a \in \text{carrier } G; b \in \text{carrier } G; c \in \text{carrier } G \rrbracket \implies a = b$ "
  and r_cancel:
    " $\bigwedge a b c. \llbracket a \otimes c = b \otimes c; a \in \text{carrier } G; b \in \text{carrier } G; c \in \text{carrier } G \rrbracket \implies a = b$ "
  shows "monoid_cancel G"
  <proof>

```

```

lemma (in monoid_cancel) is_monoid_cancel:
  "monoid_cancel G"
  <proof>

```

```

sublocale group  $\subseteq$  monoid_cancel
  <proof>

```

```

locale comm_monoid_cancel = monoid_cancel + comm_monoid

```

```

lemma comm_monoid_cancelI:
  fixes G (structure)
  assumes "comm_monoid G"
  assumes cancel:
    " $\bigwedge a b c. \llbracket a \otimes c = b \otimes c; a \in \text{carrier } G; b \in \text{carrier } G; c \in \text{carrier } G \rrbracket \implies a = b$ "
  shows "comm_monoid_cancel G"
  <proof>

```

```

lemma (in comm_monoid_cancel) is_comm_monoid_cancel:
  "comm_monoid_cancel G"
  <proof>

```

```

sublocale comm_group  $\subseteq$  comm_monoid_cancel
  <proof>

```

## 7.2 Products of Units in Monoids

```
lemma (in monoid) Units_m_closed[simp, intro]:
  assumes h1unit: "h1 ∈ Units G" and h2unit: "h2 ∈ Units G"
  shows "h1 ⊗ h2 ∈ Units G"
<proof>
```

```
lemma (in monoid) prod_unit_l:
  assumes abunit[simp]: "a ⊗ b ∈ Units G" and aunit[simp]: "a ∈ Units G"
  and carr[simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "b ∈ Units G"
<proof>
```

```
lemma (in monoid) prod_unit_r:
  assumes abunit[simp]: "a ⊗ b ∈ Units G" and bunit[simp]: "b ∈ Units G"
  and carr[simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "a ∈ Units G"
<proof>
```

```
lemma (in comm_monoid) unit_factor:
  assumes abunit: "a ⊗ b ∈ Units G"
  and [simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "a ∈ Units G"
<proof>
```

## 7.3 Divisibility and Association

### 7.3.1 Function definitions

```
definition
  factor :: "[_, 'a, 'a] ⇒ bool" (infix "dividesι" 65)
  where "a dividesG b ⟷ (∃ c ∈ carrier G. b = a ⊗G c)"
```

```
definition
  associated :: "[_, 'a, 'a] ⇒ bool" (infix "∼ι" 55)
  where "a ∼G b ⟷ a dividesG b ∧ b dividesG a"
```

```
abbreviation
  "division_rel G == (carrier = carrier G, eq = op ∼G, le = op dividesG)"
```

```
definition
  properfactor :: "[_, 'a, 'a] ⇒ bool"
  where "properfactor G a b ⟷ a dividesG b ∧ ¬(b dividesG a)"
```

```
definition
  irreducible :: "[_, 'a] ⇒ bool"
  where "irreducible G a ⟷ a ∉ Units G ∧ (∀ b ∈ carrier G. properfactor G b a ⟶ b ∈ Units G)"
```



**definition**

```

prime :: "[_, 'a] ⇒ bool" where
  "prime G p ⟷
    p ∉ Units G ∧
    (∀a∈carrier G. ∀b∈carrier G. p dividesG (a ⊗G b) ⟶ p dividesG
a ∨ p dividesG b)"

```

### 7.3.2 Divisibility

**lemma dividesI:**

```

  fixes G (structure)
  assumes carr: "c ∈ carrier G"
    and p: "b = a ⊗ c"
  shows "a divides b"
⟨proof⟩

```

**lemma dividesI' [intro]:**

```

  fixes G (structure)
  assumes p: "b = a ⊗ c"
    and carr: "c ∈ carrier G"
  shows "a divides b"
⟨proof⟩

```

**lemma dividesD:**

```

  fixes G (structure)
  assumes "a divides b"
  shows "∃c∈carrier G. b = a ⊗ c"
⟨proof⟩

```

**lemma dividesE [elim]:**

```

  fixes G (structure)
  assumes d: "a divides b"
    and elim: "⋀c. [b = a ⊗ c; c ∈ carrier G] ⟹ P"
  shows "P"
⟨proof⟩

```

**lemma (in monoid) divides\_refl[simp, intro!]:**

```

  assumes carr: "a ∈ carrier G"
  shows "a divides a"
⟨proof⟩

```

**lemma (in monoid) divides\_trans [trans]:**

```

  assumes dvds: "a divides b" "b divides c"
    and acarr: "a ∈ carrier G"
  shows "a divides c"
⟨proof⟩

```

**lemma (in monoid) divides\_mult\_lI [intro]:**

```

    assumes ab: "a divides b"
    and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
    shows "(c ⊗ a) divides (c ⊗ b)"
  <proof>

lemma (in monoid_cancel) divides_mult_l [simp]:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "(c ⊗ a) divides (c ⊗ b) = a divides b"
  <proof>

lemma (in comm_monoid) divides_mult_rI [intro]:
  assumes ab: "a divides b"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "(a ⊗ c) divides (b ⊗ c)"
  <proof>

lemma (in comm_monoid_cancel) divides_mult_r [simp]:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "(a ⊗ c) divides (b ⊗ c) = a divides b"
  <proof>

lemma (in monoid) divides_prod_r:
  assumes ab: "a divides b"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "a divides (b ⊗ c)"
  <proof>

lemma (in comm_monoid) divides_prod_l:
  assumes carr[intro]: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier
G"
  and ab: "a divides b"
  shows "a divides (c ⊗ b)"
  <proof>

lemma (in monoid) unit_divides:
  assumes uunit: "u ∈ Units G"
  and acarr: "a ∈ carrier G"
  shows "u divides a"
  <proof>

lemma (in comm_monoid) divides_unit:
  assumes udvd: "a divides u"
  and carr: "a ∈ carrier G" "u ∈ Units G"
  shows "a ∈ Units G"
  <proof>

lemma (in comm_monoid) Unit_eq_dividesone:
  assumes ucarr: "u ∈ carrier G"
  shows "u ∈ Units G = u divides 1"

```

*<proof>*

### 7.3.3 Association

```
lemma associatedI:
  fixes G (structure)
  assumes "a divides b" "b divides a"
  shows "a ~ b"
<proof>
```

```
lemma (in monoid) associatedI2:
  assumes uunit[simp]: "u ∈ Units G"
  and a: "a = b ⊗ u"
  and bcarr[simp]: "b ∈ carrier G"
  shows "a ~ b"
<proof>
```

```
lemma (in monoid) associatedI2':
  assumes a: "a = b ⊗ u"
  and uunit: "u ∈ Units G"
  and bcarr: "b ∈ carrier G"
  shows "a ~ b"
<proof>
```

```
lemma associatedD:
  fixes G (structure)
  assumes "a ~ b"
  shows "a divides b"
<proof>
```

```
lemma (in monoid_cancel) associatedD2:
  assumes assoc: "a ~ b"
  and carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "∃u∈Units G. a = b ⊗ u"
<proof>
```

```
lemma associatedE:
  fixes G (structure)
  assumes assoc: "a ~ b"
  and e: "[a divides b; b divides a] ⇒ P"
  shows "P"
<proof>
```

```
lemma (in monoid_cancel) associatedE2:
  assumes assoc: "a ~ b"
  and e: "∧u. [a = b ⊗ u; u ∈ Units G] ⇒ P"
  and carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "P"
<proof>
```

```

lemma (in monoid) associated_refl [simp, intro!]:
  assumes "a ∈ carrier G"
  shows "a ~ a"
<proof>

```

```

lemma (in monoid) associated_sym [sym]:
  assumes "a ~ b"
  and "a ∈ carrier G" "b ∈ carrier G"
  shows "b ~ a"
<proof>

```

```

lemma (in monoid) associated_trans [trans]:
  assumes "a ~ b" "b ~ c"
  and "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "a ~ c"
<proof>

```

```

lemma (in monoid) division_equiv [intro, simp]:
  "equivalence (division_rel G)"
<proof>

```

### 7.3.4 Division and associativity

```

lemma divides_antisym:
  fixes G (structure)
  assumes "a divides b" "b divides a"
  and "a ∈ carrier G" "b ∈ carrier G"
  shows "a ~ b"
<proof>

```

```

lemma (in monoid) divides_cong_l [trans]:
  assumes xx': "x ~ x'"
  and xdvdy: "x' divides y"
  and carr [simp]: "x ∈ carrier G" "x' ∈ carrier G" "y ∈ carrier
G"
  shows "x divides y"
<proof>

```

```

lemma (in monoid) divides_cong_r [trans]:
  assumes xdvdy: "x divides y"
  and yy': "y ~ y'"
  and carr[simp]: "x ∈ carrier G" "y ∈ carrier G" "y' ∈ carrier G"
  shows "x divides y'"
<proof>

```

```

lemma (in monoid) division_weak_partial_order [simp, intro!]:
  "weak_partial_order (division_rel G)"
<proof>

```

### 7.3.5 Multiplication and associativity

```
lemma (in monoid_cancel) mult_cong_r:
  assumes "b ~ b'"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "b' ∈ carrier G"
  shows "a ⊗ b ~ a ⊗ b'"
<proof>
```

```
lemma (in comm_monoid_cancel) mult_cong_l:
  assumes "a ~ a'"
  and carr: "a ∈ carrier G" "a' ∈ carrier G" "b ∈ carrier G"
  shows "a ⊗ b ~ a' ⊗ b"
<proof>
```

```
lemma (in monoid_cancel) assoc_l_cancel:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G" "b' ∈ carrier G"
  and "a ⊗ b ~ a ⊗ b'"
  shows "b ~ b'"
<proof>
```

```
lemma (in comm_monoid_cancel) assoc_r_cancel:
  assumes "a ⊗ b ~ a' ⊗ b"
  and carr: "a ∈ carrier G" "a' ∈ carrier G" "b ∈ carrier G"
  shows "a ~ a'"
<proof>
```

### 7.3.6 Units

```
lemma (in monoid_cancel) assoc_unit_l [trans]:
  assumes asc: "a ~ b" and bunit: "b ∈ Units G"
  and carr: "a ∈ carrier G"
  shows "a ∈ Units G"
<proof>
```

```
lemma (in monoid_cancel) assoc_unit_r [trans]:
  assumes aunit: "a ∈ Units G" and asc: "a ~ b"
  and bcarr: "b ∈ carrier G"
  shows "b ∈ Units G"
<proof>
```

```
lemma (in comm_monoid) Units_cong:
  assumes aunit: "a ∈ Units G" and asc: "a ~ b"
  and bcarr: "b ∈ carrier G"
  shows "b ∈ Units G"
<proof>
```

```
lemma (in monoid) Units_assoc:
  assumes units: "a ∈ Units G" "b ∈ Units G"
  shows "a ~ b"
<proof>
```

```
lemma (in monoid) Units_are_ones:
  "Units G {.=}(division_rel G) {1}"
<proof>
```

```
lemma (in comm_monoid) Units_Lower:
  "Units G = Lower (division_rel G) (carrier G)"
<proof>
```

### 7.3.7 Proper factors

```
lemma properfactorI:
  fixes G (structure)
  assumes "a divides b"
    and "¬(b divides a)"
  shows "properfactor G a b"
<proof>
```

```
lemma properfactorI2:
  fixes G (structure)
  assumes advdb: "a divides b"
    and neq: "¬(a ~ b)"
  shows "properfactor G a b"
<proof>
```

```
lemma (in comm_monoid_cancel) properfactorI3:
  assumes p: "p = a ⊗ b"
    and nunit: "b ∉ Units G"
    and carr: "a ∈ carrier G" "b ∈ carrier G" "p ∈ carrier G"
  shows "properfactor G a p"
<proof>
```

```
lemma properfactorE:
  fixes G (structure)
  assumes pf: "properfactor G a b"
    and r: "[a divides b; ¬(b divides a)] ⇒ P"
  shows "P"
<proof>
```

```
lemma properfactorE2:
  fixes G (structure)
  assumes pf: "properfactor G a b"
    and elim: "[a divides b; ¬(a ~ b)] ⇒ P"
  shows "P"
<proof>
```

```
lemma (in monoid) properfactor_unitE:
  assumes uunit: "u ∈ Units G"
    and pf: "properfactor G a u"
```

```

    and acarr: "a ∈ carrier G"
    shows "P"
  <proof>

```

```

lemma (in monoid) properfactor_divides:
  assumes pf: "properfactor G a b"
  shows "a divides b"
  <proof>

```

```

lemma (in monoid) properfactor_trans1 [trans]:
  assumes dvds: "a divides b" "properfactor G b c"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "properfactor G a c"
  <proof>

```

```

lemma (in monoid) properfactor_trans2 [trans]:
  assumes dvds: "properfactor G a b" "b divides c"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "properfactor G a c"
  <proof>

```

```

lemma properfactor_lless:
  fixes G (structure)
  shows "properfactor G = lless (division_rel G)"
  <proof>

```

```

lemma (in monoid) properfactor_cong_l [trans]:
  assumes x'x: "x' ~ x"
  and pf: "properfactor G x y"
  and carr: "x ∈ carrier G" "x' ∈ carrier G" "y ∈ carrier G"
  shows "properfactor G x' y"
  <proof>

```

```

lemma (in monoid) properfactor_cong_r [trans]:
  assumes pf: "properfactor G x y"
  and yy': "y ~ y'"
  and carr: "x ∈ carrier G" "y ∈ carrier G" "y' ∈ carrier G"
  shows "properfactor G x y'"
  <proof>

```

```

lemma (in monoid_cancel) properfactor_mult_lI [intro]:
  assumes ab: "properfactor G a b"
  and carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"
  shows "properfactor G (c ⊗ a) (c ⊗ b)"
  <proof>

```

```

lemma (in monoid_cancel) properfactor_mult_l [simp]:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G" "c ∈ carrier G"

```

shows "properfactor G (c  $\otimes$  a) (c  $\otimes$  b) = properfactor G a b"  
 $\langle proof \rangle$

lemma (in comm\_monoid\_cancel) properfactor\_mult\_rI [intro]:  
 assumes ab: "properfactor G a b"  
 and carr: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"  
 shows "properfactor G (a  $\otimes$  c) (b  $\otimes$  c)"  
 $\langle proof \rangle$

lemma (in comm\_monoid\_cancel) properfactor\_mult\_r [simp]:  
 assumes carr: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"  
 shows "properfactor G (a  $\otimes$  c) (b  $\otimes$  c) = properfactor G a b"  
 $\langle proof \rangle$

lemma (in monoid) properfactor\_prod\_r:  
 assumes ab: "properfactor G a b"  
 and carr[simp]: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"  
 shows "properfactor G a (b  $\otimes$  c)"  
 $\langle proof \rangle$

lemma (in comm\_monoid) properfactor\_prod\_l:  
 assumes ab: "properfactor G a b"  
 and carr[simp]: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"  
 shows "properfactor G a (c  $\otimes$  b)"  
 $\langle proof \rangle$

## 7.4 Irreducible Elements and Primes

### 7.4.1 Irreducible elements

lemma irreducibleI:  
 fixes G (structure)  
 assumes "a  $\notin$  Units G"  
 and " $\bigwedge b. \llbracket b \in \text{carrier G}; \text{properfactor G b a} \rrbracket \implies b \in \text{Units G}$ "  
 shows "irreducible G a"  
 $\langle proof \rangle$

lemma irreducibleE:  
 fixes G (structure)  
 assumes irr: "irreducible G a"  
 and elim: " $\llbracket a \notin \text{Units G}; \forall b. b \in \text{carrier G} \wedge \text{properfactor G b a} \rrbracket \implies P$ "  
 shows "P"  
 $\langle proof \rangle$

lemma irreducibleD:  
 fixes G (structure)  
 assumes irr: "irreducible G a"  
 and pf: "properfactor G b a"  
 and bcarr: "b  $\in$  carrier G"



```

    shows "b ∈ Units G"
  <proof>

```

```

lemma (in monoid_cancel) irreducible_cong [trans]:
  assumes irred: "irreducible G a"
    and aa': "a ~ a'"
    and carr[simp]: "a ∈ carrier G" "a' ∈ carrier G"
  shows "irreducible G a'"
  <proof>

```

```

lemma (in monoid) irreducible_prod_rI:
  assumes airr: "irreducible G a"
    and bunit: "b ∈ Units G"
    and carr[simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "irreducible G (a ⊗ b)"
  <proof>

```

```

lemma (in comm_monoid) irreducible_prod_lI:
  assumes birr: "irreducible G b"
    and aunit: "a ∈ Units G"
    and carr [simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "irreducible G (a ⊗ b)"
  <proof>

```

```

lemma (in comm_monoid_cancel) irreducible_prodE [elim]:
  assumes irr: "irreducible G (a ⊗ b)"
    and carr[simp]: "a ∈ carrier G" "b ∈ carrier G"
    and e1: "[irreducible G a; b ∈ Units G] ⇒ P"
    and e2: "[a ∈ Units G; irreducible G b] ⇒ P"
  shows "P"
  <proof>

```

#### 7.4.2 Prime elements

```

lemma primeI:
  fixes G (structure)
  assumes "p ∉ Units G"
    and "∧a b. [a ∈ carrier G; b ∈ carrier G; p divides (a ⊗ b)] ⇒
p divides a ∨ p divides b"
  shows "prime G p"
  <proof>

```

```

lemma primeE:
  fixes G (structure)
  assumes pprime: "prime G p"
    and e: "[p ∉ Units G; ∀a∈carrier G. ∀b∈carrier G.
p divides a ⊗ b ⇒ p divides a ∨ p divides
b] ⇒ P"
  shows "P"

```

*<proof>*

```
lemma (in comm_monoid_cancel) prime_divides:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
    and pprime: "prime G p"
    and pdvd: "p divides a ⊗ b"
  shows "p divides a ∨ p divides b"
<proof>
```

```
lemma (in monoid_cancel) prime_cong [trans]:
  assumes pprime: "prime G p"
    and pp': "p ∼ p'"
    and carr[simp]: "p ∈ carrier G" "p' ∈ carrier G"
  shows "prime G p'"
<proof>
```

## 7.5 Factorization and Factorial Monoids

### 7.5.1 Function definitions

**definition**

```
factors :: "[_, 'a list, 'a] ⇒ bool"
where "factors G fs a ⟷ (∀x ∈ (set fs). irreducible G x) ∧ foldr
(op ⊗G) fs 1G = a"
```

**definition**

```
wfactors :: "[_, 'a list, 'a] ⇒ bool"
where "wfactors G fs a ⟷ (∀x ∈ (set fs). irreducible G x) ∧ foldr
(op ⊗G) fs 1G ∼G a"
```

**abbreviation**

```
list_assoc :: "('a, _) monoid_scheme ⇒ 'a list ⇒ 'a list ⇒ bool" (in-
fix "[~]G" 44)
where "list_assoc G == list_all2 (op ∼G)"
```

**definition**

```
essentially_equal :: "[_, 'a list, 'a list] ⇒ bool"
where "essentially_equal G fs1 fs2 ⟷ (∃fs1'. fs1 <~~> fs1' ∧ fs1'
[~]G fs2)"
```

```
locale factorial_monoid = comm_monoid_cancel +
  assumes factors_exist:
    "[a ∈ carrier G; a ∉ Units G] ⟹ ∃fs. set fs ⊆ carrier G ∧
factors G fs a"
    and factors_unique:
    "[factors G fs a; factors G fs' a; a ∈ carrier G; a ∉ Units
G;
    set fs ⊆ carrier G; set fs' ⊆ carrier G] ⟹ essentially_equal
G fs fs'"
```

## 7.5.2 Comparing lists of elements

Association on lists

```
lemma (in monoid) listassoc_refl [simp, intro]:
  assumes "set as  $\subseteq$  carrier G"
  shows "as  $[\sim]$  as"
<proof>
```

```
lemma (in monoid) listassoc_sym [sym]:
  assumes "as  $[\sim]$  bs"
  and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"
  shows "bs  $[\sim]$  as"
<proof>
```

```
lemma (in monoid) listassoc_trans [trans]:
  assumes "as  $[\sim]$  bs" and "bs  $[\sim]$  cs"
  and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G" and "set cs  $\subseteq$ 
carrier G"
  shows "as  $[\sim]$  cs"
<proof>
```

```
lemma (in monoid_cancel) irrlist_listassoc_cong:
  assumes " $\forall a \in \text{set as. irreducible G a}$ "
  and "as  $[\sim]$  bs"
  and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"
  shows " $\forall a \in \text{set bs. irreducible G a}$ "
<proof>
```

Permutations

```
lemma perm_map [intro]:
  assumes p: "a  $\langle \sim \rangle$  b"
  shows "map f a  $\langle \sim \rangle$  map f b"
<proof>
```

```
lemma perm_map_switch:
  assumes m: "map f a = map f b" and p: "b  $\langle \sim \rangle$  c"
  shows " $\exists d. a \langle \sim \rangle d \wedge \text{map f d} = \text{map f c}$ "
<proof>
```

```
lemma (in monoid) perm_assoc_switch:
  assumes a: "as  $[\sim]$  bs" and p: "bs  $\langle \sim \rangle$  cs"
  shows " $\exists bs'. as \langle \sim \rangle bs' \wedge bs' [\sim] cs$ "
<proof>
```

```
lemma (in monoid) perm_assoc_switch_r:
  assumes p: "as  $\langle \sim \rangle$  bs" and a: "bs  $[\sim]$  cs"
  shows " $\exists bs'. as [\sim] bs' \wedge bs' \langle \sim \rangle cs$ "
<proof>
```

```

declare perm_sym [sym]

lemma perm_setP:
  assumes perm: "as <~~> bs"
  and as: "P (set as)"
  shows "P (set bs)"
<proof>

lemmas (in monoid) perm_closed =
  perm_setP[of _ _ "λas. as ⊆ carrier G"]

lemmas (in monoid) irrlist_perm_cong =
  perm_setP[of _ _ "λas. ∀a∈as. irreducible G a"]

Essentially equal factorizations

lemma (in monoid) essentially_equalI:
  assumes ex: "fs1 <~~> fs1'" "fs1' [~] fs2"
  shows "essentially_equal G fs1 fs2"
<proof>

lemma (in monoid) essentially_equalE:
  assumes ee: "essentially_equal G fs1 fs2"
  and e: "∧fs1'. [fs1 <~~> fs1'; fs1' [~] fs2] ⇒ P"
  shows "P"
<proof>

lemma (in monoid) ee_refl [simp,intro]:
  assumes carr: "set as ⊆ carrier G"
  shows "essentially_equal G as as"
<proof>

lemma (in monoid) ee_sym [sym]:
  assumes ee: "essentially_equal G as bs"
  and carr: "set as ⊆ carrier G" "set bs ⊆ carrier G"
  shows "essentially_equal G bs as"
<proof>

lemma (in monoid) ee_trans [trans]:
  assumes ab: "essentially_equal G as bs" and bc: "essentially_equal
G bs cs"
  and ascarr: "set as ⊆ carrier G"
  and bscarr: "set bs ⊆ carrier G"
  and cscarr: "set cs ⊆ carrier G"
  shows "essentially_equal G as cs"
<proof>

```

### 7.5.3 Properties of lists of elements

Multiplication of factors in a list

```

lemma (in monoid) multlist_closed [simp, intro]:
  assumes ascarr: "set fs  $\subseteq$  carrier G"
  shows "foldr (op  $\otimes$ ) fs 1  $\in$  carrier G"
<proof>

lemma (in comm_monoid) multlist_dividesI :
  assumes "f  $\in$  set fs" and "f  $\in$  carrier G" and "set fs  $\subseteq$  carrier G"
  shows "f divides (foldr (op  $\otimes$ ) fs 1)"
<proof>

lemma (in comm_monoid_cancel) multlist_listassoc_cong:
  assumes "fs [~] fs'"
  and "set fs  $\subseteq$  carrier G" and "set fs'  $\subseteq$  carrier G"
  shows "foldr (op  $\otimes$ ) fs 1  $\sim$  foldr (op  $\otimes$ ) fs' 1"
<proof>

lemma (in comm_monoid) multlist_perm_cong:
  assumes prm: "as <~~> bs"
  and ascarr: "set as  $\subseteq$  carrier G"
  shows "foldr (op  $\otimes$ ) as 1 = foldr (op  $\otimes$ ) bs 1"
<proof>

lemma (in comm_monoid_cancel) multlist_ee_cong:
  assumes "essentially_equal G fs fs'"
  and "set fs  $\subseteq$  carrier G" and "set fs'  $\subseteq$  carrier G"
  shows "foldr (op  $\otimes$ ) fs 1  $\sim$  foldr (op  $\otimes$ ) fs' 1"
<proof>

7.5.4 Factorization in irreducible elements

lemma wfactorsI:
  fixes G (structure)
  assumes " $\forall f \in \text{set fs. irreducible } G \ f$ "
  and "foldr (op  $\otimes$ ) fs 1  $\sim$  a"
  shows "wfactors G fs a"
<proof>

lemma wfactorsE:
  fixes G (structure)
  assumes wf: "wfactors G fs a"
  and e: " $\llbracket \forall f \in \text{set fs. irreducible } G \ f; \text{foldr (op } \otimes \text{) fs 1 } \sim \text{ a} \rrbracket \implies$ "
  p"
  shows "P"
<proof>

lemma (in monoid) factorsI:
  assumes " $\forall f \in \text{set fs. irreducible } G \ f$ "
  and "foldr (op  $\otimes$ ) fs 1 = a"
  shows "factors G fs a"

```

*<proof>*

```
lemma factorsE:
  fixes G (structure)
  assumes f: "factors G fs a"
    and e: "[ $\forall f \in \text{set fs. irreducible } G f; \text{foldr } (\text{op } \otimes) \text{ fs } 1 = a$ ]  $\implies P$ "
  shows "P"
<proof>
```

```
lemma (in monoid) factors_wfactors:
  assumes "factors G as a" and "set as  $\subseteq$  carrier G"
  shows "wfactors G as a"
<proof>
```

```
lemma (in monoid) wfactors_factors:
  assumes "wfactors G as a" and "set as  $\subseteq$  carrier G"
  shows " $\exists a'. \text{factors } G \text{ as } a' \wedge a' \sim a$ "
<proof>
```

```
lemma (in monoid) factors_closed [dest]:
  assumes "factors G fs a" and "set fs  $\subseteq$  carrier G"
  shows "a  $\in$  carrier G"
<proof>
```

```
lemma (in monoid) nunit_factors:
  assumes anunit: "a  $\notin$  Units G"
    and fs: "factors G as a"
  shows "length as > 0"
<proof>
```

```
lemma (in monoid) unit_wfactors [simp]:
  assumes aunit: "a  $\in$  Units G"
  shows "wfactors G [] a"
<proof>
```

```
lemma (in comm_monoid_cancel) unit_wfactors_empty:
  assumes aunit: "a  $\in$  Units G"
    and wf: "wfactors G fs a"
    and carr[simp]: "set fs  $\subseteq$  carrier G"
  shows "fs = []"
<proof>
```

Comparing wfactors

```
lemma (in comm_monoid_cancel) wfactors_listassoc_cong_1:
  assumes fact: "wfactors G fs a"
    and asc: "fs  $[\sim]$  fs'"
    and carr: "a  $\in$  carrier G" "set fs  $\subseteq$  carrier G" "set fs'  $\subseteq$  carrier G"
  shows "wfactors G fs' a"
```

*<proof>*

```
lemma (in comm_monoid) wfactors_perm_cong_l:
  assumes "wfactors G fs a"
    and "fs <~~> fs'"
    and "set fs  $\subseteq$  carrier G"
  shows "wfactors G fs' a"
<proof>
```

```
lemma (in comm_monoid_cancel) wfactors_ee_cong_l [trans]:
  assumes ee: "essentially_equal G as bs"
    and bfs: "wfactors G bs b"
    and carr: "b  $\in$  carrier G" "set as  $\subseteq$  carrier G" "set bs  $\subseteq$  carrier
G"
  shows "wfactors G as b"
<proof>
```

```
lemma (in monoid) wfactors_cong_r [trans]:
  assumes fac: "wfactors G fs a" and aa': "a  $\sim$  a'"
    and carr[simp]: "a  $\in$  carrier G" "a'  $\in$  carrier G" "set fs  $\subseteq$  carrier
G"
  shows "wfactors G fs a'"
<proof>
```

### 7.5.5 Essentially equal factorizations

```
lemma (in comm_monoid_cancel) unitfactor_ee:
  assumes uunit: "u  $\in$  Units G"
    and carr: "set as  $\subseteq$  carrier G"
  shows "essentially_equal G (as[0 := (as!0  $\otimes$  u)]) as" (is "essentially_equal
G ?as' as")
<proof>
```

```
lemma (in comm_monoid_cancel) factors_cong_unit:
  assumes uunit: "u  $\in$  Units G" and anunit: "a  $\notin$  Units G"
    and afs: "factors G as a"
    and ascarr: "set as  $\subseteq$  carrier G"
  shows "factors G (as[0 := (as!0  $\otimes$  u)]) (a  $\otimes$  u)" (is "factors G ?as'
?a'")
<proof>
```

```
lemma (in comm_monoid) perm_wfactorsD:
  assumes prm: "as <~~> bs"
    and afs: "wfactors G as a" and bfs: "wfactors G bs b"
    and [simp]: "a  $\in$  carrier G" "b  $\in$  carrier G"
    and ascarr[simp]: "set as  $\subseteq$  carrier G"
  shows "a  $\sim$  b"
<proof>
```

```

lemma (in comm_monoid_cancel) listassoc_wfactorsD:
  assumes assoc: "as [~] bs"
    and afs: "wfactors G as a" and bfs: "wfactors G bs b"
    and [simp]: "a ∈ carrier G" "b ∈ carrier G"
    and [simp]: "set as ⊆ carrier G" "set bs ⊆ carrier G"
  shows "a ~ b"
<proof>

lemma (in comm_monoid_cancel) ee_wfactorsD:
  assumes ee: "essentially_equal G as bs"
    and afs: "wfactors G as a" and bfs: "wfactors G bs b"
    and [simp]: "a ∈ carrier G" "b ∈ carrier G"
    and ascarr[simp]: "set as ⊆ carrier G" and bscarr[simp]: "set bs
⊆ carrier G"
  shows "a ~ b"
<proof>

lemma (in comm_monoid_cancel) ee_factorsD:
  assumes ee: "essentially_equal G as bs"
    and afs: "factors G as a" and bfs: "factors G bs b"
    and "set as ⊆ carrier G" "set bs ⊆ carrier G"
  shows "a ~ b"
<proof>

lemma (in factorial_monoid) ee_factorsI:
  assumes ab: "a ~ b"
    and afs: "factors G as a" and anunit: "a ∉ Units G"
    and bfs: "factors G bs b" and bnunit: "b ∉ Units G"
    and ascarr: "set as ⊆ carrier G" and bscarr: "set bs ⊆ carrier G"
  shows "essentially_equal G as bs"
<proof>

lemma (in factorial_monoid) ee_wfactorsI:
  assumes asc: "a ~ b"
    and asf: "wfactors G as a" and bsf: "wfactors G bs b"
    and acarr[simp]: "a ∈ carrier G" and bcarr[simp]: "b ∈ carrier G"
    and ascarr[simp]: "set as ⊆ carrier G" and bscarr[simp]: "set bs
⊆ carrier G"
  shows "essentially_equal G as bs"
<proof>

lemma (in factorial_monoid) ee_wfactors:
  assumes asf: "wfactors G as a"
    and bsf: "wfactors G bs b"
    and acarr: "a ∈ carrier G" and bcarr: "b ∈ carrier G"
    and ascarr: "set as ⊆ carrier G" and bscarr: "set bs ⊆ carrier G"
  shows asc: "a ~ b = essentially_equal G as bs"
<proof>

```



```

lemma (in factorial_monoid) wfactors_exist [intro, simp]:
  assumes acar: "a ∈ carrier G"
  shows "∃ fs. set fs ⊆ carrier G ∧ wfactors G fs a"
⟨proof⟩

lemma (in monoid) wfactors_prod_exists [intro, simp]:
  assumes "∀ a ∈ set as. irreducible G a" and "set as ⊆ carrier G"
  shows "∃ a. a ∈ carrier G ∧ wfactors G as a"
⟨proof⟩

lemma (in factorial_monoid) wfactors_unique:
  assumes "wfactors G fs a" and "wfactors G fs' a"
  and "a ∈ carrier G"
  and "set fs ⊆ carrier G" and "set fs' ⊆ carrier G"
  shows "essentially_equal G fs fs'"
⟨proof⟩

lemma (in monoid) factors_mult_single:
  assumes "irreducible G a" and "factors G fb b" and "a ∈ carrier G"
  shows "factors G (a # fb) (a ⊗ b)"
⟨proof⟩

lemma (in monoid_cancel) wfactors_mult_single:
  assumes f: "irreducible G a" "wfactors G fb b"
  "a ∈ carrier G" "b ∈ carrier G" "set fb ⊆ carrier G"
  shows "wfactors G (a # fb) (a ⊗ b)"
⟨proof⟩

lemma (in monoid) factors_mult:
  assumes factors: "factors G fa a" "factors G fb b"
  and acar: "set fa ⊆ carrier G" and bscarr: "set fb ⊆ carrier G"
  shows "factors G (fa @ fb) (a ⊗ b)"
⟨proof⟩

lemma (in comm_monoid_cancel) wfactors_mult [intro]:
  assumes asf: "wfactors G as a" and bsf: "wfactors G bs b"
  and acar: "a ∈ carrier G" and bcarr: "b ∈ carrier G"
  and ascarr: "set as ⊆ carrier G" and bscarr: "set bs ⊆ carrier G"
  shows "wfactors G (as @ bs) (a ⊗ b)"
⟨proof⟩

lemma (in comm_monoid) factors_dividesI:
  assumes "factors G fs a" and "f ∈ set fs"
  and "set fs ⊆ carrier G"
  shows "f divides a"
⟨proof⟩

lemma (in comm_monoid) wfactors_dividesI:
  assumes p: "wfactors G fs a"

```

```

    and fscarr: "set fs  $\subseteq$  carrier G" and acarr: "a  $\in$  carrier G"
    and f: "f  $\in$  set fs"
    shows "f divides a"
  <proof>

```

### 7.5.6 Factorial monoids and wfactors

```

lemma (in comm_monoid_cancel) factorial_monoidI:
  assumes wfactors_exists:
    " $\bigwedge a. a \in \text{carrier } G \implies \exists \text{fs. set fs} \subseteq \text{carrier } G \wedge \text{wfactors}$ 
G fs a"
    and wfactors_unique:
    " $\bigwedge a \text{ fs fs}'. \llbracket a \in \text{carrier } G; \text{set fs} \subseteq \text{carrier } G; \text{set fs}' \subseteq \text{carrier}$ 
G;
    wfactors G fs a; wfactors G fs' a  $\rrbracket \implies \text{essentially\_equal}$ 
G fs fs'"
  shows "factorial_monoid G"
  <proof>

```

## 7.6 Factorizations as Multisets

Gives useful operations like intersection

### abbreviation

```
"assocs G x == eq_closure_of (division_rel G) {x}"
```

### definition

```
"fmset G as = multiset_of (map ( $\lambda a. \text{assocs } G a$ ) as)"
```

Helper lemmas

```

lemma (in monoid) assocs_repr_independence:
  assumes "y  $\in$  assocs G x"
  and "x  $\in$  carrier G"
  shows "assocs G x = assocs G y"
  <proof>

```

```

lemma (in monoid) assocs_self:
  assumes "x  $\in$  carrier G"
  shows "x  $\in$  assocs G x"
  <proof>

```

```

lemma (in monoid) assocs_repr_independenceD:
  assumes repr: "assocs G x = assocs G y"
  and ycarr: "y  $\in$  carrier G"
  shows "y  $\in$  assocs G x"
  <proof>

```

```

lemma (in comm_monoid) assocs_assoc:
  assumes "a  $\in$  assocs G b"
  and "b  $\in$  carrier G"

```

shows "a  $\sim$  b"  
*<proof>*

lemmas (in comm\_monoid) assocs\_eqD =  
 assocs\_repr\_independenceD[THEN assocs\_assoc]

### 7.6.1 Comparing multisets

lemma (in monoid) fmset\_perm\_cong:  
 assumes prm: "as  $\sim$  bs"  
 shows "fmset G as = fmset G bs"  
*<proof>*

lemma (in comm\_monoid\_cancel) eqc\_listassoc\_cong:  
 assumes "as  $[\sim]$  bs"  
 and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "map (assocs G) as = map (assocs G) bs"  
*<proof>*

lemma (in comm\_monoid\_cancel) fmset\_listassoc\_cong:  
 assumes "as  $[\sim]$  bs"  
 and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "fmset G as = fmset G bs"  
*<proof>*

lemma (in comm\_monoid\_cancel) ee\_fmset:  
 assumes ee: "essentially\_equal G as bs"  
 and ascarr: "set as  $\subseteq$  carrier G" and bscarr: "set bs  $\subseteq$  carrier G"  
 shows "fmset G as = fmset G bs"  
*<proof>*

lemma (in monoid\_cancel) fmset\_ee\_hlp\_induct:  
 assumes prm: "cas  $\sim$  cbs"  
 and cdef: "cas = map (assocs G) as" "cbs = map (assocs G) bs"  
 shows " $\forall$  as bs. (cas  $\sim$  cbs  $\wedge$  cas = map (assocs G) as  $\wedge$   
 cbs = map (assocs G) bs)  $\longrightarrow$  ( $\exists$  as'. as  $\sim$  as'  $\wedge$  map  
 (assocs G) as' = cbs)"  
*<proof>*

lemma (in comm\_monoid\_cancel) fmset\_ee:  
 assumes mset: "fmset G as = fmset G bs"  
 and ascarr: "set as  $\subseteq$  carrier G" and bscarr: "set bs  $\subseteq$  carrier G"  
 shows "essentially\_equal G as bs"  
*<proof>*

lemma (in comm\_monoid\_cancel) ee\_is\_fmset:  
 assumes "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "essentially\_equal G as bs = (fmset G as = fmset G bs)"  
*<proof>*

### 7.6.2 Interpreting multisets as factorizations

```

lemma (in monoid) mset_fmsetEx:
  assumes elems: " $\bigwedge X. X \in \text{set\_of } Cs \implies \exists x. P\ x \wedge X = \text{assocs } G\ x$ "
  shows " $\exists cs. (\forall c \in \text{set } cs. P\ c) \wedge \text{fmset } G\ cs = Cs$ "
<proof>

lemma (in monoid) mset_wfactorsEx:
  assumes elems: " $\bigwedge X. X \in \text{set\_of } Cs$ "
   $\implies \exists x. (x \in \text{carrier } G \wedge \text{irreducible } G\ x) \wedge X =$ 
   $\text{assocs } G\ x$ 
  shows " $\exists c\ cs. c \in \text{carrier } G \wedge \text{set } cs \subseteq \text{carrier } G \wedge \text{wfactors } G\ cs\ c$ 
 $\wedge \text{fmset } G\ cs = Cs$ "
<proof>

```

### 7.6.3 Multiplication on multisets

```

lemma (in factorial_monoid) mult_wfactors_fmset:
  assumes afs: "wfactors G as a" and bfs: "wfactors G bs b" and cfs:
  "wfactors G cs (a  $\otimes$  b)"
  and carr: "a  $\in$  carrier G" "b  $\in$  carrier G"
  "set as  $\subseteq$  carrier G" "set bs  $\subseteq$  carrier G" "set cs  $\subseteq$  carrier
  G"
  shows "fmset G cs = fmset G as + fmset G bs"
<proof>

lemma (in factorial_monoid) mult_factors_fmset:
  assumes afs: "factors G as a" and bfs: "factors G bs b" and cfs: "factors
  G cs (a  $\otimes$  b)"
  and "set as  $\subseteq$  carrier G" "set bs  $\subseteq$  carrier G" "set cs  $\subseteq$  carrier
  G"
  shows "fmset G cs = fmset G as + fmset G bs"
<proof>

lemma (in comm_monoid_cancel) fmset_wfactors_mult:
  assumes mset: "fmset G cs = fmset G as + fmset G bs"
  and carr: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"
  "set as  $\subseteq$  carrier G" "set bs  $\subseteq$  carrier G" "set cs  $\subseteq$  carrier
  G"
  and fs: "wfactors G as a" "wfactors G bs b" "wfactors G cs c"
  shows "c  $\sim$  a  $\otimes$  b"
<proof>

```

### 7.6.4 Divisibility on multisets

```

lemma (in factorial_monoid) divides_fmsubset:
  assumes ab: "a divides b"
  and afs: "wfactors G as a" and bfs: "wfactors G bs b"
  and carr: "a  $\in$  carrier G" "b  $\in$  carrier G" "set as  $\subseteq$  carrier G"
  "set bs  $\subseteq$  carrier G"

```

shows "fmset G as  $\leq$  fmset G bs"  
 $\langle proof \rangle$

lemma (in comm\_monoid\_cancel) fmsubset\_divides:  
 assumes msubset: "fmset G as  $\leq$  fmset G bs"  
 and afs: "wfactors G as a" and bfs: "wfactors G bs b"  
 and acar: "a  $\in$  carrier G" and bcarr: "b  $\in$  carrier G"  
 and ascarr: "set as  $\subseteq$  carrier G" and bscarr: "set bs  $\subseteq$  carrier G"  
 shows "a divides b"  
 $\langle proof \rangle$

lemma (in factorial\_monoid) divides\_as\_fmsubset:  
 assumes "wfactors G as a" and "wfactors G bs b"  
 and "a  $\in$  carrier G" and "b  $\in$  carrier G"  
 and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "a divides b = (fmset G as  $\leq$  fmset G bs)"  
 $\langle proof \rangle$

Proper factors on multisets

lemma (in factorial\_monoid) fmset\_properfactor:  
 assumes asubb: "fmset G as  $\leq$  fmset G bs"  
 and anb: "fmset G as  $\neq$  fmset G bs"  
 and "wfactors G as a" and "wfactors G bs b"  
 and "a  $\in$  carrier G" and "b  $\in$  carrier G"  
 and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "properfactor G a b"  
 $\langle proof \rangle$

lemma (in factorial\_monoid) properfactor\_fmset:  
 assumes pf: "properfactor G a b"  
 and "wfactors G as a" and "wfactors G bs b"  
 and "a  $\in$  carrier G" and "b  $\in$  carrier G"  
 and "set as  $\subseteq$  carrier G" and "set bs  $\subseteq$  carrier G"  
 shows "fmset G as  $\leq$  fmset G bs  $\wedge$  fmset G as  $\neq$  fmset G bs"  
 $\langle proof \rangle$

## 7.7 Irreducible Elements are Prime

lemma (in factorial\_monoid) irreducible\_is\_prime:  
 assumes pirr: "irreducible G p"  
 and pcarr: "p  $\in$  carrier G"  
 shows "prime G p"  
 $\langle proof \rangle$

lemma (in factorial\_monoid) factors\_irreducible\_is\_prime:  
 assumes pirr: "irreducible G p"  
 and pcarr: "p  $\in$  carrier G"  
 shows "prime G p"  
 $\langle proof \rangle$

## 7.8 Greatest Common Divisors and Lowest Common Multiples

### 7.8.1 Definitions

**definition**

```
isgcd :: "(['a,_) monoid_scheme, 'a, 'a, 'a]  $\Rightarrow$  bool" ("(_ gcdofz _ _)" [81,81,81] 80)
where "x gcdofG a b  $\longleftrightarrow$  x dividesG a  $\wedge$  x dividesG b  $\wedge$ 
      ( $\forall y \in \text{carrier } G. (y \text{ divides}_G a \wedge y \text{ divides}_G b \longrightarrow y \text{ divides}_G x)$ )"
```

**definition**

```
islcm :: "[_ , 'a, 'a, 'a]  $\Rightarrow$  bool" ("(_ lcmofz _ _)" [81,81,81] 80)
where "x lcmofG a b  $\longleftrightarrow$  a dividesG x  $\wedge$  b dividesG x  $\wedge$ 
      ( $\forall y \in \text{carrier } G. (a \text{ divides}_G y \wedge b \text{ divides}_G y \longrightarrow x \text{ divides}_G y)$ )"
```

**definition**

```
somegcd :: "(['a,_) monoid_scheme  $\Rightarrow$  'a  $\Rightarrow$  'a  $\Rightarrow$  'a"
where "somegcd G a b = (SOME x. x  $\in$  carrier G  $\wedge$  x gcdofG a b)"
```

**definition**

```
somelcm :: "(['a,_) monoid_scheme  $\Rightarrow$  'a  $\Rightarrow$  'a  $\Rightarrow$  'a"
where "somelcm G a b = (SOME x. x  $\in$  carrier G  $\wedge$  x lcmofG a b)"
```

**definition**

```
"SomeGcd G A = inf (division_rel G) A"
```

```
locale gcd_condition_monoid = comm_monoid_cancel +
```

```
  assumes gcdof_exists:
```

```
    "[a  $\in$  carrier G; b  $\in$  carrier G]  $\Longrightarrow$   $\exists c. c \in \text{carrier } G \wedge c \text{ gcdof } a \ b$ "
```

```
locale primeness_condition_monoid = comm_monoid_cancel +
```

```
  assumes irreducible_prime:
```

```
    "[a  $\in$  carrier G; irreducible G a]  $\Longrightarrow$  prime G a"
```

```
locale divisor_chain_condition_monoid = comm_monoid_cancel +
```

```
  assumes division_wellfounded:
```

```
    "wf {(x, y). x  $\in$  carrier G  $\wedge$  y  $\in$  carrier G  $\wedge$  properfactor G x y}"
```

### 7.8.2 Connections to Lattice.thy

**lemma** gcdof\_greatestLower:

```
  fixes G (structure)
```

```
  assumes carr[simp]: "a  $\in$  carrier G" "b  $\in$  carrier G"
```

```
  shows "(x  $\in$  carrier G  $\wedge$  x gcdof a b) =
```

```
    greatest (division_rel G) x (Lower (division_rel G) {a, b})"
```

*<proof>*

```

lemma lcmof_leastUpper:
  fixes G (structure)
  assumes carr[simp]: "a ∈ carrier G" "b ∈ carrier G"
  shows "(x ∈ carrier G ∧ x lcmof a b) =
    least (division_rel G) x (Upper (division_rel G) {a, b})"
  <proof>

```

```

lemma somegcd_meet:
  fixes G (structure)
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "somegcd G a b = meet (division_rel G) a b"
  <proof>

```

```

lemma (in monoid) isgcd_divides_l:
  assumes "a divides b"
  and "a ∈ carrier G" "b ∈ carrier G"
  shows "a gcdof a b"
  <proof>

```

```

lemma (in monoid) isgcd_divides_r:
  assumes "b divides a"
  and "a ∈ carrier G" "b ∈ carrier G"
  shows "b gcdof a b"
  <proof>

```

### 7.8.3 Existence of gcd and lcm

```

lemma (in factorial_monoid) gcdof_exists:
  assumes acarr: "a ∈ carrier G" and bcarr: "b ∈ carrier G"
  shows "∃c. c ∈ carrier G ∧ c gcdof a b"
  <proof>

```

```

lemma (in factorial_monoid) lcmof_exists:
  assumes acarr: "a ∈ carrier G" and bcarr: "b ∈ carrier G"
  shows "∃c. c ∈ carrier G ∧ c lcmof a b"
  <proof>

```

## 7.9 Conditions for Factoriality

### 7.9.1 Gcd condition

```

lemma (in gcd_condition_monoid) division_weak_lower_semilattice [simp]:
  shows "weak_lower_semilattice (division_rel G)"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcdof_cong_l:
  assumes a'a: "a' ~ a"
  and agcd: "a gcdof b c"

```

```

    and a'carr: "a' ∈ carrier G" and carr': "a ∈ carrier G" "b ∈ carrier
G" "c ∈ carrier G"
    shows "a' gcdof b c"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_closed [simp]:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "somegcd G a b ∈ carrier G"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_isgcd:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "(somegcd G a b) gcdof a b"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_exists:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "∃x∈carrier G. x = somegcd G a b"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_divides_l:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "(somegcd G a b) divides a"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_divides_r:
  assumes carr: "a ∈ carrier G" "b ∈ carrier G"
  shows "(somegcd G a b) divides b"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_divides:
  assumes sub: "z divides x" "z divides y"
  and L: "x ∈ carrier G" "y ∈ carrier G" "z ∈ carrier G"
  shows "z divides (somegcd G x y)"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_cong_l:
  assumes xx': "x ∼ x'"
  and carr: "x ∈ carrier G" "x' ∈ carrier G" "y ∈ carrier G"
  shows "somegcd G x y ∼ somegcd G x' y"
  <proof>

```

```

lemma (in gcd_condition_monoid) gcd_cong_r:
  assumes carr: "x ∈ carrier G" "y ∈ carrier G" "y' ∈ carrier G"
  and yy': "y ∼ y'"
  shows "somegcd G x y ∼ somegcd G x y'"
  <proof>

```



```

lemma (in gcd_condition_monoid) gcdI:
  assumes dvd: "a divides b" "a divides c"
    and others: " $\forall y \in \text{carrier } G. y \text{ divides } b \wedge y \text{ divides } c \longrightarrow y \text{ divides } a$ "
    and acarr: "a  $\in$  carrier G" and bcarr: "b  $\in$  carrier G" and ccarr:
      "c  $\in$  carrier G"
    shows "a  $\sim$  somegcd G b c"
  <proof>

lemma (in gcd_condition_monoid) gcdI2:
  assumes "a gcdof b c"
    and "a  $\in$  carrier G" and bcarr: "b  $\in$  carrier G" and ccarr: "c  $\in$  carrier
      G"
    shows "a  $\sim$  somegcd G b c"
  <proof>

lemma (in gcd_condition_monoid) SomeGcd_ex:
  assumes "finite A" "A  $\subseteq$  carrier G" "A  $\neq$  {}"
    shows " $\exists x \in$  carrier G. x = SomeGcd G A"
  <proof>

lemma (in gcd_condition_monoid) gcd_assoc:
  assumes carr: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"
    shows "somegcd G (somegcd G a b) c  $\sim$  somegcd G a (somegcd G b c)"
  <proof>

lemma (in gcd_condition_monoid) gcd_mult:
  assumes acarr: "a  $\in$  carrier G" and bcarr: "b  $\in$  carrier G" and ccarr:
    "c  $\in$  carrier G"
    shows "c  $\otimes$  somegcd G a b  $\sim$  somegcd G (c  $\otimes$  a) (c  $\otimes$  b)"
  <proof>

lemma (in monoid) assoc_subst:
  assumes ab: "a  $\sim$  b"
    and cP: "ALL a b. a : carrier G & b : carrier G & a  $\sim$  b
      --> f a : carrier G & f b : carrier G & f a  $\sim$  f b"
    and carr: "a  $\in$  carrier G" "b  $\in$  carrier G"
    shows "f a  $\sim$  f b"
  <proof>

lemma (in gcd_condition_monoid) relprime_mult:
  assumes abrelprime: "somegcd G a b  $\sim$  1" and acrelprime: "somegcd G
    a c  $\sim$  1"
    and carr[simp]: "a  $\in$  carrier G" "b  $\in$  carrier G" "c  $\in$  carrier G"
    shows "somegcd G a (b  $\otimes$  c)  $\sim$  1"
  <proof>

lemma (in gcd_condition_monoid) primeness_condition:

```

```
"primeness_condition_monoid G"
⟨proof⟩
```

```
sublocale gcd_condition_monoid ⊆ primeness_condition_monoid
⟨proof⟩
```

### 7.9.2 Divisor chain condition

```
lemma (in divisor_chain_condition_monoid) wfactors_exist:
  assumes acarr: "a ∈ carrier G"
  shows "∃ as. set as ⊆ carrier G ∧ wfactors G as a"
⟨proof⟩
```

### 7.9.3 Primeness condition

```
lemma (in comm_monoid_cancel) multlist_prime_pos:
  assumes carr: "a ∈ carrier G" "set as ⊆ carrier G"
  and aprime: "prime G a"
  and "a divides (foldr (op ⊗) as 1)"
  shows "∃ i < length as. a divides (as!i)"
⟨proof⟩
```

```
lemma (in primeness_condition_monoid) wfactors_unique_hlp_induct:
  "∀ a as'. a ∈ carrier G ∧ set as ⊆ carrier G ∧ set as' ⊆ carrier G
  ∧
  wfactors G as a ∧ wfactors G as' a ⟶ essentially_equal G
  as as'"
⟨proof⟩
```

```
lemma (in primeness_condition_monoid) wfactors_unique:
  assumes "wfactors G as a" "wfactors G as' a"
  and "a ∈ carrier G" "set as ⊆ carrier G" "set as' ⊆ carrier G"
  shows "essentially_equal G as as'"
⟨proof⟩
```

### 7.9.4 Application to factorial monoids

Number of factors for wellfoundedness

**definition**

```
factorcount :: "_ ⇒ 'a ⇒ nat" where
  "factorcount G a =
  (THE c. (ALL as. set as ⊆ carrier G ∧ wfactors G as a ⟶ c = length
  as))"
```

```
lemma (in monoid) ee_length:
  assumes ee: "essentially_equal G as bs"
  shows "length as = length bs"
⟨proof⟩
```

```

lemma (in factorial_monoid) factorcount_exists:
  assumes carr[simp]: "a ∈ carrier G"
  shows "EX c. ALL as. set as ⊆ carrier G ∧ wfactors G as a ⟶ c = length
as"
⟨proof⟩

```

```

lemma (in factorial_monoid) factorcount_unique:
  assumes afs: "wfactors G as a"
  and acarr[simp]: "a ∈ carrier G" and ascarr[simp]: "set as ⊆ carrier
G"
  shows "factorcount G a = length as"
⟨proof⟩

```

```

lemma (in factorial_monoid) divides_fcount:
  assumes dvd: "a divides b"
  and acarr: "a ∈ carrier G" and bcarr:"b ∈ carrier G"
  shows "factorcount G a ≤ factorcount G b"
⟨proof⟩

```

```

lemma (in factorial_monoid) associated_fcount:
  assumes acarr: "a ∈ carrier G" and bcarr:"b ∈ carrier G"
  and asc: "a ∼ b"
  shows "factorcount G a = factorcount G b"
⟨proof⟩

```

```

lemma (in factorial_monoid) properfactor_fcount:
  assumes acarr: "a ∈ carrier G" and bcarr:"b ∈ carrier G"
  and pf: "properfactor G a b"
  shows "factorcount G a < factorcount G b"
⟨proof⟩

```

```

sublocale factorial_monoid ⊆ divisor_chain_condition_monoid
⟨proof⟩

```

```

sublocale factorial_monoid ⊆ primeness_condition_monoid
⟨proof⟩

```

```

lemma (in factorial_monoid) primeness_condition:
  shows "primeness_condition_monoid G"
⟨proof⟩

```

```

lemma (in factorial_monoid) gcd_condition [simp]:
  shows "gcd_condition_monoid G"
⟨proof⟩

```

```

sublocale factorial_monoid ⊆ gcd_condition_monoid
⟨proof⟩

```

```
lemma (in factorial_monoid) division_weak_lattice [simp]:
  shows "weak_lattice (division_rel G)"
  <proof>
```

## 7.10 Factoriality Theorems

```
theorem factorial_condition_one:
  shows "(divisor_chain_condition_monoid G ∧ primeness_condition_monoid
G) =
  factorial_monoid G"
  <proof>
```

```
theorem factorial_condition_two:
  shows "(divisor_chain_condition_monoid G ∧ gcd_condition_monoid G)
= factorial_monoid G"
  <proof>
```

end

```
theory Ring
imports FiniteProduct
uses ("ringsimp.ML")
begin
```

# 8 The Algebraic Hierarchy of Rings

## 8.1 Abelian Groups

```
record 'a ring = "'a monoid" +
  zero :: 'a ("0")
  add :: "'a, 'a] => 'a" (infixl "⊕" 65)
```

Derived operations.

```
definition
  a_inv :: "['a, 'm) ring_scheme, 'a] => 'a" ("⊖" [81] 80)
  where "a_inv R = m_inv (| carrier = carrier R, mult = add R, one =
zero R |)"
```

```
definition
  a_minus :: "['a, 'm) ring_scheme, 'a, 'a] => 'a" (infixl "⊖" 65)
  where "[| x ∈ carrier R; y ∈ carrier R |] ==> x ⊖R y = x ⊕R (⊖R y)"
```

```
locale abelian_monoid =
  fixes G (structure)
  assumes a_comm_monoid:
    "comm_monoid (| carrier = carrier G, mult = add G, one = zero G |)"
```

The following definition is redundant but simple to use.

```

locale abelian_group = abelian_monoid +
  assumes a_comm_group:
    "comm_group (| carrier = carrier G, mult = add G, one = zero G |)"

```

## 8.2 Basic Properties

```

lemma abelian_monoidI:
  fixes R (structure)
  assumes a_closed:
    "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==> x ⊕ y ∈ carrier
R"
  and zero_closed: "0 ∈ carrier R"
  and a_assoc:
    "!!x y z. [| x ∈ carrier R; y ∈ carrier R; z ∈ carrier R |] ==>
(x ⊕ y) ⊕ z = x ⊕ (y ⊕ z)"
  and l_zero: "!!x. x ∈ carrier R ==> 0 ⊕ x = x"
  and a_comm:
    "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==> x ⊕ y = y ⊕ x"
  shows "abelian_monoid R"
  <proof>

```

```

lemma abelian_groupI:
  fixes R (structure)
  assumes a_closed:
    "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==> x ⊕ y ∈ carrier
R"
  and zero_closed: "zero R ∈ carrier R"
  and a_assoc:
    "!!x y z. [| x ∈ carrier R; y ∈ carrier R; z ∈ carrier R |] ==>
(x ⊕ y) ⊕ z = x ⊕ (y ⊕ z)"
  and a_comm:
    "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==> x ⊕ y = y ⊕ x"
  and l_zero: "!!x. x ∈ carrier R ==> 0 ⊕ x = x"
  and l_inv_ex: "!!x. x ∈ carrier R ==> EX y : carrier R. y ⊕ x = 0"
  shows "abelian_group R"
  <proof>

```

```

lemma (in abelian_monoid) a_monoid:
  "monoid (| carrier = carrier G, mult = add G, one = zero G |)"
  <proof>

```

```

lemma (in abelian_group) a_group:
  "group (| carrier = carrier G, mult = add G, one = zero G |)"
  <proof>

```

```

lemmas monoid_record_simps = partial_object.simps monoid.simps

```

```

lemma (in abelian_monoid) a_closed [intro, simp]:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> x ⊕ y ∈ carrier G"

```

*<proof>*

**lemma** (in abelian\_monoid) zero\_closed [intro, simp]:  
 "0 ∈ carrier G"  
*<proof>*

**lemma** (in abelian\_group) a\_inv\_closed [intro, simp]:  
 "x ∈ carrier G ==>  $\ominus$  x ∈ carrier G"  
*<proof>*

**lemma** (in abelian\_group) minus\_closed [intro, simp]:  
 "[| x ∈ carrier G; y ∈ carrier G |] ==> x  $\ominus$  y ∈ carrier G"  
*<proof>*

**lemma** (in abelian\_group) a\_l\_cancel [simp]:  
 "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>  
 (x  $\oplus$  y = x  $\oplus$  z) = (y = z)"  
*<proof>*

**lemma** (in abelian\_group) a\_r\_cancel [simp]:  
 "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>  
 (y  $\oplus$  x = z  $\oplus$  x) = (y = z)"  
*<proof>*

**lemma** (in abelian\_monoid) a\_assoc:  
 "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>  
 (x  $\oplus$  y)  $\oplus$  z = x  $\oplus$  (y  $\oplus$  z)"  
*<proof>*

**lemma** (in abelian\_monoid) l\_zero [simp]:  
 "x ∈ carrier G ==> 0  $\oplus$  x = x"  
*<proof>*

**lemma** (in abelian\_group) l\_neg:  
 "x ∈ carrier G ==>  $\ominus$  x  $\oplus$  x = 0"  
*<proof>*

**lemma** (in abelian\_monoid) a\_comm:  
 "[| x ∈ carrier G; y ∈ carrier G |] ==> x  $\oplus$  y = y  $\oplus$  x"  
*<proof>*

**lemma** (in abelian\_monoid) a\_lcomm:  
 "[| x ∈ carrier G; y ∈ carrier G; z ∈ carrier G |] ==>  
 x  $\oplus$  (y  $\oplus$  z) = y  $\oplus$  (x  $\oplus$  z)"  
*<proof>*

**lemma** (in abelian\_monoid) r\_zero [simp]:  
 "x ∈ carrier G ==> x  $\oplus$  0 = x"  
*<proof>*

```

lemma (in abelian_group) r_neg:
  "x ∈ carrier G ==> x ⊕ (⊖ x) = 0"
  ⟨proof⟩

lemma (in abelian_group) minus_zero [simp]:
  "⊖ 0 = 0"
  ⟨proof⟩

lemma (in abelian_group) minus_minus [simp]:
  "x ∈ carrier G ==> ⊖ (⊖ x) = x"
  ⟨proof⟩

lemma (in abelian_group) a_inv_inj:
  "inj_on (a_inv G) (carrier G)"
  ⟨proof⟩

lemma (in abelian_group) minus_add:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> ⊖ (x ⊕ y) = ⊖ x ⊕ ⊖ y"
  ⟨proof⟩

lemma (in abelian_group) minus_equality:
  "[| x ∈ carrier G; y ∈ carrier G; y ⊕ x = 0 |] ==> ⊖ x = y"
  ⟨proof⟩

lemma (in abelian_monoid) minus_unique:
  "[| x ∈ carrier G; y ∈ carrier G; y' ∈ carrier G;
    y ⊕ x = 0; x ⊕ y' = 0 |] ==> y = y'"
  ⟨proof⟩

lemmas (in abelian_monoid) a_ac = a_assoc a_comm a_lcomm

Derive an abelian_group from a comm_group

lemma comm_group_abelian_groupI:
  fixes G (structure)
  assumes cg: "comm_group (|carrier = carrier G, mult = add G, one = zero
G|)"
  shows "abelian_group G"
  ⟨proof⟩

```

### 8.3 Sums over Finite Sets

This definition makes it easy to lift lemmas from `finprod`.

#### definition

```

finsum :: "[('b, 'm) ring_scheme, 'a => 'b, 'a set] => 'b" where
  "finsum G f A = finprod (| carrier = carrier G, mult = add G, one =
zero G |) f A"

```

syntax

```

    "_finsum" :: "index => idt => 'a set => 'b => 'b"
      ("( $3 \oplus \_ \_ \_ \_$ )" [1000, 0, 51, 10] 10)
syntax (xsymbols)
    "_finsum" :: "index => idt => 'a set => 'b => 'b"
      ("( $3 \oplus \_ \_ \in \_ \_$ )" [1000, 0, 51, 10] 10)
syntax (HTML output)
    "_finsum" :: "index => idt => 'a set => 'b => 'b"
      ("( $3 \oplus \_ \_ \in \_ \_$ )" [1000, 0, 51, 10] 10)
translations
    " $\oplus \imath i:A. b$ " == "CONST finsum  $\diamond \imath$  (%i. b) A"
    — Beware of argument permutation!

context abelian_monoid begin

lemma finsum_empty [simp]:
  "finsum G f {} = 0"
  <proof>

lemma finsum_insert [simp]:
  "[| finite F; a  $\notin$  F; f  $\in$  F -> carrier G; f a  $\in$  carrier G |]
  ==> finsum G f (insert a F) = f a  $\oplus$  finsum G f F"
  <proof>

lemma finsum_zero [simp]:
  "finite A ==> ( $\oplus i \in A. 0$ ) = 0"
  <proof>

lemma finsum_closed [simp]:
  fixes A
  assumes fin: "finite A" and f: "f  $\in$  A -> carrier G"
  shows "finsum G f A  $\in$  carrier G"
  <proof>

lemma finsum_Un_Int:
  "[| finite A; finite B; g  $\in$  A -> carrier G; g  $\in$  B -> carrier G |] ==>
    finsum G g (A Un B)  $\oplus$  finsum G g (A Int B) =
    finsum G g A  $\oplus$  finsum G g B"
  <proof>

lemma finsum_Un_disjoint:
  "[| finite A; finite B; A Int B = {};
    g  $\in$  A -> carrier G; g  $\in$  B -> carrier G |]
  ==> finsum G g (A Un B) = finsum G g A  $\oplus$  finsum G g B"
  <proof>

lemma finsum_addf:
  "[| finite A; f  $\in$  A -> carrier G; g  $\in$  A -> carrier G |] ==>

```



```

    finsum G (%x. f x  $\oplus$  g x) A = (finsum G f A  $\oplus$  finsum G g A)"
  <proof>

```

```

lemma finsum_cong':
  "[| A = B; g : B -> carrier G;
    !!i. i : B ==> f i = g i |] ==> finsum G f A = finsum G g B"
  <proof>

```

```

lemma finsum_0 [simp]:
  "f : {0::nat} -> carrier G ==> finsum G f {...0} = f 0"
  <proof>

```

```

lemma finsum_Suc [simp]:
  "f : {...Suc n} -> carrier G ==>
    finsum G f {...Suc n} = (f (Suc n)  $\oplus$  finsum G f {...n})"
  <proof>

```

```

lemma finsum_Suc2:
  "f : {...Suc n} -> carrier G ==>
    finsum G f {...Suc n} = (finsum G (%i. f (Suc i)) {...n}  $\oplus$  f 0)"
  <proof>

```

```

lemma finsum_add [simp]:
  "[| f : {...n} -> carrier G; g : {...n} -> carrier G |] ==>
    finsum G (%i. f i  $\oplus$  g i) {...n::nat} =
    finsum G f {...n}  $\oplus$  finsum G g {...n}"
  <proof>

```

```

lemma finsum_cong:
  "[| A = B; f : B -> carrier G;
    !!i. i : B ==> f i = g i |] ==> finsum G f A = finsum G g B"
  <proof>

```

Usually, if this rule causes a failed congruence proof error, the reason is that the premise  $g \in B \rightarrow \text{carrier } G$  cannot be shown. Adding `Pi_def` to the simpset is often useful.

```

lemma finsum_reindex:
  assumes fin: "finite A"
  shows "f : (h ' A)  $\rightarrow$  carrier G  $\implies$ 
    inj_on h A ==> finsum G f (h ' A) = finsum G (%x. f (h x)) A"
  <proof>

```

```

lemma finsum_singleton:
  assumes i_in_A: "i  $\in$  A" and fin_A: "finite A" and f_Pi: "f  $\in$  A  $\rightarrow$ 
    carrier G"

```

```

shows "( $\bigoplus_{j \in A} \text{if } i = j \text{ then } f \ j \text{ else } 0$ ) = f i"
  <proof>

end



### 8.4 Rings: Basic Definitions



locale ring = abelian_group R + monoid R for R (structure) +
  assumes l_distr: "[| x ∈ carrier R; y ∈ carrier R; z ∈ carrier R |]
    ==> (x ⊕ y) ⊗ z = x ⊗ z ⊕ y ⊗ z"
  and r_distr: "[| x ∈ carrier R; y ∈ carrier R; z ∈ carrier R |]
    ==> z ⊗ (x ⊕ y) = z ⊗ x ⊕ z ⊗ y"

locale cring = ring + comm_monoid R

locale "domain" = cring +
  assumes one_not_zero [simp]: "1 ~= 0"
  and integral: "[| a ⊗ b = 0; a ∈ carrier R; b ∈ carrier R |] ==>
    a = 0 | b = 0"

locale field = "domain" +
  assumes field_Units: "Units R = carrier R - {0}"

```

### 8.5 Rings

```

lemma ringI:
  fixes R (structure)
  assumes abelian_group: "abelian_group R"
    and monoid: "monoid R"
    and l_distr: "!!x y z. [| x ∈ carrier R; y ∈ carrier R; z ∈ carrier
R |]
    ==> (x ⊕ y) ⊗ z = x ⊗ z ⊕ y ⊗ z"
    and r_distr: "!!x y z. [| x ∈ carrier R; y ∈ carrier R; z ∈ carrier
R |]
    ==> z ⊗ (x ⊕ y) = z ⊗ x ⊕ z ⊗ y"
  shows "ring R"
  <proof>

lemma (in ring) is_abelian_group:
  "abelian_group R"
  <proof>

lemma (in ring) is_monoid:
  "monoid R"
  <proof>

lemma (in ring) is_ring:
  "ring R"
  <proof>

```

```
lemmas ring_record_simps = monoid_record_simps ring_simps
```

```
lemma cringI:
  fixes R (structure)
  assumes abelian_group: "abelian_group R"
    and comm_monoid: "comm_monoid R"
    and l_distr: "!!x y z. [| x ∈ carrier R; y ∈ carrier R; z ∈ carrier
R |]
    ==> (x ⊕ y) ⊗ z = x ⊗ z ⊕ y ⊗ z"
  shows "cring R"
  <proof>
```

```
lemma (in cring) is_cring:
  "cring R" <proof>
```

### 8.5.1 Normaliser for Rings

```
lemma (in abelian_group) r_neg2:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> x ⊕ (⊖ x ⊕ y) = y"
  <proof>
```

```
lemma (in abelian_group) r_neg1:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> ⊖ x ⊕ (x ⊕ y) = y"
  <proof>
```

The following proofs are from Jacobson, Basic Algebra I, pp. 88–89

```
lemma (in ring) l_null [simp]:
  "x ∈ carrier R ==> 0 ⊗ x = 0"
  <proof>
```

```
lemma (in ring) r_null [simp]:
  "x ∈ carrier R ==> x ⊗ 0 = 0"
  <proof>
```

```
lemma (in ring) l_minus:
  "[| x ∈ carrier R; y ∈ carrier R |] ==> ⊖ x ⊗ y = ⊖ (x ⊗ y)"
  <proof>
```

```
lemma (in ring) r_minus:
  "[| x ∈ carrier R; y ∈ carrier R |] ==> x ⊗ ⊖ y = ⊖ (x ⊗ y)"
  <proof>
```

```
lemma (in abelian_group) minus_eq:
  "[| x ∈ carrier G; y ∈ carrier G |] ==> x ⊖ y = x ⊕ ⊖ y"
  <proof>
```

Setup algebra method: compute distributive normal form in locale contexts

$\langle ML \rangle$

```

lemmas (in ring) ring_simplrules
  [algebra ring "zero R" "add R" "a_inv R" "a_minus R" "one R" "mult R"]
=
  a_closed zero_closed a_inv_closed minus_closed m_closed one_closed
  a_assoc l_zero l_neg a_comm m_assoc l_one l_distr minus_eq
  r_zero r_neg r_neg2 r_neg1 minus_add minus_minus minus_zero
  a_lcomm r_distr l_null r_null l_minus r_minus

lemmas (in cring)
  [algebra del: ring "zero R" "add R" "a_inv R" "a_minus R" "one R" "mult
R"] =
  -

lemmas (in cring) cring_simplrules
  [algebra add: cring "zero R" "add R" "a_inv R" "a_minus R" "one R" "mult
R"] =
  a_closed zero_closed a_inv_closed minus_closed m_closed one_closed
  a_assoc l_zero l_neg a_comm m_assoc l_one l_distr m_comm minus_eq
  r_zero r_neg r_neg2 r_neg1 minus_add minus_minus minus_zero
  a_lcomm m_lcomm r_distr l_null r_null l_minus r_minus

```

```

lemma (in cring) nat_pow_zero:
  "(n::nat) ~= 0 ==> 0 (^) n = 0"
   $\langle proof \rangle$ 

```

```

lemma (in ring) one_zeroD:
  assumes onezero: "1 = 0"
  shows "carrier R = {0}"
   $\langle proof \rangle$ 

```

```

lemma (in ring) one_zeroI:
  assumes carrzero: "carrier R = {0}"
  shows "1 = 0"
   $\langle proof \rangle$ 

```

```

lemma (in ring) carrier_one_zero:
  shows "(carrier R = {0}) = (1 = 0)"
   $\langle proof \rangle$ 

```

```

lemma (in ring) carrier_one_not_zero:
  shows "(carrier R  $\neq$  {0}) = (1  $\neq$  0)"
   $\langle proof \rangle$ 

```

Two examples for use of method algebra

```

lemma
  fixes R (structure) and S (structure)

```

```

    assumes "ring R" "cring S"
    assumes RS: "a ∈ carrier R" "b ∈ carrier R" "c ∈ carrier S" "d ∈ carrier
S"
    shows "a ⊕ ⊖ (a ⊕ ⊖ b) = b & c ⊗S d = d ⊗S c"
  <proof>
  <ML>
  <proof>

```

```

lemma
  fixes R (structure)
  assumes "ring R"
  assumes R: "a ∈ carrier R" "b ∈ carrier R"
  shows "a ⊖ (a ⊖ b) = b"
  <proof>

```

### 8.5.2 Sums over Finite Sets

```

lemma (in ring) finsum_ldistr:
  "[| finite A; a ∈ carrier R; f ∈ A -> carrier R |] ==>
  finsum R f A ⊗ a = finsum R (%i. f i ⊗ a) A"
  <proof>

```

```

lemma (in ring) finsum_rdistr:
  "[| finite A; a ∈ carrier R; f ∈ A -> carrier R |] ==>
  a ⊗ finsum R f A = finsum R (%i. a ⊗ f i) A"
  <proof>

```

## 8.6 Integral Domains

```

lemma (in "domain") zero_not_one [simp]:
  "0 ~= 1"
  <proof>

```

```

lemma (in "domain") integral_iff:
  "[| a ∈ carrier R; b ∈ carrier R |] ==> (a ⊗ b = 0) = (a = 0 | b =
0)"
  <proof>

```

```

lemma (in "domain") m_lcancel:
  assumes prem: "a ~= 0"
  and R: "a ∈ carrier R" "b ∈ carrier R" "c ∈ carrier R"
  shows "(a ⊗ b = a ⊗ c) = (b = c)"
  <proof>

```

```

lemma (in "domain") m_rcancel:
  assumes prem: "a ~= 0"
  and R: "a ∈ carrier R" "b ∈ carrier R" "c ∈ carrier R"
  shows conc: "(b ⊗ a = c ⊗ a) = (b = c)"
  <proof>

```

## 8.7 Fields

Field would not need to be derived from domain, the properties for domain follow from the assumptions of field

```
lemma (in cring) cring_fieldI:
  assumes field_Units: "Units R = carrier R - {0}"
  shows "field R"
<proof>
```

Another variant to show that something is a field

```
lemma (in cring) cring_fieldI2:
  assumes notzero: "0 ≠ 1"
  and invex: "∧a. [a ∈ carrier R; a ≠ 0] ⇒ ∃b∈carrier R. a ⊗ b = 1"
  shows "field R"
<proof>
```

## 8.8 Morphisms

definition

```
ring_hom :: "[('a, 'm) ring_scheme, ('b, 'n) ring_scheme] => ('a => 'b) set"
where "ring_hom R S =
  {h. h ∈ carrier R -> carrier S &
    (ALL x y. x ∈ carrier R & y ∈ carrier R -->
      h (x ⊗R y) = h x ⊗S h y & h (x ⊕R y) = h x ⊕S h y) &
    h 1R = 1S}"
```

```
lemma ring_hom_memI:
  fixes R (structure) and S (structure)
  assumes hom_closed: "!!x. x ∈ carrier R ==> h x ∈ carrier S"
  and hom_mult: "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==>
    h (x ⊗ y) = h x ⊗S h y"
  and hom_add: "!!x y. [| x ∈ carrier R; y ∈ carrier R |] ==>
    h (x ⊕ y) = h x ⊕S h y"
  and hom_one: "h 1 = 1S"
  shows "h ∈ ring_hom R S"
<proof>
```

```
lemma ring_hom_closed:
  "[| h ∈ ring_hom R S; x ∈ carrier R |] ==> h x ∈ carrier S"
<proof>
```

```
lemma ring_hom_mult:
  fixes R (structure) and S (structure)
  shows
    "[| h ∈ ring_hom R S; x ∈ carrier R; y ∈ carrier R |] ==>
      h (x ⊗ y) = h x ⊗S h y"
<proof>
```

```

lemma ring_hom_add:
  fixes R (structure) and S (structure)
  shows
    "[| h ∈ ring_hom R S; x ∈ carrier R; y ∈ carrier R |] ==>
     h (x ⊕ y) = h x ⊕S h y"
    ⟨proof⟩

lemma ring_hom_one:
  fixes R (structure) and S (structure)
  shows "h ∈ ring_hom R S ==> h 1 = 1S"
  ⟨proof⟩

locale ring_hom_cring = R: cring R + S: cring S
  for R (structure) and S (structure) +
  fixes h
  assumes homh [simp, intro]: "h ∈ ring_hom R S"
  notes hom_closed [simp, intro] = ring_hom_closed [OF homh]
    and hom_mult [simp] = ring_hom_mult [OF homh]
    and hom_add [simp] = ring_hom_add [OF homh]
    and hom_one [simp] = ring_hom_one [OF homh]

lemma (in ring_hom_cring) hom_zero [simp]:
  "h 0 = 0S"
  ⟨proof⟩

lemma (in ring_hom_cring) hom_a_inv [simp]:
  "x ∈ carrier R ==> h (⊖ x) = ⊖S h x"
  ⟨proof⟩

lemma (in ring_hom_cring) hom_finsum [simp]:
  "[| finite A; f ∈ A -> carrier R |] ==>
   h (finsum R f A) = finsum S (h o f) A"
  ⟨proof⟩

lemma (in ring_hom_cring) hom_finprod:
  "[| finite A; f ∈ A -> carrier R |] ==>
   h (finprod R f A) = finprod S (h o f) A"
  ⟨proof⟩

declare ring_hom_cring.hom_finprod [simp]

lemma id_ring_hom [simp]:
  "id ∈ ring_hom R R"
  ⟨proof⟩

end

```

```
theory AbelCoset
imports Coset Ring
begin
```

## 8.9 More Lifting from Groups to Abelian Groups

### 8.9.1 Definitions

Hiding  $\langle + \rangle$  from Sum\_Type until I come up with better syntax here

```
no_notation Plus (infixr "<+>" 65)
```

**definition**

```
a_r_coset      :: "[_, 'a set, 'a]  $\Rightarrow$  'a set"      (infixl "<+>" 60)
  where "a_r_coset G = r_coset ( $\lfloor$ carrier = carrier G, mult = add G, one
= zero G $\rfloor$ )"
```

**definition**

```
a_l_coset      :: "[_, 'a, 'a set]  $\Rightarrow$  'a set"      (infixl "<+>" 60)
  where "a_l_coset G = l_coset ( $\lfloor$ carrier = carrier G, mult = add G, one
= zero G $\rfloor$ )"
```

**definition**

```
A_RCOSSETS    :: "[_, 'a set]  $\Rightarrow$  ('a set)set"      ("a'_rcosets" [81] 80)
  where "A_RCOSSETS G H = RCOSSETS ( $\lfloor$ carrier = carrier G, mult = add G,
one = zero G $\rfloor$ ) H"
```

**definition**

```
set_add       :: "[_, 'a set, 'a set]  $\Rightarrow$  'a set" (infixl "<+>" 60)
  where "set_add G = set_mult ( $\lfloor$ carrier = carrier G, mult = add G, one
= zero G $\rfloor$ )"
```

**definition**

```
A_SET_INV     :: "[_, 'a set]  $\Rightarrow$  'a set"      ("a'_set'_inv" [81] 80)
  where "A_SET_INV G H = SET_INV ( $\lfloor$ carrier = carrier G, mult = add G,
one = zero G $\rfloor$ ) H"
```

**definition**

```
a_r_congruent :: "[('a,'b) ring_scheme, 'a set]  $\Rightarrow$  ('a*'a)set" ("racong"
_)
  where "a_r_congruent G = r_congruent ( $\lfloor$ carrier = carrier G, mult = add
G, one = zero G $\rfloor$ )"
```

**definition**

```
A_FactGroup   :: "[('a,'b) ring_scheme, 'a set]  $\Rightarrow$  ('a set) monoid" (in-
fixl "A'_Mod" 65)
  — Actually defined for groups rather than monoids
  where "A_FactGroup G H = FactGroup ( $\lfloor$ carrier = carrier G, mult = add
G, one = zero G $\rfloor$ ) H"
```



**definition**

`a_kernel :: "('a, 'm) ring_scheme  $\Rightarrow$  ('b, 'n) ring_scheme  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  'a set"`

— the kernel of a homomorphism (additive)

`where "a_kernel G H h =`

`kernel ( $\lambda$ carrier = carrier G, mult = add G, one = zero G)`

`( $\lambda$ carrier = carrier H, mult = add H, one = zero H) h"`

`locale abelian_group_hom = G: abelian_group G + H: abelian_group H`

`for G (structure) and H (structure) +`

`fixes h`

`assumes a_group_hom: "group_hom ( $\lambda$  carrier = carrier G, mult = add G, one = zero G  $\lambda$ )`

`( $\lambda$  carrier = carrier H, mult = add H,`

`one = zero H  $\lambda$ ) h"`

`lemmas a_r_coset_defs =`

`a_r_coset_def r_coset_def`

`lemma a_r_coset_def':`

`fixes G (structure)`

`shows "H  $\rightarrow$  a  $\equiv \bigcup_{h \in H}. \{h \oplus a\}"$`

*<proof>*

`lemmas a_l_coset_defs =`

`a_l_coset_def l_coset_def`

`lemma a_l_coset_def':`

`fixes G (structure)`

`shows "a  $\leftarrow$  H  $\equiv \bigcup_{h \in H}. \{a \oplus h\}"$`

*<proof>*

`lemmas A_RCOSETS_defs =`

`A_RCOSETS_def RCOSETS_def`

`lemma A_RCOSETS_def':`

`fixes G (structure)`

`shows "a_rcosets H  $\equiv \bigcup_{a \in \text{carrier } G}. \{H \rightarrow a\}"$`

*<proof>*

`lemmas set_add_defs =`

`set_add_def set_mult_def`

`lemma set_add_def':`

`fixes G (structure)`

`shows "H  $\leftrightarrow$  K  $\equiv \bigcup_{h \in H}. \bigcup_{k \in K}. \{h \oplus k\}"$`

*<proof>*

```

lemmas A_SET_INV_defs =
  A_SET_INV_def SET_INV_def

```

```

lemma A_SET_INV_def':
  fixes G (structure)
  shows "a_set_inv H  $\equiv \bigcup_{h \in H}. \{\ominus h\}$ "
  <proof>

```

### 8.9.2 Cosets

```

lemma (in abelian_group) a_coset_add_assoc:
  "[| M  $\subseteq$  carrier G; g  $\in$  carrier G; h  $\in$  carrier G |]
  ==> (M +> g) +> h = M +> (g  $\oplus$  h)"
  <proof>

```

```

lemma (in abelian_group) a_coset_add_zero [simp]:
  "M  $\subseteq$  carrier G ==> M +> 0 = M"
  <proof>

```

```

lemma (in abelian_group) a_coset_add_inv1:
  "[| M +> (x  $\oplus$  ( $\ominus$  y)) = M; x  $\in$  carrier G; y  $\in$  carrier G;
  M  $\subseteq$  carrier G |] ==> M +> x = M +> y"
  <proof>

```

```

lemma (in abelian_group) a_coset_add_inv2:
  "[| M +> x = M +> y; x  $\in$  carrier G; y  $\in$  carrier G; M  $\subseteq$  carrier
  G |]
  ==> M +> (x  $\oplus$  ( $\ominus$  y)) = M"
  <proof>

```

```

lemma (in abelian_group) a_coset_join1:
  "[| H +> x = H; x  $\in$  carrier G; subgroup H (|carrier = carrier G,
  mult = add G, one = zero G|) |] ==> x  $\in$  H"
  <proof>

```

```

lemma (in abelian_group) a_solve_equation:
  "[subgroup H (|carrier = carrier G, mult = add G, one = zero G|);
  x  $\in$  H; y  $\in$  H]  $\implies \exists h \in H. y = h \oplus x$ "
  <proof>

```

```

lemma (in abelian_group) a_repr_independence:
  "[y  $\in$  H +> x; x  $\in$  carrier G; subgroup H (|carrier = carrier G, mult
  = add G, one = zero G|) ]  $\implies H +> x = H +> y$ "
  <proof>

```

```

lemma (in abelian_group) a_coset_join2:
  "[x  $\in$  carrier G; subgroup H (|carrier = carrier G, mult = add G,
  one = zero G|); x  $\in$  H]  $\implies H +> x = H$ "
  <proof>

```

```

lemma (in abelian_monoid) a_r_coset_subset_G:
  "[| H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> H +> x  $\subseteq$  carrier G"
<proof>

```

```

lemma (in abelian_group) a_rcosI:
  "[| h  $\in$  H; H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> h  $\oplus$  x  $\in$  H +> x"
<proof>

```

```

lemma (in abelian_group) a_rcosetsI:
  "[| H  $\subseteq$  carrier G; x  $\in$  carrier G |] ==> H +> x  $\in$  a_rcosets H"
<proof>

```

Really needed?

```

lemma (in abelian_group) a_transpose_inv:
  "[| x  $\oplus$  y = z; x  $\in$  carrier G; y  $\in$  carrier G; z  $\in$  carrier G |]
  ==> ( $\ominus$  x)  $\oplus$  z = y"
<proof>

```

### 8.9.3 Subgroups

```

locale additive_subgroup =
  fixes H and G (structure)
  assumes a_subgroup: "subgroup H ( $\langle$ carrier = carrier G, mult = add G,
one = zero G $\rangle$ )"

```

```

lemma (in additive_subgroup) is_additive_subgroup:
  shows "additive_subgroup H G"
<proof>

```

```

lemma additive_subgroupI:
  fixes G (structure)
  assumes a_subgroup: "subgroup H ( $\langle$ carrier = carrier G, mult = add G,
one = zero G $\rangle$ )"
  shows "additive_subgroup H G"
<proof>

```

```

lemma (in additive_subgroup) a_subset:
  "H  $\subseteq$  carrier G"
<proof>

```

```

lemma (in additive_subgroup) a_closed [intro, simp]:
  "[| x  $\in$  H; y  $\in$  H |] ==> x  $\oplus$  y  $\in$  H"
<proof>

```

```

lemma (in additive_subgroup) zero_closed [simp]:
  "0  $\in$  H"
<proof>

```

```

lemma (in additive_subgroup) a_inv_closed [intro,simp]:
  "x ∈ H ⇒ ⊖ x ∈ H"
⟨proof⟩

```

#### 8.9.4 Additive subgroups are normal

Every subgroup of an `abelian_group` is normal

```

locale abelian_subgroup = additive_subgroup + abelian_group G +
  assumes a_normal: "normal H (carrier = carrier G, mult = add G, one
    = zero G)"

```

```

lemma (in abelian_subgroup) is_abelian_subgroup:
  shows "abelian_subgroup H G"
⟨proof⟩

```

```

lemma abelian_subgroupI:
  assumes a_normal: "normal H (carrier = carrier G, mult = add G, one
    = zero G)"
    and a_comm: "!!x y. [| x ∈ carrier G; y ∈ carrier G |] ==> x ⊕G
    y = y ⊕G x"
  shows "abelian_subgroup H G"
⟨proof⟩

```

```

lemma abelian_subgroupI2:
  fixes G (structure)
  assumes a_comm_group: "comm_group (carrier = carrier G, mult = add
    G, one = zero G)"
    and a_subgroup: "subgroup H (carrier = carrier G, mult = add G,
    one = zero G)"
  shows "abelian_subgroup H G"
⟨proof⟩

```

```

lemma abelian_subgroupI3:
  fixes G (structure)
  assumes asg: "additive_subgroup H G"
    and ag: "abelian_group G"
  shows "abelian_subgroup H G"
⟨proof⟩

```

```

lemma (in abelian_subgroup) a_coset_eq:
  "(∀x ∈ carrier G. H +> x = x <+ H)"
⟨proof⟩

```

```

lemma (in abelian_subgroup) a_inv_op_closed1:
  shows "[x ∈ carrier G; h ∈ H] ⇒ (⊖ x) ⊕ h ⊕ x ∈ H"
⟨proof⟩

```

```

lemma (in abelian_subgroup) a_inv_op_closed2:
  shows "[x ∈ carrier G; h ∈ H] ⇒ x ⊕ h ⊕ (⊖ x) ∈ H"

```

*<proof>*

Alternative characterization of normal subgroups

```
lemma (in abelian_group) a_normal_inv_iff:
  "(N <| carrier = carrier G, mult = add G, one = zero G |) =
   (subgroup N (|carrier = carrier G, mult = add G, one = zero G|) &
    (∀ x ∈ carrier G. ∀ h ∈ N. x ⊕ h ⊕ (⊖ x) ∈ N))"
  (is "_ = ?rhs")
```

*<proof>*

```
lemma (in abelian_group) a_lcos_m_assoc:
  "[| M ⊆ carrier G; g ∈ carrier G; h ∈ carrier G |]
   ==> g <+ (h <+ M) = (g ⊕ h) <+ M"
```

*<proof>*

```
lemma (in abelian_group) a_lcos_mult_one:
  "M ⊆ carrier G ==> 0 <+ M = M"
```

*<proof>*

```
lemma (in abelian_group) a_lcoset_subset_G:
  "[| H ⊆ carrier G; x ∈ carrier G |] ==> x <+ H ⊆ carrier G"
```

*<proof>*

```
lemma (in abelian_group) a_lcoset_swap:
  "[| y ∈ x <+ H; x ∈ carrier G; subgroup H (|carrier = carrier G, mult
= add G, one = zero G|) |] ==> x ∈ y <+ H"
```

*<proof>*

```
lemma (in abelian_group) a_lcoset_carrier:
  "[| y ∈ x <+ H; x ∈ carrier G; subgroup H (|carrier = carrier G,
mult = add G, one = zero G|) |] ==> y ∈ carrier G"
```

*<proof>*

```
lemma (in abelian_group) a_lrepr_imp_subset:
  assumes y: "y ∈ x <+ H" and x: "x ∈ carrier G" and sb: "subgroup H
(|carrier = carrier G, mult = add G, one = zero G|)"
  shows "y <+ H ⊆ x <+ H"
```

*<proof>*

```
lemma (in abelian_group) a_lrepr_independence:
  assumes y: "y ∈ x <+ H" and x: "x ∈ carrier G" and sb: "subgroup H
(|carrier = carrier G, mult = add G, one = zero G|)"
  shows "x <+ H = y <+ H"
```

*<proof>*

```
lemma (in abelian_group) setadd_subset_G:
  "[| H ⊆ carrier G; K ⊆ carrier G |] ==> H <+> K ⊆ carrier G"
```

*<proof>*

**lemma** (in abelian\_group) subgroup\_add\_id: "subgroup H ( $\text{carrier} = \text{carrier } G, \text{mult} = \text{add } G, \text{one} = \text{zero } G$ )  $\implies H <+> H = H$ "  
*<proof>*

**lemma** (in abelian\_subgroup) a\_rcos\_inv:  
 assumes x: "x  $\in$  carrier G"  
 shows "a\_set\_inv (H +> x) = H +> ( $\ominus$  x)"  
*<proof>*

**lemma** (in abelian\_group) a\_setmult\_rcos\_assoc:  
 "[H  $\subseteq$  carrier G; K  $\subseteq$  carrier G; x  $\in$  carrier G]  
 $\implies H <+> (K +> x) = (H <+> K) +> x$ "  
*<proof>*

**lemma** (in abelian\_group) a\_rcos\_assoc\_lcos:  
 "[H  $\subseteq$  carrier G; K  $\subseteq$  carrier G; x  $\in$  carrier G]  
 $\implies (H +> x) <+> K = H <+> (x <+ K)$ "  
*<proof>*

**lemma** (in abelian\_subgroup) a\_rcos\_sum:  
 "[x  $\in$  carrier G; y  $\in$  carrier G]  
 $\implies (H +> x) <+> (H +> y) = H +> (x \oplus y)$ "  
*<proof>*

**lemma** (in abelian\_subgroup) rcosets\_add\_eq:  
 "M  $\in$  a\_rcosets H  $\implies H <+> M = M$ "  
 — generalizes subgroup\_mult\_id  
*<proof>*

### 8.9.5 Congruence Relation

**lemma** (in abelian\_subgroup) a\_equiv\_rcong:  
 shows "equiv (carrier G) (racong H)"  
*<proof>*

**lemma** (in abelian\_subgroup) a\_l\_coset\_eq\_rcong:  
 assumes a: "a  $\in$  carrier G"  
 shows "a <+ H = racong H `` {a}"  
*<proof>*

**lemma** (in abelian\_subgroup) a\_rcos\_equation:  
 shows  
 "[ha  $\oplus$  a = h  $\oplus$  b; a  $\in$  carrier G; b  $\in$  carrier G;  
 h  $\in$  H; ha  $\in$  H; hb  $\in$  H]  
 $\implies hb \oplus a \in (\bigcup_{h \in H}. \{h \oplus b\})$ "  
*<proof>*

```

lemma (in abelian_subgroup) a_rcos_disjoint:
  shows "[a ∈ a_rcosets H; b ∈ a_rcosets H; a ≠ b] ⇒ a ∩ b = {}"
  <proof>

lemma (in abelian_subgroup) a_rcos_self:
  shows "x ∈ carrier G ⇒ x ∈ H +> x"
  <proof>

lemma (in abelian_subgroup) a_rcosets_part_G:
  shows "⋃ (a_rcosets H) = carrier G"
  <proof>

lemma (in abelian_subgroup) a_cosets_finite:
  "[c ∈ a_rcosets H; H ⊆ carrier G; finite (carrier G)] ⇒ finite
  c"
  <proof>

lemma (in abelian_group) a_card_cosets_equal:
  "[c ∈ a_rcosets H; H ⊆ carrier G; finite(carrier G)]
  ⇒ card c = card H"
  <proof>

lemma (in abelian_group) rcosets_subset_PowG:
  "additive_subgroup H G ⇒ a_rcosets H ⊆ Pow(carrier G)"
  <proof>

theorem (in abelian_group) a_lagrange:
  "[finite(carrier G); additive_subgroup H G]
  ⇒ card(a_rcosets H) * card(H) = order(G)"
  <proof>

```

### 8.9.6 Factorization

```

lemmas A_FactGroup_defs = A_FactGroup_def FactGroup_def

```

```

lemma A_FactGroup_def':
  fixes G (structure)
  shows "G A_Mod H ≡ (carrier = a_rcosets_G H, mult = set_add G, one =
  H)"
  <proof>

```

```

lemma (in abelian_subgroup) a_setmult_closed:
  "[K1 ∈ a_rcosets H; K2 ∈ a_rcosets H] ⇒ K1 <+> K2 ∈ a_rcosets H"
  <proof>

```

```

lemma (in abelian_subgroup) a_setinv_closed:
  "K ∈ a_rcosets H ⇒ a_set_inv K ∈ a_rcosets H"
  <proof>

```

```

lemma (in abelian_subgroup) a_rcosets_assoc:
  "[[M1 ∈ a_rcosets H; M2 ∈ a_rcosets H; M3 ∈ a_rcosets H]]
   ⇒ M1 <+> M2 <+> M3 = M1 <+> (M2 <+> M3)"
<proof>

```

```

lemma (in abelian_subgroup) a_subgroup_in_rcosets:
  "H ∈ a_rcosets H"
<proof>

```

```

lemma (in abelian_subgroup) a_rcosets_inv_mult_group_eq:
  "M ∈ a_rcosets H ⇒ a_set_inv M <+> M = H"
<proof>

```

```

theorem (in abelian_subgroup) a_factorgroup_is_group:
  "group (G A_Mod H)"
<proof>

```

Since the Factorization is based on an *abelian* subgroup, it results in a commutative group

```

theorem (in abelian_subgroup) a_factorgroup_is_comm_group:
  "comm_group (G A_Mod H)"
<proof>

```

```

lemma add_A_FactGroup [simp]: "X ⊗(G A_Mod H) X' = X <+>G X'"
<proof>

```

```

lemma (in abelian_subgroup) a_inv_FactGroup:
  "X ∈ carrier (G A_Mod H) ⇒ invG A_Mod H X = a_set_inv X"
<proof>

```

The coset map is a homomorphism from  $G$  to the quotient group  $G \text{ Mod } H$

```

lemma (in abelian_subgroup) a_r_coset_hom_A_Mod:
  "(λa. H +> a) ∈ hom (|carrier = carrier G, mult = add G, one = zero G|)
  (G A_Mod H)"
<proof>

```

The isomorphism theorems have been omitted from lifting, at least for now

### 8.9.7 The First Isomorphism Theorem

The quotient by the kernel of a homomorphism is isomorphic to the range of that homomorphism.

```

lemmas a_kernel_defs =
  a_kernel_def kernel_def

```

```

lemma a_kernel_def':
  "a_kernel R S h = {x ∈ carrier R. h x = 0S}"
<proof>

```



### 8.9.8 Homomorphisms

```
lemma abelian_group_homI:
  assumes "abelian_group G"
  assumes "abelian_group H"
  assumes a_group_hom: "group_hom (| carrier = carrier G, mult = add
    G, one = zero G |)
    (| carrier = carrier H, mult = add H,
    one = zero H |) h"
  shows "abelian_group_hom G H h"
<proof>
```

```
lemma (in abelian_group_hom) is_abelian_group_hom:
  "abelian_group_hom G H h"
<proof>
```

```
lemma (in abelian_group_hom) hom_add [simp]:
  "[| x : carrier G; y : carrier G |]
   ==> h (x  $\oplus_G$  y) = h x  $\oplus_H$  h y"
<proof>
```

```
lemma (in abelian_group_hom) hom_closed [simp]:
  "x  $\in$  carrier G  $\implies$  h x  $\in$  carrier H"
<proof>
```

```
lemma (in abelian_group_hom) zero_closed [simp]:
  "h 0  $\in$  carrier H"
<proof>
```

```
lemma (in abelian_group_hom) hom_zero [simp]:
  "h 0 = 0H"
<proof>
```

```
lemma (in abelian_group_hom) a_inv_closed [simp]:
  "x  $\in$  carrier G  $\implies$  h ( $\ominus$ x)  $\in$  carrier H"
<proof>
```

```
lemma (in abelian_group_hom) hom_a_inv [simp]:
  "x  $\in$  carrier G  $\implies$  h ( $\ominus$ x) =  $\ominus_H$  (h x)"
<proof>
```

```
lemma (in abelian_group_hom) additive_subgroup_a_kernel:
  "additive_subgroup (a_kernel G H h) G"
<proof>
```

The kernel of a homomorphism is an abelian subgroup

```
lemma (in abelian_group_hom) abelian_subgroup_a_kernel:
  "abelian_subgroup (a_kernel G H h) G"
<proof>
```

```

lemma (in abelian_group_hom) A_FactGroup_nonempty:
  assumes X: "X ∈ carrier (G A_Mod a_kernel G H h)"
  shows "X ≠ {}"
<proof>

```

```

lemma (in abelian_group_hom) FactGroup_contents_mem:
  assumes X: "X ∈ carrier (G A_Mod (a_kernel G H h))"
  shows "contents (h'X) ∈ carrier H"
<proof>

```

```

lemma (in abelian_group_hom) A_FactGroup_hom:
  "(λX. contents (h'X)) ∈ hom (G A_Mod (a_kernel G H h))
  (|carrier = carrier H, mult = add H, one = zero H)"
<proof>

```

```

lemma (in abelian_group_hom) A_FactGroup_inj_on:
  "inj_on (λX. contents (h ' X)) (carrier (G A_Mod a_kernel G H h))"
<proof>

```

If the homomorphism  $h$  is onto  $H$ , then so is the homomorphism from the quotient group

```

lemma (in abelian_group_hom) A_FactGroup_onto:
  assumes h: "h ' carrier G = carrier H"
  shows "(λX. contents (h ' X)) ' carrier (G A_Mod a_kernel G H h) =
  carrier H"
<proof>

```

If  $h$  is a homomorphism from  $G$  onto  $H$ , then the quotient group  $G \text{ Mod } \text{kernel } G \ H \ h$  is isomorphic to  $H$ .

```

theorem (in abelian_group_hom) A_FactGroup_iso:
  "h ' carrier G = carrier H
  ⇒ (λX. contents (h'X)) ∈ (G A_Mod (a_kernel G H h)) ≅
  (| carrier = carrier H, mult = add H, one = zero H |)"
<proof>

```

### 8.9.9 Cosets

Not everything from `CosetExt.thy` is lifted here.

```

lemma (in additive_subgroup) a_Hcarr [simp]:
  assumes hH: "h ∈ H"
  shows "h ∈ carrier G"
<proof>

```

```

lemma (in abelian_subgroup) a_elemrcos_carrier:
  assumes acarr: "a ∈ carrier G"
  and a': "a' ∈ H +> a"
  shows "a' ∈ carrier G"

```

*<proof>*

```
lemma (in abelian_subgroup) a_rcos_const:
  assumes hH: "h ∈ H"
  shows "H +> h = H"
<proof>
```

```
lemma (in abelian_subgroup) a_rcos_module_imp:
  assumes xcarr: "x ∈ carrier G"
  and x'cos: "x' ∈ H +> x"
  shows "(x' ⊕ ⊖x) ∈ H"
<proof>
```

```
lemma (in abelian_subgroup) a_rcos_module_rev:
  assumes "x ∈ carrier G" "x' ∈ carrier G"
  and "(x' ⊕ ⊖x) ∈ H"
  shows "x' ∈ H +> x"
<proof>
```

```
lemma (in abelian_subgroup) a_rcos_module:
  assumes "x ∈ carrier G" "x' ∈ carrier G"
  shows "(x' ∈ H +> x) = (x' ⊕ ⊖x ∈ H)"
<proof>
```

```
lemma (in abelian_subgroup) a_rcos_module_minus:
  assumes "ring G"
  assumes carr: "x ∈ carrier G" "x' ∈ carrier G"
  shows "(x' ∈ H +> x) = (x' ⊖ x ∈ H)"
<proof>
```

```
lemma (in abelian_subgroup) a_repr_independence':
  assumes y: "y ∈ H +> x"
  and xcarr: "x ∈ carrier G"
  shows "H +> x = H +> y"
<proof>
```

```
lemma (in abelian_subgroup) a_repr_independenceD:
  assumes ycarr: "y ∈ carrier G"
  and repr: "H +> x = H +> y"
  shows "y ∈ H +> x"
<proof>
```

```
lemma (in abelian_subgroup) a_rcosets_carrier:
  "X ∈ a_rcosets H ⇒ X ⊆ carrier G"
<proof>
```

### 8.9.10 Addition of Subgroups

```
lemma (in abelian_monoid) set_add_closed:
```

```

    assumes Acarr: "A  $\subseteq$  carrier G"
    and Bcarr: "B  $\subseteq$  carrier G"
    shows "A  $\langle + \rangle$  B  $\subseteq$  carrier G"
  <proof>

lemma (in abelian_group) add_additive_subgroups:
  assumes subH: "additive_subgroup H G"
  and subK: "additive_subgroup K G"
  shows "additive_subgroup (H  $\langle + \rangle$  K) G"
  <proof>

end

```

```

theory Ideal
imports Ring AbelCoset
begin

```

## 9 Ideals

### 9.1 Definitions

#### 9.1.1 General definition

```

locale ideal = additive_subgroup I R + ring R for I and R (structure) +
  assumes I_l_closed: "[a  $\in$  I; x  $\in$  carrier R]  $\implies$  x  $\otimes$  a  $\in$  I"
  and I_r_closed: "[a  $\in$  I; x  $\in$  carrier R]  $\implies$  a  $\otimes$  x  $\in$  I"

```

```

sublocale ideal  $\subseteq$  abelian_subgroup I R
  <proof>

```

```

lemma (in ideal) is_ideal:
  "ideal I R"
  <proof>

```

```

lemma idealI:
  fixes R (structure)
  assumes "ring R"
  assumes a_subgroup: "subgroup I ( $\langle$ carrier = carrier R, mult = add R,
one = zero R $\rangle$ )"
  and I_l_closed: " $\bigwedge$ a x. [a  $\in$  I; x  $\in$  carrier R]  $\implies$  x  $\otimes$  a  $\in$  I"
  and I_r_closed: " $\bigwedge$ a x. [a  $\in$  I; x  $\in$  carrier R]  $\implies$  a  $\otimes$  x  $\in$  I"
  shows "ideal I R"
  <proof>

```

#### 9.1.2 Ideals Generated by a Subset of carrier R

definition

```

genideal :: "('a, 'b) ring_scheme  $\Rightarrow$  'a set  $\Rightarrow$  'a set"  ("Idlz _" [80]
79)
  where "genideal R S = Inter {I. ideal I R  $\wedge$  S  $\subseteq$  I}"

```

### 9.1.3 Principal Ideals

```

locale principalideal = ideal +
  assumes generate: " $\exists i \in \text{carrier } R. I = \text{Idl } \{i\}$ "

```

```

lemma (in principalideal) is_principalideal:
  shows "principalideal I R"
<proof>

```

```

lemma principalidealI:
  fixes R (structure)
  assumes "ideal I R"
  assumes generate: " $\exists i \in \text{carrier } R. I = \text{Idl } \{i\}$ "
  shows "principalideal I R"
<proof>

```

### 9.1.4 Maximal Ideals

```

locale maximalideal = ideal +
  assumes I_notcarr: "carrier R  $\neq$  I"
    and I_maximal: " $\llbracket \text{ideal } J \text{ R}; I \subseteq J; J \subseteq \text{carrier } R \rrbracket \Longrightarrow J = I \vee J$ 
= carrier R"

```

```

lemma (in maximalideal) is_maximalideal:
  shows "maximalideal I R"
<proof>

```

```

lemma maximalidealI:
  fixes R
  assumes "ideal I R"
  assumes I_notcarr: "carrier R  $\neq$  I"
    and I_maximal: " $\bigwedge J. \llbracket \text{ideal } J \text{ R}; I \subseteq J; J \subseteq \text{carrier } R \rrbracket \Longrightarrow J = I$ 
 $\vee J = \text{carrier } R$ "
  shows "maximalideal I R"
<proof>

```

### 9.1.5 Prime Ideals

```

locale primeideal = ideal + cring +
  assumes I_notcarr: "carrier R  $\neq$  I"
    and I_prime: " $\llbracket a \in \text{carrier } R; b \in \text{carrier } R; a \otimes b \in I \rrbracket \Longrightarrow a \in$ 
I  $\vee b \in I$ "

```

```

lemma (in primeideal) is_primeideal:
  shows "primeideal I R"
<proof>

```

```

lemma primeidealI:
  fixes R (structure)
  assumes "ideal I R"
  assumes "cring R"
  assumes I_notcarr: "carrier R  $\neq$  I"
    and I_prime: " $\bigwedge a b. \llbracket a \in \text{carrier } R; b \in \text{carrier } R; a \otimes b \in I \rrbracket$ "
 $\implies a \in I \vee b \in I$ 
  shows "primeideal I R"
<proof>

```

```

lemma primeidealI2:
  fixes R (structure)
  assumes "additive_subgroup I R"
  assumes "cring R"
  assumes I_l_closed: " $\bigwedge a x. \llbracket a \in I; x \in \text{carrier } R \rrbracket \implies x \otimes a \in I$ "
    and I_r_closed: " $\bigwedge a x. \llbracket a \in I; x \in \text{carrier } R \rrbracket \implies a \otimes x \in I$ "
    and I_notcarr: "carrier R  $\neq$  I"
    and I_prime: " $\bigwedge a b. \llbracket a \in \text{carrier } R; b \in \text{carrier } R; a \otimes b \in I \rrbracket$ "
 $\implies a \in I \vee b \in I$ 
  shows "primeideal I R"
<proof>

```

## 9.2 Special Ideals

```

lemma (in ring) zeroideal:
  shows "ideal {0} R"
<proof>

```

```

lemma (in ring) oneideal:
  shows "ideal (carrier R) R"
<proof>

```

```

lemma (in "domain") zeroprimeideal:
  shows "primeideal {0} R"
<proof>

```

## 9.3 General Ideal Properties

```

lemma (in ideal) one_imp_carrier:
  assumes I_one_closed: "1  $\in$  I"
  shows "I = carrier R"
<proof>

```

```

lemma (in ideal) Icarr:
  assumes iI: "i  $\in$  I"
  shows "i  $\in$  carrier R"
<proof>

```

## 9.4 Intersection of Ideals

**Intersection of two ideals** The intersection of any two ideals is again an ideal in  $R$

```
lemma (in ring) i_intersect:
  assumes "ideal I R"
  assumes "ideal J R"
  shows "ideal (I  $\cap$  J) R"
<proof>
```

The intersection of any Number of Ideals is again an Ideal in  $R$

```
lemma (in ring) i_Intersect:
  assumes Sideals: " $\bigwedge I. I \in S \implies \text{ideal } I \text{ } R$ "
  and notempty: " $S \neq \{\}$ "
  shows "ideal (Inter S) R"
<proof>
```

## 9.5 Addition of Ideals

```
lemma (in ring) add_ideals:
  assumes idealI: "ideal I R"
  and idealJ: "ideal J R"
  shows "ideal (I  $\lt+>$  J) R"
<proof>
```

## 9.6 Ideals generated by a subset of carrier $R$

genideal generates an ideal

```
lemma (in ring) genideal_ideal:
  assumes Scarr: " $S \subseteq \text{carrier } R$ "
  shows "ideal (Idl S) R"
<proof>
```

```
lemma (in ring) genideal_self:
  assumes " $S \subseteq \text{carrier } R$ "
  shows " $S \subseteq \text{Idl } S$ "
<proof>
```

```
lemma (in ring) genideal_self':
  assumes carr: " $i \in \text{carrier } R$ "
  shows " $i \in \text{Idl } \{i\}$ "
<proof>
```

genideal generates the minimal ideal

```
lemma (in ring) genideal_minimal:
  assumes a: "ideal I R"
  and b: " $S \subseteq I$ "
  shows " $\text{Idl } S \subseteq I$ "
```

*<proof>*

Generated ideals and subsets

```
lemma (in ring) Idl_subset_ideal:
  assumes Ideal: "ideal I R"
    and Hcarr: "H  $\subseteq$  carrier R"
  shows "(Idl H  $\subseteq$  I) = (H  $\subseteq$  I)"
<proof>
```

```
lemma (in ring) subset_Idl_subset:
  assumes Icarr: "I  $\subseteq$  carrier R"
    and HI: "H  $\subseteq$  I"
  shows "Idl H  $\subseteq$  Idl I"
<proof>
```

```
lemma (in ring) Idl_subset_ideal':
  assumes acar: "a  $\in$  carrier R" and bcarr: "b  $\in$  carrier R"
  shows "(Idl {a}  $\subseteq$  Idl {b}) = (a  $\in$  Idl {b})"
<proof>
```

```
lemma (in ring) genideal_zero:
  "Idl {0} = {0}"
<proof>
```

```
lemma (in ring) genideal_one:
  "Idl {1} = carrier R"
<proof>
```

Generation of Principal Ideals in Commutative Rings

**definition**

```
cgenideal :: "('a, 'b) monoid_scheme  $\Rightarrow$  'a  $\Rightarrow$  'a set" ("PIDlz _" [80]
79)
  where "cgenideal R a = {x  $\otimes_R$  a | x. x  $\in$  carrier R}"
```

genhideal (?) really generates an ideal

```
lemma (in cring) cgenideal_ideal:
  assumes acar: "a  $\in$  carrier R"
  shows "ideal (PIDl a) R"
<proof>
```

```
lemma (in ring) cgenideal_self:
  assumes icarr: "i  $\in$  carrier R"
  shows "i  $\in$  PIDl i"
<proof>
```

cgenideal is minimal

```
lemma (in ring) cgenideal_minimal:
  assumes "ideal J R"
```



```

    assumes aJ: "a ∈ J"
    shows "PIdl a ⊆ J"
  <proof>

```

```

lemma (in cring) cgenideal_eq_genideal:
  assumes icarr: "i ∈ carrier R"
  shows "PIdl i = Idl {i}"
  <proof>

```

```

lemma (in cring) cgenideal_eq_rcos:
  "PIdl i = carrier R #> i"
  <proof>

```

```

lemma (in cring) cgenideal_is_principalideal:
  assumes icarr: "i ∈ carrier R"
  shows "principalideal (PIdl i) R"
  <proof>

```

## 9.7 Union of Ideals

```

lemma (in ring) union_genideal:
  assumes idealI: "ideal I R"
    and idealJ: "ideal J R"
  shows "Idl (I ∪ J) = I <+> J"
  <proof>

```

## 9.8 Properties of Principal Ideals

0 generates the zero ideal

```

lemma (in ring) zero_genideal:
  shows "Idl {0} = {0}"
  <proof>

```

1 generates the unit ideal

```

lemma (in ring) one_genideal:
  shows "Idl {1} = carrier R"
  <proof>

```

The zero ideal is a principal ideal

```

corollary (in ring) zeropideal:
  shows "principalideal {0} R"
  <proof>

```

The unit ideal is a principal ideal

```

corollary (in ring) onepideal:
  shows "principalideal (carrier R) R"
  <proof>

```

Every principal ideal is a right coset of the carrier

```

lemma (in principalideal) rcos_generate:
  assumes "cring R"
  shows " $\exists x \in I. I = \text{carrier } R \#> x$ "
<proof>

```

## 9.9 Prime Ideals

```

lemma (in ideal) primeidealCD:
  assumes "cring R"
  assumes notprime: " $\neg \text{primeideal } I \text{ } R$ "
  shows " $\text{carrier } R = I \vee (\exists a \ b. a \in \text{carrier } R \wedge b \in \text{carrier } R \wedge a \otimes b \in I \wedge a \notin I \wedge b \notin I)$ "
<proof>

```

```

lemma (in ideal) primeidealCE:
  assumes "cring R"
  assumes notprime: " $\neg \text{primeideal } I \text{ } R$ "
  obtains "carrier R = I"
    | " $\exists a \ b. a \in \text{carrier } R \wedge b \in \text{carrier } R \wedge a \otimes b \in I \wedge a \notin I \wedge b \notin I$ "
<proof>

```

If  $\{0\}$  is a prime ideal of a commutative ring, the ring is a domain

```

lemma (in cring) zeroprimeideal_domainI:
  assumes pi: "primeideal {0} R"
  shows "domain R"
<proof>

```

```

corollary (in cring) domain_eq_zeroprimeideal:
  shows "domain R = primeideal {0} R"
<proof>

```

## 9.10 Maximal Ideals

```

lemma (in ideal) helper_I_closed:
  assumes carr: " $a \in \text{carrier } R$ " " $x \in \text{carrier } R$ " " $y \in \text{carrier } R$ "
  and axI: " $a \otimes x \in I$ "
  shows " $a \otimes (x \otimes y) \in I$ "
<proof>

```

```

lemma (in ideal) helper_max_prime:
  assumes "cring R"
  assumes acar: " $a \in \text{carrier } R$ "
  shows "ideal {x ∈ carrier R. a ⊗ x ∈ I} R"
<proof>

```

In a cring every maximal ideal is prime

```

lemma (in cring) maximalideal_is_prime:
  assumes "maximalideal I R"

```

```

    shows "primeideal I R"
  <proof>

```

### 9.11 Derived Theorems

— A non-zero cring that has only the two trivial ideals is a field

```

lemma (in cring) trivialideals_fieldI:
  assumes carrnzero: "carrier R  $\neq$  {0}"
  and haveideals: "{I. ideal I R} = {{0}, carrier R}"
  shows "field R"
  <proof>

```

```

lemma (in field) all_ideals:
  shows "{I. ideal I R} = {{0}, carrier R}"
  <proof>
lemma (in cring) trivialideals_eq_field:
  assumes carrnzero: "carrier R  $\neq$  {0}"
  shows "{I. ideal I R} = {{0}, carrier R} = field R"
  <proof>

```

Like zeroprimeideal for domains

```

lemma (in field) zeromaximalideal:
  "maximalideal {0} R"
  <proof>

```

```

lemma (in cring) zeromaximalideal_fieldI:
  assumes zeromax: "maximalideal {0} R"
  shows "field R"
  <proof>

```

```

lemma (in cring) zeromaximalideal_eq_field:
  "maximalideal {0} R = field R"
  <proof>

```

end

```

theory RingHom
imports Ideal
begin

```

## 10 Homomorphisms of Non-Commutative Rings

Lifting existing lemmas in a ring\_hom\_ring locale

```

locale ring_hom_ring = R: ring R + S: ring S
  for R (structure) and S (structure) +
  fixes h

```

```

    assumes homh: "h ∈ ring_hom R S"
    notes hom_mult [simp] = ring_hom_mult [OF homh]
    and hom_one [simp] = ring_hom_one [OF homh]

sublocale ring_hom_cring ⊆ ring: ring_hom_ring
  ⟨proof⟩

sublocale ring_hom_ring ⊆ abelian_group: abelian_group_hom R S
  ⟨proof⟩

lemma (in ring_hom_ring) is_ring_hom_ring:
  "ring_hom_ring R S h"
  ⟨proof⟩

lemma ring_hom_ringI:
  fixes R (structure) and S (structure)
  assumes "ring R" "ring S"
  assumes
    hom_closed: "!!x. x ∈ carrier R ==> h x ∈ carrier S"
    and compatible_mult: "!!x y. [| x : carrier R; y : carrier R |]
==> h (x ⊗ y) = h x ⊗S h y"
    and compatible_add: "!!x y. [| x : carrier R; y : carrier R |] ==>
h (x ⊕ y) = h x ⊕S h y"
    and compatible_one: "h 1 = 1S"
  shows "ring_hom_ring R S h"
  ⟨proof⟩

lemma ring_hom_ringI2:
  assumes "ring R" "ring S"
  assumes h: "h ∈ ring_hom R S"
  shows "ring_hom_ring R S h"
  ⟨proof⟩

lemma ring_hom_ringI3:
  fixes R (structure) and S (structure)
  assumes "abelian_group_hom R S h" "ring R" "ring S"
  assumes compatible_mult: "!!x y. [| x : carrier R; y : carrier R |]
==> h (x ⊗ y) = h x ⊗S h y"
    and compatible_one: "h 1 = 1S"
  shows "ring_hom_ring R S h"
  ⟨proof⟩

lemma ring_hom_cringI:
  assumes "ring_hom_ring R S h" "cring R" "cring S"
  shows "ring_hom_cring R S h"
  ⟨proof⟩

```

## 10.1 The Kernel of a Ring Homomorphism

— the kernel of a ring homomorphism is an ideal

```
lemma (in ring_hom_ring) kernel_is_ideal:
  shows "ideal (a_kernel R S h) R"
<proof>
```

Elements of the kernel are mapped to zero

```
lemma (in abelian_group_hom) kernel_zero [simp]:
  "i ∈ a_kernel R S h ⇒ h i = 0_S"
<proof>
```

## 10.2 Cosets

Cosets of the kernel correspond to the elements of the image of the homomorphism

```
lemma (in ring_hom_ring) rcos_imp_homeq:
  assumes acarr: "a ∈ carrier R"
  and xrcos: "x ∈ a_kernel R S h +> a"
  shows "h x = h a"
<proof>
```

```
lemma (in ring_hom_ring) homeq_imp_rcos:
  assumes acarr: "a ∈ carrier R"
  and xcarr: "x ∈ carrier R"
  and hx: "h x = h a"
  shows "x ∈ a_kernel R S h +> a"
<proof>
```

```
corollary (in ring_hom_ring) rcos_eq_homeq:
  assumes acarr: "a ∈ carrier R"
  shows "(a_kernel R S h) +> a = {x ∈ carrier R. h x = h a}"
<proof>
```

end

```
theory QuotRing
imports RingHom
begin
```

# 11 Quotient Rings

## 11.1 Multiplication on Cosets

definition

```
rcoset_mult :: "[( 'a, _) ring_scheme, 'a set, 'a set, 'a set] ⇒ 'a
set"
```

```

    ("[mod _:] _  $\otimes$  _" [81,81,81] 80)
  where "rcoset_mult R I A B = ( $\bigcup_{a \in A} \bigcup_{b \in B} I +>_R (a \otimes_R b)$ )"

```

rcoset\_mult fulfils the properties required by congruences

```

lemma (in ideal) rcoset_mult_add:
  "[x  $\in$  carrier R; y  $\in$  carrier R]  $\implies$  [mod I:] (I +> x)  $\otimes$  (I +> y) =
  I +> (x  $\otimes$  y)"
<proof>

```

## 11.2 Quotient Ring Definition

**definition**

```

FactRing :: "[('a,'b) ring_scheme, 'a set]  $\Rightarrow$  ('a set) ring" (infixl
"Quot" 65)
  where "FactRing R I =
    (carrier = a_rcosets_R I, mult = rcoset_mult R I, one = (I +>_R 1_R),
    zero = I, add = set_add R)"

```

## 11.3 Factorization over General Ideals

The quotient is a ring

```

lemma (in ideal) quotient_is_ring:
  shows "ring (R Quot I)"
<proof>

```

This is a ring homomorphism

```

lemma (in ideal) rcos_ring_hom:
  "(op +> I)  $\in$  ring_hom R (R Quot I)"
<proof>

```

```

lemma (in ideal) rcos_ring_hom_ring:
  "ring_hom_ring R (R Quot I) (op +> I)"
<proof>

```

The quotient of a cring is also commutative

```

lemma (in ideal) quotient_is_cring:
  assumes "cring R"
  shows "cring (R Quot I)"
<proof>

```

Cosets as a ring homomorphism on crings

```

lemma (in ideal) rcos_ring_hom_cring:
  assumes "cring R"
  shows "ring_hom_cring R (R Quot I) (op +> I)"
<proof>

```

## 11.4 Factorization over Prime Ideals

The quotient ring generated by a prime ideal is a domain

```
lemma (in primeideal) quotient_is_domain:
  shows "domain (R Quot I)"
  <proof>
```

Generating right cosets of a prime ideal is a homomorphism on commutative rings

```
lemma (in primeideal) rcos_ring_hom_cring:
  shows "ring_hom_cring R (R Quot I) (op +> I)"
  <proof>
```

## 11.5 Factorization over Maximal Ideals

In a commutative ring, the quotient ring over a maximal ideal is a field. The proof follows “W. Adkins, S. Weintraub: Algebra – An Approach via Module Theory”

```
lemma (in maximalideal) quotient_is_field:
  assumes "cring R"
  shows "field (R Quot I)"
  <proof>

end
```

```
theory IntRing
imports QuotRing Lattice Int "~/src/HOL/Old_Number_Theory/Primes"
begin
```

# 12 The Ring of Integers

## 12.1 Some properties of int

```
lemma dvds_eq_abseq:
  "(l dvd k ∧ k dvd l) = (abs l = abs (k::int))"
  <proof>
```

## 12.2 $\mathbb{Z}$ : The Set of Integers as Algebraic Structure

```
definition
  int_ring :: "int ring" ("Z") where
    "int_ring = (⌊carrier = UNIV, mult = op *, one = 1, zero = 0, add = op
  +⌋)"
```

```
lemma int_Zcarr [intro!, simp]:
  "k ∈ carrier Z"
```

*<proof>*

```
lemma int_is_cring:
  "cring  $\mathcal{Z}$ "
<proof>
```

### 12.3 Interpretations

Since definitions of derived operations are global, their interpretation needs to be done as early as possible — that is, with as few assumptions as possible.

```
interpretation int: monoid  $\mathcal{Z}$ 
  where "carrier  $\mathcal{Z}$  = UNIV"
        and "mult  $\mathcal{Z}$  x y = x * y"
        and "one  $\mathcal{Z}$  = 1"
        and "pow  $\mathcal{Z}$  x n = x^n"
<proof>
```

```
interpretation int: comm_monoid  $\mathcal{Z}$ 
  where "finprod  $\mathcal{Z}$  f A = (if finite A then setprod f A else undefined)"
<proof>
```

```
interpretation int: abelian_monoid  $\mathcal{Z}$ 
  where "zero  $\mathcal{Z}$  = 0"
        and "add  $\mathcal{Z}$  x y = x + y"
        and "finsum  $\mathcal{Z}$  f A = (if finite A then setsum f A else undefined)"
<proof>
```

```
interpretation int: abelian_group  $\mathcal{Z}$ 
  where "a_inv  $\mathcal{Z}$  x = - x"
        and "a_minus  $\mathcal{Z}$  x y = x - y"
<proof>
```

```
interpretation int: "domain"  $\mathcal{Z}$ 
<proof>
```

Removal of occurrences of UNIV in interpretation result — experimental.

```
lemma UNIV:
  "x ∈ UNIV = True"
  "A ⊆ UNIV = True"
  "(ALL x : UNIV. P x) = (ALL x. P x)"
  "(EX x : UNIV. P x) = (EX x. P x)"
  "(True --> Q) = Q"
  "(True ==> PROP R) == PROP R"
<proof>
```

```
interpretation int :
  partial_order "(| carrier = UNIV::int set, eq = op =, le = op ≤ |)"
  where "carrier (| carrier = UNIV::int set, eq = op =, le = op ≤ |)
    = UNIV"
```



```

    and "le (| carrier = UNIV::int set, eq = op =, le = op ≤ |) x y =
(x ≤ y)"
    and "lless (| carrier = UNIV::int set, eq = op =, le = op ≤ |) x
y = (x < y)"
  <proof>

```

```

interpretation int :
  lattice "(| carrier = UNIV::int set, eq = op =, le = op ≤ |)"
  where "join (| carrier = UNIV::int set, eq = op =, le = op ≤ |) x y
= max x y"
    and "meet (| carrier = UNIV::int set, eq = op =, le = op ≤ |) x y
= min x y"
  <proof>

```

```

interpretation int :
  total_order "(| carrier = UNIV::int set, eq = op =, le = op ≤ |)"
  <proof>

```

## 12.4 Generated Ideals of $\mathcal{Z}$

```

lemma int_Idl:
  "Idl $\mathcal{Z}$  {a} = {x * a | x. True}"
  <proof>

```

```

lemma multiples_principalideal:
  "principalideal {x * a | x. True }  $\mathcal{Z}$ "
  <proof>

```

```

lemma prime_primeideal:
  assumes prime: "prime (nat p)"
  shows "primeideal (Idl $\mathcal{Z}$  {p})  $\mathcal{Z}$ "
  <proof>

```

## 12.5 Ideals and Divisibility

```

lemma int_Idl_subset_ideal:
  "Idl $\mathcal{Z}$  {k} ⊆ Idl $\mathcal{Z}$  {l} = (k ∈ Idl $\mathcal{Z}$  {l})"
  <proof>

```

```

lemma Idl_subset_eq_dvd:
  "(Idl $\mathcal{Z}$  {k} ⊆ Idl $\mathcal{Z}$  {l}) = (l dvd k)"
  <proof>

```

```

lemma dvds_eq_Idl:
  "(l dvd k ∧ k dvd l) = (Idl $\mathcal{Z}$  {k} = Idl $\mathcal{Z}$  {l})"
  <proof>

```

```

lemma Idl_eq_abs:
  "(Idl $\mathcal{Z}$  {k} = Idl $\mathcal{Z}$  {l}) = (abs l = abs k)"

```

*<proof>*

## 12.6 Ideals and the Modulus

**definition**

```
ZMod :: "int => int => int set"
where "ZMod k r = (IdlZ {k}) +>Z r"
```

**lemmas** ZMod\_defs =

```
ZMod_def genideal_def
```

**lemma** rcos\_zfact:

```
assumes k1l: "k ∈ ZMod l r"
shows "EX x. k = x * l + r"
```

*<proof>*

**lemma** ZMod\_imp\_zmod:

```
assumes zmods: "ZMod m a = ZMod m b"
shows "a mod m = b mod m"
```

*<proof>*

**lemma** ZMod\_mod:

```
shows "ZMod m a = ZMod m (a mod m)"
```

*<proof>*

**lemma** zmod\_imp\_ZMod:

```
assumes modeq: "a mod m = b mod m"
shows "ZMod m a = ZMod m b"
```

*<proof>*

**corollary** ZMod\_eq\_mod:

```
shows "(ZMod m a = ZMod m b) = (a mod m = b mod m)"
```

*<proof>*

## 12.7 Factorization

**definition**

```
ZFact :: "int ⇒ int set ring"
where "ZFact k = Z Quot (IdlZ {k})"
```

**lemmas** ZFact\_defs = ZFact\_def FactRing\_def

**lemma** ZFact\_is\_cring:

```
shows "cring (ZFact k)"
```

*<proof>*

**lemma** ZFact\_zero:

```
"carrier (ZFact 0) = (⋃ a. {{a}})"
```

*<proof>*

```
lemma ZFact_one:
  "carrier (ZFact 1) = {UNIV}"
  <proof>
```

```
lemma ZFact_prime_is_domain:
  assumes pprime: "prime (nat p)"
  shows "domain (ZFact p)"
  <proof>
```

```
end
```

```
theory Module
imports Ring
begin
```

## 13 Modules over an Abelian Group

### 13.1 Definitions

```
record ('a, 'b) module = "'b ring" +
  smult :: "'a, 'b] => 'b" (infixl "⊙" 70)

locale module = R: cring + M: abelian_group M for M (structure) +
  assumes smult_closed [simp, intro]:
    "[| a ∈ carrier R; x ∈ carrier M |] ==> a ⊙M x ∈ carrier M"
  and smult_l_distr:
    "[| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊕ b) ⊙M x = a ⊙M x ⊕M b ⊙M x"
  and smult_r_distr:
    "[| a ∈ carrier R; x ∈ carrier M; y ∈ carrier M |] ==>
    a ⊙M (x ⊕M y) = a ⊙M x ⊕M a ⊙M y"
  and smult_assoc1:
    "[| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊗ b) ⊙M x = a ⊙M (b ⊙M x)"
  and smult_one [simp]:
    "x ∈ carrier M ==> 1 ⊙M x = x"

locale algebra = module + cring M +
  assumes smult_assoc2:
    "[| a ∈ carrier R; x ∈ carrier M; y ∈ carrier M |] ==>
    (a ⊙M x) ⊗M y = a ⊙M (x ⊗M y)"

lemma moduleI:
  fixes R (structure) and M (structure)
  assumes cring: "cring R"
    and abelian_group: "abelian_group M"
    and smult_closed:
```

```

    "!!a x. [| a ∈ carrier R; x ∈ carrier M |] ==> a ⊙M x ∈ carrier
M"
  and smult_l_distr:
    "!!a b x. [| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊕ b) ⊙M x = (a ⊙M x) ⊕M (b ⊙M x)"
  and smult_r_distr:
    "!!a x y. [| a ∈ carrier R; x ∈ carrier M; y ∈ carrier M |] ==>
    a ⊙M (x ⊕M y) = (a ⊙M x) ⊕M (a ⊙M y)"
  and smult_assoc1:
    "!!a b x. [| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊗ b) ⊙M x = a ⊙M (b ⊙M x)"
  and smult_one:
    "!!x. x ∈ carrier M ==> 1 ⊙M x = x"
  shows "module R M"
  <proof>

```

```

lemma algebraI:
  fixes R (structure) and M (structure)
  assumes R_cring: "cring R"
  and M_cring: "cring M"
  and smult_closed:
    "!!a x. [| a ∈ carrier R; x ∈ carrier M |] ==> a ⊙M x ∈ carrier
M"
  and smult_l_distr:
    "!!a b x. [| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊕ b) ⊙M x = (a ⊙M x) ⊕M (b ⊙M x)"
  and smult_r_distr:
    "!!a x y. [| a ∈ carrier R; x ∈ carrier M; y ∈ carrier M |] ==>
    a ⊙M (x ⊕M y) = (a ⊙M x) ⊕M (a ⊙M y)"
  and smult_assoc1:
    "!!a b x. [| a ∈ carrier R; b ∈ carrier R; x ∈ carrier M |] ==>
    (a ⊗ b) ⊙M x = a ⊙M (b ⊙M x)"
  and smult_one:
    "!!x. x ∈ carrier M ==> (one R) ⊙M x = x"
  and smult_assoc2:
    "!!a x y. [| a ∈ carrier R; x ∈ carrier M; y ∈ carrier M |] ==>
    (a ⊙M x) ⊗M y = a ⊙M (x ⊗M y)"
  shows "algebra R M"
  <proof>

```

```

lemma (in algebra) R_cring:
  "cring R"
  <proof>

```

```

lemma (in algebra) M_cring:
  "cring M"
  <proof>

```

```

lemma (in algebra) module:

```

```
"module R M"
⟨proof⟩
```

### 13.2 Basic Properties of Algebras

```
lemma (in algebra) smult_l_null [simp]:
  "x ∈ carrier M ==> 0 ∘M x = 0M"
⟨proof⟩
```

```
lemma (in algebra) smult_r_null [simp]:
  "a ∈ carrier R ==> a ∘M 0M = 0M"
⟨proof⟩
```

```
lemma (in algebra) smult_l_minus:
  "[| a ∈ carrier R; x ∈ carrier M |] ==> (⊖ a) ∘M x = ⊖M (a ∘M x)"
⟨proof⟩
```

```
lemma (in algebra) smult_r_minus:
  "[| a ∈ carrier R; x ∈ carrier M |] ==> a ∘M (⊖M x) = ⊖M (a ∘M x)"
⟨proof⟩
```

```
end
```

```
theory UnivPoly
imports Module RingHom
begin
```

## 14 Univariate Polynomials

Polynomials are formalised as modules with additional operations for extracting coefficients from polynomials and for obtaining monomials from coefficients and exponents (record `up_ring`). The carrier set is a set of bounded functions from `Nat` to the coefficient domain. Bounded means that these functions return zero above a certain bound (the degree). There is a chapter on the formalisation of polynomials in the PhD thesis [1], which was implemented with axiomatic type classes. This was later ported to `Locales`.

### 14.1 The Constructor for Univariate Polynomials

Functions with finite support.

```
locale bound =
  fixes z :: 'a
  and n :: nat
  and f :: "nat => 'a"
  assumes bound: "!!m. n < m ==> f m = z"
```

```

declare bound.intro [intro!]
  and bound.bound [dest]

lemma bound_below:
  assumes bound: "bound z m f" and nonzero: "f n  $\neq$  z" shows "n  $\leq$  m"
  <proof>

record ('a, 'p) up_ring = "('a, 'p) module" +
  monom :: "[ 'a, nat ] => 'p"
  coeff :: "[ 'p, nat ] => 'a"

definition
  up :: "('a, 'm) ring_scheme => (nat => 'a) set"
  where "up R = {f. f  $\in$  UNIV -> carrier R & (EX n. bound 0R n f)}"

definition UP :: "('a, 'm) ring_scheme => ('a, nat => 'a) up_ring"
  where "UP R = (|
    carrier = up R,
    mult = (%p:up R. %q:up R. %n.  $\bigoplus_{R} i \in \{..n\}. p\ i \otimes_R q\ (n-i)$ ),
    one = (%i. if i=0 then 1R else 0R),
    zero = (%i. 0R),
    add = (%p:up R. %q:up R. %i. p i  $\oplus_R$  q i),
    smult = (%a:carrier R. %p:up R. %i. a  $\otimes_R$  p i),
    monom = (%a:carrier R. %n i. if i=n then a else 0R),
    coeff = (%p:up R. %n. p n) |)"

Properties of the set of polynomials up.

lemma mem_upI [intro]:
  "[| !n. f n  $\in$  carrier R; EX n. bound (zero R) n f |] ==> f  $\in$  up R"
  <proof>

lemma mem_upD [dest]:
  "f  $\in$  up R ==> f n  $\in$  carrier R"
  <proof>

context ring
begin

lemma bound_upD [dest]: "f  $\in$  up R ==> EX n. bound 0 n f" <proof>

lemma up_one_closed: "(%n. if n = 0 then 1 else 0)  $\in$  up R" <proof>

lemma up_smult_closed: "[| a  $\in$  carrier R; p  $\in$  up R |] ==> (%i. a  $\otimes$  p i)  $\in$  up R" <proof>

lemma up_add_closed:
  "[| p  $\in$  up R; q  $\in$  up R |] ==> (%i. p i  $\oplus$  q i)  $\in$  up R"
  <proof>

```

```

lemma up_a_inv_closed:
  "p ∈ up R ==> (%i. ⊖ (p i)) ∈ up R"
  <proof>

lemma up_minus_closed:
  "[| p ∈ up R; q ∈ up R |] ==> (%i. p i ⊖ q i) ∈ up R"
  <proof>

lemma up_mult_closed:
  "[| p ∈ up R; q ∈ up R |] ==>
  (%n. ⊕ i ∈ {...n}. p i ⊗ q (n-i)) ∈ up R"
  <proof>

end

```

## 14.2 Effect of Operations on Coefficients

```

locale UP =
  fixes R (structure) and P (structure)
  defines P_def: "P == UP R"

locale UP_ring = UP + R: ring R

locale UP_cring = UP + R: cring R

sublocale UP_cring < UP_ring
  <proof>

locale UP_domain = UP + R: "domain" R

sublocale UP_domain < UP_cring
  <proof>

context UP
begin

  Temporarily declare  $P \equiv UP\ R$  as simp rule.

  declare P_def [simp]

  lemma up_eqI:
    assumes prem: "!!n. coeff P p n = coeff P q n" and R: "p ∈ carrier
    P" "q ∈ carrier P"
    shows "p = q"
    <proof>

  lemma coeff_closed [simp]:
    "p ∈ carrier P ==> coeff P p n ∈ carrier R" <proof>

```

end

context UP\_ring  
begin

```
lemma coeff_monom [simp]:
  "a ∈ carrier R ==> coeff P (monom P a m) n = (if m=n then a else 0)"
  <proof>

lemma coeff_zero [simp]: "coeff P 0_P n = 0" <proof>

lemma coeff_one [simp]: "coeff P 1_P n = (if n=0 then 1 else 0)"
  <proof>

lemma coeff_smult [simp]:
  "[| a ∈ carrier R; p ∈ carrier P |] ==> coeff P (a ⊙_P p) n = a ⊗ coeff
P p n"
  <proof>

lemma coeff_add [simp]:
  "[| p ∈ carrier P; q ∈ carrier P |] ==> coeff P (p ⊕_P q) n = coeff
P p n ⊕ coeff P q n"
  <proof>

lemma coeff_mult [simp]:
  "[| p ∈ carrier P; q ∈ carrier P |] ==> coeff P (p ⊗_P q) n = (⊕ i ∈
{..n}. coeff P p i ⊗ coeff P q (n-i))"
  <proof>
```

end

### 14.3 Polynomials Form a Ring.

context UP\_ring  
begin

Operations are closed over P.

```
lemma UP_mult_closed [simp]:
  "[| p ∈ carrier P; q ∈ carrier P |] ==> p ⊗_P q ∈ carrier P" <proof>

lemma UP_one_closed [simp]:
  "1_P ∈ carrier P" <proof>

lemma UP_zero_closed [intro, simp]:
  "0_P ∈ carrier P" <proof>

lemma UP_a_closed [intro, simp]:
```



```

" [| p ∈ carrier P; q ∈ carrier P |] ==> p ⊕P q ∈ carrier P" <proof>

lemma monom_closed [simp]:
  "a ∈ carrier R ==> monom P a n ∈ carrier P" <proof>

lemma UP_smult_closed [simp]:
  " [| a ∈ carrier R; p ∈ carrier P |] ==> a ⊙P p ∈ carrier P" <proof>

end

declare (in UP) P_def [simp del]

Algebraic ring properties

context UP_ring
begin

lemma UP_a_assoc:
  assumes R: "p ∈ carrier P" "q ∈ carrier P" "r ∈ carrier P"
  shows "(p ⊕P q) ⊕P r = p ⊕P (q ⊕P r)" <proof>

lemma UP_l_zero [simp]:
  assumes R: "p ∈ carrier P"
  shows "0P ⊕P p = p" <proof>

lemma UP_l_neg_ex:
  assumes R: "p ∈ carrier P"
  shows "EX q : carrier P. q ⊕P p = 0P"
  <proof>

lemma UP_a_comm:
  assumes R: "p ∈ carrier P" "q ∈ carrier P"
  shows "p ⊕P q = q ⊕P p" <proof>

lemma UP_m_assoc:
  assumes R: "p ∈ carrier P" "q ∈ carrier P" "r ∈ carrier P"
  shows "(p ⊗P q) ⊗P r = p ⊗P (q ⊗P r)"
  <proof>

lemma UP_r_one [simp]:
  assumes R: "p ∈ carrier P" shows "p ⊗P 1P = p"
  <proof>

lemma UP_l_one [simp]:
  assumes R: "p ∈ carrier P"
  shows "1P ⊗P p = p"
  <proof>

lemma UP_l_distr:
  assumes R: "p ∈ carrier P" "q ∈ carrier P" "r ∈ carrier P"

```

```

shows "(p ⊕P q) ⊗P r = (p ⊗P r) ⊕P (q ⊗P r)"
⟨proof⟩

lemma UP_r_distr:
  assumes R: "p ∈ carrier P" "q ∈ carrier P" "r ∈ carrier P"
  shows "r ⊗P (p ⊕P q) = (r ⊗P p) ⊕P (r ⊗P q)"
  ⟨proof⟩

theorem UP_ring: "ring P"
  ⟨proof⟩

end

14.4 Polynomials Form a Commutative Ring.

context UP_cring
begin

lemma UP_m_comm:
  assumes R1: "p ∈ carrier P" and R2: "q ∈ carrier P" shows "p ⊗P q
= q ⊗P p"
  ⟨proof⟩

14.5 Polynomials over a commutative ring for a commutative
      ring

theorem UP_cring:
  "cring P" ⟨proof⟩

end

context UP_ring
begin

lemma UP_a_inv_closed [intro, simp]:
  "p ∈ carrier P ==> ⊖P p ∈ carrier P"
  ⟨proof⟩

lemma coeff_a_inv [simp]:
  assumes R: "p ∈ carrier P"
  shows "coeff P (⊖P p) n = ⊖ (coeff P p n)"
  ⟨proof⟩

end

sublocale UP_ring < P: ring P ⟨proof⟩
sublocale UP_cring < P: cring P ⟨proof⟩

```

## 14.6 Polynomials Form an Algebra

context UP\_ring  
begin

lemma UP\_smult\_l\_distr:

"[| a ∈ carrier R; b ∈ carrier R; p ∈ carrier P |] ==>  
(a ⊕ b) ⊙<sub>P</sub> p = a ⊙<sub>P</sub> p ⊕ b ⊙<sub>P</sub> p"  
⟨proof⟩

lemma UP\_smult\_r\_distr:

"[| a ∈ carrier R; p ∈ carrier P; q ∈ carrier P |] ==>  
a ⊙<sub>P</sub> (p ⊕<sub>P</sub> q) = a ⊙<sub>P</sub> p ⊕<sub>P</sub> a ⊙<sub>P</sub> q"  
⟨proof⟩

lemma UP\_smult\_assoc1:

"[| a ∈ carrier R; b ∈ carrier R; p ∈ carrier P |] ==>  
(a ⊗ b) ⊙<sub>P</sub> p = a ⊙<sub>P</sub> (b ⊙<sub>P</sub> p)"  
⟨proof⟩

lemma UP\_smult\_zero [simp]:

"p ∈ carrier P ==> 0 ⊙<sub>P</sub> p = 0<sub>P</sub>"  
⟨proof⟩

lemma UP\_smult\_one [simp]:

"p ∈ carrier P ==> 1 ⊙<sub>P</sub> p = p"  
⟨proof⟩

lemma UP\_smult\_assoc2:

"[| a ∈ carrier R; p ∈ carrier P; q ∈ carrier P |] ==>  
(a ⊙<sub>P</sub> p) ⊗<sub>P</sub> q = a ⊙<sub>P</sub> (p ⊗<sub>P</sub> q)"  
⟨proof⟩

end

Interpretation of lemmas from algebra.

lemma (in cring) cring:

"cring R" ⟨proof⟩

lemma (in UP\_cring) UP\_algebra:

"algebra R P" ⟨proof⟩

sublocale UP\_cring < algebra R P ⟨proof⟩

## 14.7 Further Lemmas Involving Monomials

context UP\_ring  
begin

lemma monom\_zero [simp]:

"monom P 0 n = 0<sub>P</sub>" *<proof>*

**lemma** monom\_mult\_is\_smult:  
 assumes R: "a ∈ carrier R" "p ∈ carrier P"  
 shows "monom P a 0 ⊗<sub>P</sub> p = a ⊙<sub>P</sub> p"  
*<proof>*

**lemma** monom\_one [simp]:  
 "monom P 1 0 = 1<sub>P</sub>"  
*<proof>*

**lemma** monom\_add [simp]:  
 "[| a ∈ carrier R; b ∈ carrier R |] ==>  
 monom P (a ⊕ b) n = monom P a n ⊕<sub>P</sub> monom P b n"  
*<proof>*

**lemma** monom\_one\_Suc:  
 "monom P 1 (Suc n) = monom P 1 n ⊗<sub>P</sub> monom P 1 1"  
*<proof>*

**lemma** monom\_one\_Suc2:  
 "monom P 1 (Suc n) = monom P 1 1 ⊗<sub>P</sub> monom P 1 n"  
*<proof>*

The following corollary follows from lemmas monom P 1 (Suc ?n) = monom P 1 ?n ⊗<sub>P</sub> monom P 1 1 and monom P 1 (Suc ?n) = monom P 1 1 ⊗<sub>P</sub> monom P 1 ?n, and is trivial in UP\_cring

**corollary** monom\_one\_comm: shows "monom P 1 k ⊗<sub>P</sub> monom P 1 1 = monom P 1 1 ⊗<sub>P</sub> monom P 1 k"  
*<proof>*

**lemma** monom\_mult\_smult:  
 "[| a ∈ carrier R; b ∈ carrier R |] ==> monom P (a ⊗ b) n = a ⊙<sub>P</sub> monom P b n"  
*<proof>*

**lemma** monom\_one\_mult:  
 "monom P 1 (n + m) = monom P 1 n ⊗<sub>P</sub> monom P 1 m"  
*<proof>*

**lemma** monom\_one\_mult\_comm: "monom P 1 n ⊗<sub>P</sub> monom P 1 m = monom P 1 m ⊗<sub>P</sub> monom P 1 n"  
*<proof>*

**lemma** monom\_mult [simp]:  
 assumes a\_in\_R: "a ∈ carrier R" and b\_in\_R: "b ∈ carrier R"  
 shows "monom P (a ⊗ b) (n + m) = monom P a n ⊗<sub>P</sub> monom P b m"  
*<proof>*

```

lemma monom_a_inv [simp]:
  "a ∈ carrier R ==> monom P (⊖ a) n = ⊖P monom P a n"
  ⟨proof⟩

```

```

lemma monom_inj:
  "inj_on (%a. monom P a n) (carrier R)"
  ⟨proof⟩

```

```

end

```

## 14.8 The Degree Function

```

definition
  deg :: "('a, 'm) ring_scheme, nat => 'a] => nat"
  where "deg R p = (LEAST n. bound 0R n (coeff (UP R) p))"

```

```

context UP_ring
begin

```

```

lemma deg_aboveI:
  "[| (!m. n < m ==> coeff P p m = 0); p ∈ carrier P |] ==> deg R p <= n"
  ⟨proof⟩

```

```

lemma deg_aboveD:
  assumes "deg R p < m" and "p ∈ carrier P"
  shows "coeff P p m = 0"
  ⟨proof⟩

```

```

lemma deg_belowI:
  assumes non_zero: "n ~≠ 0 ==> coeff P p n ~≠ 0"
  and R: "p ∈ carrier P"
  shows "n <= deg R p"
  — Logically, this is a slightly stronger version of deg_aboveD
  ⟨proof⟩

```

```

lemma lcoeff_nonzero_deg:
  assumes deg: "deg R p ~≠ 0" and R: "p ∈ carrier P"
  shows "coeff P p (deg R p) ~≠ 0"
  ⟨proof⟩

```

```

lemma lcoeff_nonzero_nonzero:
  assumes deg: "deg R p = 0" and nonzero: "p ~≠ 0P" and R: "p ∈ carrier P"
  shows "coeff P p 0 ~≠ 0"
  ⟨proof⟩

```

```

lemma lcoeff_nonzero:
  assumes neq: "p ~= 0P" and R: "p ∈ carrier P"
  shows "coeff P p (deg R p) ~= 0"
<proof>

lemma deg_eqI:
  "[| !!m. n < m ==> coeff P p m = 0;
    !!n. n ~= 0 ==> coeff P p n ~= 0; p ∈ carrier P |] ==> deg R p =
n"
<proof>

```

Degree and polynomial operations

```

lemma deg_add [simp]:
  "p ∈ carrier P ==> q ∈ carrier P ==>
  deg R (p ⊕P q) <= max (deg R p) (deg R q)"
<proof>

lemma deg_monom_le:
  "a ∈ carrier R ==> deg R (monom P a n) <= n"
<proof>

lemma deg_monom [simp]:
  "[| a ~= 0; a ∈ carrier R |] ==> deg R (monom P a n) = n"
<proof>

lemma deg_const [simp]:
  assumes R: "a ∈ carrier R" shows "deg R (monom P a 0) = 0"
<proof>

lemma deg_zero [simp]:
  "deg R 0P = 0"
<proof>

lemma deg_one [simp]:
  "deg R 1P = 0"
<proof>

lemma deg_uminus [simp]:
  assumes R: "p ∈ carrier P" shows "deg R (⊖P p) = deg R p"
<proof>

```

The following lemma is later *overwritten* by the most specific one for domains, `deg_smult`.

```

lemma deg_smult_ring [simp]:
  "[| a ∈ carrier R; p ∈ carrier P |] ==>
  deg R (a ⊙P p) <= (if a = 0 then 0 else deg R p)"
<proof>

end

```

```
context UP_domain
begin
```

```
lemma deg_smult [simp]:
  assumes R: "a ∈ carrier R" "p ∈ carrier P"
  shows "deg R (a ⊙P p) = (if a = 0 then 0 else deg R p)"
  <proof>

end
```

```
context UP_ring
begin
```

```
lemma deg_mult_ring:
  assumes R: "p ∈ carrier P" "q ∈ carrier P"
  shows "deg R (p ⊗P q) ≤ deg R p + deg R q"
  <proof>

end
```

```
context UP_domain
begin
```

```
lemma deg_mult [simp]:
  "[| p ~ 0P; q ~ 0P; p ∈ carrier P; q ∈ carrier P |] ==>
  deg R (p ⊗P q) = deg R p + deg R q"
  <proof>

end
```

The following lemmas also can be lifted to UP\_ring.

```
context UP_ring
begin
```

```
lemma coeff_finsum:
  assumes fin: "finite A"
  shows "p ∈ A -> carrier P ==>
  coeff P (finsum P p A) k = (⊕ i ∈ A. coeff P (p i) k)"
  <proof>
```

```
lemma up_repr:
  assumes R: "p ∈ carrier P"
  shows "(⊕P i ∈ {..deg R p}. monom P (coeff P p i) i) = p"
  <proof>
```

```
lemma up_repr_le:
  "[| deg R p ≤ n; p ∈ carrier P |] ==>
  (⊕P i ∈ {..n}. monom P (coeff P p i) i) = p"
```

*<proof>*

end

## 14.9 Polynomials over Integral Domains

lemma domainI:

assumes cring: "cring R"  
 and one\_not\_zero: "one R  $\neq$  zero R"  
 and integral: "!!a b. [| mult R a b = zero R; a  $\in$  carrier R;  
 b  $\in$  carrier R |] ==> a = zero R | b = zero R"  
 shows "domain R"  
*<proof>*

context UP\_domain

begin

lemma UP\_one\_not\_zero:

"1<sub>P</sub>  $\neq$  0<sub>P</sub>"  
*<proof>*

lemma UP\_integral:

"[| p  $\otimes_P$  q = 0<sub>P</sub>; p  $\in$  carrier P; q  $\in$  carrier P |] ==> p = 0<sub>P</sub> | q = 0<sub>P</sub>"  
*<proof>*

theorem UP\_domain:

"domain P"  
*<proof>*

end

Interpretation of theorems from domain.

sublocale UP\_domain < "domain" P

*<proof>*

## 14.10 The Evaluation Homomorphism and Universal Property

lemma (in abelian\_monoid) boundD\_carrier:

"[| bound 0 n f; n < m |] ==> f m  $\in$  carrier G"  
*<proof>*

context ring

begin

theorem diagonal\_sum:

"[| f  $\in$  {.. $n$  +  $m$  : nat}  $\rightarrow$  carrier R; g  $\in$  {.. $n$  +  $m$ }  $\rightarrow$  carrier R |] ==>  
 $(\bigoplus_{k \in \{.. $n$  +  $m\}}$ .  $\bigoplus_{i \in \{.. $k\}}$ . f i  $\otimes$  g (k - i)) =  
 $(\bigoplus_{k \in \{.. $n$  +  $m\}}$ .  $\bigoplus_{i \in \{.. $n$  +  $m$  - k\}}$ . f k  $\otimes$  g i)"$$$



*<proof>*

**theorem** cauchy\_product:

assumes bf: "bound 0 n f" and bg: "bound 0 m g"  
 and Rf: "f ∈ {...n} -> carrier R" and Rg: "g ∈ {...m} -> carrier R"  
 shows " $(\bigoplus_{k \in \{..n+m\}} \bigoplus_{i \in \{..k\}} f\ i \otimes g\ (k - i)) =$   
 $(\bigoplus_{i \in \{..n\}} f\ i) \otimes (\bigoplus_{i \in \{..m\}} g\ i)$ "

*<proof>*

**end**

**lemma** (in UP\_ring) const\_ring\_hom:

"(%a. monom P a 0) ∈ ring\_hom R P"

*<proof>*

**definition**

eval :: "[('a, 'm) ring\_scheme, ('b, 'n) ring\_scheme,  
 'a => 'b, 'b, nat => 'a] => 'b"  
 where "eval R S phi s = (λp ∈ carrier (UP R).  
 $\bigoplus_{i \in \{..deg\ R\ p\}} \text{phi}\ (\text{coeff}\ (\text{UP}\ R)\ p\ i) \otimes_S s\ (^)_{S\ i})$ "

**context** UP

**begin**

**lemma** eval\_on\_carrier:

fixes S (structure)  
 shows "p ∈ carrier P ==>  
 eval R S phi s p =  $(\bigoplus_{i \in \{..deg\ R\ p\}} \text{phi}\ (\text{coeff}\ P\ p\ i) \otimes_S s\ (^)_{S\ i})$ "

*<proof>*

**lemma** eval\_extensional:

"eval R S phi p ∈ extensional (carrier P)"

*<proof>*

**end**

The universal property of the polynomial ring

**locale** UP\_pre\_univ\_prop = ring\_hom\_cring + UP\_cring

**locale** UP\_univ\_prop = UP\_pre\_univ\_prop +

fixes s and Eval

assumes indet\_img\_carrier [simp, intro]: "s ∈ carrier S"

defines Eval\_def: "Eval == eval R S h s"

JE: I have moved the following lemma from Ring.thy and lifted then to the  
 locale ring\_hom\_ring from ring\_hom\_cring.

JE: I was considering using it in `eval_ring_hom`, but that property does not hold for non commutative rings, so maybe it is not that necessary.

```
lemma (in ring_hom_ring) hom_finsum [simp]:
  "[| finite A; f ∈ A -> carrier R |] ==>
   h (finsum R f A) = finsum S (h o f) A"
  <proof>
```

```
context UP_pre_univ_prop
begin
```

```
theorem eval_ring_hom:
  assumes S: "s ∈ carrier S"
  shows "eval R S h s ∈ ring_hom P S"
  <proof>
```

The following lemma could be proved in `UP_cring` with the additional assumption that `h` is closed.

```
lemma (in UP_pre_univ_prop) eval_const:
  "[| s ∈ carrier S; r ∈ carrier R |] ==> eval R S h s (monom P r 0) =
  h r"
  <proof>
```

Further properties of the evaluation homomorphism.

The following proof is complicated by the fact that in arbitrary rings one might have  $1 = 0$ .

```
lemma (in UP_pre_univ_prop) eval_monom1:
  assumes S: "s ∈ carrier S"
  shows "eval R S h s (monom P 1 1) = s"
  <proof>
```

```
end
```

Interpretation of ring homomorphism lemmas.

```
sublocale UP_univ_prop < ring_hom_cring P S Eval
  <proof>
```

```
lemma (in UP_cring) monom_pow:
  assumes R: "a ∈ carrier R"
  shows "(monom P a n) (^)P m = monom P (a (^) m) (n * m)"
  <proof>
```

```
lemma (in ring_hom_cring) hom_pow [simp]:
  "x ∈ carrier R ==> h (x (^) n) = h x (^)S (n::nat)"
  <proof>
```

```
lemma (in UP_univ_prop) Eval_monom:
  "r ∈ carrier R ==> Eval (monom P r n) = h r ⊗S s (^)S n"
```

*<proof>*

```
lemma (in UP_pre_univ_prop) eval_monom:
  assumes R: "r ∈ carrier R" and S: "s ∈ carrier S"
  shows "eval R S h s (monom P r n) = h r ⊗S s (^)S n"
<proof>
```

```
lemma (in UP_univ_prop) Eval_smult:
  "[| r ∈ carrier R; p ∈ carrier P |] ==> Eval (r ⊙P p) = h r ⊗S Eval
p"
<proof>
```

```
lemma ring_hom_cringI:
  assumes "cring R"
  and "cring S"
  and "h ∈ ring_hom R S"
  shows "ring_hom_cring R S h"
<proof>
```

```
context UP_pre_univ_prop
begin
```

```
lemma UP_hom_unique:
  assumes "ring_hom_cring P S Phi"
  assumes Phi: "Phi (monom P 1 (Suc 0)) = s"
  "!!r. r ∈ carrier R ==> Phi (monom P r 0) = h r"
  assumes "ring_hom_cring P S Psi"
  assumes Psi: "Psi (monom P 1 (Suc 0)) = s"
  "!!r. r ∈ carrier R ==> Psi (monom P r 0) = h r"
  and P: "p ∈ carrier P" and S: "s ∈ carrier S"
  shows "Phi p = Psi p"
<proof>
```

```
lemma ring_homD:
  assumes Phi: "Phi ∈ ring_hom P S"
  shows "ring_hom_cring P S Phi"
<proof>
```

```
theorem UP_universal_property:
  assumes S: "s ∈ carrier S"
  shows "EX! Phi. Phi ∈ ring_hom P S ∩ extensional (carrier P) &
Phi (monom P 1 1) = s &
(ALL r : carrier R. Phi (monom P r 0) = h r)"
<proof>
```

```
end
```

JE: The following lemma was added by me; it might be even lifted to a simpler locale

```
context monoid
begin
```

```
lemma nat_pow_eone[simp]: assumes x_in_G: "x ∈ carrier G" shows "x
(^) (1::nat) = x"
  ⟨proof⟩
```

```
end
```

```
context UP_ring
begin
```

```
abbreviation lcoeff :: "(nat =>'a) => 'a" where "lcoeff p == coeff P
p (deg R p)"
```

```
lemma lcoeff_nonzero2: assumes p_in_R: "p ∈ carrier P" and p_not_zero:
"p ≠ 0P" shows "lcoeff p ≠ 0"
  ⟨proof⟩
```

#### 14.11 The long division algorithm: some previous facts.

```
lemma coeff_minus [simp]:
  assumes p: "p ∈ carrier P" and q: "q ∈ carrier P" shows "coeff P (p
⊖P q) n = coeff P p n ⊖ coeff P q n"
  ⟨proof⟩
```

```
lemma lcoeff_closed [simp]: assumes p: "p ∈ carrier P" shows "lcoeff
p ∈ carrier R"
  ⟨proof⟩
```

```
lemma deg_smult_decr: assumes a_in_R: "a ∈ carrier R" and f_in_P: "f
∈ carrier P" shows "deg R (a ⊙P f) ≤ deg R f"
  ⟨proof⟩
```

```
lemma coeff_monom_mult: assumes R: "c ∈ carrier R" and P: "p ∈ carrier
P"
  shows "coeff P (monom P c n ⊗P p) (m + n) = c ⊗ (coeff P p m)"
  ⟨proof⟩
```

```
lemma deg_lcoeff_cancel:
  assumes p_in_P: "p ∈ carrier P" and q_in_P: "q ∈ carrier P" and r_in_P:
"r ∈ carrier P"
  and deg_r_nonzero: "deg R r ≠ 0"
  and deg_R_p: "deg R p ≤ deg R r" and deg_R_q: "deg R q ≤ deg R r"

  and coeff_R_p_eq_q: "coeff P p (deg R r) = ⊖R (coeff P q (deg R r))"
  shows "deg R (p ⊕P q) < deg R r"
  ⟨proof⟩
```

```

lemma monom_deg_mult:
  assumes f_in_P: "f ∈ carrier P" and g_in_P: "g ∈ carrier P" and deg_le:
    "deg R g ≤ deg R f"
  and a_in_R: "a ∈ carrier R"
  shows "deg R (g ⊗P monom P a (deg R f - deg R g)) ≤ deg R f"
  ⟨proof⟩

```

```

lemma deg_zero_impl_monom:
  assumes f_in_P: "f ∈ carrier P" and deg_f: "deg R f = 0"
  shows "f = monom P (coeff P f 0) 0"
  ⟨proof⟩

```

end

## 14.12 The long division proof for commutative rings

```

context UP_cring
begin

```

```

lemma exI3: assumes exist: "Pred x y z"
  shows "∃ x y z. Pred x y z"
  ⟨proof⟩

```

Jacobson's Theorem 2.14

```

lemma long_div_theorem:
  assumes g_in_P [simp]: "g ∈ carrier P" and f_in_P [simp]: "f ∈ carrier
P"
  and g_not_zero: "g ≠ 0P"
  shows "∃ q r (k::nat). (q ∈ carrier P) ∧ (r ∈ carrier P) ∧ (lcoeff
g)(^)Rk ⊙P f = g ⊗P q ⊕P r ∧ (r = 0P | deg R r < deg R g)"
  ⟨proof⟩

```

end

The remainder theorem as corollary of the long division theorem.

```

context UP_cring
begin

```

```

lemma deg_minus_monom:
  assumes a: "a ∈ carrier R"
  and R_not_trivial: "(carrier R ≠ {0})"
  shows "deg R (monom P 1R 1 ⊖P monom P a 0) = 1"
  (is "deg R ?g = 1")
  ⟨proof⟩

```

```

lemma lcoeff_monom:
  assumes a: "a ∈ carrier R" and R_not_trivial: "(carrier R ≠ {0})"
  shows "lcoeff (monom P 1R 1 ⊖P monom P a 0) = 1"
  ⟨proof⟩

```

```

lemma deg_nzero_nzero:
  assumes deg_p_nzero: "deg R p  $\neq$  0"
  shows "p  $\neq$  0P"
  <proof>

lemma deg_monom_minus:
  assumes a: "a  $\in$  carrier R"
  and R_not_trivial: "carrier R  $\neq$  {0}"
  shows "deg R (monom P 1R 1  $\ominus_P$  monom P a 0) = 1"
  (is "deg R ?g = 1")
  <proof>

lemma eval_monom_expr:
  assumes a: "a  $\in$  carrier R"
  shows "eval R R id a (monom P 1R 1  $\ominus_P$  monom P a 0) = 0"
  (is "eval R R id a ?g = _")
  <proof>

lemma remainder_theorem_exist:
  assumes f: "f  $\in$  carrier P" and a: "a  $\in$  carrier R"
  and R_not_trivial: "carrier R  $\neq$  {0}"
  shows " $\exists$  q r. (q  $\in$  carrier P)  $\wedge$  (r  $\in$  carrier P)  $\wedge$  f = (monom P 1R
1  $\ominus_P$  monom P a 0)  $\otimes_P$  q  $\oplus_P$  r  $\wedge$  (deg R r = 0)"
  (is " $\exists$  q r. (q  $\in$  carrier P)  $\wedge$  (r  $\in$  carrier P)  $\wedge$  f = ?g  $\otimes_P$  q  $\oplus_P$  r  $\wedge$ 
(deg R r = 0)")
  <proof>

lemma remainder_theorem_expression:
  assumes f [simp]: "f  $\in$  carrier P" and a [simp]: "a  $\in$  carrier R"
  and q [simp]: "q  $\in$  carrier P" and r [simp]: "r  $\in$  carrier P"
  and R_not_trivial: "carrier R  $\neq$  {0}"
  and f_expr: "f = (monom P 1R 1  $\ominus_P$  monom P a 0)  $\otimes_P$  q  $\oplus_P$  r"
  (is "f = ?g  $\otimes_P$  q  $\oplus_P$  r" is "f = ?gq  $\oplus_P$  r")
  and deg_r_0: "deg R r = 0"
  shows "r = monom P (eval R R id a f) 0"
  <proof>

corollary remainder_theorem:
  assumes f [simp]: "f  $\in$  carrier P" and a [simp]: "a  $\in$  carrier R"
  and R_not_trivial: "carrier R  $\neq$  {0}"
  shows " $\exists$  q r. (q  $\in$  carrier P)  $\wedge$  (r  $\in$  carrier P)  $\wedge$ 
f = (monom P 1R 1  $\ominus_P$  monom P a 0)  $\otimes_P$  q  $\oplus_P$  monom P (eval R R id a
f) 0"
  (is " $\exists$  q r. (q  $\in$  carrier P)  $\wedge$  (r  $\in$  carrier P)  $\wedge$  f = ?g  $\otimes_P$  q  $\oplus_P$  monom
P (eval R R id a f) 0")
  <proof>

end

```

### 14.13 Sample Application of Evaluation Homomorphism

```
lemma UP_pre_univ_propI:
  assumes "cring R"
    and "cring S"
    and "h ∈ ring_hom R S"
  shows "UP_pre_univ_prop R S h"
  ⟨proof⟩
```

**definition**

```
INTEG :: "int ring"
where "INTEG = (| carrier = UNIV, mult = op *, one = 1, zero = 0, add
= op + |)"
```

```
lemma INTEG_cring: "cring INTEG"
  ⟨proof⟩
```

```
lemma INTEG_id_eval:
  "UP_pre_univ_prop INTEG INTEG id"
  ⟨proof⟩
```

Interpretation now enables to import all theorems and lemmas valid in the context of homomorphisms between INTEG and UP INTEG globally.

```
interpretation INTEG: UP_pre_univ_prop INTEG INTEG id "UP INTEG"
  ⟨proof⟩
```

```
lemma INTEG_closed [intro, simp]:
  "z ∈ carrier INTEG"
  ⟨proof⟩
```

```
lemma INTEG_mult [simp]:
  "mult INTEG z w = z * w"
  ⟨proof⟩
```

```
lemma INTEG_pow [simp]:
  "pow INTEG z n = z ^ n"
  ⟨proof⟩
```

```
lemma "eval INTEG INTEG id 10 (monom (UP INTEG) 5 2) = 500"
  ⟨proof⟩
```

**end**

## References

- [1] C. Ballarin. *Computer Algebra and Theorem Proving*. PhD thesis, University of Cambridge, 1999. <http://www4.in.tum.de/~ballarin/>

[publications.html](#).

- [2] N. Jacobson. *Basic Algebra I*. Freeman, 1985.
- [3] F. Kammüller and L. C. Paulson. A formal proof of sylow's theorem: An experiment in abstract algebra with Isabelle HOL. *J. Automated Reasoning*, (23):235–264, 1999.